

Paths and distances

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December 14, 2022

Outline

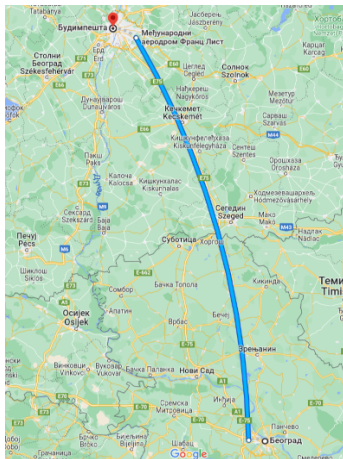
- 1 Introduction
- 2 Definition
- 3 Centrality measures

Importance of distance

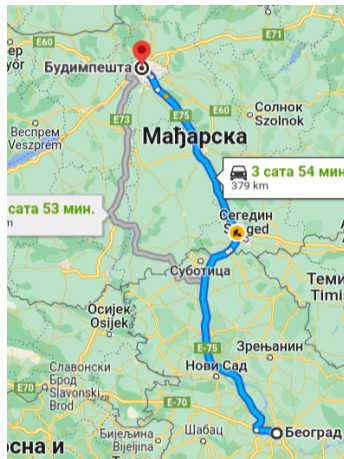
- **Physics:** strength of the force decreases with distance
- **Geography:** everything is related to everything else, but near things are more related than distant things

Distance in physical world

Euclidian

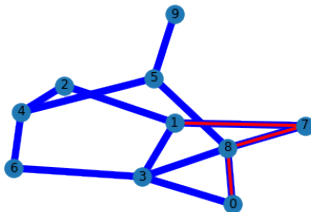
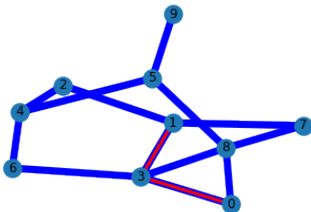


Non-Euclidian



Distance in networks

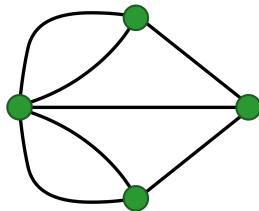
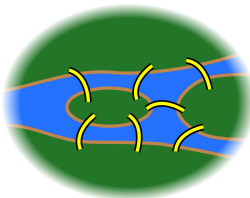
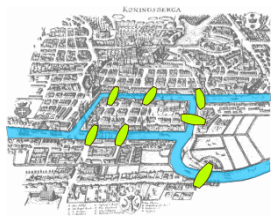
What is the distance between two webpages or two people?



In complex networks theory distance==length of path

Everything started with a path

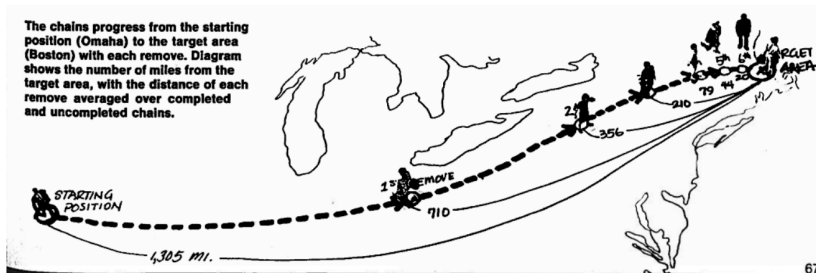
Seven Bridges of Königsberg



Euler set the foundations of network theory

It continued with paths

Small-world experiment in 1967



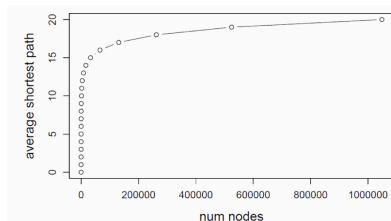
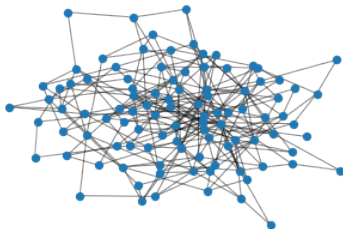
67

Six degrees of separation

Source: Migram's article in Psychology Today

It continued with paths

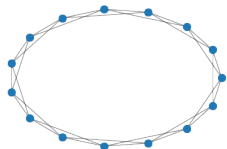
Erdos-Renyi networks - N nodes; p - probability to have a link between any pair of nodes;



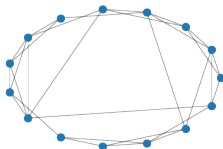
$$\langle l \rangle = \frac{\log(n)}{\log(p(n-1))}$$

It continued with paths

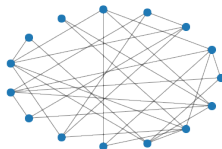
N nodes; k -neighbours; p - rewiring probability;



$$T = 15, \langle l \rangle = 2.29$$



$$T = 11, \langle l \rangle = 2.03$$

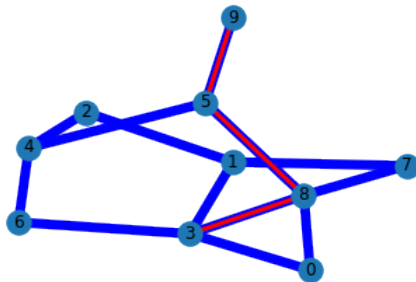


$$T = 9, \langle l \rangle = 1.9$$

Path

- **Path** is a route that runs along the links of the network
- **Path's length** l is the number of links the path contains
- Path P between nodes i_0 and i_n : $P = \{(i_0, i_1), (i_1, i_2), \dots, (i_{n-1}, i_n)\}$
- Length of path P is $l = n$

Undirected networkss: path

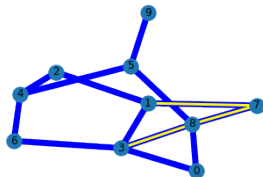
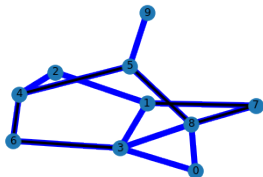
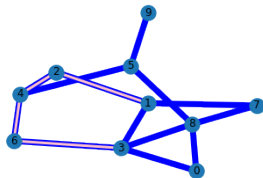
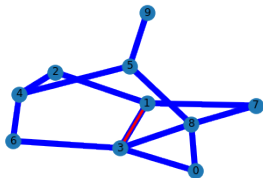


path from node 3 to node 9: $P_{3,9} = \{(3, 8), (8, 5), (5, 9)\}$

path's length $l = 3$

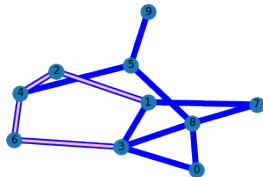
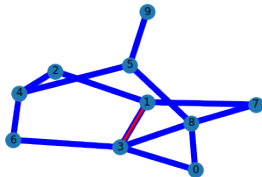
Undirected networks: paths

There is more than one path between a pair of nodes



Undirected networks: paths

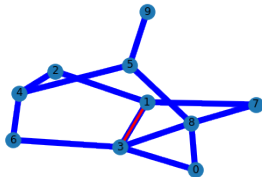
There is more than one path between a pair of nodes



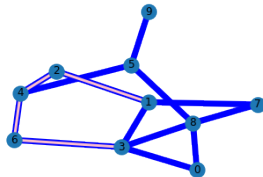
$$P_{3,1} = \{(3,1)\}$$

Undirected networks: paths

There is more than one path between a pair of nodes



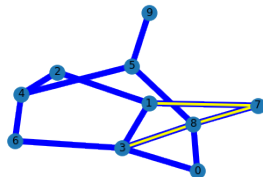
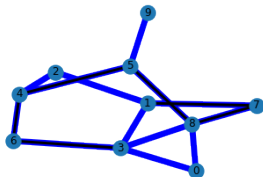
$$P_{3,1} = \{(3, 1)\}; l = 1$$



$$P_{3,1} = \{(3, 6), (6, 4), (4, 2), (2, 1)\}; l = 4$$

Undirected networks: paths

There is more than one path between a pair of nodes

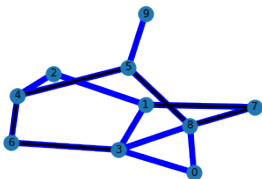


$$P_{3,1} = \{(3,6), (6,4), (4,5), (5,8), (8,7), (7,1)\}$$

$$l = 6$$

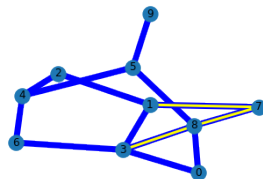
Undirected networks: paths

There is more than one path between a pair of nodes



$$P_{3,1} = \{(3, 6), (6, 4), (4, 5), (5, 8), (8, 7), (7, 1)\}$$

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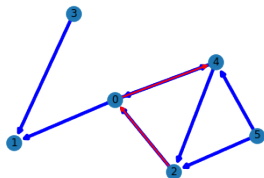
$$P_{3,1} = \{(3, 8), (8, 7), (7, 1)\}$$

$$l = 3$$

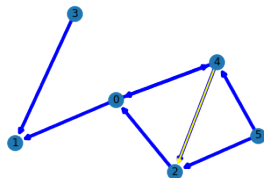
In undirected networks $P_{1,3} \equiv P_{3,1}$

Directed networks: paths

In directed networks the existence of a path from node i to node j does not guarantee the existence of a path from j to i



$$P_{2,4} = \{(2, 0), (0, 4)\}; l = 2$$



$$P_{4,2} = \{(4, 2)\}; l = 1$$

Weighted networks: paths

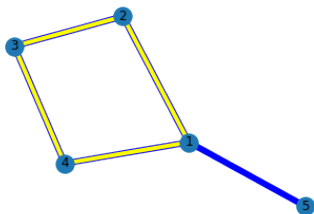
- We can have 2 paths - weighted and topological (unweighted) paths
- These paths are different
- If link weights are distances or strengths - weighted path

$$P_{i_0 i_n}^w \{(i_0, i_1), (i_1, i_2), \dots, (i_{n-1}, i_n)\}, l_{i_0 i_n}^w = \sum_{(i,j) \in P_{i_0 i_n}^w} W_{ij}$$

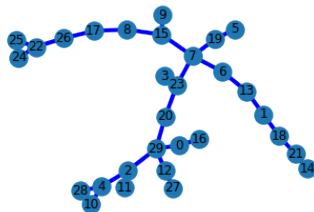
Shortest path

- **Shortest path** between nodes i and j is the path with the fewest number of links
- Shortest path between i and j is called **distance** between these nodes d_{ij}
- There can be multiple shortest paths between the same pair of nodes
- Shortest path cannot contain loops and cannot intersect itself
- undirected $d_{ij} = d_{ji}$; directed $d_{ij} \neq d_{ji}$

Path types



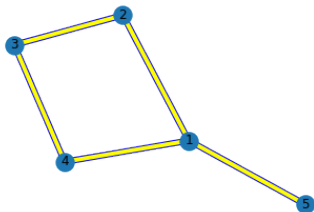
Cycle - a path with the same start and end node



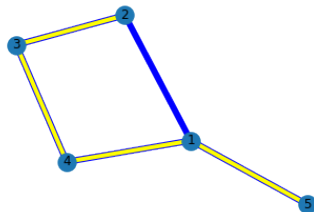
Networks without cycles - **acyclic network**

Connected acyclic network is a tree

Path types



Eulerian Path - a path that traverses each link exactly once



Hamiltonian Path - a path that visits each node exactly once

Average path length and diameter

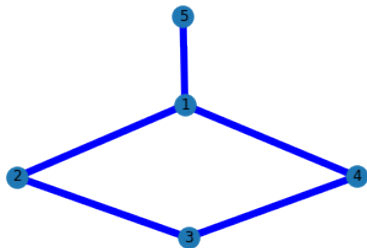
Average path length $\langle l \rangle$ - the average of the shortest paths between all pairs of nodes

$$\langle l \rangle = \frac{1}{\frac{1}{2}N(N-1)} \sum_{i>j} d_{ij}$$

Diameter - the longest shortest path in a network, or the distance between the two furthest nodes

$$D = \max_{i>j} (d_{ij})$$

Average path length and diameter



$$d_{12} = 1, d_{13} = 2, d_{14} = 1, \\ d_{15} = 1, d_{23} = 1, d_{2,4} = 2, \\ d_{25} = 2, d_{34} = 1, d_{35} = 3, \\ d_{45} = 2$$

$$\langle l \rangle = 1.6$$

$$D = 3$$

Diameter of disconnected network is $D = \infty$.

Number of shortest paths

Number of shortest paths N_{ij} and distance d_{ij} between nodes i and j can be calculated using adjacency matrix A

$d_{ij} = 1$: if $A_{ij} = 1$ $N_{ij} = 1$ otherwise it is 0

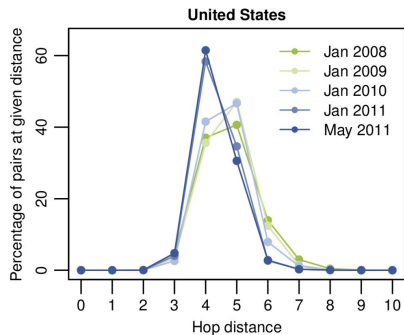
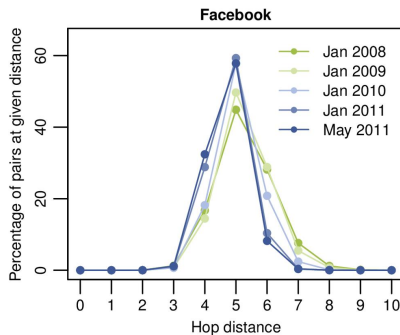
$d_{ij} = 2$: $A_{ij}^2 > 1$, $N_{ij} = A_{ij}^2$

$d_{ij} = 3$: $A_{ij}^3 > 1$, $N_{ij} = A_{ij}^3$

$d_{ij} = d$: $A_{ij}^d > 1$, $N_{ij} = A_{ij}^d$

Algorithm for finding shortest paths: Breadth-First Search (BFS)
Algorithm. See section 2.8 in the Network Science book for details.

Shortest path distribution



Source: <https://www.facebook.com/notes/10158927855913415/>

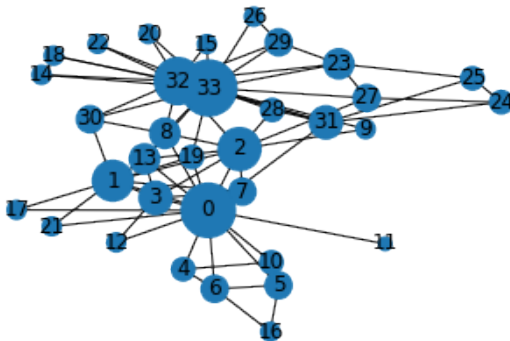
Why do we care?

- In complex networks not all nodes are equally important: airport New York is more important than some local airport
- The same is true for edges: New York - London connection
- Importance depends on the process - different centrality measures

Hubs and degree centrality

Node degree $q_i = \sum_{j=1}^N A_{ij}$ - degree centrality

Nodes with high number of neighbours are more important



Path based centrality measures

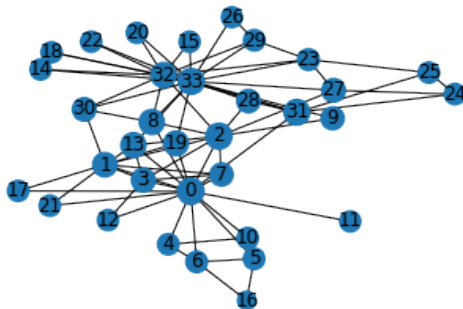
- Closeness centrality - how close is the node to other nodes
- Betweenness centrality of
 - a node - a fraction of shortest paths that pass through a node
 - an edge - a fraction of shortest paths that pass through an edge

Closeness centrality

Closeness centrality of node i - inverse of the sum of the distances of node i to all other nodes in the network

$$G_i = \frac{1}{\sum_{j \neq i} d_{ij}}$$

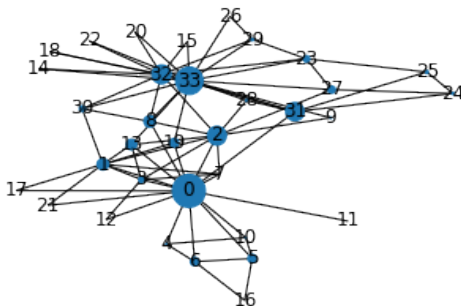
Normalized - $g_i = \frac{N-1}{\sum_{j \neq i} d_{ij}}$



Betweenness centrality of a node

σ_{hj} - number of shortest paths from h to j , and $\sigma_{hj}(i)$ the ones that pass through node i

Betweenness centrality of node i $b_i = \sum_{h \neq i, j \neq i} \frac{\sigma_{hj}(i)}{\sigma_{hj}}$



Betweenness centrality of an edge

σ_{hj} - number of shortest paths from h to j , and $\sigma_{hj}(e_{ik})$ the ones that pass through an edge e_{ik}

Betweenness centrality of an edge e_{ik} $b_{e_{ik}} = \sum_{h \neq i, k, j \neq i, k} \frac{\sigma_{hj}(e_{ik})}{\sigma_{hj}}$

