

# Queue Implementation Using Array

Generate DSA Notes for implement QUEUE using Array

Here are the DSA Notes for implementing a Queue using an Array:

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## Queue Data Structure

A **Queue** is a linear data structure that follows the **FIFO (First In, First Out)** principle, where the element inserted first is the first to be removed.

### Operations in a Queue

1. **Enqueue (Insert):** Adds an element to the end of the queue.
  2. **Dequeue (Remove):** Removes an element from the front of the queue.
  3. **Peek/Front:** Retrieves the element at the front without removing it.
  4. **IsEmpty:** Checks if the queue is empty.
  5. **IsFull:** Checks if the queue is full (for array-based queues).
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## Array Implementation of Queue

An array-based queue implementation maintains the following:

- **Front:** Index of the first element in the queue.
- **Rear:** Index of the last element in the queue.
- **Size:** Maximum capacity of the queue.

### Challenges in Array Implementation

- After repeated enqueue and dequeue operations, elements may shift, and space may remain unused. To resolve this, we use **circular queues**.
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# Steps to Implement a Queue Using Array

1. **Define the Array:** Create an array to store elements of the queue.
2. **Initialize Front and Rear:** Initially set to `-1` to indicate an empty queue.
3. **Enqueue Operation:**
  - Check if the queue is full ( `rear == size - 1` for a simple queue).
  - Increment the `rear` and insert the element at `rear`.
  - If the queue was empty, update `front` to `0`.
4. **Dequeue Operation:**
  - Check if the queue is empty ( `front == -1` or `front > rear` ).
  - Retrieve the element at `front` and increment `front`.
  - If `front` surpasses `rear`, reset both to `-1` (queue is empty).
5. **Peek Operation:**
  - Return the element at the `front` if the queue is not empty.
6. **Handle Edge Cases:**
  - Queue overflow and underflow.
  - Reset indices when the queue becomes empty.

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## Pseudocode

### Initialization:

java

```
class Queue {  
    int[] arr; // Array to hold queue elements  
    int front, rear, size, capacity;  
  
    Queue(int capacity) {  
        this.capacity = capacity;  
        arr = new int[capacity];  
        front = rear = -1;  
    }  
}
```

```
        size = 0;
    }
}
```

### Enqueue Operation:

java

```
void enqueue(int value) {
    if (rear == capacity - 1) { // Queue is full
        System.out.println("Queue Overflow");
        return;
    }
    if (front == -1) front = 0; // First element
    arr[++rear] = value;
    size++;
}
```

### Dequeue Operation:

java

```
int dequeue() {
    if (front == -1 || front > rear) { // Queue is empty
        System.out.println("Queue Underflow");
        return -1;
    }
    int dequeuedValue = arr[front++];
    size--;
    if (front > rear) { // Reset queue after last element is dequeued
        front = rear = -1;
    }
    return dequeuedValue;
}
```

### Peek Operation:

java

```
int peek() {
    if (front == -1) { // Queue is empty
        System.out.println("Queue is Empty");
        return -1;
    }
}
```

```
    }  
    return arr[front];  
}
```

### IsEmpty:

java

```
boolean isEmpty() {  
    return front == -1;  
}
```

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## Circular Queue Optimization

In a circular queue:

- When `rear` reaches the last index, it wraps around to `0` (if there's space at the front).
- Formula for the next position: `(rear + 1) % capacity`.

### Key Changes for Circular Queue:

#### 1. Enqueue:

java

```
void enqueue(int value) {  
    if ((rear + 1) % capacity == front) { // Queue is full  
        System.out.println("Queue Overflow");  
        return;  
    }  
    if (front == -1) front = 0; // First element  
    rear = (rear + 1) % capacity;  
    arr[rear] = value;  
    size++;  
}
```

#### 2. Dequeue:

java

```
int dequeue() {
    if (front == -1) { // Queue is empty
        System.out.println("Queue Underflow");
        return -1;
    }
    int dequeuedValue = arr[front];
    if (front == rear) { // Last element
        front = rear = -1;
    } else {
        front = (front + 1) % capacity;
    }
    size--;
    return dequeuedValue;
}
```

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## Time Complexity

- Enqueue:  $O(1)$
- Dequeue:  $O(1)$
- Peek:  $O(1)$

## Space Complexity

- Array:  $O(n)$ , where  $n$  is the queue capacity.

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## Advantages of Array-based Queue

- Easy to implement.
- Predictable memory usage.

## Disadvantages

- Fixed size (requires resizing or circular implementation for efficiency).
- Shifting elements (in a simple queue, not circular).

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Let me know if you'd like to explore linked-list-based implementation for a queue or need deeper explanations for circular queues.