

Q.1.

$$a. f(m) = (m^3)^3 ; g(m) = m^{(m^3)}$$

$$f(m) = m^{3m}$$

$$\text{Taking log both sides. } g(m) = m^{m^3}$$

$$f(m) = 3m \log m$$

$$g(m) = m^3 \log m$$

$$\lim_{m \rightarrow 0} \frac{f(m)}{g(m)} = \frac{3m \log m}{m^3 \log m} = \frac{3}{m^2} = 0$$

Since  $\lim_{m \rightarrow 0} \frac{f(m)}{g(m)} = 0$  (We can conclude)

Answer.

$$\textcircled{1} f(m) = o(g(m)) \quad \& \quad f(m) = O(g(m)) \text{ (Little \& Big Oh..)}$$

$$b) f(m) = m! ; g(m) = (m+1)!$$

$$f(m) = 1 \times 2 \times 3 \times \dots \times (m-1) \times m ; g(m) = 1 \times 2 \times 3 \times \dots \times (m-1) \times m \times (m+1)$$

$$f(m) \approx m^m ; g(m) \approx m^{m+1}$$

$$\lim_{m \rightarrow 0} \frac{f(m)}{g(m)} = \frac{m^m}{m^{m+1}} = m^{(m-m-1)} = \frac{1}{m} = 0$$

We can conclude the following.

Answer.

$$\textcircled{1} f(m) = o(g(m)) \quad \& \quad \textcircled{2} f(m) = O(g(m)) \text{ (Little \& Big Oh..)}$$

c.  $f(n) = n^{0.99} + 15 (\log_2 n)^{100}$  ;  $g(n) = \frac{1}{6} \log_6 n^2$

$f(n) = n^{0.99} + 15 (\log_2 n)^{100}$  ;  $g(n) = n^{2 \cdot \log_6 \frac{1}{6}} = n^{2 \times \frac{1}{2}} = n$

$\therefore$   $\boxed{\text{Asymptotically equal}} = f(n) \Theta g(n)$   $g(n) = n$

Alternatively we can take log both sides

$f(n) = 0.99 \log n + 1500 \log \log n$  ;  $g(n) = \log n$

$\therefore$  Conclude that  $\boxed{\begin{array}{l} \textcircled{1} f(n) = \Theta g(n) \\ \textcircled{2} f(n) = O g(n) \\ \textcircled{3} f(n) = \Omega g(n) \end{array}}$   $\leftarrow$  Answer

4.  $f(n) = n^{0.000001}$  ;  $g(n) = (\log n)^{1000001}$

Take logs both the sides

$f(n) = 0.000001 \log n$  ;  $g(n) = 1000001 \log \log n$

$f(n) \geq g(n)$

We can conclude  $\boxed{\begin{array}{l} \textcircled{1} f(n) = \Omega g(n) \\ \textcircled{2} f(n) = \omega g(n) \end{array}}$   $\leftarrow$  Answer

5.  $f(n) = n^{5/\log_2 n}$  ;  $g(n) = 10000$

Take logs both sides

$f(n) = \frac{5}{\log_2 n} \times \log n$  ;  $g(n) = \log 10000$

$f(n) = 5$  ;  $g(n) = \log 10000$

$\therefore$  Asymptotically equal.

$$f(n) \approx g(n)$$

$$\begin{array}{|l} \textcircled{1} f(n) \Theta g(n) \\ \textcircled{2} f(n) O g(n) \\ \textcircled{3} f(n) \Omega g(n) \end{array}$$

→ Answer.

Problem 1-2.

$$1. \frac{n^2 + 2}{1 + n^3 \cdot 2^{-n}} \approx \frac{n^2}{n^3 \cdot 2^{-n}} = \frac{2^n}{n} \approx \Theta(2^n) \leftarrow \text{Answer.}$$

$$2. \log(\sqrt{\log(5n)}) \cdot \log(37n^3 + 45) = \log_2 n^{1/3} \log_2 n^3 \\ \approx \frac{1}{3} \times 3 \times \log n \times \log n = (\log n)^2 \approx \Theta(\log n)^2 \leftarrow \text{Answer.}$$

$$3. \log(n!) + 10^{205} n \approx \log(n^n) + n \approx n \log n + n \\ \approx \Theta(n \log n) \leftarrow \text{Answer.}$$

$$4. \frac{6^n - 1000}{2^n + 1} \approx \frac{6^n}{2^n} \approx 3^n \approx \Theta(3^n) \leftarrow \text{Answer.}$$

$$5. n^{2/\log n} + 1000 \approx n^{2 \cdot 1/\log n} \approx \sqrt[n]{n^2} \\ \approx \Theta(1) \leftarrow \text{Answer.}$$

$$6. 2^{n + \log n} \approx 2^n \cdot 2^{\log n} \approx 2^n \cdot n \cdot \log 2 \\ \approx 2^n \cdot n \approx \Theta(n 2^n) \leftarrow \text{Answer.}$$

$$7. 2^m + 10 (\log m)^{\log m} \approx 2^m = \Theta(2^m) \leftarrow \text{Answer.}$$

$$8. 2^{3m} + 4^m \approx 2^{3m} + 2^{(m+1)} \approx 2^{3m} + 2^m \\ \approx \Theta(2^{3m}) \leftarrow \text{Answer.}$$

b) Functions per asymptotic growth. (least to maximum.)

$$\cancel{1, (\log m)^2, m \log m, 2^m, 2^m, m 2^m, 3^m,}$$

$$1, (\log m)^2, m \log m, 2^m, 2^m, m 2^m, 2^{3m}, 3^m.$$