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8 5	4 19	118	16	62	13		for (int i=0; i <n; i++)<="" td=""></n;>
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						0	min = 0;
		10	7	02	13		
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5 52	9	18	38	62	13		E dearches for amaller elemon
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5 9	13	13	38	52	62		
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Insertion Sort

It works the way we don't playing cords in our hand we choose one card and indert st in its parties (autending or descending)

1 2 3 4 5 3 3

Working Of Insertion Sort

1. We start by making the second element of the given array, i.e. element at index 1, the key. The key element here is the new card that we need to add to our existing sorted set of cards (as explained cards example before).

2. We compare the key element with the element(s) before it, in this case, element at index 0.

2. If the start element is less than the first element, we insert the key element before the 2.2 If the key element is greater than the first element, then we insert it after the first element.

3. There element.

4. There element.

5. There element.

6. The element of the array as key and will compare it with elements to it's left and insert if at the right position.

6. And we go on repeating this, until the array is sorted.

into temps;;

for (into l=1; i< n; i++)

temp= arr [i];

defects the temp variable

temp= arr [i];

catile (j > 0 & 8 arr [j-1] > temp)

arr [j] = arr [j-1];

j-j-1; // j--;

arr [j] = temp;

arr [j] = temp;