

ToC Question Bank – 3

1. Construct DPDA for
 - a. $L = \{a^n b^n \mid n > 0\}$
 - b. $L = \{wcw^r \mid w \in \{a, b\}^*\}$
 - c. $L = \{n_a(w) = n_b(w) \mid w \in \{a, b\}^*\}$
 - d. $L = \{a^n b^n c^m \mid n > 0, m > 0\}$
 - e. $L = \{a^{n+m} b^n c^m \mid n > 0, m > 0\}$
 - f. $L = \{a^n b m c^{m+n} \mid n > 0, m > 0\}$
 - g. $L = \{a^n b^{2n} \mid n > 0\}$
 - h. $L = \{a^{2n} b^{3n} \mid n > 0\}$
2. Construct PDA for the following languages. Check whether each one is a DPDA.
 - a. $L = \{a^n b^n \mid n \geq 1\} \cup \{a^n b^{2n} \mid n \geq 1\}$
 - b. $\{ww^R \mid w = \{a, b\}^*\}$
 - c. $L = \{n_a(w) > n_b(w) \mid w \in \{a, b\}^*\}$
 - d. $L = \{w \in \{0, 1\}^* \mid w = w^R \text{ and the length of } w \text{ is odd}\}$
 - e. $L = \{w \in \{0, 1\}^* \mid w = w^R \text{ and the length of } w \text{ is any}\}$
3. Consider the following CFG $G = (V, \Sigma, R, S)$, where $V = \{S, T, X\}$, $\Sigma = \{a, b\}$, the start variable is S, and the rules R are

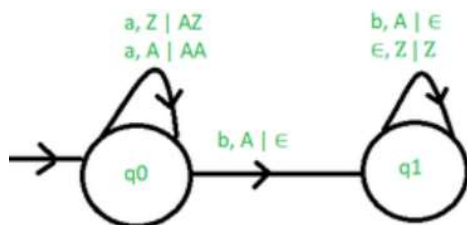
$S \rightarrow aT Xb$

$T \rightarrow XT S \mid \epsilon$

$X \rightarrow a \mid b$

Convert G to an equivalent PDA

4. Convert PDA to CFG



5. Design Turing Machine to recognize
 - a. $L = \{0^{2n} \mid n \geq 0\}$.
 - b. $L = \{0^n 1^n 2^n\}$ where $n \geq 1$
 - c. All strings over $\{0, 1\}$ in which the number "01" pairs is odd.
 - d. $q_0 w \vdash q_f w w^R$
 - e. To check the palindrome of the string of any length.
 - f. To add 2 unary numbers separated by a 0.
 - g. To get the function $c = f(a-b)$ where a and b are two unary numbers separated by a 0 and a is always greater than b.
 - h. $L = \{0^i 1^j 2^k\}$ where $i \times j = k$, $i, j, k \geq 1$
 - i. Given a list of strings over $\{0, 1\}$ separated by #, determine if all strings are different.
 - j. $ww \vdash w\#w$, where $w \in \{1\}^*$
 - k. $w \vdash w^R$ where $w \in \{1\}^*$

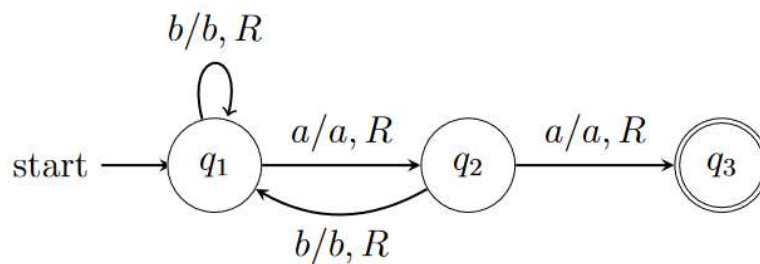
- l. $w \vdash w \bmod 2$ where $w \in \{1\}^*$
 - m. $L = \{(a + b)^*b(a + b)^*\}$.
 - n. $L = \{w\#w' : w \in \{a, b\}^* \text{ and } w \text{ is a substring of } w'\}$.
 - o. To accept a string with 'aba' as its substring
 - p. $L = \{b^n a^n b^n \mid n \geq 0\}$.
 - q. $w \in \{a, b, c\}^* \mid n_a(w) = n_b(w) \text{ and } n_a(w) > n_c(w)\}$
 - r. that changes all the a's on its tape to b's and vice versa
 - s. that doubles each character in its input string. For example, if the input is 0100, then the machine should change its tape so it contains 00110000
6. Consider the TM $M = (Q, \Sigma, \Gamma, \delta, q_0, q_f)$ where $Q = \{q_0, q_1, q_f\}$, $\Sigma = \{0, 1\}$, $\Gamma = \{0, 1, B\}$, and δ is given by:

$\delta(q_0, 0) = (q_0, 0, R)$, $\delta(q_0, 1) = (q_0, 1, R)$, $\delta(q_0, B) = (q_1, B, L)$, $\delta(q_1, 0) = (q_f, 1, R)$,

$\delta(q_1, 1) = (q_1, 0, L)$, $\delta(q_1, B) = (q_f, B, L)$

What is the function computed by M ?

7. What is the regular expression accepted by the below TM?



8. $L = \{wcw \text{ where } w \in \{a, b\}^*\}$ using 2 tape Turing Machine.