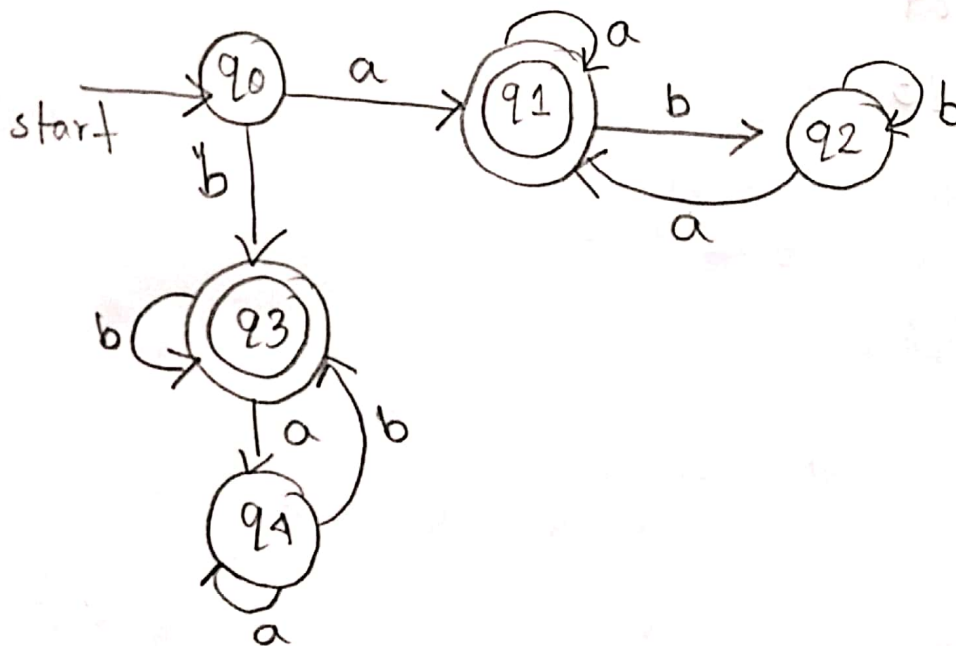


Q.1

a) The process of reducing a given DFA to its minimal form is called as minimization of DFA. It contains minimum number of states.

b) $L = \{ \text{strings that start and end with same symbol} \}$
 $= \{ aa, aba, abbaba, \dots, bb, baab, \dots \}$



$Q = \{ q_0, q_1, q_2, q_3, q_4 \}$

$\Sigma \rightarrow \{ a, b \}$

$q_0 \rightarrow \{ q_0 \}$

$F \rightarrow \{ q_1, q_3 \}$

6 - transition table.

pg-2

$q_0, a \rightarrow q_1$

$q_0, b \rightarrow q_3$

$q_1, a \rightarrow q_1$

$q_1, b \rightarrow q_2$

$q_2, a \rightarrow q_1$

$q_2, b \rightarrow q_2$

$q_3, a \rightarrow q_4$

$q_3, b \rightarrow q_3$

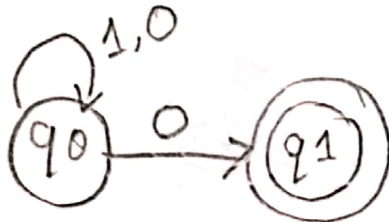
$q_4, a \rightarrow q_4$

$q_4, b \rightarrow q_3.$

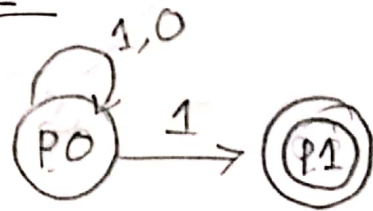
2

- c) R1 for even numbers (the last no. must be zero)
 R2 for odd numbers. (last no. must be 1)

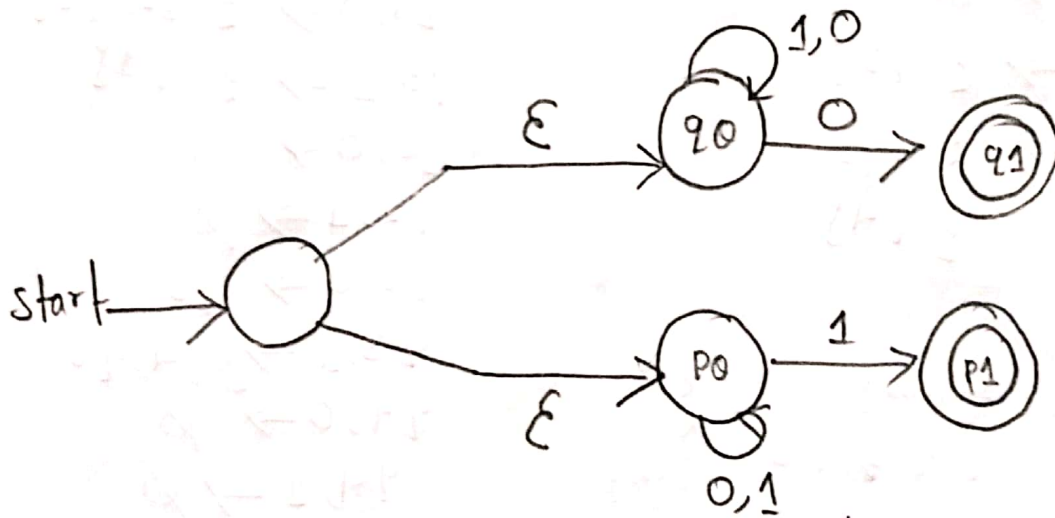
R1



R2



$$RE = R1 \mid R2$$

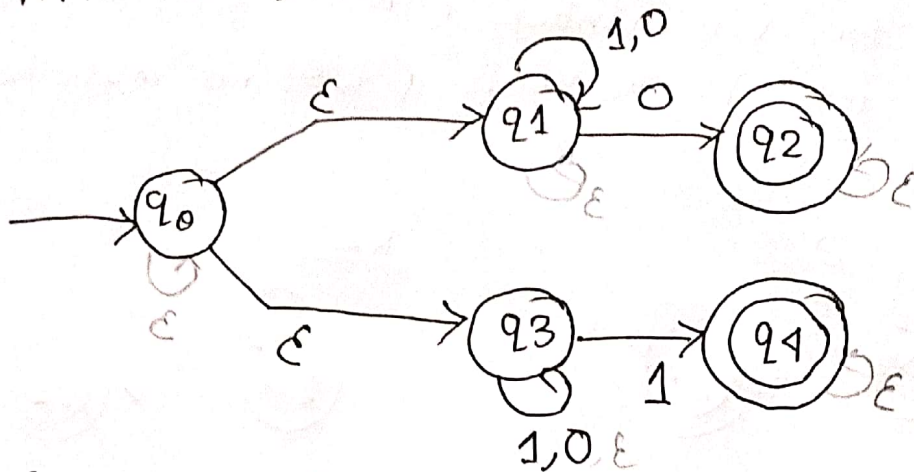


2. a) $RE = (0+1)^* 0$

b) $RE = (0+1)^+ 1$

d) NFA to DFA.

pg-4



first ϵ NFA \rightarrow NFA.

ϵ closure

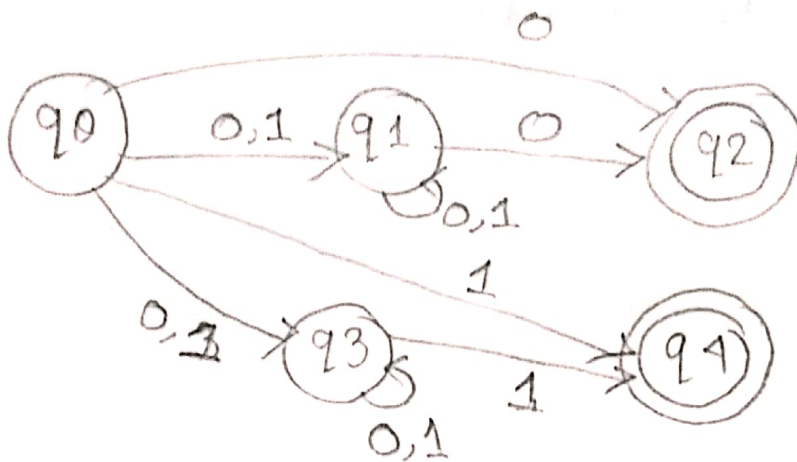
ϵ^+ of $q_0 \rightarrow \{q_1, q_3\}$
 ϵ^+ of $q_1 \rightarrow \{q_1\}$
 ϵ^+ of $q_2 \rightarrow \{q_2\}$
 ϵ^+ of $q_3 \rightarrow \{q_3\}$
 ϵ^+ of $q_4 \rightarrow \{q_4\}$

$q_0, 0 \rightarrow \{q_2, q_1, q_3\}$
 $q_0, 1 \rightarrow \{q_3, q_4, q_1\}$
 $q_1, 0 \rightarrow \{q_2, q_1\}$
 $q_1, 1 \rightarrow \{q_1\}$
 $q_2, 0 \rightarrow \emptyset$
 $q_2, 1 \rightarrow \emptyset$
 $q_3, 0 \rightarrow q_3$
 $q_3, 1 \rightarrow \{q_4, q_3\}$
 $q_4, 0 \rightarrow \emptyset$
 $q_4, 1 \rightarrow \emptyset$

states	ϵ^+	symbol	ϵ^*
$q_0, 0$	q_1	$\{q_2, q_1\}$	q_2, q_1
	q_3	$\{q_3\}$	q_3
$q_0, 1$	q_1	$\{q_1\}$	q_1
	q_3	$\{q_3, q_4\}$	q_3, q_4
$q_1, 0$	q_1	$\{q_2, q_1\}$	q_2, q_1
$q_1, 1$	q_1	$\{q_1\}$	q_1
$q_2, 0$	q_2	\emptyset	\emptyset
$q_2, 1$	q_2	\emptyset	\emptyset
$q_3, 0$	q_3	q_3	q_3
$q_3, 1$	q_3	q_4, q_3	q_4, q_3
$q_4, 0$	q_4	\emptyset	\emptyset
$q_4, 1$	q_4	\emptyset	\emptyset

→ NFA

pg-5

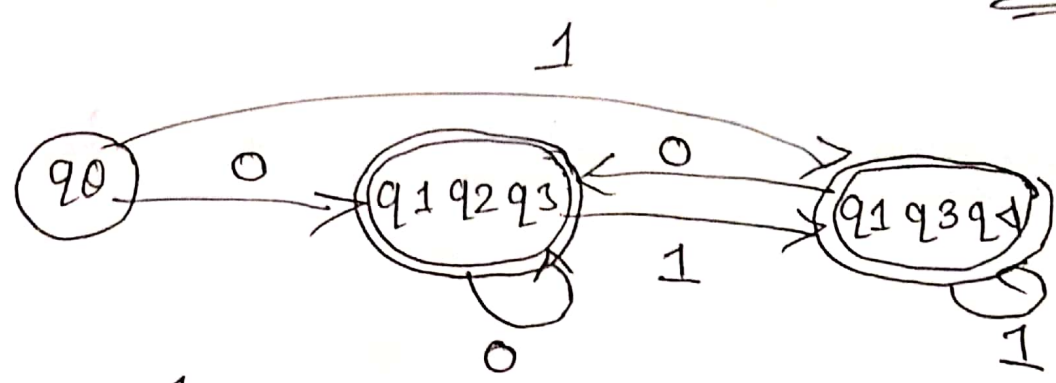


Transition table for NFA

states	0	1
q0	{q2, q1, q3}	{q1, q4, q3}
q1	{q1, q2}	{q1}
q2	{∅}	{∅}
q3	{q3, ∅}	{q3, q4}
q4	{∅}	{∅}

DFA

states	0	1
q0	[q1 q2 q3]	[q1 q3 q4]
[q1 q2 q3]	[q1, q2, q3]	[q1, q3 q4]
[q1, q3 q4]	[q1, q2 q3]	[q3 q4]



DFA ↗

2. e) Minimize -

	<u>final</u>	<u>non-final</u>
0 equivalent	$[q1q2q3, q1q3q4]$	$[q0]$
1 "	$[q1q2q3, q1q3q4]$	$[q0]$
2 <u> </u>		
"		

no minimization.

~~q4~~

Q3

a) A PDA can only access the top of its stack, whereas TM can access any position on an infinite tape. The infinite tape cannot be simulated with a single stack, so a PDA is less powerful. ~~There is another difference is the~~

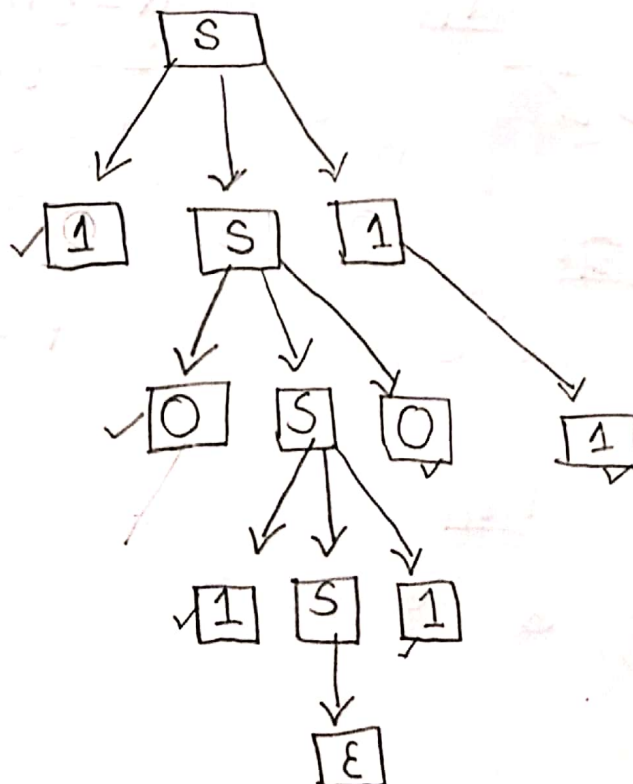
b) $S \rightarrow 0S0 \mid 1S1 \mid \epsilon$

for string 101101 :

$S \rightarrow 0S0$

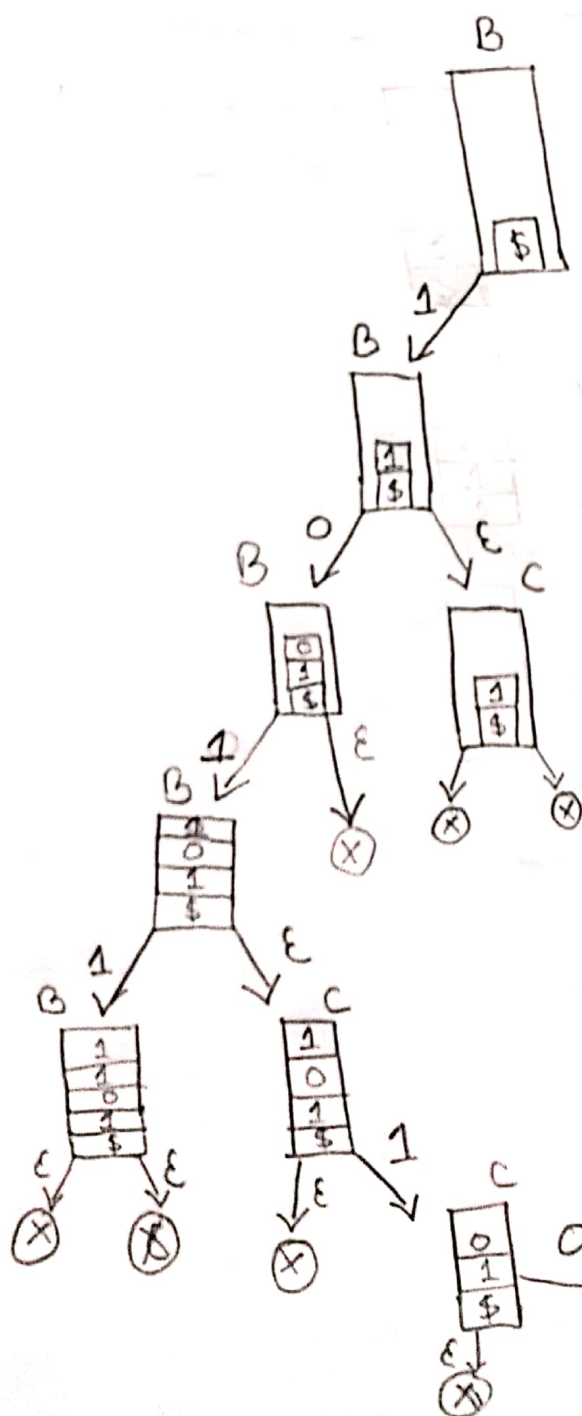
$S \rightarrow 1S1$

$S \rightarrow \epsilon$



9

for: $\varepsilon 1 \varepsilon 0 \varepsilon 1 \varepsilon 1 \varepsilon 0 \varepsilon 1$


$$Q = \{A, B, C, D\}$$
$$\Sigma = \{0, 1\}$$
$$r = \{ \frac{1}{2}, 0, 1 \}$$
$$90 = \angle A$$
$$t_0 = 2 \text{ sf}$$
$$F = \{A, D\}$$

Q.4

a) If the binary number is 100111, so to
1st start from the rightmost digit.

~~If the digit is 1 flip it unit found 0.~~

~~If the digit is 0, flip it make it 1~~
and move left and halt.

If the digit is 1, flip it and move to the
digit to the left.

b) for: 100111 \rightarrow 101000

