

0/1 knapsack problem using Dynamic Programming

$$M=8 \quad P=\{1,2,5,6\}$$

$$n=4 \quad W=\{2,3,4,5\}$$

This we can solve with table and set method

⇒ Here we can take solid object, we are not suppose to take fraction.

$$x_i = 0/1$$

$$\max \sum p_i x_i$$

- It demands maximum profit so it is optimization problem.

- It takes sequence of decision yes all items are require to check.

DP allow try for all possible decision or solution and pick up the best.

$$2^4 = 16 \Rightarrow \text{no. of solutions.}$$

$$\text{It is for } n \text{ is } 2^n \Rightarrow \cancel{O(2^n)} O(2^n)$$

Tabulation Method.

		0	1	2	3	4	5	6	7	8	← weights
P_i	W_i	0	0	0	0	0	0	0	0	0	
1	2	1	0	0	1	1	1	1	1	1	
2	3	2	0	0	1	2	3	3	3	3	
5	4	3	0	0	1	2	5	5	6	7	
6	5	4	0	0	1	2	5	6	6	7	

$$V[i, w] = \max \left\{ \overset{\substack{\uparrow \text{row} \\ \rightarrow \text{column}}}{V[i-1, w]}, \overset{\substack{\uparrow \text{column}}}{V[i-1, w-w[i]]} + \overset{\substack{\uparrow \text{column} \quad \uparrow \text{objective}}}{P[i]} \right\}$$

$$V[4, 1] = \max \left\{ V[3, 1], V[3, 1-5] + 6 \right\}$$

$$= \max \left\{ 0, V[3, -4] + 6 \right\}$$

→ value undefined ~~to~~ upto 5th weight-
 value is negative
 so fill weight

it is

	0	1	2	3	4	5	6	7	8
0	0	0	0	0	0	0	0	0	0
1	0	0	1	1	1	1	1	1	1
2	0	0	1	2	3	3	3	3	3
3	0	0	1	2	5	5	6	7	7
4	0	0	1	2	5	6			
5	0	0	1	2	5	6			

$$V[i, w] = \max \{ \underset{\substack{\downarrow \text{col} \\ \text{row}}}{V[i-1, w]}, V[i-1, \underset{\substack{\uparrow \text{row}}}{w - w[i]}] + P[i] \}$$

$$V[4, 5] = \max \{ V[3, 5], V[3, 5-5] + 6 \}$$

$$= \max \{ 5, V[3, 0] + 6 \}$$

$$= \max \{ 5, 0 + 6 \}$$

$$= \max \{ 5, 6 \}$$

$$\boxed{V[4, 5] = 6}$$

$$V[4, 6] = \max \{ V[3, 6], V[3, \overset{\substack{\uparrow \text{row}}}{6-5}}{5} + 6 \}$$

$$= \max \{ 5, 0 + 6 \} \quad V[3, \underset{6-5}{1}] + 6$$

$$= \max \{ 6 \}$$

$$\boxed{V[4, 6] = 6}$$

$$V[4, 7] = \max \{ V[3, 7], V[3, 7-5] + 6 \}$$

$$= \max \{ 7, \overset{\substack{\uparrow \\ 1}}{0} + 6 \}$$

$$= \max \{ 7 \}$$

$$\boxed{V[4, 7] = 7}$$

$$\begin{aligned}
 v[4, 8] &= \max \left\{ v[3, 8], v[3, 8-5] + 5 \right\} \\
 &= \max \left\{ 7, 2 + 5 \right\} \\
 &= \max \{ 8 \}
 \end{aligned}$$

	0	1	2	3	4	5	6	7	8
0	0	0	0	0	0	0	0	0	0
1	0	0	1	1	1	1	1	1	1
2	0	0	1	2	3	3	3	3	3
3	0	0	1	2	5	5	6	7	7
4	0	0	1	2	5	6	6	7	8

$x_1, x_2, x_3, x_4 \leftarrow$ sequence of decisions.

maximum profit = 8 including 4th object.

x_1, x_2, x_3, x_4

8

$8 - 6 = 2$
 \uparrow profit of 4th object.

x_1, x_2, x_3, x_4

$\{ 0, 1, 0, 1 \}$

$2 - 2 = 0$

solve the problem using set's method.

$$s = (P, w)$$

$$M = 8 \quad P = \{1, 2, 5, 6\}$$

$$n = 4 \quad w = \{2, 3, 4, 5\}$$

$$s^0 = \{(0, 0)\} \quad \text{No profit \& No weight.}$$

$$s_1^0 = \{(1, 2)\}$$

$$s^1 = \{(0, 0), (1, 2)\}$$

$$s_1^1 = \{(2, 3), (3, 5)\}$$

} Merge

$$s^2 = \{(0, 0), (1, 2), (2, 3), (3, 5)\}$$

$$s_1^2 = \{(5, 4), (6, 6), (7, 7), (8, 9)\}$$

} Merge

↳ don't include

discard one with lesser profit

$$s^3 = \{(0, 0), (1, 2), (2, 3), (3, 5), (5, 4), (6, 6), (7, 7)\}$$

↳ Dominance Rule

$$s_1^3 = \{(6, 5), (7, 7), (8, 8), (11, 9), (12, 10), (13, 12)\}$$

$$s^4 = \{(0, 0), (1, 2), (2, 3), (5, 4), (6, 6), (7, 7), (8, 8)\}$$

It is the easy method as compare to this tabular method.

$$\begin{aligned} (8, 8) &\in s^4 \\ \text{but } (8, 8) &\notin s^3 \quad \therefore x_4 = 1 \\ (8-6, 8-5) &= (2, 3) \\ (2, 3) &\notin s^3 \\ \& \quad (2, 3) &\notin s^2 \quad \therefore x_3 = 0 \end{aligned}$$

$$\begin{aligned} (2, 2) &\in s^2 \\ \text{but } (2, 3) &\notin s^1 \quad \therefore x_2 = 1 \\ (2-2, 3-3) &= (0, 0) \\ (0, 0) &\in s^1 \& (0, 0) \in s^0 \\ x_1 &= 0 \end{aligned}$$

(3)