

TD2:

Exercise 1: /

Exercise 2: 🌟

① • $\llbracket P(a) \rrbracket_p^I = T$

$\llbracket \exists x. P(x) \rrbracket_p^I = T$

#2 eme
façon

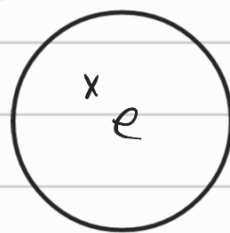
$D_I = \mathbb{N}$

$I(a) = 0$

$I(p) = \{n \in \mathbb{N} \mid n \bmod 2 = 0\}$

#1

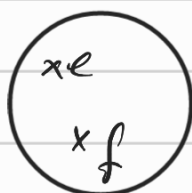
$D_I = \{e\}$



$I(a) = e$
 $I(p) = \{(e, T)\}$

• $\llbracket P(a) \rrbracket_p^I = F$

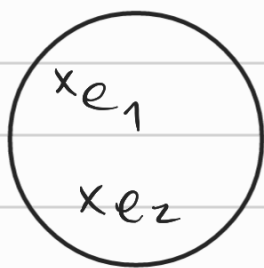
$\llbracket \exists x. P(x) \rrbracket_p^I = T$



$I(p) = \{(e, F), (f, T)\}$
 $I(a) = e$

• $\llbracket P(a) \rrbracket_p^I = T$

$\llbracket \exists x. P(x) \rrbracket_p^I = F$



$I(a) = e_1$

$I(p) = \{(e_1, T), (e_2, F)\}$

\downarrow
 $I(p)(I(a)) = T$

$\vee I(p)(e_2) = F$

$$\begin{aligned}
 & \bigvee_{v \in D} I(P)(v) = \dots \\
 & = I(P)(v_a) \bigvee_{v \in D \setminus \{v_a\}} I(P)(v) \\
 & = T \bigvee_B \dots \\
 & = T
 \end{aligned}$$

Dans D il y a forcément la valeur v_a qui interprète a : $I(a) = v_a \in D_I$

- $\text{---} = F$
 $\text{---} = f$
- \bigcirc^{D_I}
- $I(a) = e$
 $I(p) = \{(e, F)\}$
- $D_I = \{e\}$

②

- $D_I = \{e\}$ $F_1 = (\exists x. P(x)) \Rightarrow P(a)$
 $I(a) = e$
 $I(p) = \{(a, T)\}$
- $\bullet \llbracket F_1 \rrbracket_p^I = T$

- $\bullet \llbracket \forall x. P(x) \rrbracket_p^I = I(p)(e) = T$

- $D_I = \{e, f\}$ $\bullet \llbracket F_2 \rrbracket_p^I = T \Rightarrow F = F$
 $I(a) = e$
 $I(p) = \{(e, F), (f, T)\}$

- $\bullet \llbracket \forall x. P(x) \rrbracket_p^I = I(p)(e) \wedge I(p)(f) = F \wedge T = F$

- $D_I = \{e\}$ $\bullet \llbracket F_3 \rrbracket_p^I = F \Rightarrow F = T$

$$I(a) = e$$

$$I(P) = \{e, F\}$$

$$\bullet \llbracket \forall x. P(x) \rrbracket_P^I = I(P)(e) = F$$

$\llbracket P(a) \Rightarrow \exists x. P(x) \rrbracket_P^I$ Soit I une interprétation quelconque du vocabulaire on a forcément un élément v_a dans D_I qui interprète a , c'est à dire:
 $I(a) = v_a \in D_I$

$$I(P)(v_a) \Rightarrow_\beta \left(I(P)(v_a) \vee_\beta \bigvee_{v \in D_I \setminus \{v_a\}} I(P)(v) \right)$$

Par cas:

- Soit $I(P)(v_a) = F$ alors $(i) = F \Rightarrow_\beta \text{---} = T$
- Soit $I(P)(v_a) = T$ alors $(i) = T \Rightarrow_\beta \left(T \vee_\beta \bigvee_{v \in D_I \setminus \{v_a\}} I(P)(v) \right)$

Donc $P(a) \Rightarrow \exists x. P(x)$ est valide.

Exercice 3: ∇

$$\textcircled{4} (\forall x. P(x) \wedge Q(x)) \Rightarrow (\forall x. P(x)) \wedge (\forall x. Q(x)) = A$$

$$\llbracket A \rrbracket_P^I = \llbracket \forall x. P(x) \wedge Q(x) \rrbracket_P^I \Rightarrow_\beta \left(\llbracket \forall x. P(x) \rrbracket_P^I \wedge_\beta \llbracket \forall x. Q(x) \rrbracket_P^I \right)$$

Soit il existe $v_0 \in D_I$ tel que $I(P)(v_0) = F$
 ou $I(Q)(v_0) = F$

$$\llbracket A \rrbracket_P^I = \bigwedge_{v \in D_I} \left(I(P)(v) \wedge_\beta I(Q)(v) \right) \Rightarrow_\beta \dots$$

$$= \bigwedge_{v \in D_I \setminus \{v_0\}} \left(I(P)(v) \wedge_\beta I(Q)(v) \right) \wedge_\beta \left(I(P)(v_0) \wedge_\beta I(Q)(v_0) \right) \Rightarrow_\beta \dots$$

$$\llbracket A \rrbracket_P^I = F \Rightarrow \dots = T$$

Sinon, pour tout $v \in D_I$ on a : $I(P)(v) = I(Q)(v) = T$

$$\llbracket A \rrbracket_P^I = T \Rightarrow \left(\bigwedge_{v \in D_I} I(P)(v) \right) \wedge_P \left(\bigwedge_{v \in D_I} I(Q)(v) \right)$$

$$= \left(\bigwedge_{v \in D_I} I(P)(v) \right) \wedge_P \left(\bigwedge_{v \in D_I} I(Q)(v) \right)$$

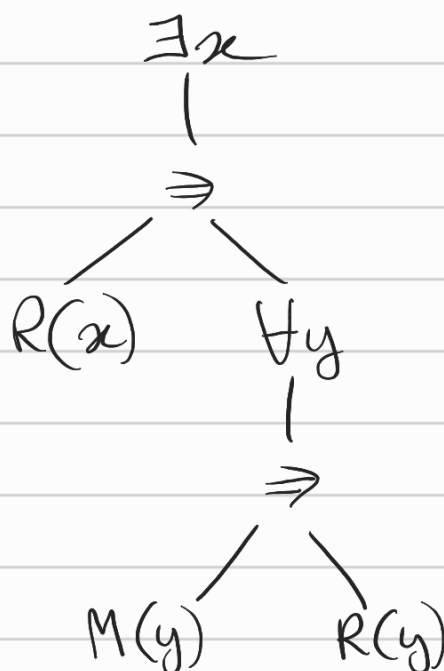
$$= \left(\bigwedge_{v \in D_I} T \right) \wedge_P \left(\bigwedge_{v \in D_I} T \right) = T$$

Exercice 4: 

$R(x)$ = "x résout ce problème"

$M(x)$ = "x est mathématicien"

C = "Cabot"



$$(\exists x. R(x)) \Rightarrow \forall y. M(y) \Rightarrow R(y)$$

$$(\forall x. R(x) \Rightarrow \forall y. M(y) \Rightarrow R(y)) \Rightarrow M(c) \wedge \neg R(c) \Rightarrow \forall z. R(z)$$

valide?

