DAY 9

```
def maxCoins(piles):
  # Sort the piles array in ascending order
  piles.sort()
  # Initialize the variable to keep track of your total coins
  total coins = 0
  # Since the total number of piles is 3n, we can process them in chunks of 3.
  # We will pick the second largest pile from each triplet
  for i in range(len(piles) // 3):
     total_coins += piles[-2 - 2*i] # This is the second largest in the triplet
  return total_coins
def minPatches(coins, target):
  coins.sort() # Sort the coins to make the greedy strategy work
  current_sum = 0 # This tracks the maximum sum we can form so far
  patches = 0 # The number of new coins we need to add
  i = 0 # Pointer to the coins array
  while current_sum < target:
     # If the current coin can be used to form the sum `current sum + 1`
     if i < len(coins) and coins[i] <= current_sum + 1:
       current sum += coins[i]
       i += 1
     else:
       # Otherwise, we need to add a coin of value `current_sum + 1`
       current_sum += current_sum + 1
       patches += 1
  return patches
def minimumTimeRequired(jobs, k):
  def canAssignJobs(maxTime):
     # Function to check if it's possible to assign jobs such that no worker has more than maxTime work
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workers = [0] * k # Initialize all workers' workloads to 0
     def backtrack(i):
       if i == len(jobs): # All jobs have been assigned
          return True
       for j in range(k):
          if workers[j] + jobs[i] <= maxTime:
            workers[j] += jobs[i] # Assign job to worker
            if backtrack(i + 1): # Recur to assign the next job
            workers[j] -= jobs[i] # Undo assignment if it didn't work
          # If a worker is idle (workers[j] == 0) and we already tried this, skip to avoid duplicate work
          if workers[j] == 0:
             break
       return False
     return backtrack(0)
  # Binary search for the minimum possible maximum working time
  low, high = max(jobs), sum(jobs)
  while low < high:
     mid = (low + high) // 2
     if canAssignJobs(mid):
       high = mid # Try to minimize the max working time
     else:
       low = mid + 1 # Increase the max working time
  return low
import bisect
def jobScheduling(startTime, endTime, profit):
  # Create a list of jobs with start, end, and profit
  jobs = list(zip(startTime, endTime, profit))
  # Sort jobs by their end time
  jobs.sort(key=lambda x: x[1])
  # dp[i] will store the maximum profit considering the first i jobs
  n = len(jobs)
  dp = [0] * (n + 1)
  # Helper function to find the rightmost job that doesn't overlap with job i
  def findLastNonConflictingJob(i):
     # We need to find the job whose end time is <= the current job's start time
     # Binary search on the end times
     low, high = 0, i - 1
     while low <= high:
       mid = (low + high) // 2
       if jobs[mid][1] \le jobs[i][0]:
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if jobs[mid + 1][1] <= jobs[i][0]:
             low = mid + 1
          else:
            return mid
       else:
          high = mid - 1
     return -1
  # Fill dp array
  for i in range(1, n + 1):
     # Option 1: Do not take this job, profit remains same as previous job
     dp[i] = dp[i - 1]
     # Option 2: Take this job
     prevJobIndex = findLastNonConflictingJob(i - 1)
     profitIncludingCurrent = jobs[i - 1][2] + (dp[prevJobIndex + 1] if prevJobIndex != -1 else 0)
     # Maximize profit by taking or not taking the job
     dp[i] = max(dp[i], profitIncludingCurrent)
  # The answer will be in dp[n], the maximum profit
  return dp[n]
# Example 1:
startTime1 = [1, 2, 3, 3]
endTime1 = [3, 4, 5, 6]
profit1 = [50, 10, 40, 70]
print(jobScheduling(startTime1, endTime1, profit1)) # Output: 120
# Example 2:
startTime2 = [1, 2, 3, 4, 6]
endTime2 = [3, 5, 10, 6, 9]
profit2 = [20, 20, 100, 70, 60]
print(jobScheduling(startTime2, endTime2, profit2)) # Output: 150
import heapq
def dijkstra(n, graph, source):
  # Initialize distances as infinity for all vertices except the source
  dist = [float('inf')] * n
  dist[source] = 0
  # Priority queue to store (distance, vertex) pairs
  pq = [(0, source)] # (distance, vertex)
  # While there are vertices to process
  while pg:
     # Get the vertex with the smallest tentative distance
     current_dist, u = heapq.heappop(pq)
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# If the current distance is greater than the stored distance, continue
     if current_dist > dist[u]:
        continue
     # Explore neighbors of vertex u
     for v in range(n):
        if graph[u][v] != float('inf'): # Check if there's an edge
           weight = graph[u][v]
           # If the new calculated distance is shorter, update it
           if dist[u] + weight < dist[v]:
              dist[v] = dist[u] + weight
              heapq.heappush(pq, (dist[v], v))
  return dist
# Test Case 1
n1 = 5
graph1 = [
  [0, 10, 3, float('inf'), float('inf')],
  [float('inf'), 0, 1, 2, float('inf')],
  [float('inf'), 4, 0, 8, 2],
  [float('inf'), float('inf'), float('inf'), 0, 7],
  [float('inf'), float('inf'), float('inf'), 9, 0]
source1 = 0
print(dijkstra(n1, graph1, source1)) # Output: [0, 7, 3, 9, 5]
# Test Case 2
n2 = 4
graph2 = [
  [0, 5, float('inf'), 10],
  [float('inf'), 0, 3, float('inf')],
  [float('inf'), float('inf'), 0, 1],
  [float('inf'), float('inf'), float('inf'), 0]
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source2 = 0
print(dijkstra(n2, graph2, source2)) # Output: [0, 5, 8, 9]
6.
import heapq
def dijkstra(n, edges, source, target):
  # Step 1: Convert the edge list to an adjacency list
  graph = {i: [] for i in range(n)}
  for u, v, w in edges:
     graph[u].append((v, w))
     graph[v].append((u, w)) # For undirected graph, add reverse edge too
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# Step 2: Initialize the distance array and priority queue
  dist = [float('inf')] * n
  dist[source] = 0
  pq = [(0, source)] # Priority queue stores (distance, vertex)
  # Step 3: Process the vertices
  while pq:
     current_dist, u = heapq.heappop(pq)
     # If we reached the target, return the distance
     if u == target:
       return current_dist
     # If the current distance is greater than the known distance, skip it
     if current dist > dist[u]:
       continue
     # Explore the neighbors of the current vertex
     for v, weight in graph[u]:
        new_dist = current_dist + weight
       if new dist < dist[v]:
          dist[v] = new dist
          heapq.heappush(pq, (new_dist, v))
  # If we reach here, no path exists to the target
  return -1
# Test Case 1
n1 = 6
edges1 = [
  (0, 1, 7), (0, 2, 9), (0, 5, 14), (1, 2, 10), (1, 3, 15),
  (2, 3, 11), (2, 5, 2), (3, 4, 6), (4, 5, 9)
source1 = 0
target1 = 4
print(dijkstra(n1, edges1, source1, target1)) # Output: 20
# Test Case 2
n2 = 5
edges2 = [
  (0, 1, 10), (0, 4, 3), (1, 2, 2), (1, 4, 4), (2, 3, 9),
  (3, 2, 7), (4, 1, 1), (4, 2, 8), (4, 3, 2)
source2 = 0
target2 = 3
print(dijkstra(n2, edges2, source2, target2)) # Output: 8
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class Node:
  def __init__(self, char, freq):
     self.char = char
     self.freq = freq
     self.left = None
     self.right = None
  # To make the node comparable based on frequency
  def It (self, other):
     return self.freq < other.freq
def build_huffman_tree(characters, frequencies):
  # Step 1: Create a priority queue (min-heap) and insert all the nodes
  priority queue = []
  for char, freg in zip(characters, frequencies):
     heapq.heappush(priority_queue, Node(char, freq))
  # Step 2: Build the Huffman Tree
  while len(priority queue) > 1:
     # Extract the two nodes with the lowest frequency
     left = heapq.heappop(priority queue)
     right = heapq.heappop(priority queue)
     # Create a new internal node with the sum of frequencies
     merged_node = Node(None, left.freq + right.freq)
     merged node.left = left
     merged node.right = right
     # Insert the new internal node back into the priority queue
     heapq.heappush(priority_queue, merged_node)
  # The remaining node is the root of the Huffman tree
  return priority queue[0]
def generate huffman codes(root, current code=""):
  # Step 3: Generate the Huffman codes by traversing the tree
  if root is None:
     return {}
  if root.char is not None: # It's a leaf node
     return {root.char: current code}
  # Recursively get the codes for the left and right subtrees
  codes = {}
  codes.update(generate huffman codes(root.left, current code + "0"))
  codes.update(generate_huffman_codes(root.right, current_code + "1"))
  return codes
def huffman coding(characters, frequencies):
  # Build the Huffman Tree
  root = build huffman tree(characters, frequencies)
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# Generate the Huffman Codes
  codes = generate huffman codes(root)
  # Return the codes sorted by the characters
  return sorted(codes.items())
# Test Case 1
characters1 = ['a', 'b', 'c', 'd']
frequencies 1 = [5, 9, 12, 13]
print(huffman coding(characters1, frequencies1))
# Expected Output: [('a', '110'), ('b', '10'), ('c', '0'), ('d', '111')]
# Test Case 2
characters2 = ['f', 'e', 'd', 'c', 'b', 'a']
frequencies2 = [5, 9, 12, 13, 16, 45]
print(huffman coding(characters2, frequencies2))
# Expected Output: [('a', '0'), ('b', '101'), ('c', '100'), ('d', '111'), ('e', '1101'), ('f', '1100')]
import heapq
class Node:
  def __init__(self, char, freq):
     self.char = char
     self.freq = freq
     self.left = None
     self.right = None
  # To make the node comparable based on frequency
  def It (self, other):
     return self.freq < other.freq
def build_huffman_tree(characters, frequencies):
  # Step 1: Create a priority queue (min-heap) and insert all the nodes
  priority queue = []
  for char, freq in zip(characters, frequencies):
     heapq.heappush(priority_queue, Node(char, freq))
  # Step 2: Build the Huffman Tree
  while len(priority queue) > 1:
     # Extract the two nodes with the lowest frequency
     left = heapq.heappop(priority queue)
     right = heapq.heappop(priority_queue)
     # Create a new internal node with the sum of frequencies
     merged_node = Node(None, left.freq + right.freq)
     merged node.left = left
     merged node.right = right
     # Insert the new internal node back into the priority queue
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heapq.heappush(priority_queue, merged_node)
  # The remaining node is the root of the Huffman tree
  return priority queue[0]
def decode huffman tree(root, encoded string):
  # Step 3: Decode the encoded string
  decoded message = []
  node = root # Start at the root of the Huffman tree
  for bit in encoded_string:
     # Traverse the tree based on the bit: '0' means left, '1' means right
     if bit == '0':
       node = node.left
     else:
       node = node.right
     # If we reach a leaf node, add the character to the result and reset to the root
     if node.char is not None:
       decoded message.append(node.char)
       node = root # Reset to the root
  # Join the decoded characters to form the decoded message
  return ".join(decoded message)
def huffman_decoding(characters, frequencies, encoded_string):
  # Build the Huffman Tree
  root = build huffman tree(characters, frequencies)
  # Decode the string using the Huffman Tree
  decoded_message = decode_huffman_tree(root, encoded_string)
  return decoded message
# Test Case 1
characters1 = ['a', 'b', 'c', 'd']
frequencies1 = [5, 9, 12, 13]
encoded string1 = '1101100111110'
print(huffman decoding(characters1, frequencies1, encoded string1))
# Expected Output: "abacd"
# Test Case 2
characters2 = ['f', 'e', 'd', 'c', 'b', 'a']
frequencies2 = [5, 9, 12, 13, 16, 45]
encoded_string2 = '110011011100101111001011'
print(huffman decoding(characters2, frequencies2, encoded string2))
# Expected Output: "fcbade"
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# Step 1: Sort the weights in descending order
  sorted weights = sorted(weights, reverse=True)
  # Step 2: Initialize the total weight loaded into the container
  total weight = 0
  # Step 3: Try to add items to the container
  for weight in sorted weights:
     if total weight + weight <= max capacity:
       total weight += weight
     else:
       break # Stop if adding this item exceeds the capacity
  return total weight
# Test Case 1
weights1 = [10, 20, 30, 40, 50]
max capacity 1 = 60
print(max_weight_loaded(weights1, max_capacity1)) # Expected Output: 50
# Test Case 2
weights2 = [5, 10, 15, 20, 25, 30]
max capacity2 = 50
print(max_weight_loaded(weights2, max_capacity2)) # Expected Output: 50
10.
def min_containers(weights, max_capacity):
  # Step 1: Sort the weights in descending order
  weights.sort(reverse=True)
  # Step 2: Initialize the number of containers and current container's remaining capacity
  container count = 0
  current capacity = 0
  # Step 3: Try to fit items into containers
  for weight in weights:
     if current capacity + weight <= max capacity:
       # Add the item to the current container
       current capacity += weight
     else:
       # The current container can't fit this item, so use a new container
       container count += 1
       current_capacity = weight # Start a new container with this item
  # Account for the last container being used
  if current capacity > 0:
     container_count += 1
  return container_count
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# Test Case 1
weights1 = [5, 10, 15, 20, 25, 30, 35]
max capacity 1 = 50
print(min_containers(weights1, max_capacity1)) # Expected Output: 4
# Test Case 2
weights2 = [10, 20, 30, 40, 50, 60, 70, 80]
max_capacity2 = 100
print(min_containers(weights2, max_capacity2)) # Expected Output: 6
11.
class UnionFind:
  def __init__(self, n):
     self.parent = list(range(n))
     self.rank = [0] * n
  def find(self, u):
     if self.parent[u] != u:
       self.parent[u] = self.find(self.parent[u]) # Path compression
     return self.parent[u]
  def union(self, u, v):
     root u = self.find(u)
     root_v = self.find(v)
     if root_u != root_v:
       # Union by rank
       if self.rank[root_u] > self.rank[root_v]:
          self.parent[root v] = root u
       elif self.rank[root_u] < self.rank[root_v]:
          self.parent[root_u] = root_v
       else:
          self.parent[root v] = root u
          self.rank[root u] += 1
       return True
     return False
def kruskal(n, edges):
  # Step 1: Sort the edges by weight
  edges.sort(key=lambda x: x[2]) # Sort by the third element in each tuple (weight)
  uf = UnionFind(n)
  mst = ∏
  mst_weight = 0
  # Step 2: Iterate through sorted edges
  for u, v, weight in edges:
     if uf.union(u, v): # If u and v were not connected
       mst.append((u, v, weight))
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mst_weight += weight
     # If we've added n-1 edges, we can stop
     if len(mst) == n - 1:
        break
  return mst, mst_weight
# Test Case 1
n1 = 4
m1 = 5
edges1 = [(0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4)]
mst1, mst_weight1 = kruskal(n1, edges1)
print("Edges in MST:", mst1)
print("Total weight of MST:", mst_weight1)
# Test Case 2
n2 = 5
m2 = 7
edges2 = [(0, 1, 2), (0, 3, 6), (1, 2, 3), (1, 3, 8), (1, 4, 5), (2, 4, 7), (3, 4, 9)]
mst2, mst_weight2 = kruskal(n2, edges2)
print("Edges in MST:", mst2)
print("Total weight of MST:", mst_weight2)
12.
class UnionFind:
  def __init__(self, n):
     self.parent = list(range(n))
     self.rank = [0] * n
  def find(self, u):
     if self.parent[u] != u:
        self.parent[u] = self.find(self.parent[u]) # Path compression
     return self.parent[u]
  def union(self, u, v):
     root_u = self.find(u)
     root_v = self.find(v)
     if root_u != root_v:
       # Union by rank
       if self.rank[root_u] > self.rank[root_v]:
          self.parent[root_v] = root_u
       elif self.rank[root_u] < self.rank[root_v]:
          self.parent[root_u] = root_v
       else:
          self.parent[root v] = root u
          self.rank[root u] += 1
        return True
     return False
```

```
def kruskal(n, edges):
  # Step 1: Sort the edges by weight
  edges.sort(key=lambda x: x[2]) # Sort by the third element in each tuple (weight)
  uf = UnionFind(n)
  mst = []
  mst weight = 0
  # Step 2: Iterate through sorted edges
  for u, v, weight in edges:
     if uf.union(u, v): # If u and v were not connected
       mst.append((u, v, weight))
       mst weight += weight
     # If we've added n-1 edges, we can stop
     if len(mst) == n - 1:
       break
  return mst, mst weight
def is mst unique(n, edges, given mst):
  # Step 1: Get the MST from Kruskal's algorithm
  mst, mst_weight = kruskal(n, edges)
  # Step 2: Check if the given MST matches the one found by Kruskal's
  given mst weight = sum(weight for u, v, weight in given mst)
  if given mst weight!= mst weight:
    return True, None # If the weights are not equal, the given MST is incorrect
  # Step 3: Check for alternative MSTs
  # Since the MST weight is the same, check if we can swap any edges without increasing the total weight
  # We will attempt to find another valid MST by considering edges with the same weight.
  # We will try to construct a different MST
  uf = UnionFind(n)
  alternative_mst = []
  # Start with the edges in the original MST
  mst set = set(given mst)
  # Add edges to the alternative MST from the sorted list
  for u, v, weight in edges:
     if uf.union(u, v) and (u, v, weight) not in mst set:
       alternative mst.append((u, v, weight))
     # If we have found a valid alternative MST, break
     if len(alternative mst) == n - 1:
       break
  # If we find a valid alternative MST, then the given MST is not unique
  if len(alternative mst) == n - 1:
     return False, alternative mst
  else:
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return True, None
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```
# Test Case 1
n1 = 4
m1 = 5
edges1 = [(0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4)]
given_mst1 = [(2, 3, 4), (0, 3, 5), (0, 1, 10)]
is_unique1, alt_mst1 = is_mst_unique(n1, edges1, given_mst1)
if is_unique1:
  print("Is the given MST unique?", True)
else:
  print("Is the given MST unique?", False)
  print("Another possible MST:", alt_mst1)
# Test Case 2
n2 = 5
m2 = 6
edges2 = [(0, 1, 1), (0, 2, 1), (1, 3, 2), (2, 3, 2), (3, 4, 3), (4, 2, 3)]
given_mst2 = [(0, 1, 1), (0, 2, 1), (1, 3, 2), (3, 4, 3)]
is_unique2, alt_mst2 = is_mst_unique(n2, edges2, given_mst2)
if is_unique2:
  print("Is the given MST unique?", True)
else:
  print("Is the given MST unique?", False)
  print("Another possible MST:", alt_mst2)
```