

## INTEL Registers: Names, sizes and bit positions

63	31	16	7	0
rax	eax	ax	al	
rbx	ebx	bx	bl	
rcx	ecx	cx	cl	
rdx	edx	dx	dl	
rsi	esi	si	sil	
rdi	edi	di	dil	
rbp	ebp	bp	bpl	
rsp	esp	sp	spl	
r8	r8d	r8w	r8b/r8l	
r9	r9d	r9w	r9b/r9l	
r10	r10d	r10w	r10b/r10l	
r11	r11d	r11w	r11b/r11l	
r12	r12d	r12w	r12b/r12l	
r13	r13d	r13w	r13b/r13l	
r14	r14d	r14w	r14b/r14l	
r15	r15d	r15w	r15b/r15l	

INTEL Address modes: The ways you can specify the source or destinations for an instruction : Allowed combinations are:

**Immediate:** the source value is stored in the instruction  
eg. ADD EAX, 14 # Add 14 into the 32 bit EAX register  
MOV RAX, 0xdeadbeef # set RAX

**Register to register:**  
eg. ADD R8B, AL # add 8 bit AL value to R8B register  
MOV RAX, R8 # copy the value from R8 into RAX

**Memory operands:**  
[BaseReg + scale \* IndexReg + Displacement]  
Where BaseReg and IndexReg can be any general purpose register  
scale is a numeric value of 1,2,4,8  
Displacement is 8, 16 or 32 bit value. Often this will be a symbolic label  
MOV RAX, QWORD PTR [RBX + 8\*RDI + XARRAY]  
Note: In general you can omit various terms to meet your needs

When doing moves only one operand can be a memory operand.

## INTEL EFLAGS Single bit flags that we are concerned with

ZF	Zero Flag set if result was zero
SF	Sign Flag set if result was negative (most significant bit set)
CF	Carry Flag set if operation generates a carry or borrow
OF	Overflow Flag set if overflow on signed operation

Typical Intel GNU Assembly Instruction formats:

mnemonic : mnemonic with no operands  
eg. int3  
mnemonic <dst> : mnemonic with one destination operand:  
eg. inc RAX. # RAX=RAX+1  
mnemonic <dst>, <src> : mnemonic with one destination and one src  
eg. add RAX, RBX # RAX = RAX + RBX  
mnemonic <dst>, <srcA>, <srcB> : mnemonic with one destination and two srcs  
eg. imul rax, rbx, 42 # rax = rbx \* 42

Common Intel mnemonics (instructions):

Data Transfer :

MOV : move to/from locations : mov dest, src : dst = src;

ALU:

ADD <dst>, <src> : add integers (signed or unsigned) : dst = dst + src;  
SUB <dst>, <src> : subtract integers (signed or unsigned) : dst = dst - src;  
IMUL <dst>, <srcA>, <srcB> : multiplies integers (signed): dst = srcA \* srcB;  
INC. : Increment : inc dst : dst = dst + 1  
DEC. : Decrement : dec dst : dst = dst - 1  
AND <dst>, <src> : Bitwise boolean and : dst = dst & src  
OR <dst>, <src> : Bitwise boolean or : dst = dst | src  
XOR <dst>, <src> : Bitwise boolean xor : dst = dst ^ src  
NOT <dst>. : Bitwise boolean not : dst = ~dst  
SHR <dst>, imm. : logical shift right : dst = dst >> imm (zero extends)  
SAR <dst>, imm : arithmetic shift right : dst = dst >> imm (sign extends)  
SHL <dst>, imm : logical shift left: dst = dst << imm (zero fill)  
SAL <dst>, imm : arithmetic shift left: dst = dst << imm (zero fill same as SHL)  
CMP <dst>, <src> : set eflags based on subtraction: dst - src  
TEST <dst>, <src> : set eflags based on bitwise and of dst & src

Control flow:

JMP <dst> : Unconditional jump : jmp <dst> : pc = dst : dst is usually a label  
eg. jmp loop\_begin  
but can also be indirect:  
1) a register (which contains the address to jump to)  
jmp rax  
2) ptr to a memory location which contains the address to jump to  
jmp qword ptr [myjumptable]  
JE <dst> : jump if equal : jmp if zero eflag is set (1) same as JZ  
JNE <dst> : jump if not equal: jmp if zero eflag is NOT set (0)  
JZ <dst> : jump if zero : same as JE  
JNZ <dst> : jump not zero: same as JNE  
JG <dst> : jump if greater (signed)  
JGE <dst> : jump if greater or equal (signed)  
JL <dst> : jump if less (signed)  
JLE <dst> : jump if less or equal (signed)  
JA <dst> : jump if above (unsigned)  
JAE <dst> : jump if above or equal (unsigned)  
JB <dst> : jump if below (unsigned)  
JBE <dst> : jump if below or equal (unsigned)

STACK:

PUSH : stack push : push src : rsp=rsp - len(src); M[rsp] = src;  
POP : stack pop : pop dst : dst = M[rsp]; rsp = rsp + len(src);  
CALL/RET : call and return from subroutine : call pushes address of the following instruction on the stack and then sets pc to the specified target address. ret pops the top value from the stack and sets the pc to this address

MISC:

NOP : no operation  
INT3 : hand control back to debugger  
SYSCALL : request operating system call routine

## Byte Vector Sizes and Names

1 Byte : INTEL BYTE : GAS directive .byte. : C unsigned char and char (signed)  
2 Bytes : INTEL WORD : GAS directive .short : C unsigned short and short (signed)  
4 Bytes : INTEL DWORD : GAS directive .long : C unsigned int and int (signed)  
8 Bytes : INTEL QWORD : GAS directive .quad : C unsigned long long and long long (signed)  
NOTE: On INTEL 64 bit machines all pointer types (char \*, short \*, int \*, long long \* and void \*) are 8 bytes in size

## GDB Commands:

file <binary> : opens a new binary replacing the current one eg. file empty  
run : creates a process from the current open binary and initiates the cpu's execution within it  
b <symbol> : sets a breakpoint to stop execution when the PC equals the address of symbol eg. b \_start  
c : continue execution from current PC address until execution terminates or a break point is hit  
si : single step a cpu instruction eg. unfreeze the cpu so it can do one execution loop  
p /x \$<REG> : print the current value of the specified register in hex  
x/<n>bx <address> : print/examine n memory bytes sized values start at the specified address in hex notation  
x/<n>hx <address> : same as above but n memory 2-byte sized values  
x/<n>wx <address> : same as above but n memory 4-byte sized values  
x/<n>gx <address> : same as above but n memory 8-byte sized values  
set \$<REG>=<value>: sets the value of the specified registers. Value can be specified in notations by using the right prefix eg. 0x for hex, 0b for binary. The default is signed two-complement integers.  
set {CType}(address)=<value>: set in memory at the address specified. CType is one of the C programming type names for bytes sized quantities. See notes below for a list

## NOTES:

1. When using x to display multi-bytes sized (eg. x/1hx <addr>) gdb will reorder to account for endianness of the computer. For example if the bytes at address \_start, on a little endian computer, are 0xFA 0x10 and we use the command x/1hx & \_start gdb will display something like  
(gdb) x/1hx &\_start  
0x401000 <\_start>: 0x10FA  
This is true for all the other multi-bytes sizes (h,w,g)
2. For the p and x command the following Format letters can be used  
o(octal), x(hex), d(decimal), u(unsigned decimal), t(binary), f(float), a(address), i(instruction), c(char), s(string) and z(hex, zero padded on the left).
3. CType names: "unsigned char" : 1 byte, "unsigned short" : 2 byte,  
"unsigned int" : 4 byte, "unsigned long long": 8 byte

## INTEL C Linux Calling Conventions :

Defines how registers should be used by caller and callee code. It also defines how arguments and the return value for a C function should be assigned to registers and the stack. The First 6 integer arguments are passed in registers as follows

Argument 0 : rdi

Argument 1 : rsi

Argument 2 : rdx

Argument 3 : rcx

Argument 4 : r8

Argument 5 : r9

Return value : rax

If more than 6 arguments are required the remainder are pushed on the stack in reverse order (last pushed first). A functions return value must be place in rax.

The function code (callee) is free to overwrite any of the 7 above registers along with r10 and r11.

Calling code (caller) needs to save and restore these registers if it wants to rely on their values. Thus they are called volatile and caller saved. The values of the remaining general purpose registers (rbx, rsp, rbp, r12-r15) must not be affected by a function as such they are called non-volatile and callee saved. Eg. if a function writes them it must restore their value before returning to the caller

## ASCII Hex Table (Hex Character)

00 nul	01 soh	02 stx	03 etx	04 eot	05 enq	06 ack	07 bel
08 bs	09 ht	0a nl	0b vt	0c np	0d cr	0e so	0f si
10 dle	11 dc1	12 dc2	13 dc3	14 dc4	15 nak	16 syn	17 etb
18 can	19 em	1a sub	1b esc	1c fs	1d gs	1e rs	1f us
20 sp	21 !	22 "	23 #	24 \$	25 %	26 &	27 '
28 (	29 )	2a *	2b +	2c ,	2d -	2e .	2f /
30 0	31 1	32 2	33 3	34 4	35 5	36 6	37 7
38 8	39 9	3a :	3b ;	3c <	3d =	3e >	3f ?
40 @	41 A	42 B	43 C	44 D	45 E	46 F	47 G
48 H	49 I	4a J	4b K	4c L	4d M	4e N	4f O
50 P	51 Q	52 R	53 S	54 T	55 U	56 V	57 W
58 X	59 Y	5a Z	5b [	5c \	5d ]	5e ^	5f _
60 `	61 a	62 b	63 c	64 d	65 e	66 f	67 g
68 h	69 i	6a j	6b k	6c l	6d m	6e n	6f o
70 p	71 q	72 r	73 s	74 t	75 u	76 v	77 w
78 x	79 y	7a z	7b {	7c	7d }	7e ~	7f del