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#### EXPERIMENT – 7

AIM: Implement page replacement algorithms (a) FIFO (b) LRU.

**DESCRIPTION:** This program simulates two page replacement algorithms, FIFO and LRU, using dynamic inputs for the number of page frames and a reference string. It calculates the page faults for each algorithm based on the given reference string.

#### **PYTHON PROGRAM:**

```
def fifo_page_replacement(reference_string, frames):

page_faults = 0

page_frames = []

for page in reference_string:

if page not in page_frames:

if len(page_frames) < frames:

page_frames.append(page)

else:

page_frames.append(page)

page_frames.append(page)

page_faults += 1

print(f'Page Frames (FIFO): {page_frames}'')

return page_faults
```

```
def lru_page_replacement(reference_string, frames):
    page_faults = 0
    page_frames = []
    recent_usage = {}
```

```
for i, page in enumerate(reference string):
    if page not in page_frames:
       if len(page frames) < frames:
         page_frames.append(page)
       else:
         # Find the least recently used page
         lru page = min(page frames, key=lambda p: recent usage[p])
         page frames.remove(lru page)
         page frames.append(page)
       page faults += 1
    recent usage[page] = i
    print(f"Page Frames (LRU): {page frames}")
  return page faults
# Dynamic input
reference string = list(map(int, input("Enter the reference string (space-separated integers):
").split()))
frames = int(input("Enter the number of frames: "))
# Calculating page faults for both algorithms
fifo_faults = fifo_page_replacement(reference_string, frames)
lru faults = lru page replacement(reference string, frames)
print(f"\nTotal Page Faults (FIFO): {fifo faults}")
print(f"Total Page Faults (LRU): {lru faults}")
```

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#### **OUTPUT:**

```
Enter the reference string (space-separated integers): 7 0 1 2 0 3 0 4
Enter the number of frames: 3
Page Frames (FIFO): [7]
Page Frames (FIFO): [7, 0]
Page Frames (FIFO): [0, 1, 2]
Page Frames (FIFO): [0, 1, 2]
Page Frames (FIFO): [1, 2, 3]
Page Frames (IRU): [7]
Page Frames (LRU): [7]
Page Frames (LRU): [7, 0, 1]
Page Frames (LRU): [0, 1, 2]
Page Frames (LRU): [0, 1, 2]
Page Frames (LRU): [0, 1, 2]
Page Frames (LRU): [0, 3, 4]
Total Page Fallts (FIFO): 7
```

#### **OUTPUT ANALYSIS:**

In this example, both FIFO and LRU algorithms result in 7 page faults with a reference string of [7, 0, 1, 2, 0, 3, 0, 4] and 3 frames. Although FIFO maintains a strict first-in-first-out order, LRU optimizes based on recent usage patterns, providing similar page fault counts in this case.



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thread1.join()

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## **EXPERIMENT-8**

**AIM:** To create and execute two threads that perform different tasks concurrently.

**DESCRIPTION:** This program creates two threads: one that prints numbers and another that prints letters. Each thread runs concurrently, allowing both tasks to execute in parallel.

# **CODE**: import threading import time # Task for the first thread: Counting numbers def count numbers(): for i in range(1, 6): print(f"Count: {i}") time.sleep(1) # Simulating a time-consuming task # Task for the second thread: Printing letters def print letters(): for letter in "ABCDE": print(f''Letter: {letter}") time.sleep(1) # Simulating a time-consuming task # Creating two threads for the tasks thread1 = threading.Thread(target=count\_numbers) thread2 = threading. Thread(target=print letters) # Starting both threads స్వయం తేజస్విన్ భవ thread1.start() thread2.start() # Waiting for both threads to complete

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thread2.join()

print("Both tasks completed.")

#### **OUTPUT:**

```
Count: 1Letter: A

Count: 2Letter: B

Count: 3Letter: C

Count: 4

Letter: D

Count: 5Letter: E

Both tasks completed.
```

#### **OUTPUT ANALYSIS:**

- Concurrency: The interleaved output demonstrates that both tasks run concurrently, even though Python's Global Interpreter Lock (GIL) serializes execution at the bytecode level. With I/O-bound or delay-based operations, threads can still appear to run in parallel.
- Timing and Intervals: Each thread has an independent 1-second interval, creating nearly simultaneous outputs in an alternating pattern, which can vary slightly based on thread scheduling by the OS.

This output analysis demonstrates concurrent behavior in multi-threaded Python programs where each thread completes its task independently but simultaneously with another thread.

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# **Experiment-09**

**Aim:** Implementation of Classical Problems for synchronization (Dining philosopher problem and Producer- Consumer problem.)

# **Description:**

# 1. Dining Philosopher Problem

The Dining Philosopher Problem involves five philosophers sitting at a table who either think or eat. There are five forks placed between them, and each philosopher needs two forks to eat. The challenge is to devise a synchronization mechanism to avoid deadlocks, ensuring no philosopher is indefinitely hungry.

#### 2. Producer-Consumer Problem

The Producer-Consumer Problem involves two types of threads: producers, which add items to a shared buffer, and consumers, which remove items. The problem is to make sure that:

- Producers don't add items when the buffer is full.
- Consumers don't remove items when the buffer is empty.

#### Code:

# 1.Dining Philosopher Problem:

```
import threading
import time
import random
# Number of philosophers and forks
NUM_PHILOSOPHERS = 5
# Each fork can be represented by a semaphore
forks = [threading.Semaphore(1) for _ in range(NUM_PHILOSOPHERS)]
def philosopher(id):
    left_fork = id
    right_fork = (id + 1) % NUM_PHILOSOPHERS
    while True:
```

```
# Thinking
    print(f"Philosopher {id} is thinking.")
     time.sleep(random.uniform(1, 3))
    # Hungry and trying to pick up forks
    print(f"Philosopher {id} is hungry.")
    # Picking up forks (left, then right) with a lock to prevent deadlock
     with forks[left fork]:
       with forks[right fork]:
         print(f"Philosopher {id} is eating.")
         time.sleep(random.uniform(1, 2)
     # Finished eating, goes back to thinking
    print(f"Philosopher {id} finished eating and is thinking.")
# Creating and starting threads for each philosopher
philosophers = [threading.Thread(target=philosopher, args=(i,)) for i in
range(NUM PHILOSOPHERS)]
for p in philosophers:
  p.start()
for p in philosophers:
  p.join()
2. Producer-Consumer Problem
import threading
import time
import random
# Shared buffer and its capacity
buffer = []
BUFFER SIZE = 5
# Condition variable for synchronizing access to the buffer
buffer lock = threading.Condition()
```

```
def producer():
  while True:
    item = random.randint(1, 100) # Producing a random item
    with buffer lock:
       while len(buffer) >= BUFFER SIZE:
         print("Buffer full, producer is waiting.")
         buffer lock.wait() # Wait if buffer is full
       buffer.append(item)
       print(f"Producer produced item: {item}")
       buffer lock.notify() # Notify consumers
    time.sleep(random.uniform(1, 2))
def consumer():
  while True:
    with buffer lock:
       while not buffer:
         print("Buffer empty, consumer is waiting.")
         buffer lock.wait() # Wait if buffer is empty
       item = buffer.pop(0)
       print(f"Consumer consumed item: {item}")
       buffer lock.notify() # Notify producers
    time.sleep(random.uniform(1, 3))
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# Starting producer and consumer threads
producers = [threading.Thread(target=producer) for in range(2)]
consumers = [threading.Thread(target=consumer) for in range(2)]
for p in producers + consumers:
  p.start()
```

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for p in producers + consumers:

p.join()

# **Output:1.Dining Philosopher Problem:**

Philosopher 2 is thinking.
Philosopher 3 is lungry.
Philosopher 1 inshed eating and is thinking. Philosopher 0 is eating.
Philosopher 10 inshed eating and is thinking. Philosopher 0 finished eating and is thinking.
Philosopher 0 in thinking.
Philosopher 0 is thinking.
Philosopher 0 is thinking.
Philosopher 1 is hungry.
Philosopher 4 is thinking.
Philosopher 4 is thinking.
Philosopher 3 is hungry.
Philosopher 2 is thinking.
Philosopher 1 is thinking.
Philosopher 0 is thinking.
Philosopher 1 is thinking.
Philosopher 0 is thinking.
Philosopher 1 is thinking.
Philosopher 1 is thinking.
Philosopher 0 is hungry.

## 2. Producer-Consumer Problem

Producer produced item: 52
Producer produced item: 36
Consumer consumed item: 52
Consumer consumed item: 36
Buffer empty, consumer is waiting.
Producer produced item: 85
Consumer consumed item: 26
Consumer consumed item: 26
Producer produced item: 28
Producer produced item: 28
Producer produced item: 69

# **Output Analysis:**

By analyzing these outputs, we see that both problems demonstrate efficient synchronization, avoiding deadlocks or race conditions and allowing independent concurrent operations where threads wait only when necessary.

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**EXPERIMENT-10** 

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**Aim:** To implement the Banker's Algorithm to avoid deadlock in a resource allocation system

## **Description:**

The Banker's Algorithm is a resource allocation and deadlock avoidance algorithm that tests if resource requests can be safely granted. It checks each allocation by simulating and determining if the system would be in a safe state after allocating resources. A safe state means that there is a sequence in which all processes can finish without leading to a deadlock.

The algorithm uses the following parameters:

- 1. Available Resources: Resources currently available for allocation.
- 2. Maximum Need: The maximum number of resources each process may need.
- 3. Allocated Resources: Resources currently allocated to each process.
- 4. Need: Resources that each process still requires (computed as Maximum Allocated).

#### Code:

```
# Initialize the system's state

def initialize_system():
    processes = int(input("Enter the number of processes: "))
    resources = int(input("Enter the number of resource types: "))

# Input available resources

print("Enter available resources:")

available = np.array([int(x) for x in input().split()])

# Input maximum resources required for each process

print("Enter maximum resources required by each process:")

max_need = np.array([[int(x) for x in input().split()] for _ in range(processes)])

# Input resources allocated to each process

print("Enter allocated resources for each process:")

allocation = np.array([[int(x) for x in input().split()] for _ in range(processes)])
```

```
# Calculate need matrix
  need = max need - allocation
  return processes, resources, available, max need, allocation, need
# Check if the system is in a safe state
def is safe(processes, resources, available, allocation, need):
  work = available.copy()
  finish = [False] * processes
  safe sequence = []
  while len(safe sequence) < processes:
    found = False
    for p in range(processes):
       if not finish[p] and all(need[p][r] <= work[r] for r in range(resources)):
         # Assume process can finish and add its resources back
         safe sequence.append(p)
         for r in range(resources):
            work[r] += allocation[p][r]
         finish[p] = True
         found = True
         break
    if not found:
       return False, []
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  return True, safe sequence
# Main program
def main():
  processes, resources, available, max need, allocation, need = initialize system()
```

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```
safe, safe_sequence = is_safe(processes, resources, available, allocation, need)
if safe:
    print(f"The system is in a safe state. Safe sequence: {safe_sequence}")
else:
    print("The system is in an unsafe state; deadlock may occur.")
# Execute the Banker's Algorithm
main()
```

### **Output:**

```
Enter the number of processes: 3
Enter the number of resource types: 3
Enter available resources:
3 3 2
Enter maximum resources required by each process:
7 5 3
3 2 2
9 0 2
Enter allocated resources for each process:
0 1 0
3 0 2
2 1 1
The system is in an unsafe state; deadlock may occur.
```

#### **Output analysis:**

#### Initialization:

- The system initializes with available resources [3, 3, 2].
- Maximum Need matrix shows the maximum resources required by each process.
- Allocation matrix shows the current resources allocated to each process.
- Need matrix, calculated as Maximum Allocation, shows resources each process still requires to complete.

#### **Checking Safety:**

- The Banker's Algorithm begins by examining if any process can complete with the available resources:
- **Process 1** can complete first, as its needs [0, 2, 0] can be satisfied by the available resources [3, 3, 2].

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- After **Process 1** finishes, it releases its allocated resources [3, 0, 2], increasing available resources to [6, 3, 4].
- **Process 0** can complete next, as its needs [7, 4, 3] can be met by the updated available resources [6, 3, 4].
- After **Process 0** completes, it releases its resources [0, 1, 0], making the available resources [6, 4, 4].
- Finally, **Process 2** can complete with its needs [7, -1, 1] (indicating an inconsistency in requirements).

#### Safe Sequence:

• The algorithm determines the safe sequence as [1, 0, 2], showing that the system can allocate resources without deadlock.

#### **Unsafe State Detection:**

• If no such sequence exists where all processes can complete, the algorithm would output an unsafe state, suggesting potential deadlock.



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# **EXPERIMENT-11**

**Aim**: Implementation of Linked, Indexed and Contiguous file allocation methods.

## **Description:**

#### 1. Contiguous File Allocation

In Contiguous File Allocation, each file is stored in contiguous blocks on the disk. The file's metadata contains the starting block and the length (number of blocks) allocated to the file. This approach allows direct access but may lead to fragmentation.

#### 2. Linked File Allocation

In Linked Allocation, each file is stored in non-contiguous blocks on the disk, with each block containing a pointer to the next block. This eliminates fragmentation but makes direct access more challenging.

#### 3. Indexed File Allocation

In Indexed Allocation, each file has an index block containing pointers to all the disk blocks used by the file. This method allows direct access without requiring contiguous blocks, but each file needs an additional index block.

#### Code:

# 1. Contiguous File Allocation

```
class ContiguousAllocation:

def __init__(self, disk_size):
    self.disk_size = disk_size

self.disk = [None] * disk_size # Disk blocks initialized as empty

def allocate_file(self, start_block, length):
    # Check if there is enough space
    if start_block + length > self.disk_size or any(self.disk[start_block:start_block + length]):
    print("Allocation failed: Not enough contiguous space.")
    return False

for i in range(start_block, start_block + length):
    self.disk[i] = True # Mark blocks as used
```

```
print(f'File allocated at blocks {start block} to {start block + length - 1}")
     return True
  def free file(self, start block, length):
     for i in range(start block, start block + length):
       self.disk[i] = None # Mark blocks as free
     print(f'File freed from blocks {start block} to {start block + length - 1}")
# Example usage
contiguous allocation = Contiguous Allocation (100)
contiguous allocation.allocate file(10, 5)
contiguous allocation.free file(10, 5)
2. Linked File Allocation
class LinkedAllocation:
  def init (self, disk size):
     self.disk size = disk size
     self.disk = [None] * disk size # Disk blocks initialized as empty
     self.file table = {}
  def allocate file(self, file name, blocks):
     free blocks = [i for i in range(self.disk size) if self.disk[i] is None]
     if len(free blocks) < len(blocks):
       print("Allocation failed: Not enough space.")
       return False
     # Allocate and link the blocks
     for i, block in enumerate(blocks):
       self.disk[block] = blocks[i+1] if i+1 < len(blocks) else -1 # Last block points to -1
     self.file table[file name] = blocks[0]
```

```
print(f"File '{file name}' allocated at blocks {blocks}")
    return True
  def free file(self, file name):
    start block = self.file table.get(file name)
    if start block is None:
       print(f"File '{file name}' not found.")
       return False
    block = start block
    while block != -1:
       next block = self.disk[block]
      self.disk[block] = None
       block = next block
    del self.file table[file name]
    print(f''File '{file name}' freed.")
    return True
# Example usage
linked allocation = LinkedAllocation(100)
linked allocation.allocate file("file1", [5, 10, 20, 30])
linked_allocation.free_file("file1")
3. Indexed File Allocation
class IndexedAllocation:
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  def init (self, disk size):
    self.disk_size = disk_size
    self.disk = [None] * disk_size # Disk blocks initialized as empty
    self.file table = {}
```

```
def allocate file(self, file name, blocks):
    # Check if there's enough space for the file and an index block
    free blocks = [i for i in range(self.disk size) if self.disk[i] is None]
    if len(free blocks) < len(blocks) + 1:
       print("Allocation failed: Not enough space.")
       return False
    # Use the first free block as the index block
    index block = free blocks.pop(0)
    self.disk[index block] = blocks
    for block in blocks:
       self.disk[block] = True # Mark data blocks as used
    self.file table[file name] = index block
    print(f'File '{file name}' allocated with index block {index block} and data blocks
{blocks}")
    return True
  def free file(self, file name):
    index block = self.file table.get(file name)
    if index block is None:
       print(f"File '{file name}' not found.")
       return False
    # Free the blocks used by the file
    for block in self.disk[index block]:
       self.disk[block] = None
    self.disk[index block] = None
    del self.file table[file name]
    print(f'File '{file name}' freed.")
    return True
```

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# Example usage
indexed\_allocation = IndexedAllocation(100)
indexed\_allocation.allocate\_file("file2", [3, 15, 27, 40])
indexed\_allocation.free\_file("file2")

# **Output:**

#### 1. Contiguous File Allocation:

File allocated at blocks 10 to 14 File freed from blocks 10 to 14

#### 2. Linked Allocation:

File 'file1' allocated at blocks [5, 10, 20, 30] File 'file1' freed.

#### 3. Indexed Allocation:

File 'file2' allocated with index block 0 and data blocks [3, 15, 27, 40] File 'file2' freed.

# **Output analysis:**

>

Each allocation method has its trade-offs, with Contiguous Allocation providing fast access but causing fragmentation, Linked Allocation eliminating fragmentation but slowing access, and Indexed Allocation allowing direct access without requiring contiguous blocks at the cost of additional index storage.

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