

# **15-213 Recitation 5: Attack Lab**

**26 Sept 2016**

# Agenda

- **Reminders**
- **Stacks**
- **Attack Lab Activities**

# Reminders

- **Bomb lab is due tomorrow!**

- “But if you wait until the last minute, it only takes a minute!” - ***NOT!***
- Don't waste your grace days on this assignment

- **Attack lab will be released tomorrow!**

# Stacks

## ■ Last-In, First-Out

- just like a stack of plates
- pushes and pops to prese

## ■ x86 stack grows down

- lowest address is “top”

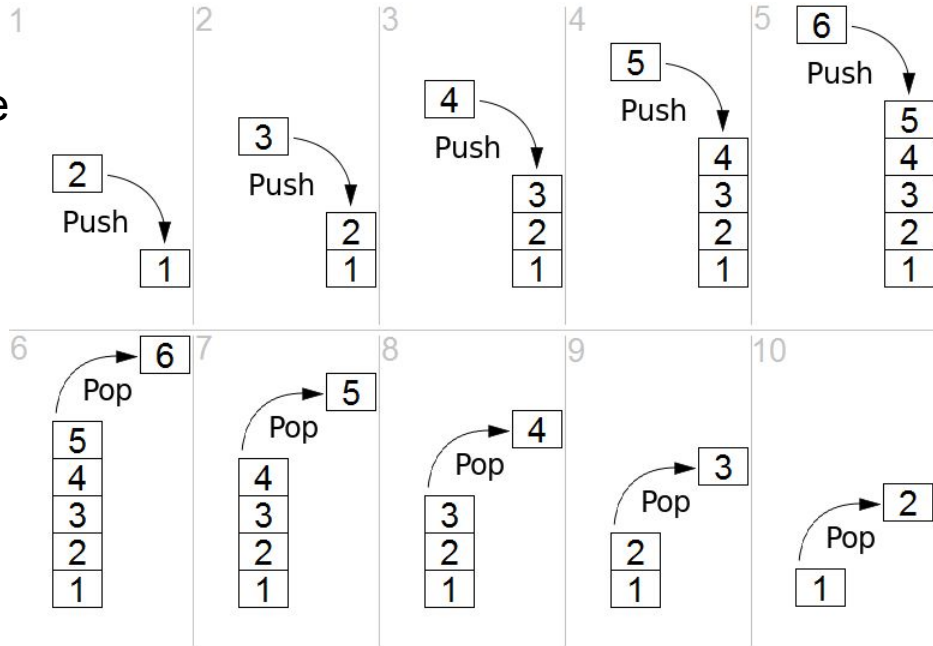


Image credit: Wikimedia Commons

# Stack

- **Stack space is allocated in “frames”**
  - Represents the state of a single function invocation
- **Used primarily for two things:**
  - Storing callee save registers
  - Storing the return address of a function
- **Can also store:**
  - Local variables that don't fit in registers
  - Function arguments 7+

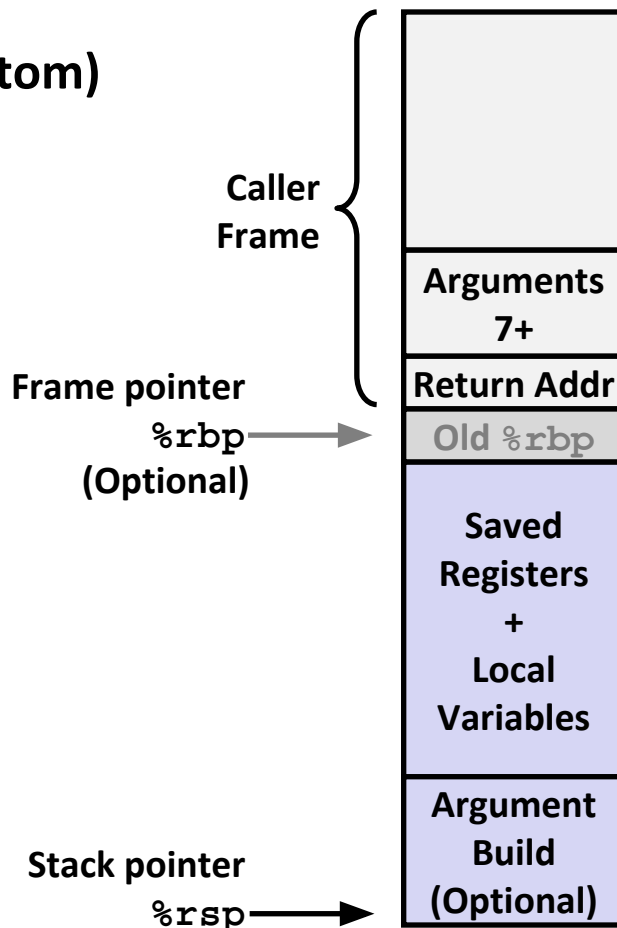
# x86-64/Linux Stack Frame

## ■ Current Stack Frame (“Top” to Bottom)

- “Argument build:”  
Parameters for function about to call
- Local variables  
If can’t keep in registers
- Saved register context
- Old frame pointer (optional)

## ■ Caller Stack Frame

- Return address
  - Pushed by `call` instruction
- Arguments for this call



# Stack Maintenance

- **Functions free their frame before returning**
- **Return instruction looks for the return address at the top of the stack**
  - What if the return address has been changed?

# Attack Lab Activities

- **Three activities**

- Each relies on a specially crafted assembly sequence to purposefully overwrite the stack

- **Activity 1 – Overwrites the return addresses**

- **Activity 2 – Writes an assembly sequence onto the stack**

- **Activity 3 – Uses byte sequences in libc as the instructions**



# Form pairs

- One student needs a laptop
- Login to a shark machine

```
$ wget http://www.cs.cmu.edu/~213/activities/rec5.tar
```

```
$ tar xf rec5.tar
```

```
$ cd rec5
```

```
$ make
```

```
$ gdb act1
```

# Activity 1

**(gdb) break clobber**

**(gdb) run**

**(gdb) x \$rsp**

**(gdb) backtrace**

**Q. Does the value at the top of the stack match any frame?**

**(gdb) x /2gx \$rdi                // Here are the two key values**

**(gdb) stepi                    // Keep doing this until**

```
(gdb)
clobber () at support.s:16
16                               ret
```

**(gdb) x \$rsp**

**Q. Has the return address changed?**

**(gdb) fin**

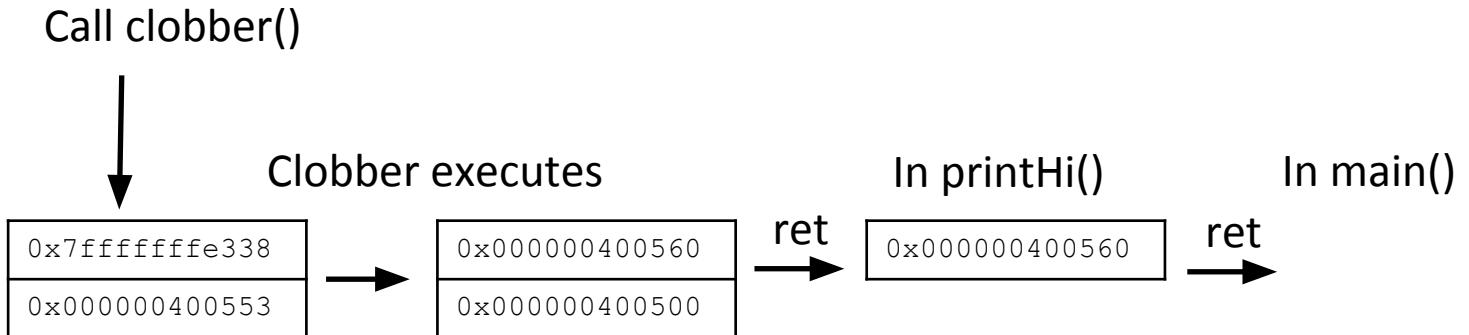


Should exit normally,  
May segfault

# Activity 1 Post

## ■ In this activity, we overwrote part of the stack

- Placing two return addresses onto the stack
- Return to printHi()
- Return to main



# Activity 2

**\$gdb act2**

**(gdb) break clobber**

**(gdb) run**

**(gdb) x \$rsp**

**Q. What is the address of the stack and the return address?**

**(gdb) x /4gx \$rdi**

**Q. What will the new return address be?  
(i.e., what is the first value?)**

**(gdb) x/5i \$rdi + 8    // Display as instructions**

**Q. Why rdi + 8?**

**Q. What are the three addresses?**

**(gdb) break puts**

**(gdb) break exit**

**Q. Do these addresses look familiar?**

# Activity 2 Post

- **Normally programs cannot execute instructions on the stack**
  - Main used `mprotect` to change the memory protection for this activity
- **Clobber wrote a return address of the stack to the stack**
  - And a sequence of instructions
  - Three addresses: `"Hi\n"`, `puts()`, `exit()`
- **Why callq `*%rsi`?**
  - As the attacklab writeup notes, calling functions is hard.
  - Return oriented programming is much easier.

# Activity 3

**\$gdb act3**

**(gdb) break clobber**

**(gdb) run**

**(gdb) x /5gx \$rdi**

**Q. Which value will be first on the stack?**

**Q. At the end of clobber, where will it return?**

**(gdb) x /2i <return address>**

**Q. What does this sequence do?**

**Q. Do the same for the other addresses. Note that some are return addresses and some are for data. When you continue, what will the code now do?**

# How was it constructed?

- **Think of possible executions**
- **What are the bytes of the instructions?**
  - Write short assembly into `foo.s`
  - `gcc -c foo.s`
  - `objdump -d foo.o`
  - OR: Convert them to byte sequences (Attacklab write-up has a table)
    - Also important so you can switch between register names
- **After determining the desired instruction(s)**
  - Use the Linux tool `xxd` to dump the raw bytes to a file
  - Or: `Objdump -d rtarget` (or `act3` or ...)
  - Search the file

# If You Get Stuck

- Please read the writeup. *Please read the writeup. Please read the writeup. **Please read the writeup!***
- CS:APP Chapter 3
- View lecture notes and course FAQ at <http://www.cs.cmu.edu/~213>
- Office hours Sunday through Thursday 5:00-9:00pm in WeH 5207
- Post a private question on Piazza
- `man gdb`, `gdb's help command`



# Remember...



# Appendix

# Attack Lab Tools

💡 **gcc -c file.s**

💡 convert the assembly code in file.s to object code in file.o

💡 **objdump -d file.o**

💡 disassemble the code in file.o; shows the actual bytes for the instructions

💡 **./hex2raw**

💡 convert hex codes into raw ASCII strings to pass to targets

💡 **gdb**

💡 determine stack addresses

💡 **paper and pencil**

💡 for drawing stack diagrams

## More Useful GDB Commands

<b><u>x</u>/[<i>n</i>]i &lt;address&gt;</b>	<b>disassemble <i>n</i> instructions at &lt;address&gt;</b>
<b>b &lt;loc&gt; if &lt;cond&gt;</b>	<b>conditional breakpoint, stop only if &lt;cond&gt;</b>
<b>cond &lt;bp&gt; &lt;cond&gt;</b>	<b>true</b>
<b><u>commands</u> &lt;bp&gt;</b>	<b>add condition to existing breakpoint &lt;bp&gt;</b>
<b><u>tbreak</u> &lt;loc&gt;</b>	<b>execute commands when breakpoint &lt;bp&gt;</b>
<b><u>finish</u></b>	<b>hit</b>
	<b>set temporary breakpoint – auto-deletes</b>
<b><u>layout</u> <u>asm</u></b>	<b>when hit!</b>
	<b>run until current frame (function) returns,</b>
<b><u>layout</u> <u>reg</u></b>	<b>and print return value</b>
	<b>split the screen into separate disassembly</b>
	<b>and command windows</b>
	<b>show register window as well (after layout</b>