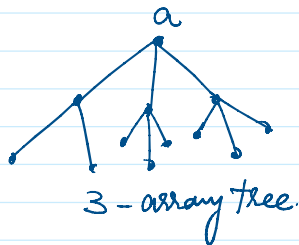
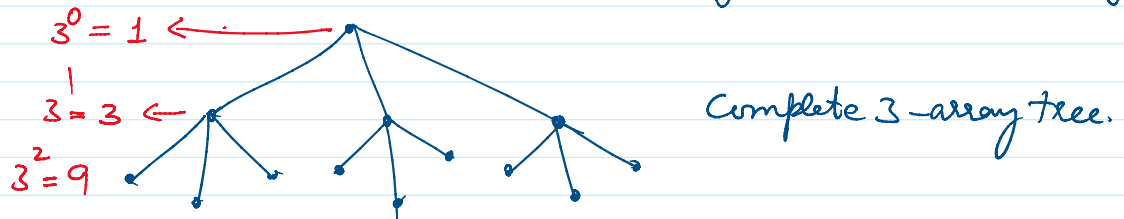


m-array Tree \rightarrow A tree is called m-array tree if every vertex of the tree has at most m-offsprings (childs)

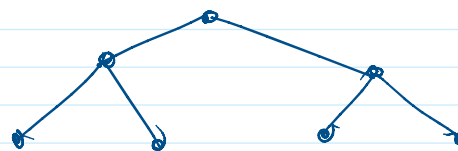
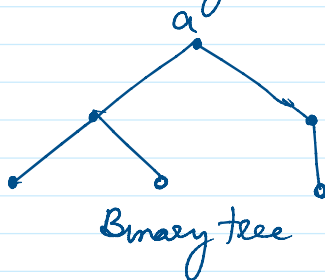


Complete m-array tree \rightarrow if Every Vertex has Exactly m-offspring

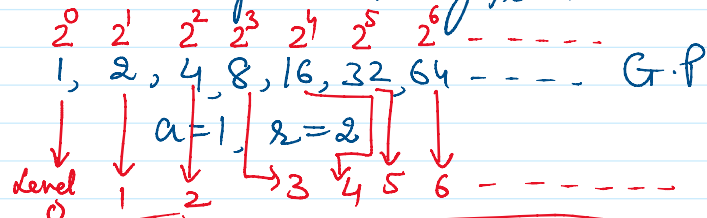


Binary-Tree \rightarrow 2-array tree

Complete Binary Tree \rightarrow Complete 2-array tree.



Complete Binary Tree.



Note: No. of Vertices in a Binary Tree at n^{th} level = 2^n

The Maximum of Vertices in a Binary Tree at n^{th} level = 2^n

(1) The Max. No. of Vertices/Node in a m-array tree at the n^{th} level = m^n

(2) No. of Nodes/Vertices in a Complete m-array tree at the n^{th} level = m^n

(Ex) ^{Max.} No. of nodes in a 4-array tree at 3rd level are m^n $4^3 = 64$

(1) (2) (3) (4) (5) (6) (7) (8) (9) (10) (11) (12) (13) (14) (15) (16) (17) (18) (19) (20) (21) (22) (23) (24) (25) (26) (27) (28) (29) (30) (31) (32) (33) (34) (35) (36) (37) (38) (39) (40) (41) (42) (43) (44) (45) (46) (47) (48) (49) (50) (51) (52) (53) (54) (55) (56) (57) (58) (59) (60) (61) (62) (63) (64) (65) (66) (67) (68) (69) (70) (71) (72) (73) (74) (75) (76) (77) (78) (79) (80) (81) (82) (83) (84) (85) (86) (87) (88) (89) (90) (91) (92) (93) (94) (95) (96) (97) (98) (99) (100)

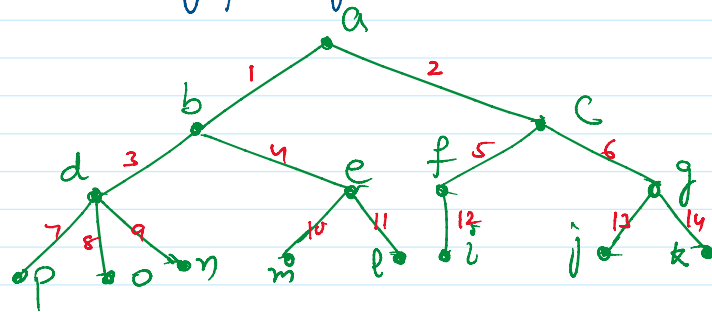
(Ex) No. of nodes in a 4-ary tree at 3rd level are $m^i \quad 4^3 = 64$
 (a) 16 (b) 64 (c) 128 (d) 4

① There is One and only One path between every pair of Vertices in a Tree.
Unique.

$a \rightarrow b \rightarrow d \rightarrow n$

$a \rightarrow c \rightarrow f \rightarrow i$

$d \rightarrow b \rightarrow a \rightarrow c \rightarrow g \rightarrow k$



If in a graph G , there is a Unique path b/w every pair of Vertices, G is a Tree.

② A Tree with n -Vertices has $(n-1)$ -edges.

$n = \text{No. of Vertices} = 15$
 $\text{No. of edges} = 14 = (n-1)$

No. of Edges = One less than no. of Vertices

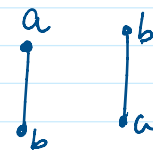
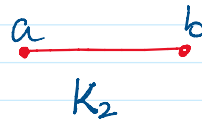
No. of Vertices = One more than no. of Edges

③ In a full/complete m -ary tree with i -Internal Vertices, then No. of Vertices is given by $n = mi + 1$

m - m-ary
 i = No. of Internal Vertices
 n = Total No. of Vertices

④ In any non-trivial tree, there are atleast two pendent Vertices.

Trivial \rightarrow No Edge
 (a) Trivial graph



\rightarrow degree one

- # Tree
- ① Connected
 - ② No Cycle
 - ③ No. of edges = $n-1$

① Tree

② Not Tree

③ Tree

