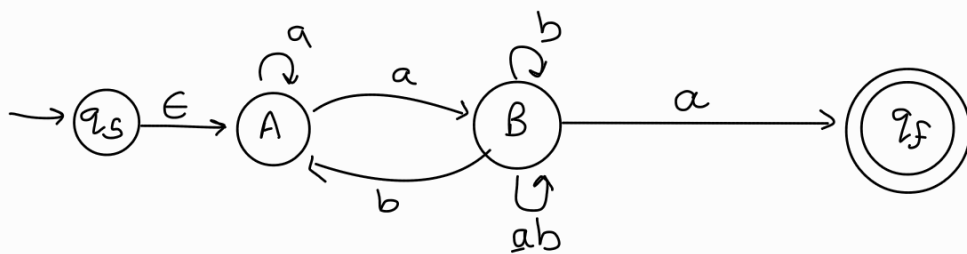
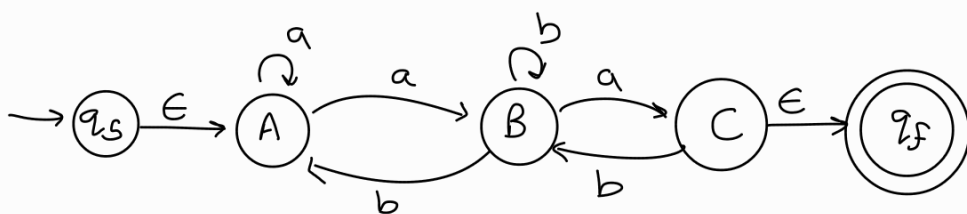
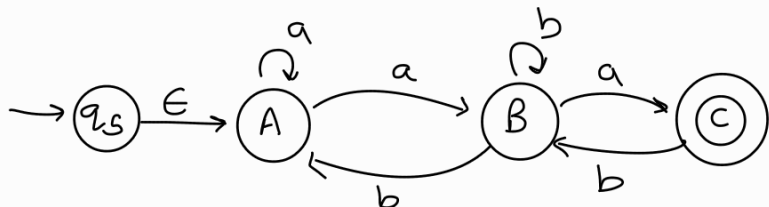
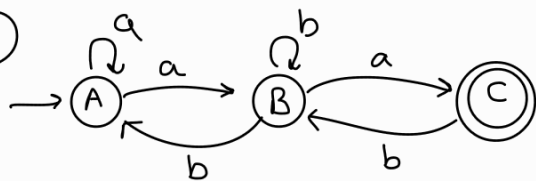


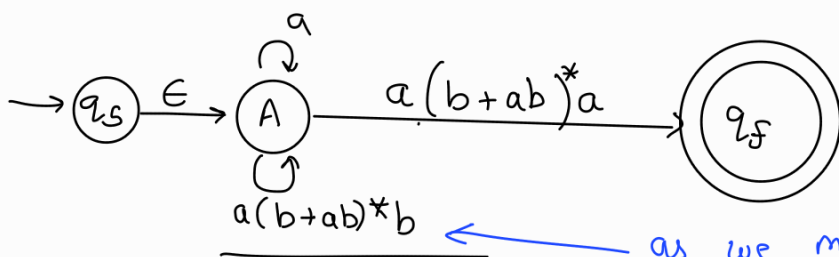
8/05/23

Regular expression using state elimination.

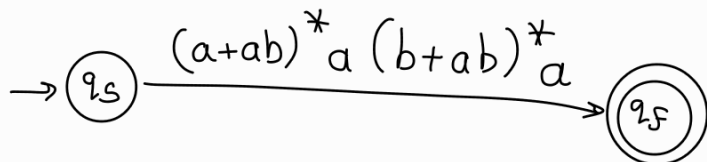
(1)



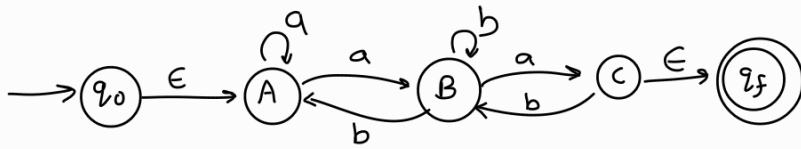
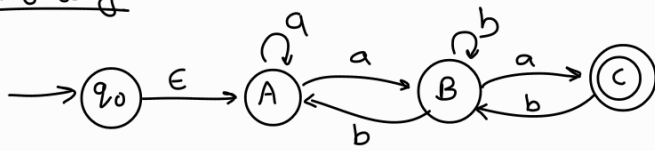
// elim state C.



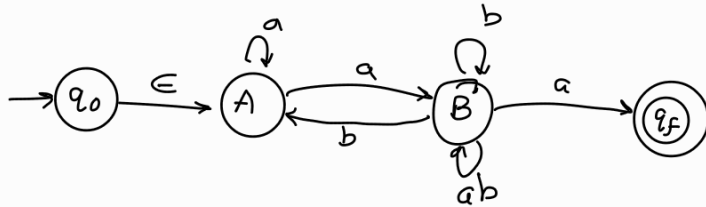
as we might go to (B) have the loops * times and go back to A.



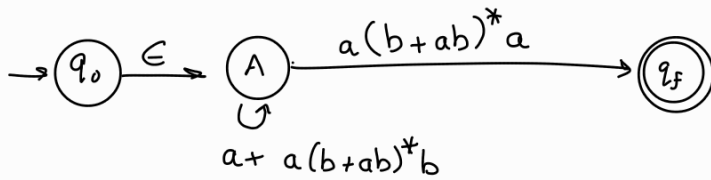
Ma'am's way



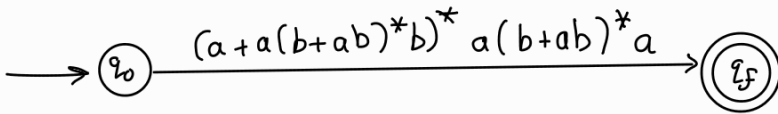
Eliminate C:

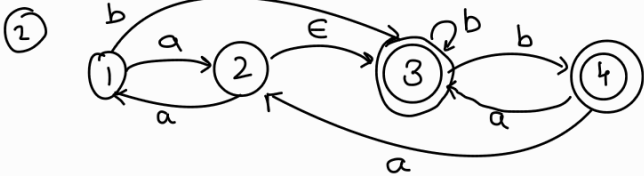


Elim B:

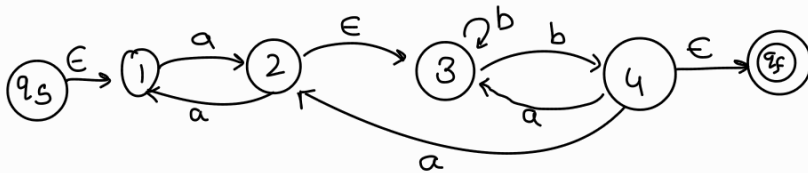
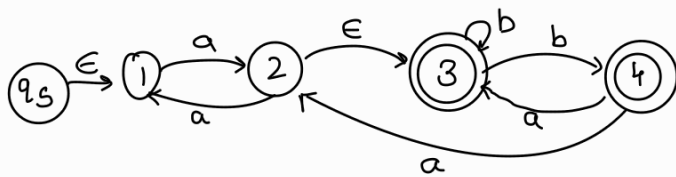


Elim A:

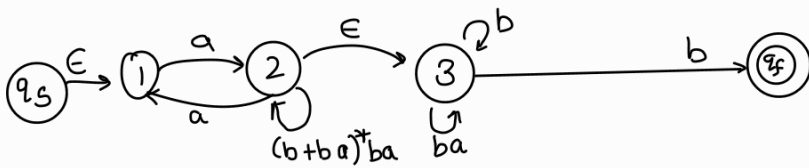




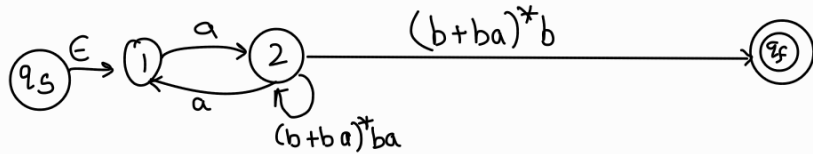
Solved at the end



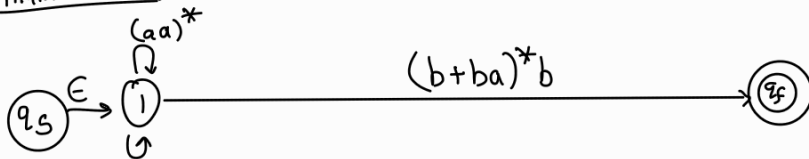
Eliminate 4 :



Eliminate 3

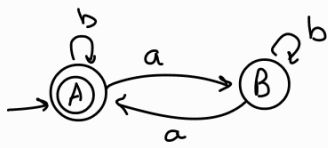


Eliminate 2 :



$$b + a(aa)^*b^* + (b + (b + a(aa)^*b^*b(a + a(aa)^*b^*b))^*(\epsilon + (a + a(aa)^*b^*)))$$

① Find RE using Arden's theorem:



$$A = \epsilon + Ab + Ba$$

$$B = Aa + Bb$$

$$B = Aab^* \dots \textcircled{1}$$

$$A = \epsilon + Ab + Ba$$

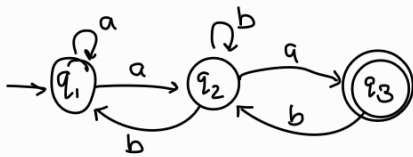
$$= \epsilon + Ab + Aab^*a$$

$$A = \epsilon + A(b + ab^*a)$$

$$A = \epsilon (b + ab^*a)^*$$

$$\therefore \text{Regex} = (b + ab^*a)^*$$

②



$$q_1 = \epsilon + q_1a + q_2b$$

$$q_2 = q_1a + q_2b + q_3b$$

$$q_3 = q_2a$$

$$q_2 = q_1a + q_2b + q_2ab$$

$$q_2 = q_1a + q_2(b + ab)$$

$$q_2 = q_1a(b + ab)^*$$

$$q_1 = \epsilon + q_1a + q_1a(b + ab)^*b$$

$$q_1 = \epsilon + q_1(a + a(b + ab)^*b)$$

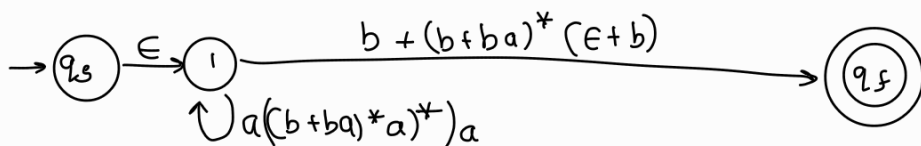
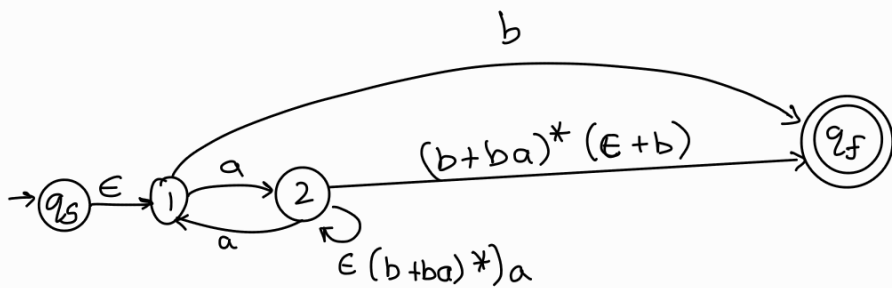
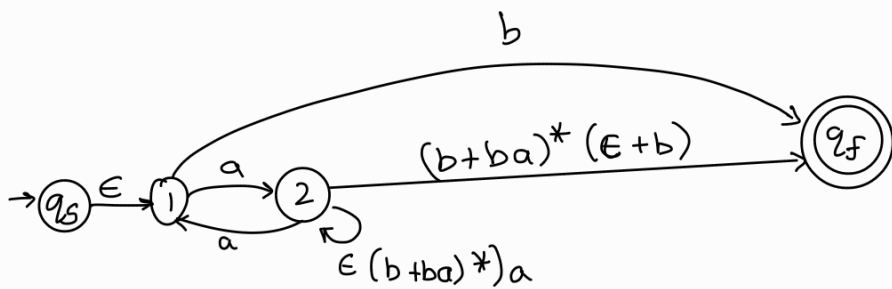
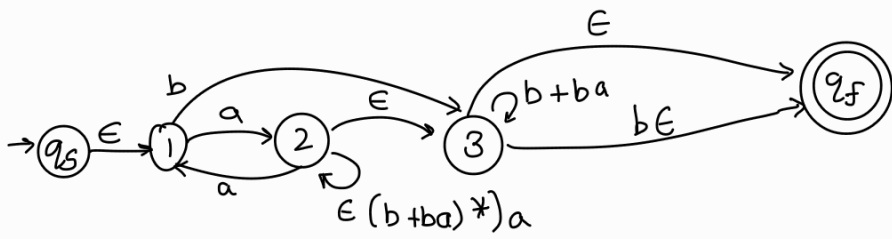
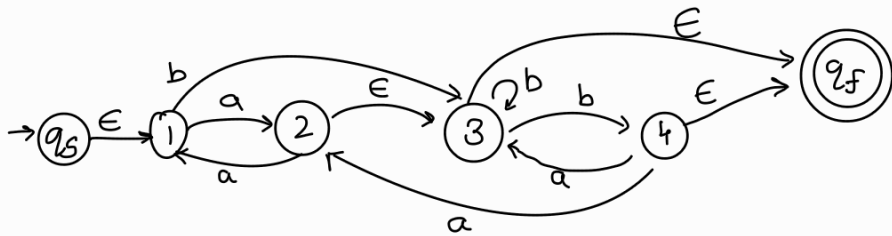
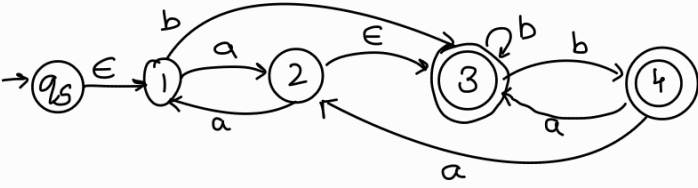
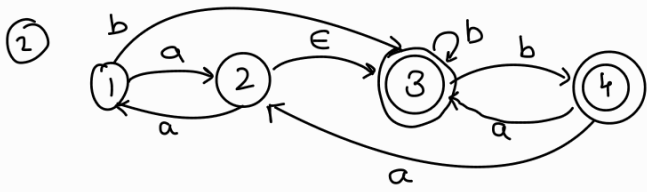
$$q_1 = \epsilon(a + a(b + ab)^*b)^*$$

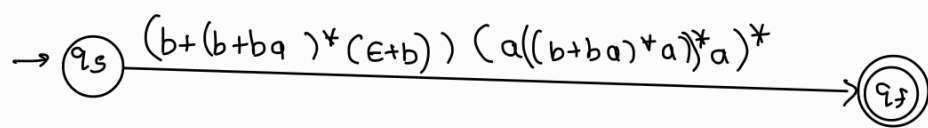
$$q_1 = (a + a(b + ab)^*b)^*$$

$$q_2 = q_1a(b + ab)^*$$

$$q_2 = (a + a(b + ab)^*b)^*a(b + ab)^*$$

$$q_3 = (a + a(b + ab)^*b)^*a(b + ab)^*a$$





$$b + a(aa)^*b^* + (b + (b + a(aa)^*b^*b)(a + a(aa)^*b^*b)^*(\epsilon + (a + a(aa)^*b^*b)))$$

Design CFG for w s.t. $w \in \{0,1\}^*$ | w contains at least three 1's;

$$S \rightarrow x1x1x1x$$

$$x \rightarrow 0x | 1x | \epsilon$$

Design CFG for $a^n b^n$ s.t. $n > 1$

$$S \rightarrow aSb | ab$$

$$G = (\{S\}, \{a, b\}, P, S)$$

$$P \rightarrow \{S \rightarrow aSb | ab\}$$

③ Design CFG for $a^n b^{2^n} | n \geq 0$.

$$S \rightarrow aSbb | \epsilon$$

$$G = (\{S\}, \{a, b\}, P, S)$$

$$P = \{S \rightarrow aSbb | \epsilon\}$$

$$(4) a^m b^n c^{m+n} \quad | m \geq 0, n \geq 0$$

$$a^m b^n c^n c^m$$

$$S \rightarrow aSc \mid B$$

$$B \rightarrow bBc \mid \epsilon$$

$$a^3 b^5 c^4$$

$$a a a \underline{b b b} c c c c$$

$$\underline{b b} c c c$$

Q PDA by empty stack.

$$S \rightarrow xS \mid \epsilon$$

$$X \rightarrow aXb \mid ab$$

$$\delta(q, \epsilon, s) \rightarrow \{ (q, xs), (q, \epsilon) \}$$

$$\delta(q, \epsilon, x) \rightarrow \{ (q, axb), (q, ab) \}$$

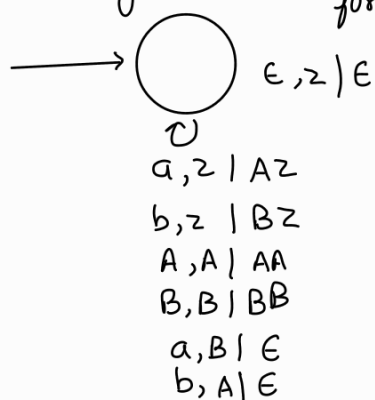
We Pop Here

$$\delta(q, a, a) \rightarrow \{ (q, \epsilon) \}$$

$$\delta(q, b, b) \rightarrow \{ (q, \epsilon) \}$$

$$\delta(q, aabb, s)$$

Design a PDA with single state for $n_a(w) = n_b(w)$



(A) $S \rightarrow qzq$

she will give only
EMPTY stack mostly.

(B) $[qzq] \rightarrow a [qA^{\textcircled{B}}z] [qAq]$

$[qzq] \rightarrow b [qB^{\textcircled{C}}z] [qBq]$

$[qAq] \rightarrow a [qAq] [qAq]$

$[qBq] \rightarrow b [qBq] [qBq]$

$[qBq] \rightarrow a$

$[qAq] \rightarrow b$

$[qzq] \rightarrow \epsilon$

$S \rightarrow abab$

$(q, abab, S) \Rightarrow (q, bab, BS)$

$S \rightarrow A$

$A \rightarrow aBA \mid b \in A \mid b$

$B \rightarrow aBB \mid b$

$C \rightarrow bcc \mid a$