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**Lecture – 19**  
**Dual Simplex Method**

Good morning students, today we will learn a method called the dual simplex method. This method is a modification of the simplex method for a particular type of problem, that is, which has a relationship with the dual. So as in how we proceed, we will see how the dual simplex method can be applied to obtain the solution of a linear programming problem.

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The outline of today's lecture is as follows: background, then the dual simplex method with the help of an exercise and finally a question for you to solve. Now let us look at the definition of the primal and the dual once again.

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## Background

### Primal:

Minimize  $f(X) = \sum_{j=1}^n c_j x_j$   
subject to  $\sum_{j=1}^n a_{ij} x_j \geq b_j, i = 1, 2, \dots, m,$   
 $x_j \geq 0, j = 1, 2, \dots, n.$

### Dual:

Maximize  $w(Y) = \sum_{i=1}^m b_i y_i$   
subject to  $\sum_{i=1}^m a_{ij} y_i \leq c_j, j = 1, 2, \dots, n$   
 $y_i \geq 0, i = 1, 2, \dots, m.$

As you know, that if the primal is in the form of minimization of the objective function subject to the inequality constraints of the type  $\geq$  and the dual is of the maximization type and the constraints are of the  $\leq$  type. Then the standard form of the primal can be written as minimization of  $f(X) = \sum_{j=1}^n c_j x_j$  subject to  $\sum_{j=1}^n a_{ij} x_j \geq b_j$ , where  $i$  goes from 1, 2 up to  $m$  and  $x_j$  are the decision variables of the primal, they are  $\geq 0$  for  $j = 1, 2$  to up to  $n$ . The corresponding dual is maximization of a function  $w(Y) = \sum_{i=1}^m b_i y_i$  subject to  $\sum_{i=1}^m a_{ij} y_i \leq c_j$  where  $j$  goes to 1 to  $n$  and the decision variables of the dual are  $y_i$  are  $\geq 0$  for  $i = 1$  to up to  $m$ .

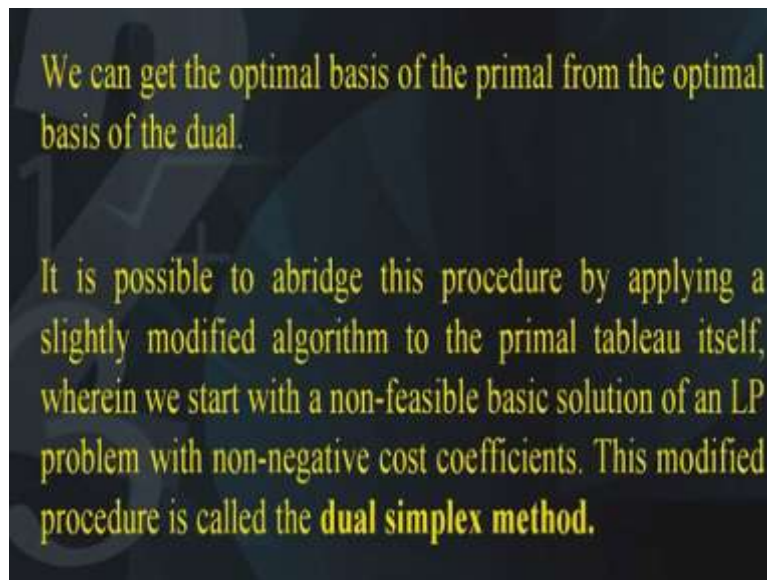
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If, however, some or all  $b_i > 0$  and all  $c_j \geq 0$ , then the basis  $x_{n+1}, x_{n+2}, \dots, x_{n+m}$  is not feasible for the primal, but the basis  $y_{m+1}, y_{m+2}, \dots, y_{m+n}$  is feasible for the dual. This situation is called: **primal infeasible and dual feasible.**

Now suppose all the  $c$  that is the  $c_j \geq 0$  and the  $b_i \leq 0$ . This is a special case, we are assuming. A special case then the basis consisting of the basic variables that is  $x_{n+1}, x_{n+2}, \dots, x_{n+m}$  which are nothing but the slack variables is feasible and also optimum. Similarly the corresponding basis of the dual is feasible and also optimum. Now suppose there is a situation that some or all of the  $b_i$ 's are strictly  $> 0$  not  $\geq$ , they are strictly  $> 0$  and all the  $c_j$ 's are  $\geq 0$ . Then the basis  $x_{n+1}, x_{n+2}, \dots, x_{n+m}$  is not feasible for the primal. But the basis  $y_{m+1}, y_{m+2}, \dots, y_{m+n}$  is feasible for the dual. Now this will hold only if some or all the  $b_i$ 's are strictly  $> 0$ . This kind of a special situation is called primal infeasible and dual feasible. This is called primal infeasible and dual feasible.

Now suppose we start our calculations of the simplex algorithm on the dual then we shall be moving through a successive iterations, where the basic feasible solution of the dual which occurs at the basic variables which means that the  $c_j$ 's are  $\geq 0$ , till the final relative cost coefficient or the deviation coefficients that is the  $b_i$ 's of the dual are all non-positive, then we would have then arrived in the optimal solution of the dual. So this is a special situation and under this condition we would have arrived at the optimal solution of the dual.

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We can get the optimum basis of the primal from the optimal basis of the dual. Therefore it is possible to abridge this procedure, by applying a slightly modified version of the simplex algorithm to the primal table itself, wherein we start with a non-feasible basic solution of an LP

with non-negative cost coefficients and this kind of a procedure is termed as the dual simplex method.

So basically in the dual simplex method, we are starting from infeasible solution which is optimum and at the successive iterations we are going towards optimality and feasibility both.

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Write the simplex tableau for the primal problem with a basis consisting of slack variables.

Basis	Values	$x_1$	$x_2$	..	$x_p$	..	$x_n$	$x_{n+1}$	..	$x_{n+m}$
$x_{n+1}$	$-b_1$	$-a_{11}$	$-a_{12}$	..	$-a_{1p}$	..	$-a_{1n}$	1	..	0
$x_{n+2}$	$-b_2$	$-a_{21}$	$-a_{22}$	..	$-a_{2p}$	..	$-a_{2n}$	0	..	0
..	..	..	..	..	..	..	..	..	..	..
$x_{n+i}$	$-b_i$	$-a_{i1}$	$-a_{i2}$	..	$-a_{ip}$	..	$-a_{in}$	0	..	0
..	..	..	..	..	..	..	..	..	..	..
$x_{n+m}$	$-b_m$	$-a_{m1}$	$-a_{m2}$	..	$-a_{mp}$	..	$-a_{mn}$	0	..	1
$-f$	0	$c_1$	$c_2$	..	$c_p$	..	$c_n$	0	..	0

So let us write the simplex table for the primal with the basis consisting of the slack variables now this can be written in this tabular form on the first column we have the basis then in the second column we have the right hand side and under every column  $x_1, x_2, x_3$  etc we have their corresponding entries and in the end we have the basis that is  $x_{n+1}, x_{n+2}, \dots, x_{n+m}$ . So these entries are the unit vectors 1, 0, 0 etc is 0, 1, 0, 0, etc and finally 0, 0, 0, 1.

So the basis has been excluded towards the right.

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The dual simplex method consists in changing a negative basic variable in such a way that the value of the new basic variable in its place would be positive, and the relative cost coefficients for the changed basis still remain non-negative.

The dual simplex method consists of changing a negative basic variable in such a way that the value of the new basis variable in its place would be positive and the relative cost coefficients or the deviations for the changed basis still remains non-negative. Because, that is the condition for optimality if you remember, we must have all the deviation entries non-negative.

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Assume that some  $b_j > 0$  (that is, values of some basic variables are negative) and all  $c_j \geq 0$ .

For example, let  $b_r > 0$  so that corresponding basic variable  $x_{n+r}$  is negative.

Also let some coefficients  $-a_{rj}$  be negative. Let, in particular,  $-a_{rp} < 0$ .

We may replace  $x_{n+r}$  by  $x_p$  in basis by dividing the  $r$ th equation by  $-a_{rp}$  and eliminating  $x_p$  from all other equations and also from the last row giving the expression for  $f$  in terms of nonbasic variables and relative cost coefficients.

Now assume, that some  $b_j > 0$  that is, it is not  $\geq$  it is strictly  $> 0$  that is value of some basic variable are negative and all  $c_j$  are  $\geq 0$ . For example let  $b_r$  be strictly  $> 0$  so that the corresponding basic variable  $x_{n+r}$  is negative. Also let some coefficient of the coefficient matrix let us say,  $-a_{rj}$  be negative. So  $a_{rj}$  is negative. Let, in particular it be called as  $-a_{rp}$ . This is strictly  $< 0$ .

Now we may replace  $x_{n+r}$  by let us say  $x_p$  in the basis by dividing the  $r$ th equation by  $-a_{rp}$  and then eliminating  $x_p$  from all other equations and also from the last row giving the expression for the objective function  $f$  in terms of the non-basic variables and the deviation entries that is the relative cost coefficients.

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This change should be such that no relative cost coefficient becomes negative. This will be so when:

$$c_j - \frac{a_{ij}}{a_{rp}} c_p \geq 0, \quad j = 1, 2, \dots, n+m,$$

Or  $\frac{c_j}{a_{rj}} \geq \frac{c_p}{a_{rp}}$   
over all those  $j$  for which  $-a_{rj} < 0$   
Or  $\min_j \frac{c_j}{a_{rj}} = \frac{c_p}{a_{rp}}, \quad -a_{rj} < 0$

Now this change should be made very carefully such that no relative cost coefficient becomes negative, this will be so only if the following conditions hold that is  $c_j - (a_{ij} / a_{rp}) c_p \geq 0$  for all  $j = 1, 2$  up to  $n+m$ . In other words, this inequality should be satisfied over all those  $j$  for which  $-a_{rj}$  is strictly  $< 0$  that is the minimum over  $j$  for this condition holds that is  $-a_{rj}$  is  $< 0$ .

Now this leads to the determination of  $p$ . So, we know what is  $p$  the value of the new basic variable  $x_p$  would be  $(-b_r) / (-a_{rp})$  and as you know negative multiplied by negative is positive. So the whole expression will turn out to be positive if for  $-b_r < 0$  then there is no  $-a_{rp} < 0$  and the problem is infeasible. Now, we may change the basis in this way step by step iteration after iteration such that one basic variable in each iteration till all the basic variables comes to have non-negative values. Thus we shall arrive at a basis which is a feasible solution and which is also an optimal solution as all the relative cost coefficients that is the  $c_j \geq 0$ .

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It may be noted that in this method we move through a set of points which are not primal feasible taking care all the time that the relative cost coefficients remain non-negative so that the moment we arrive at a feasible basis, we find ourselves at the feasible optimal basis.

Now it may be noted that in this method we move through a set of points which are not primal feasible taking care all the time that the relative cost coefficients or the deviations remain non-negative so that the moment we arrive at a feasible basis we find ourselves at the feasible optimum basis.

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**Example:**

Minimize  $Z = x_1 + 4x_2 + 3x_4$   
 Subject to  $x_1 + 2x_2 - x_3 + x_4 \geq 3$   
 $-2x_1 - x_2 + 4x_3 + x_4 \geq 2$   
 $x_1, x_2, x_3, x_4 \geq 0$   
 LP in standard form is:  
 Minimize  $Z = x_1 + 4x_2 + 3x_4$   
 Subject to  $x_1 + 2x_2 - x_3 + x_4 - x_5 = 3$   
 $-2x_1 - x_2 + 4x_3 + x_4 - x_6 = 2$   
 $x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$

Now in order to understand this procedure let us look at an example, it says we have the primal as minimization of  $Z = x_1 + 4x_2 + 3x_4$  there is no  $x_3$  term subject to  $x_1 + 2x_2 - x_3 + x_4 \geq 3$  and the second constraint is  $-2x_1 - x_2 + 4x_3 + x_4 \geq 2$  and all the four decision variables  $x_1, x_2, x_3, x_4 \geq 0$ . Now the LP in the standard form looks like this the minimization is the same of the

objective function. But we have to subtract surplus variables in both the equations and the surplus variable in the first equation is called as  $x_5$  and the surplus variable in the second equation is called as  $x_6$ . So we have the two conditions  $x_1 + 2x_2 - x_3 + x_4 - x_5 = 3$  and the second condition is  $-2x_1 - x_2 + 4x_3 + x_4 - x_6 = 2$ . Of course all the decision variables  $x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$ .

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Multiplying both equations by  $-1$ , we get:

$$-x_1 - 2x_2 + x_3 - x_4 + x_5 = -3$$

$$2x_1 + x_2 - 4x_3 - x_4 + x_6 = -2$$

The sol. Is  $x_5 = -3, x_6 = -2$ , all other zero.

This is optimal but primal infeasible, and dual feasible.

Now let us multiply both the equations with the negative sign and what do we find, we find that we get a basic variable in both the equations. So in the first equation we have  $x_5$  as the basic variable and in the second equation we have  $x_6$  as the basic variables because they were negative and since we have multiplied both the equations with the negative sign therefore we get  $x_5$  and  $x_6$  as positive and this canonical form gives us a solution  $x_5 = -3, x_6 = -2$ . Now you will wonder that how can the values of the right-hand side be negative. Yes that is true, this is what is the beauty about this method, that is the basis that is the basic variables, they are negative, both of them are negative and of course all others are 0. So what is this BFS corresponding to, this BFS is optimum but it is primal infeasible why is it infeasible because we want that all the  $x_i$ 's should be  $\geq 0$ . But unfortunately both these variables are negative. So they are not feasible, i.e., they are not primal feasible that is why they are written to be primal infeasible both these variables are primal infeasible and of course they are dual feasible. As we have seen in the explanation that they will be dual feasible because this satisfy the dual constraints so they are dual feasible but they are primal infeasible and they are optimum.



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Table 1								
		1	4	0	3	0	0	
$C_B$	Basis	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	RHS
0	$x_5$	-1	-2	1	-1	1	0	-3
0	$x_6$	2	1	-4	-1	0	1	-2
dev row		1	4	0	3	0	0	

So now let us record all this information in this table. Now as I said the basis is nothing but  $x_5$  and  $x_6$  and the coefficients of the objective function are 0 and 0 corresponding to  $x_5$  and  $x_6$ . In the  $x_1$  column we will write the entries -1 and 2; under  $x_2$ , -2 and 1; under  $x_3$ , 1 and -4. Similarly -1 and 1 under  $x_4$ ; 1 0 under  $x_5$  and under  $x_6$  0, 1 and of course the right-hand side is -3 and -2.

On the top row over here we have to write the coefficients of the objective function corresponding to each of the variables. So what do we find? Corresponding to  $x_1$  we have 1,  $x_2$  4, as you remember there was no  $x_3$  term in the objective function so it is 0;  $x_4$  is 3 and of course for  $x_5$  and  $x_6$  the entries are 0. Next we will calculate the deviation row that is the cost coefficients and they are nothing but as before  $1 - (0, 0) \cdot (-1, 2)^t$  that comes out to be 1. Similarly for the others  $4 - (0, 0) \cdot (-2, 1)^t$  that comes out to be 4. Similarly 0, 3, 0 and 0. So, I hope you have understood how the initial table has been prepared the peculiarity in this table is that the right hand side entries are negative which is indicating that although this BFS is optimum but it is infeasible as far as the primal is concerned.

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Criteria for leaving variable is:

most negative RHS

So,  $x_5$  leaves the basis

Criteria for entering non basic variable is: maximum ratio  
between dev ratio and only negative co-efficients in pivot  
row:

$$\max \left\{ \frac{1}{-1}, \frac{4}{-2}, \frac{3}{-1} \right\} = -1$$

So  $x_1$  enters the basis

Now we have to change the basis so we need to decide the pivot and we need to decide which variable should enter and which variable should leave. Now the difference between the simplex method and the dual simplex method is that in the simplex method we first find the entering variable and then find the leaving variable, however in the case of the dual simplex method it is the other way round. That is, first of all we find the leaving variable that is that variable which has to leave and then we find the entering variable. So what is the criteria for the leaving variable and the criteria is most negative right-hand side. So if you look at the right hand side what were the two entries let us go back yeah -3 and -2, you find that -3 and -2 are the right-hand side entries and the most negative is -3. Therefore the variable corresponding to -3 should be leaving.

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Criteria for leaving variable is:

most negative RHS

So,  $x_5$  leaves the basis

Criteria for entering non basic variable is: maximum ratio between dev ratio and only negative co-efficients in pivot row:

$$\max\left\{\frac{1}{-1}, \frac{4}{-2}, \frac{3}{-1}\right\} = -1$$

So  $x_1$  enters the basis

So,  $x_5$  is the variable corresponding to -3, that is,  $x_5$  that has to leave. So  $x_5$  leaves the basis that is the first step that has to be taken. Next, we have to look at the criteria for the entering non-basic variable into the basis. So, we have to look at the criteria is the maximum ratio test has to be performed. Again, if you remember in the simplex method, we perform the minimum ratio test. However, in the dual simplex method, we perform the maximum ratio test and the maximum ratio test has to be performed between the deviation row and only the negative coefficients in the pivot row. So for the non-basic variables, what are those entries? Let us look at it 1/-1 let me go back; 1/-1 that is corresponding to  $x_1$ , similarly corresponding to  $x_2$ , 4/-2, and the 3rd one is 3/-1. Now, out of these 3 we have to select the one that is maximum and you can see that the maximum is = -1. So what does this indicate? This indicates that  $x_1$  should enter the basis and that is what is to be done.

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From Table 1 we get Table 2								
		1	4	0	3	0	0	
$C_B$	Basis	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	RHS
0	$x_5$	-1	-2	1	-1	1	0	-3
0	$x_6$	2	1	-4	-1	0	1	-2
	dev row	1	4	0	3	0	0	
$R_1$ is replaced by $-R_1$ and $R_2$ is replaced by $R_2-2R_1$								
		1	4	0	3	0	0	
$C_B$	Basis	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	RHS
1	$x_1$	1	2	-1	1	-1	0	3
0	$x_6$	0	-3	-2	-3	2	1	-8
	dev row	0	2	1	2	1	0	

So from the initial table that is the table number 1, we obtained the table number 2 by applying the elementary row operations in such a way that  $x_1$  variable becomes the basis. Here you can see in the second table under the  $x_1$  column we have the entries 1 and 0, as you know that the elementary row operations have to be performed like this,  $R_1$  is replaced by  $-R_1$  because this entry we have to make as 1, it is -1 we have to make it as 1. So in order to do that we need to multiply  $R_1$  with the negative sign. Secondly we have to replace the  $R_2$  by  $R_2-2R_1$ . If we do that, then the entry corresponding to the second row under  $x_1$  column becomes 0. Therefore, applying these two elementary row operations we get the table number 2 which is shown below and after that we need to calculate the deviation entries. So the deviation entries are to be calculated by making sure that in this first column. We change the coefficient of the basis so the coefficient of basis for  $x_1$  variable is 1 and for the  $x_6$  is 0, that will remain as before and just as before. We calculate the deviation entries like this  $1 - (1, 0) (1, 0)^t$  that is 0. Anyway this is a basic variable so automatically this entry will become 0 similarly  $4 - (1, 0) (2, -3)^t$  which comes out to be 2 and similarly  $0 - (1, 0) (-1, -2)^t$  which comes out to be 1. Similarly  $3 - (1, 0) (1, -3)^t$  which comes out to be 2 and then  $0 - (1, 0) (-1, 2)^t$  which is 1 and finally  $x_6$  variable is the basic variable, So the entry is 0 that is how we have obtained table number 2 from table number 1.

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Most negative RHS is -8

So,  $x_6$  leaves the basis

Maximum ratio between dev ratio and only negative coefficients of non-basic variables in pivot row is:

$$\max\left\{\frac{2}{-3}, \frac{1}{-2}, \frac{2}{-3}\right\} = -1/2$$

So  $x_3$  enters the basis

Again, the next iteration has to be performed by finding out the most negative right hand side and we find that the most negative right hand side is -8 and this corresponds to the basic variable  $x_6$  and this indicates that the basic variable  $x_6$  should leave the basis and we have to now perform the maximum ratio test to determine which variable should enter the basis. So the maximum ratio test between the deviation row. Only the negative coefficients of the non-basic variable is in the pivot row is maximization of  $2/-3$ ,  $1/-2$ ,  $2/-3$  which comes out to be  $-1/2$  this has to be the maximum and this indicates that this is corresponding to the entering variable.

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**From Table 2 we get Table 3**

		1	4	0	3	0	0	
$C_B$	Basis	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	RHS
1	$x_1$	1	2	-1	1	-1	0	3
0	$x_6$	0	-3	-2	-3	2	1	-8
	dev row	0	2	1	2	1	0	

$R_2$  is replaced by  $(-R_2)/2$  and  $R_1$  is replaced by  $R_1 + R_2$

		1	4	0	3	0	0	
$C_B$	Basis	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	RHS
1	$x_1$	1	7/2	0	5/2	-2	-1/2	7
0	$x_3$	0	3/2	1	3/2	-1	-1/2	4
	dev row	0	1/2	0	1/2	2	1/2	

Therefore from table number 2 we get table number 3 as follows, under the  $x_1$  column we have 1,0,0. So now this we have to apply the elementary row operations in such a way that  $R_2$  is

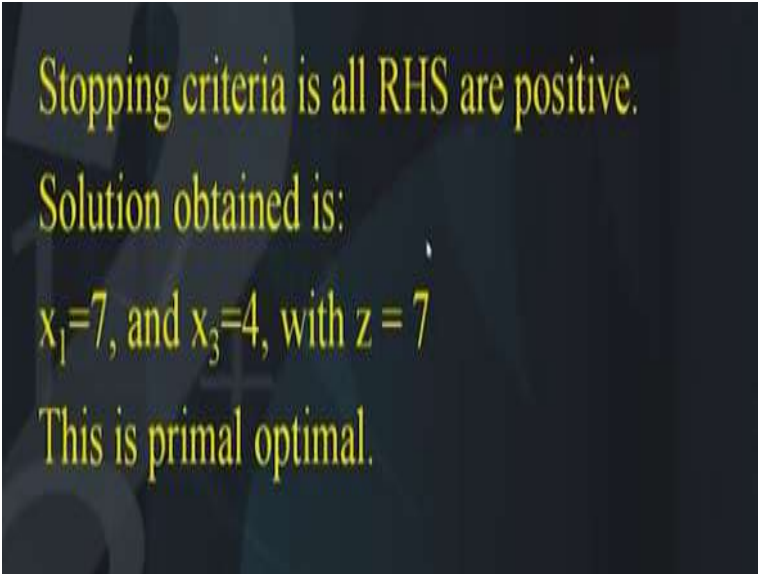


replaced by  $-R_2/2$  and  $R_1$  is replaced by  $R_1+R_2$ . So the elementary row operations are applied like this,  $R_2$  is replaced by  $-R_2/2$  and  $R_1$  is replaced by  $R_1+R_2$ .

So, we get table number 3 from table number 2 like this under the  $x_1$  heading we have 1,0 under  $x_2$  we have  $7/2$ ,  $3/2$  under  $x_3$  we have 0 1, under  $x_4$  we have  $5/2$  and  $3/2$   $x_5$  is -2,-1  $x_3$  is -1/2 and -1/2 and the right hand side is 7 and 4. Now you will observe that we will calculate the last row that is the deviation rows it is calculated as usual that is  $1 - (1,0) (1,0)^t$  that is 0, then the next one is  $1/2$ , 0,  $1/2$ , 0 and finally  $1/2$ .

Now you will find that this is optimum as well as feasible because the right hand side has become  $>0$  and as you know for feasibility it is necessary that the right hand side should be  $>0$  and that is what is happening here that right hand side has become  $>0$ .

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Stopping criteria is all RHS are positive.  
Solution obtained is:  
 $x_1=7$ , and  $x_3=4$ , with  $z = 7$   
This is primal optimal.

Now the stopping criteria has to be defined for the dual simplex method and the stopping criteria is that all the right hand side should be positive as you know that is the requirement for feasibility and since the stopping criteria has been satisfied in the second iteration itself, therefore the solution obtained is  $x_1=7$ , and  $x_3=4$ , and of course the objective function value that is  $z$  is 7. This is the primal optimum it is feasible as well as optimum.

So what do we find in the simplex method we move from non-optimum to the optimum. Whereas, in the dual simplex method we move from the primal infeasible to the primal feasible.

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Primal is Minimize  $Z = x_1 + 4x_2 + 3x_4$

Subject to  $x_1 + 2x_2 - x_3 + x_4 \geq 3$

$-2x_1 - x_2 + 4x_3 + x_4 \geq 2$

$x_1, x_2, x_3, x_4 \geq 0$

Dual is: Maximize  $W = 3y_1 + 2y_2$

Subject to  $y_1 - 2y_2 \leq 1$

$2y_1 - y_2 \leq 4$

$-y_1 + 4y_2 \leq 0$

$y_1 + y_2 \leq 3$

$y_1 \geq 0, y_2 \geq 0$

So the primal is in the problem that we had was given by the equation  $x_1 + 4x_2 + 3x_4$  subject to  $x_1 + 2x_2 - x_3 + x_4 \geq 3$ ,  $-2x_1 - x_2 + 4x_3 + x_4 \geq 2$ , all  $x_i$ 's  $\geq 0$  and the dual was maximization of  $3y_1 + 2y_2$  subject to  $y_1 - 2y_2 \leq 1$ ,  $2y_1 - y_2 \leq 4$ ,  $-y_1 + 4y_2 \leq 0$ ,  $y_1 + y_2 \leq 3$  and all the  $y_1$  and the  $y_2 \geq 0$ . Now we have seen that we have solved this primal okay and we started with a infeasible primal and we reached at a feasible primal.

Now I have written the dual of this problem and now let us see what happens if you look at this matrix B.

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$$[B]_{2 \times 2} = (P_5 \quad P_6)$$

$$\text{Hence } B^{-1} = \begin{pmatrix} -2 & -1/2 \\ -1 & -1/2 \end{pmatrix}$$

$$\bar{b} = B^{-1}b = \begin{pmatrix} -2 & -1/2 \\ -1 & -1/2 \end{pmatrix} \begin{pmatrix} -3 \\ -2 \end{pmatrix} = \begin{pmatrix} 7 \\ 4 \end{pmatrix}$$

Simplex multipliers are  $(\pi_1 \quad \pi_2)$

$$= c_B B^{-1}$$

$$= (1 \quad 0) \begin{pmatrix} -2 & -1/2 \\ -1 & -1/2 \end{pmatrix}$$

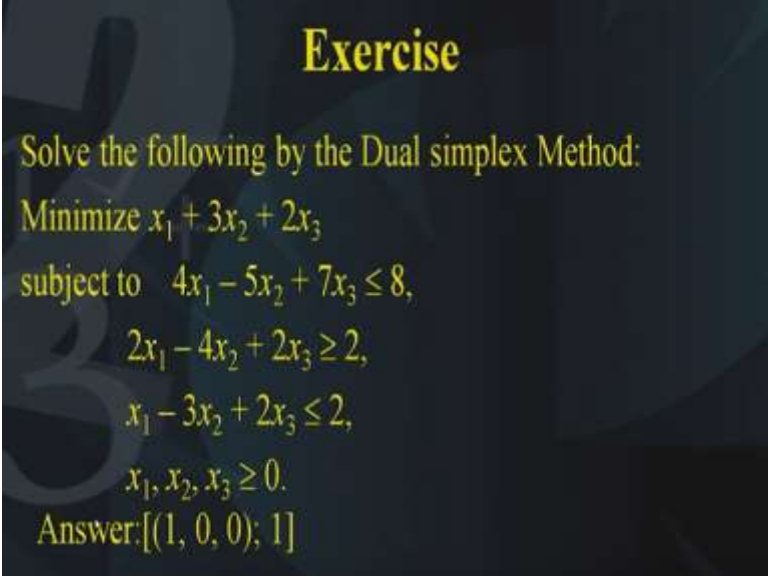
$$= \begin{pmatrix} -2 \\ -1/2 \end{pmatrix}$$

This is the solution of the dual

That is the basis the column basis that is the  $P_5$  and the  $P_6$ . So what is  $B^{-1}$ ?  $B^{-1}$  is nothing but -2, -1; -1/2, -1/2 and the  $b$  bar is nothing but  $B^{-1}b$  which is nothing but this matrix (-2, -1; -1/2, -1/2) (-3, -2) which comes out to be (7, 4). You can just check these calculations, then the simplex multipliers  $\pi$  are nothing but  $\pi_1$  and  $\pi_2$  which is  $= c_B B^{-1}$  and this is  $=(1, 0) (-2, -1; -1/2, -1/2)$ , this nothing but -2 and -1/2 and this is the solution of the dual. This is the solution of the dual, because as you know that when we find the simplex multipliers from the primal final table we can read the value of the solution of the dual. So, the actually the simplex multipliers are the solution of the dual, so this solution -2 and -1/2 is actually the solution of that dual. So let me go back.

Yeah this is the solution to the dual of this problem you can just check that it should satisfy all the constraints and it should be optimum to the dual. So -2, -1/2 is the simplex multipliers and therefore they are the solution of the dual.

**(Refer Slide Time: 35:34)**



**Exercise**

Solve the following by the Dual simplex Method:

Minimize  $x_1 + 3x_2 + 2x_3$

subject to  $4x_1 - 5x_2 + 7x_3 \leq 8,$

$2x_1 - 4x_2 + 2x_3 \geq 2,$

$x_1 - 3x_2 + 2x_3 \leq 2,$

$x_1, x_2, x_3 \geq 0.$

Answer:  $[(1, 0, 0), 1]$

So in the end let us conclude this lecture with a question which I want you to do at your place. Solve the following LPP by the dual simplex method. So you have LPP as follows, minimization of  $x_1 + 3x_2 + 2x_3$  subject to  $4x_1 - 5x_2 + 7x_3 \leq 8$ , second constraint is  $2x_1 - 4x_2 + 2x_3 \geq 2$ ,  $x_1 - 3x_2 + 2x_3 \leq 2$  and of course  $x_1, x_2, x_3 \geq 0$ .

So you have to solve this problem with the dual simplex method and I have already told you how to do that, you have to make the initial table with the help of multiplying those equations

where you have added surplus variable or rather subtracted surplus variables, so that the right hand side becomes negative the answer to this problem is also given. It is  $(1, 0, 0)$  with an objective function value 1. So I hope you will be able to complete this exercise. Thank you.