

Sliding Window

Algorithm

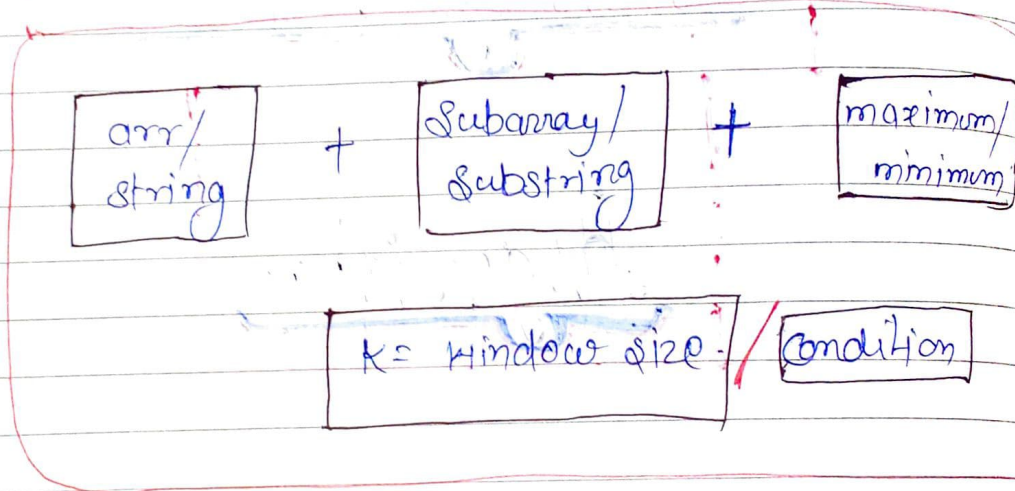
Sliding window technique is useful for solving problems in array or string, especially it is considered as a technique that could reduce the time complexity from $O(n^2)$ to $O(n)$.

Sliding Window Introduction

Identification And Types.

It's a continuous block of array which doesn't break.

Identification



Type

Fixed

Variable size.

Given: arr + window size



min/max (target)

Given: arr + target



min/max (window size)

2.

Sliding Window Problems

problems

Fixed

- (1.) max/min sub array of size k .
- (2.) 1st -ve in every min size of k .
- (3.) Count occurrence of anagram.
- (4.) Max. of all subarray of size k .
- (5.) max of min for every min size.

Variable

- (1.) largest/smallest subarray with sum k .
- (2.) largest sum-string with k distinct character.
- (3.) length of largest substring with no repetitive characters.
- (4.) Pick Toy.
- (5.) minimum window substring.

Q.

max. sum sub-array of size k

PS-IP-OP →

IP

size = 7

arr[] = [2, 5, 1, 8, 2, 9, 1]

Window size k = 3

PS

return max of
all window
sum

OP

19

[CODE] (Sliding-Window-Maximum)

int i = 0;

int me = INT_MIN;

int j = 0

int sum = 0;

while (j < size())
{

sum += arr[j];

→ Calculation

If window
size is less
than given
window
size

if (j - i + 1 < k)
j++;

When window
reached
its form

```
else if (j - i + 1 == k)
```

```
max = max(max, sum);
```

```
sum = sum - arr[i];
```

// excluding first
element

maintaining
the window
size

```
i++;  
j++;
```

```
}
```

```
return max;
```

// answer

```
}
```

Q1

First Negative Number in Every Window of Size k

Given an arr.

arr() = [12, -1, -7, 8, -15, 30, 16, 28]

k = 3

~~o/p~~

o/p = [-1, -1, -7, -15, -15, 0]

=====

CODE

int i = 0;
int j = 0;

← Initializing variables

resultant array →

vector<int> res;
deque<int> list;

← for storing useful element

while (j < arr.size())

{

← traversal

calculation →

if (arr[j] < 0)
list.push-back(arr[j]);

matching ←
subarray
with size k

```
if (j - i + 1 < k)
{
    j++;
    continue;
}
```

when it's
matched

→ else if (j - i + 1 == k)

```
{
    if (!list.empty())
        res.push_back(list.front());
    else
        res.push_back(0);
}
```

Calculating →
ans

```
if (arr[i] < 0)
    list.pop_front();
    i++;
}
j++;
}
```

Sliding →
Window

```
return res;
```

← returning ans

6.

Maximum of all Subarray.
of size K.

Given.

arr[] = [1, 2, 3, 1, 4, 5, 2, 3, 6]

K = 3

o/p → [3, 3, 4, 5, 5, 5, 6]

Sliding Window - Generic Structure.

```
int i = 0;
```

```
int j = 0;
```

```
while (j < arr.size())
```

```
{
```

```
    [calculation]
```

```
    if (j - i + 1 < K)
```

```
        j++;
```

```
    else if (j - i + 1 == K)
```

```
    {
```

```
        [ans]
```

```
        [slid the window]
```

```
    }
```


Running Code

```
int i = 0;
```

```
int j = 0;
```

```
vector<int> res;
```

```
deque<int> dq;
```

```
while (j < arr.size())
```

```
{
```

deleting

```
while (!dq.empty() && arr[j] >= arr[dq.back()])
```

```
    dq.pop_back();
```

```
    dq.push_back(j);
```

```
    if (j - i + 1 < k) {
```

```
        j++;  
        continue;
```

```
    }  
    if (j - i + 1 == k)
```

```
{
```

maximum element

```
    res.push_back(arr[dq.front()]);
```

```
    while (!dq.empty() && dq.front() <= j - k)
```

```
        dq.pop_front();
```

```
    i++; j++;
```

```
}
```

```
}
```

```
return res;
```

```
}
```

sliding window

(2nd variation)

Variable Size Sliding Window

Longest Subarray of sum k

problem statement:-

arr() = [4, 1, 1, 1, 2, 3, 5]

k = 5

sum = 5
(k)

o/p 4

CODE

```
int i = 0;
```

```
int j = 0;
```

```
int sum = 0;
```

```
int mx = INT_MIN;
```

```
while (j < arr.size())
```

```
{
```

```
    sum += arr[j];
```

→ calculation

1st case \leftarrow `if (sum < k)`
`j++;`

not case \leftarrow `else if (sum == k)`
`{`
`max = max(max, j - i + 1);` \rightarrow calculating ans
`j++;`
`}`

3rd case \leftarrow `else if (sum > k)`
`{`
`while (sum > k)` \rightarrow removing from front
`{`
`sum = sum - arr[i];`
`i++;`
`}`
`j++;`
`}`

`}`
`return max;` \rightarrow returning ans
`}`

Note \Rightarrow This approach will only work for positive integer array.
 \Rightarrow For negative and +ve integers use unordered_map

Q.

Longest Substring With K Unique Characters

PP: Q

s = aabacbebebe

K = 3

O/P = 7

CODE

```
unordered_map<char, int> mp;
```

```
int i = 0, j = 0, mx = 0;
```

```
while (j < arr.size())  
{
```

```
    mp[arr[j]]++;
```

```
    if (mp.size() < K)  
        j++;
```

```
    else if (mp.size() == K)  
    {  
        mx = max(mx, j - i + 1);  
        j++;  
    }
```

```
else if (mp.size() > K)
{
```

```
    while (mp.size() > K)
```

```
    {
```

```
        mp[arr[i]]--;
```

```
        if (mp[arr[i]] == 0)
```

```
            mp.erase(arr[i]);
```

```
        i++;
```

```
    }
```

```
    j++;
```

```
}
```

```
if (mx == 0 || arr.size() < K)
```

```
    return -1;
```

```
return mx;
```

```
}
```

Q.1

Longest Substring With Without Repeating Characters.

Given ↪

s = saturday

O/p ↪ 4

CODE

```
unordered_map<char, int> mp;  
int i = 0; j = 0; mx = INT_MIN;
```

```
while (j < s.length())  
{
```

```
    mp[s[j]]++;
```

```
    if (mp.size() == j - i + 1)
```

```
    {
```

```
        mx = max(mx, j - i + 1);
```

```
        j++;
```

```
    }
```



```
else if( mp.size() < j-i+1 )
```

```
{
```

```
    while( mp.size < j-i+1 )
```

```
    { mp[s(i)]--;
```

```
      if( mp[s(i)] == 0 )
```

```
        mp.erase(s(i));
```

```
      i++;
```

```
    }
```

```
    j++;
```

```
}
```

```
return m;
```

```
}
```