

# Big O Runtime Analysis

## Part (a)

```
void f1(int n)
{
    for (int i = 2; i < n; i = i * i)
    {
        int i = 2;
        while (i < n) {
            /* do something that takes O(1) time */
            i = i * i;
        }
    }
}
```

i	K
2	0
4	1
16	2
256	3
⋮	⋮

$$\begin{aligned}i &= 2 \\i &= 2 \times 2 = 2^2 \\2^2 \times 2^2 &= 2^4 \\2^4 \cdot 2^4 &= 2^8 \\&\vdots\end{aligned}$$

$$2^{2^k} \geq n$$

When is  $2^{2^k} = n$ ?

$$\log_2(i) = \log_2 2^{2^k} = \log_2(n)$$

$$2^k = \log_2(n)$$

$$\log_2 2^k = \log(\log_2(n))$$

$$k \cdot \log_2 2 = \log(\log_2(n))$$

$$k = \log(\log_2(n))$$

$\Theta(\log(\log_2(n)))$  since no data dependency

## Part (b)



```
void f2(int n)
{
    for(int i=1; i <= n; i++){
        if( (i % (int)sqrt(n)) == 0){
            for(int k=0; k < pow(i,3); k++) {
                /* do something that takes O(1) time */
            }
        }
    }
}
```

Outer loop runs  $n$  times, inner runs  $\sqrt{n}$

$$\begin{aligned} & \sum_{i=1}^n \left( \theta(1) + \sum_{k=0}^{\sqrt{n}} i^3 \right) \\ &= \sum_{i=1}^n \theta(1) + \sum_{i=1}^n \sum_{k=0}^{\sqrt{n}} (k\sqrt{n})^3 \\ &= \theta(n) + (\sqrt{n})^3 \sum_{k=0}^{\sqrt{n}} k^3 \\ &= \theta(n) + (\sqrt{n})^3 \left( \frac{n + \sqrt{n}}{2} \right)^2 \\ &= \theta(n) + n^{\frac{3}{2}} \left( \frac{n^2 + 2\sqrt{n} + n}{4} \right) \\ &= \theta(n) + \frac{n^{\frac{5}{2}} + 2n^{\frac{3}{2}} + n^{\frac{5}{2}}}{4} \Rightarrow \theta(n^{5/2}) \end{aligned}$$

$i$	$i^3$
1	1
2	8
3	27
$\vdots$	$\vdots$

$i = k\sqrt{n}$

Part(c)

```

    for(int i=1; i <= n; i++){
        for(int k=1; k <= n; k++){
            if( A[k] == i){
                for(int m=1; m <= n; m=m+m){
                    // do something that takes O(1) time
                    // Assume the contents of the A[] array are not changed
                }
            }
        }
    }

```

$O(n^2)$   
 $m=2^n$   
 $\log_2 n = m$

$$\begin{aligned} T(n) &= \sum_{i=1}^n \sum_{k=1}^n (\theta(1) + O(\sum_{m=1}^n \theta(1))) \\ &= \theta(n^2) + n \sum_{m=1}^n O(1) \\ &= \theta(n^2) + O(n \cdot \log_2 n) \\ &= \theta(n^2) \end{aligned}$$

Part d



```
int f (int n)
{
    int *a = new int [10];
    int size = 10;
    for (int i = 0; i < n; i ++)
    {
        if (i == size)
        {
            int newsize = 3*size/2;
            int *b = new int [newsize];
            for (int j = 0; j < size; j ++) b[j] = a[j];
            delete [] a;
            a = b;
            size = newsize;
        }
        a[i] = i*i;
    }
}
```

$$i = 10 \times \frac{3}{2} s$$

$$n = 10 \times \frac{3}{2} s$$

$$s = \log_{\frac{3}{2}} \frac{n}{10}$$

$$\sum_{i=1}^n (\theta(1) + O(\sum_{j=1}^n \theta(1)))$$

$$= \sum_{i=0}^n \theta(1) + \sum_i 10 \times \frac{3}{2}^k \theta(1)$$

$$= \theta(n) + \sum_{k=0}^{\log_{\frac{3}{2}} \frac{n}{10}} \theta(10 \times \frac{3}{2}^k)$$

$$= \theta(n) + 10 \left(\frac{3}{2}\right)^{\log_{\frac{3}{2}} \frac{n}{10}} = \theta(n)$$