

LAB 2 discussion

Simulation based Process Scheduling

- Explore multiple Schedulers and what OS tracks
- You will learn how to write a Discrete Event Simulator
- Specification of Process and Random Behavior

Process-1	0	200	40	90
	40	100	10	40
	50	20	10	10
Process-4	60	200	5	20

process
arrival/start

Total
CPU-time

CPU
Burst

IO
Burst

- Cpu-burst = random (range [1.. Proc.cpu_burst])
- Ioburst = random (range [1.. Proc.io_burst])

What you will generate ?

```
$ cat input4
```

```
0 200 40 90
40 100 10 40
50 20 10 10
60 200 5 20
```

```
$ ./sched -sF input4 rfile
```

FCFS

0000:	0	200	40	90	2		575	575	360	15
0001:	40	100	10	40	4		532	492	300	92
0002:	50	20	10	10	2		82	32	10	2
0003:	60	200	5	20	4		1174	1114	773	141
SUM:	1174	44.29	81.26	553.25	62.50	0.341				

static-prio

Finish Time

Turnaround time

I/O Time

CPU Wait time

Total Sim time

CPU Utilization

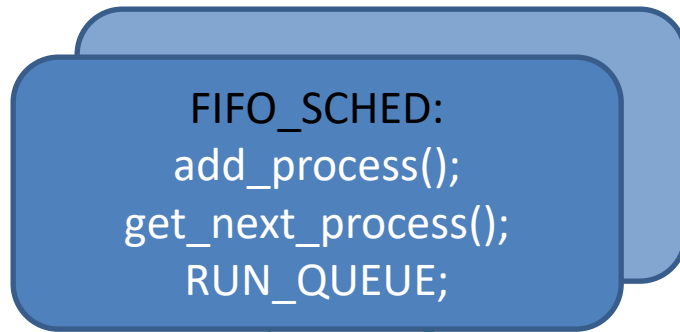
I/O Utilization

Avg Turnaround time

Avg Wait time

Throughput

Generic Structure of your Lab-2 infrastructure



Generic Interface
Allows for plugging
Different schedulers

(*add_process());

(*get_next_process());

Simulation and Scheduler Interaction

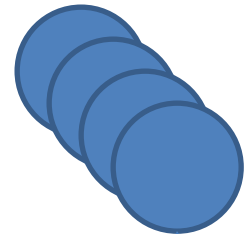
```
While (get_event()) {  
    decode_event() and run action(); // see next page  
}
```

get_event()

put_event()

rm_event()

Processes

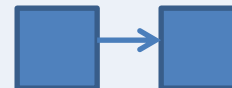


Event

DES-Layer



Time ordered



Timestamp
Process
oldstate
newstate

```

void Simulation() {
    EVENT* evt;
    bool CALL_SCHEDULER = false;
    while( (evt = get_event()) ) {
        Process *proc    = evt->evtProcess;    // this is the process the event works on
        CURRENT_TIME     = evt->evtTimeStamp;  // time jumps discretely
        int timeInPrevState = CURRENT_TIME - proc->state_ts; // good for accounting
        int transition = evt->transition;
        delete evt; evt = nullptr;    //remove current event object from Memory and don't touch anymore

        switch(transition) { // which state to transition to?
            case TRANS_TO_READY:
                // must come from BLOCKED or from PREEMPTION
                // must add to run queue
                CALL_SCHEDULER = true; // conditional on whether something is run
                break;

            case TRANS_TO_RUN:
                // create event for either preemption or blocking
                break;

            case TRANS_TO_BLOCK:
                //create an event for when process becomes READY again
                CALL_SCHEDULER = true;
                break;

            case TRANS_TO_PREEMPT:
                // add to runqueue (no event is generated)
                CALL_SCHEDULER = true;
                break;
        }

        if(CALL_SCHEDULER) {
            if (get_next_event_time() == CURRENT_TIME) {
                continue; //process all events at the same time before scheduling
            }
            CALL_SCHEDULER = false;
            if (CURRENT_RUNNING_PROCESS == nullptr) {
                CURRENT_RUNNING_PROCESS = THE_SCHEDULER->get_next_process();
                if (CURRENT_RUNNING_PROCESS == nullptr)
                    continue;
                // create event to make process runnable for same time.
            }
        }
    }
}

```

./sched -v -e input_show rfile

ShowEventQ: 0:0 20:1

0 0 0: CREATED -> READY

AddEvent(0:0:RUNNG): 20:1:READY ==> 0:0:RUNNG 20:1:READY

0 0 0: READY -> RUNNG cb=8 rem=100 prio=1

AddEvent(8:0:BLOCK): 20:1:READY ==> 8:0:BLOCK 20:1:READY

8 0 8: RUNNG -> BLOCK ib=2 rem=92

AddEvent(10:0:READY): 20:1:READY ==> 10:0:READY 20:1:READY

10 0 2: BLOCK -> READY

AddEvent(10:0:RUNNG): 20:1:READY ==> 10:0:RUNNG 20:1:READY

10 0 0: READY -> RUNNG cb=10 rem=92 prio=1

AddEvent(20:0:BLOCK): 20:1:READY ==> 20:1:READY 20:0:BLOCK

20 1 0: CREATED -> READY

20 0 10: RUNNG -> BLOCK ib=7 rem=82

AddEvent(27:0:READY): ==> 27:0:READY

AddEvent(20:1:RUNNG): 27:0:READY ==> 20:1:RUNNG 27:0:READY

20 1 0: READY -> RUNNG cb=7 rem=100 prio=3

AddEvent(27:1:BLOCK): 27:0:READY ==> 27:0:READY 27:1:BLOCK

Input file

0 100 10 10

20 100 20 10

RESULTS OF SIMULATION

FCFS

0000: 0 100 10 10 2 | 234 234 89 45

0001: 20 100 20 10 4 | 226 206 77 29

SUM: 234 85.47 57.26 220.00 37.00 0.855

./sched -v -e input_show rfile

Input file

0 100 10 10
20 100 20 10

ShowEventQ: 0:0 20:1

line triggered by "-v"
Timestamp pid howlong:
FROM -> TO

0 0 0: CREATED -> READY

AddEvent(0:0:RUNNG): 20:1:READY ==> 0:0:RUNNG 20:1:READY

0 0 0: READY -> RUNNG cb=8 rem=100 prio=1

AddEvent(8:0:BLOCK): 20:1:READY ==> 8:0:BLOCK 20:1:READY

One Event

8 0 8: RUNNG -> BLOCK ib=2 rem=92

AddEvent(10:0:READY): 20:1:READY ==> 10:0:READY 20:1:READY

10 0 2: BLOCK -> READY

AddEvent(10:0:RUNNG): 20:1:READY ==> 10:0:RUNNG 20:1:READY

line triggered by "-e"
event added: (time,pid,transition)
EventQ-Before → EventQ-After
(should be time ordered)

10 0 0: READY -> RUNNG cb=10 rem=92 prio=1

AddEvent(20:0:BLOCK): 20:1:READY ==> 20:1:READY 20:0:BLOCK

20 1 0: CREATED -> READY

20 0 10: RUNNG -> BLOCK ib=7 rem=82

AddEvent(27:0:READY): ==> 27:0:READY

AddEvent(20:1:RUNNG): 27:0:READY ==> 20:1:RUNNG 27:0:READY

20 1 0: READY -> RUNNG cb=7 rem=100 prio=3

AddEvent(27:1:BLOCK): 27:0:READY ==> 27:0:READY 27:1:BLOCK

./sched -v -t input_show rfile

Input file

0 100 10 10
20 100 20 10

```
0 0 0: CREATED -> READY
SCHED (1): 0:0
0 0 0: READY -> RUNNG cb=8 rem=100 prio=1
8 0 8: RUNNG -> BLOCK ib=2 rem=92
SCHED (0):
10 0 2: BLOCK -> READY
SCHED (1): 0:10
10 0 0: READY -> RUNNG cb=10 rem=92 prio=1
20 1 0: CREATED -> READY
20 0 10: RUNNG -> BLOCK ib=7 rem=82
SCHED (1): 1:20
20 1 0: READY -> RUNNG cb=7 rem=100 prio=3
27 0 7: BLOCK -> READY
27 1 7: RUNNG -> BLOCK ib=9 rem=93
SCHED (1): 0:27
27 0 0: READY -> RUNNG cb=7 rem=82 prio=1
34 0 7: RUNNG -> BLOCK ib=1 rem=75
SCHED (0):
35 0 1: BLOCK -> READY
SCHED (1): 0:35
35 0 0: READY -> RUNNG cb=9 rem=75 prio=1
36 1 9: BLOCK -> READY
44 0 9: RUNNG -> BLOCK ib=9 rem=66
SCHED (1): 1:36
44 1 8: READY -> RUNNG cb=11 rem=93 prio=3
53 0 9: BLOCK -> READY
55 1 11: RUNNG -> BLOCK ib=4 rem=82
```

line triggered by “-t” to show scheduler runqueue

SCHED(len): { “pid:timestamp” }*

where len is length of RQ followed by RQ entries

cb=9

Current cpu-burst

ib=9

Current io-burst

prio=3

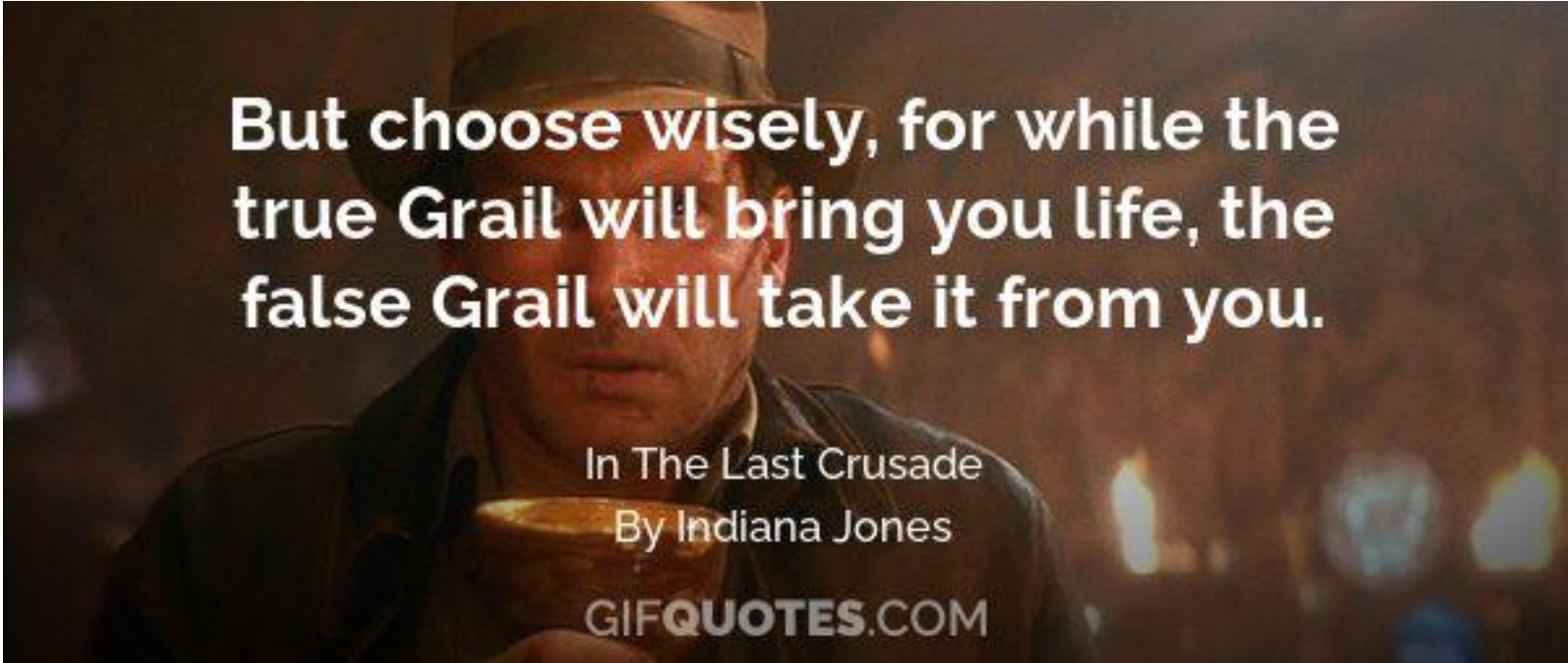
Dynamic priority

rem=82

Remaining cputime

Data Structures

- Schedulers are often about pooling requests and selecting:
 - Queueing , searching , indexing, mapping, priorities
 - Efficiency matters, not just effectiveness
 - Underlying data structures matter

A still from the movie 'The Last Crusade' showing Indiana Jones (Harrison Ford) wearing his signature fedora and holding a golden chalice (the Grail). The background is dark and atmospheric, with some light flares.

**But choose wisely, for while the
true Grail will bring you life, the
false Grail will take it from you.**

In The Last Crusade
By Indiana Jones

GIFQUOTES.COM

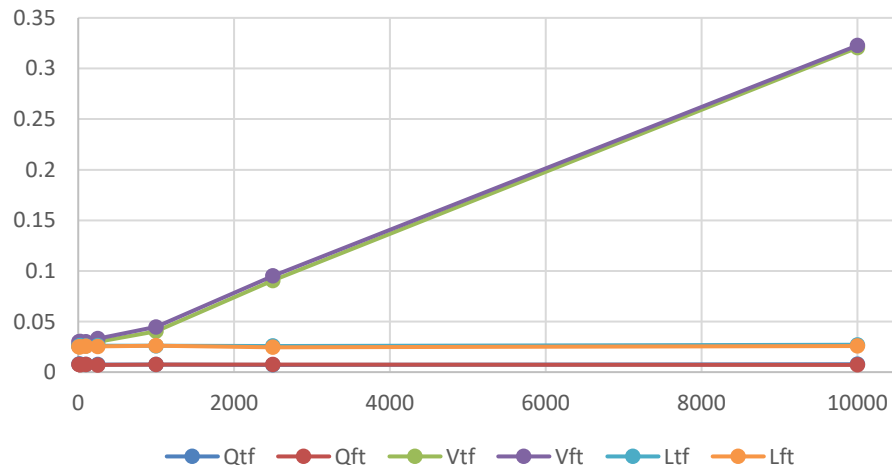
Example: Queue Semantics

- Assume existing queue size N
- Consider the following operations
 - Add front / delete tail (FT)
 - Add Tail / delete front (TF)
 - Iterate through queue with
 - C++ iterator (Iter) or C++ [] operator
- Time these operations
- Consider the following container classes:
 - Deque (Q)
 - Vector (V)
 - List (L)

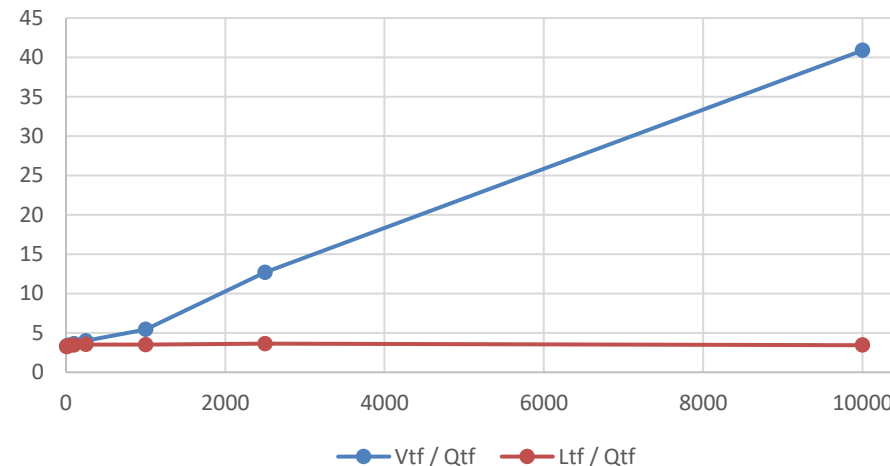
Add / Delete Opposite ends

- Scale the size of the queue (x-axis = N)
- TF and FT semantics
- Why are vectors so expensive ?

Add / Delete at opposite ends



Relative Overhead

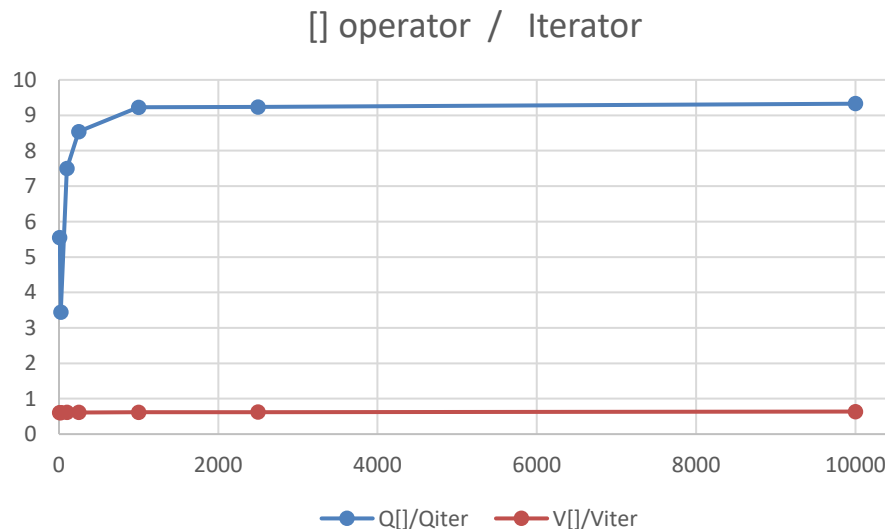


Iterations

- Iterate by 2 approaches:

iter: `for (x : objs) { x } ;`

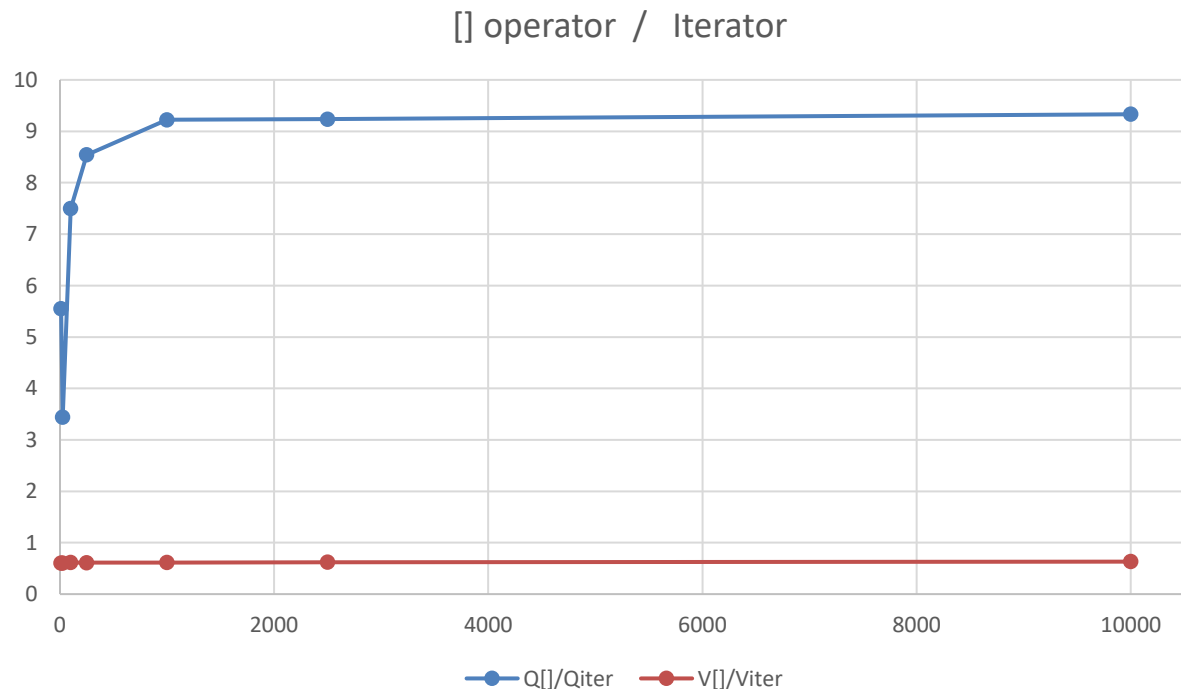
[]: `for (i = 0; i<obj.size(); i++) { x }`



note: list[] does not exist

Iteration vector vs deque

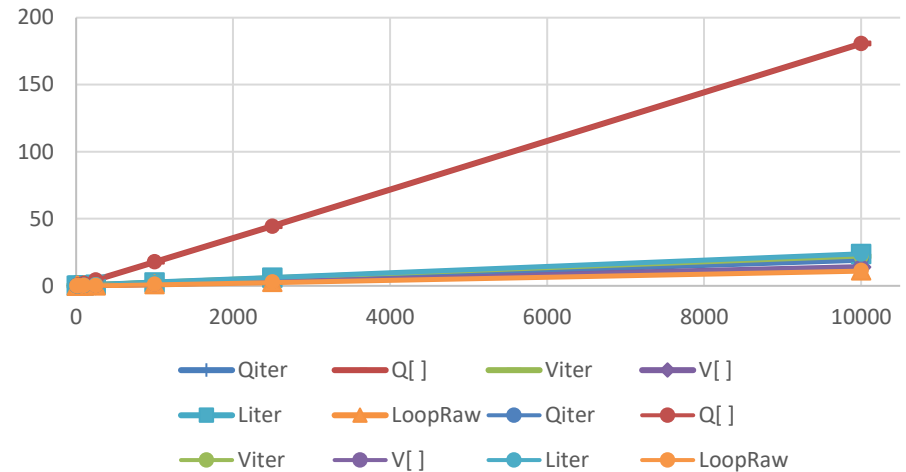
- `[]` operator relatively expensive for deque over iterator, but not for vector



Iteration vector vs deque vs list

- Avoid [] operator on queues
- Think about underlying data structures how something would be implemented and think about cost of operations.

Cost Iterations



Cost Iterations

