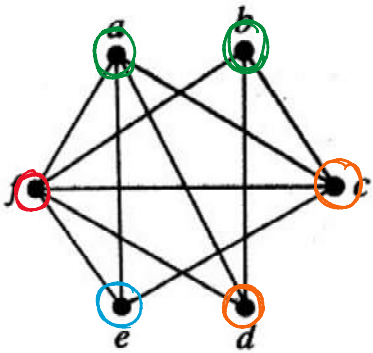


Welch Powell Algorithm:-

- (i) Order (Arrange) all the Vertices of graph G according to decreasing Order of their degrees.
- (ii) We assign the first Color C_1 to the vertex which has maximum degree. then in order assign C_1 to each vertex which is not adjacent to this vertex.
- (iii) Repeat step-2 with the second Color C_2 and the subsequence of non-colored Vertices.
- (iv) Repeat Step-3 with the 3rd Color C_3 and so-on untill all the Vertices are colored.



Find the Chromatic Number of the graph by using Welch Powell Algo.

Sol

$$\deg(a) = 4$$

$$\deg(f) = 5$$

$$\deg(b) = 3$$

$$\deg(c) = 4$$

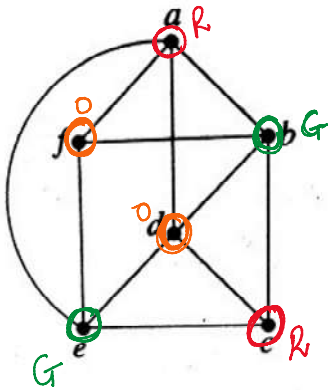
$$\deg(d) = 3$$

$$\deg(e) = 3$$

Degree	5	4	4	3	3	3
Vertex	f	a	c	b	d	e
Color	R	G	O	G	O	B

\therefore Chromatic Number = 4

$$\boxed{\chi(G) = 4}$$

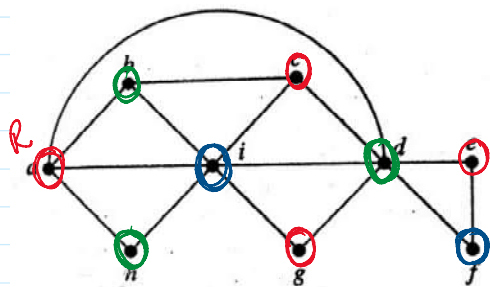


(a) 2

(b) \checkmark 3

(c) 4

(d) 5



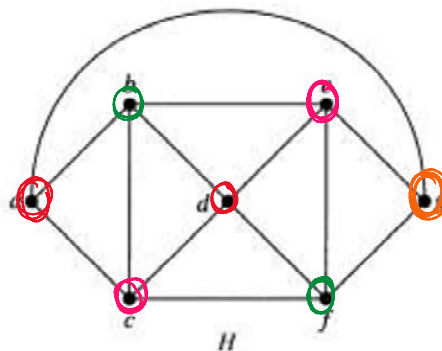
(a) — 2

(b) — 3 ✓

(c) — 4

(d) — 5

The chromatic number of the following graph is



(a) 6

(b) 2

✓
(c) 4

(d) 3

The minimum number of colours required to colour the vertices of a cycle with n nodes in such a way that no two adjacent nodes have the same colour is

✓
a) 2

✓
(b) 3

(c) 5

Ⓐ
(d) (a) or (b)

C_n $\begin{cases} 2 & \text{if } n \text{ is Even} \\ 3 & \text{if } n \text{ is odd} \end{cases}$

Shortest path Problems :→

Dijkstra's Algorithm :→

Let S be the source set

- (i) Initially there is no vertex in S
- (ii) Include a source vertex V_s in S . Find all the path from V_s to all other vertices without going through any other vertex
- (iii) Include that vertex in S which is nearest to V_s and find shortest path to all the vertices through this vertex and update the values
- (iv) Repeat the step (iii) until $(n-1)$ vertices are not included in S if there are n vertices in the graph

After completion of process we get the shortest path to all the vertices from the source vertex.

' from 'the source' vertex. "

'