

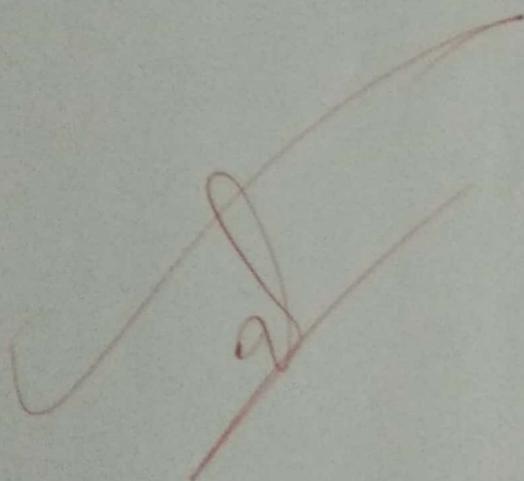
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Week 5

Given an array of non negative integers, design a linear algorithm & implement it using a program to search for element in array. Also find no of comparisons - T.C. = $O(n)$

Algorithm →

Take input array from user

Take element that is to be searched from user

Set flag = 0

Loop = ~~a~~ from i=0 to size-1

if no == array element

increment flag

After loop check if flag = 0 print "Not found"
else print "Element found"

STOP

Given an array (sorted) design an algorithm to search for a key in $O(\log n)$.

Algorithm →

Take Input array from user and key

~~Loop~~ Initialise l=0, r = size-1

Loop - while ($l \leq r$)

mid = $(l+r)/2$

if ($a[mid] == \text{key}$) then
return mid

else if ($a[mid] < \text{key}$) then

~~left~~ = mid + 1

else $r = mid - 1$

End Loop

return -1

> Given a sorted array of integers, search for a key using
JumpSort TC. $\leq O(n)$

Algorithm →

Start

$$c = \sqrt{\text{size}}$$

s~~t~~ = 0, e = ~~s~~ blocksize

while ($a[e]$ <= key) and $e < \text{size}$):

s~~t~~ = e ;

e = e + blocksize

if ($e > \text{size} - 1$) then

end = size

End Loop

FOR (i = s~~t~~ to e - 1) :

if ($a[i] == \text{key}$) then

return i

End Loop

return -1

STOP

Topic 2

Given sorted array, design an algorithm to find if an element is present. Also find its duplicates. $T.C = O(\log n)$

Algorithm -

use Binary Search to get index of first occurrence of key element present in arr[].

Let index = l.

use Binary Search to get index of last occurrence of key element occurrence of arr[].

Let index of last occurrence be j.

return (j - l + 1).

Given sorted array of positive integer, Find 3 indices i, j, k so that $a[i] + a[j] = a[k]$.

Algorithm -

Enter a sorted array

For (i = 0 to n, i++) :

Set k = 1 & for every j > 1,

then increment k until $a[k] >= a[i] + a[j]$

and increment count if equality is achieved.

Stop.

Given an array of non negative integers, design an algorithm to count number of pairs such that their difference equals k.

Algorithm -

Start

Enter key & #

Sort the array

i \leftarrow 0, j \leftarrow 0

while (i < size of array)

if (arr[i] - arr[j] == k)

then c++, i++, j++

else if arr[i] - arr[j] > k

then j++

else i++

return c

Stop

Ques 3
 Given an unsorted array. Design an algo to sort it using Insertion Sort. Also find Comparisons & Shifts.

Algorithm -

If its first element, return 1
 Pick next element

Compare with all elements in sorted sub array.

Shift all the elements in sorted sub array

for which value is greater than value to be sorted.

Insert value

Repeat until array is sorted

> Given an unsorted array, sort it using Selection Sort.
 Also find Comparisons & Shifts.

Algorithm -

selection sort(a, n)

Repeat step 2 & 3 for $i = 0$ to $n-1$

call smallest(a, i, n, pos)

swap $a[i]$ & $a[pos]$

End Loop

End

smallest(a, i, n, pos)

set small = $a[1]$

$pos = i$

Repeat for $j = i+1$ to n

if ($small > a[j]$) set $small = a[j]$

Set $pos = j$

End if

End Loop

return pos

Given unsorted array. Design an algorithm to find if array has duplicate elements or not. TC($O(n \log n)$) -

Algorithm →

Start

Input array α

Sort the array using Insert sort

$c \leftarrow 0$

for ($i = 1$ to n , $i++$)

if ($a[i-1] == a[i]$) print "Duplicate found"

return

print "No duplicate found"

Stop

WEEK 4

Given unsorted array. Design an algorithm & sort the array using merge sort. Also find no of comparisons & shifts.

Algorithm

Find middle index of array as:

$$\text{mid} = (\text{lb} + \text{ub}) / 2$$

Divide the array from middle by recursive call

Call merge sort for first half of array:

mergesort (a , lb , mid)

Call merge sort for second half of array:

mergesort (a , $\text{mid}+1$, ub)

Merge two sorted halves into single array

> Given an unsorted array, Sort it using Quicksort. Also find comparisons & shifts.

Algorithm

choose highest index as pivot

Take 2 variables to point first & last location.

~~left~~ value at left < pivot move right

~~right~~ value at right > pivot move left

If both step 5 or step 6 don't match, swap left & right

If $\text{left} \geq \text{right}$, the point where next is, is the new pivot.

Swap element so that pivot reaches its original location

Recursively call the function again

Given an unsorted array, find kth smallest element
(TC. = O(N))

Algorithm →

Start

Find middle index of the array

$$\text{mid} = (\text{lb} + \text{ub}) / 2$$

divide the array from middle

call merge sort for first half of array & then for second half

MergeSort (a, first, mid)

MergeSort (a, mid+1, last)

Merge the 2 sorted halves into single array

return element at index kth in sorted array.

Ques 5

Given unsorted array of alphabets. Find the alphabet with maximum occurrence $TC = O(N)$.

Algorithm

Start
Create function count(a , size, k)

count[$k+1$] = 0 & flag = 0

For ($i = 0$ to $n-1$; $i++$):

Increment count [$a[i] - 'a'$] by 1

End loop

For ($ch = 97$; $ch \leq 122$; $ch++$):

If count [$ch - 97$] ≥ 1 then
flag++

End Loop

Max = count[0]

For ($i = 0$ to $k-1$; $i++$):

if (count[i] $>$ max) then

max = count[i]

key = i

End if

End loop

If flag = 0 print "No Duplicate present"

Else print $\star (char)(key + 97)$, print max

Stop

> Given unsorted array of integer . Design an algorithm
of find 2 elements whose sum is more than given
element $Tc = O(n \log n)$.

Algorithm

Start

sort the array using ~~Quick Sort~~ Quick Sort

if $k \leq \text{size} - 1$, $j = 0$, $\text{flag} = 0$

while ($j < n/2$ and $k \geq n/2$):

if $(a[j] + a[k]) \geq \text{num}$)

then $k--$ ~~flag++~~

else $j++$;

if $(a[j] + a[k]) \geq \text{num}$)

then print ("a[j] + a[k]

flag++;

End Loop

If $\text{flag} \geq 0$ then

print ("No such pair exist")

Stop

Given 2 sorted integer arrays of size m & n. Design an algorithm to find element which are common to both.

$$T.C. = O(m+n)$$

Algorithm

Start
For ($i=0$ to $m-1$, $i++$) :

 if $b[i] = 0$

 For ($j=0$ to $n-1$; $j++$)

 if ($a[i] == b[j]$) then

 print ($a[i]$)

 End if

 End loop

End Loop

Stop

WEEK - 6

Given a graph, design an algorithm & implement it
using a program to find path, if exists, between 2
given nodes.

START

Input v

If $v \geq n$ go to step 7

Input temp

adj[i].push-back(temp)

If $i < v$, go to step 3

Input s, d

res = checkpath(adj, v, s-1, d-1)

res = checkpath(adj, v, s-1, d-1)

If ($res \geq 1$) print "path exists"
else print "Path does not exist"

STOP

Design an algorithm to find if a graph is Bipartite or not

Start

Input v

If $i \geq v$ go to step 7

input temp

input $g[i].push_back(temp)$

If $j < v$ go to step 3

res = Bipartite (G_i, v)

If ($res == 1$) print "Yes , Bipartite"

else print "No not Bipartite"

Stop

Given
whether

a directed graph, design an algorithm to find 47
cycle exists or not.

Start

> Input V

> Input u, v in adj[v]

> res = DFS (adj, v)

> if (res == 1) print "Cycle exists"

> else print "No cycle exists"

> Stop

WEEK - 7

> Apply Dijkstra's Algorithm to print shortest path & distance from given source to destination.

Start

> Input V

> If $v \geq 0$ go to step 8

> If $j \geq v$ go to step 7

> ~~Input~~ adj[i][j]

> If $j < v$ go to step 4

> If $s < v$ go to step 3

> Input S

> dijkstra adj, v, s-1

> Stop

6.1

Design an algorithm to find shortest path & distance
between 2 nodes using Bellmanford algorithm.

Start

Input m

If $i \geq m$ go to step 9

If $j \geq m$ go to step 7

Input graph[1][G]

If $j \leq m$ go to step 4

If $i \leq m$ go to step 3

Input s

findpath(graph, m, s=1)

Stop

Given a directed graph with two vertices, design an algorithm to find weight of shortest path from source to destination with exactly k edges on the path

- > Start
- > Input var
- > If $i \geq ver$ go to step 6
- > If $j \geq ver$ go to step 7
- > Input $graph[i][j]$
- > If $j < ver$ go to step 4
- > If $i < ver$ go to step 3
- > Input u, v, k
- > $ans = shortest_wt(graph, ver, u-1, v-1, k)$
- > Print ans
- > Stop

WEEK - 8

design an algorithm which will find maximum cost required
to connect cities (use Prim's Algo)

start

Input n

if $i \geq n$ go to step 8

if $j \geq n$ go to step 7

Input $a[i][j]$

if $j < n$ go to step 4

if $i < n$ go to step 3

prim's(a, n)

stop

2) Design an algorithm to find minimum cost required
to connect cities (Kruskal's algorithm)

> Start
> Input n
> If $i >= n$ go to step 8
> If $j >= n$ go to step 7
> Input graph[i][j]
> If $j < n$ go to Step 4
> If $i < n$ go to step 5
> kruskal(graph, n)
> Stop

Design an algorithm & implement it using program
for finding maximum budget required for a project

START

Input n

If $i \geq n$ go to step 8

If $j \geq n$ go to step 7

Input graph[][]

If $j < n$ go to step 4

If $i < n$ go to step 3

Kruskals(graph, n)

Stop

Week 9

Design an algorithm to implement Floyd-Warshall's all pair shortest path algorithm.

> Start

> Input n

> If $i >= n$ go to step 11

> If $j >= n$ go to step 10

> Input domp

> $y[domp] = "INF"$

> $graph[i][j] = \text{infinity}$

graph[i][j] = ∞

> If $j < n$ go to step 4

If $i < n$ go to step 3

If $k >= n$ go to step 19

> If $j >= n$ go to step 18

> If $j >= n$ go to step 17

> If $graph[i][k] + graph[k][j] < graph[i][j]$

graph[i][j] = $graph[i][k] + graph[k][j]$

> If $j < n$ go to step 13

> If $j < n$ go to step 12

> If $k < n$ go to step 11

> Print shortest path matrix

> If $i >= n$ go to step 26

> If $j >= n$ go to step 25

> If $(graph[i][j] >= \infty)$ print "INF"

> Else print $graph[i][j]$

> If $j < n$ go to step 21

> If $i < n$ go to step 20

26>Stop

14

> Given a knapsack of maximum capacity w . N items are provided each having its own value & weight.
> Design an algorithm to find list of selected items such that total selected list has weight w & has maximum value.

Start

> Input n

> if $i \geq n$ go to step 6

> Input items

> if $i < n$ go to step 3

> if $i > n$ go to step 11

> Input $val[i]$

> Job.push_back($val[i]$, item[i])

> double $(i+1)$

> if $i < n$ go to step 6

> Input k

> Sort (job.begin(), job.end())

> profit = 0

> if $j \geq n$ go to step 22

> if ($job[i][1] \geq k$)

> profit += $k * job[i][0]$

> Job.push_back(~~($k, job[i][2]$)~~)

> break

> else profit += $job[i][2] * job[i][0]$

> Job.push_back(~~($job[i][1], job[i][2]$)~~)

> $k = k - job[i][1]$

> print profit

> print "Item-weight"

> for every element, it . of LS

> print it . second - it . first

> Stop

Given an array of elements. Assume unit size
file i. Write an algorithm to merge all these files
into single file with minimum computation. For given 2
files A & B with sizes m & n, computation cost of merging
them is $O(m+n)$.

Start

Input m
If $i \geq n$ go to step 6

Input a[i]
If $i < n$ go to step 3

If $i \geq n$ go to step 10

minheap.push(a[i])

If $i < n$ go to step 7

ans = 0

while (minheap.size() > 1)

i) e1 = minheap.top(), minheap.pop()

ii) e2 = minheap.top(), minheap.pop()

iii) ans += e1 + e2

iv) minheap.push(e1 + e2)

print ans

Stop

WEEK 10

Given a list of activities with their starting time & finishing time. Select maximum no of activities that can be performed by single person such that selected activities must be non conflicting.

Step 1

Input n

$i, s[n], f[n]$

If $i > n$ go to step 7

Input s[i]

If $i < n$ go to step 4

If $i = 0$

If $i >= n$ go to Step 11

Input f[i]

If $i < n$ go to step 11

> vector<vector<int>> a, vector<int> act

> If $i >= n$ go to step 14

> a.push_back({f[i], {s[i], i+1}})

> sort(a.begin(), a.end())

> e = INT-MIN, c = 0

> If $i >= n$, go to step 13

> If $(a[i][1] >= e)$

> e = a[i][0]

(i) $c++$

(ii) act.pushback(a[i][2])

> If $i < n$ go to step 16

> Print "No of non conflicting", c

> Print "No of activities selected"

> If $i >= n$ go to step 17

else act[i]

> Given a long list of tasks. Each task takes specific time to accomplish it & each task has a deadline associated with it.
 +. Design an algorithm to find maximum no of tasks that can be completed without crossing their deadlines &
 b. find list of selected tasks.

Start

Input n

i, t[n], f[n]

if $i \geq n$ go to step 7

Input t[i]

if $i < n$ go to step 4

$i = 0$

> if $i > n$ go to step 11

> Input f[i]

> if $i < n$ go to step 8

> vector<vector<int>> a, vector<int> act

> if $i \geq n$ go to step 14

> a.push_back({y[i], y[i] - t[i], i + 1})

> sort(a.begin(), a.end())

> c = INT_MIN, e = 0

> if $i \geq n$ go to step 12

7) if $a[i][1] \geq e$
 $e = a[i][0]$

8) $e = a[i][0]$

9) $e++$

10) act.push_back(a[i][0])

11) if $i < n$ go to step 16

12) sort(act.begin(), act.end())

13) Print c

14) Print "Selected Task Numbers"

15) if $i \geq n$ go to step 28

16) Print act[i]

17) if $i < n$ go to step 25

28) Stop

Given an unsorted array, design an algorithm to find the majority element (one which occurs more than $\frac{n}{2}$ times).

Start

Input n

$i, a[n], i, j$

$i > n$ go to step 6

Input $a[i]$

$j = 0$

$\text{int}(a, a + n)$

$i > n$ go to step 14

$c = 1$

$j++$
while($j < n$ $\&$ $a[j + 1] == a[i]$)

$c++$
 $c > n/2$ Print "Yes"

(d) $j = 1$
(e) break

$i = j - 1$

$j = 0$ Print "No"

$j = n/2$ Print $a[n/2]$

$n/2$ Print $a[n/2] + a[n/2 - 1]/2$

STOP

WEEK 11

Given a sequence of matrices. Write an algorithm to find
an efficient way to multiply them.

Start

Input m

$b[n+1]$

If $i >= n$ go to step 6

Input $b[i]$, $b[i+1]$

Print . matching order ($p, n+1$)

Stop

match chain order (p, n)

$m[n][n], i, j, k, l, v;$

for (iz 1 to n ; $i++$)
 $m[1][i] = 0$

for ($j = 2$; $j < n$; $j++$)

{ for ($j = 1$; $j < n-l+1$; $j++$)

{ $j = i+l-1$;

$m[i][j] = INT_MAX;$

for ($k = i$; $k \leq j-1$; $k++$)

{ $v = m[i][k] + m[k+1][j] + b[i-1] * b[k]^* b[j]$

if ($v < m[i][j]$)

$m[i][j] = v$

3

3

return $m[1][n-1];$

a set of available types of coins. Let suppose you have infinite supply of each type of coin. For a given ~~set~~ N , you have to design an algorithm to find no of ways in which these coins can be added to make sum equals to N .

Start

Input n

$i, j, a[n]$

if $i >= n$ go to step 7

Input $a[i]$

Input amt

$i = 1$

if $i >= amt$ go to step 11

$ans[i] = 0$

$ans[0] = 1$

if $j >= n$ go to step 14

if $i >= amt$ go to step 11

if $(a[j] \leq i)$

$ans[i] += ans[i - a[j]]$

Print $ans[amt]$

Stop

Given a set of problem elements, you have to partition the set into 2 subsets such that the sum of the elements in both is same. Design an algorithm for it.

Start

Input n

$i, j, a[n]$

If $i > n$ go to Step 6

Input $a[i]$

Sum = 0, $i = 0$

If $i > n$ go to Step 8

If ($sum \% 2 != 0$)

Print "No", return 0

Sum = Sum / 2

bool s[n+1][sum+1]

If $i > n$ go to Step 16

If $j > sum$ go to Step 11

If ($j = 0$) $s[i][j] = 1$

else If ($i = 0$) $s[i][j] = 0$

else

(i) If ($a[i-1] > j$)

$s[i][j] = s[i-1][j]$

(ii) else $s[i][j] = s[i-1][j] || s[i-1][j-a[i-1]]$

If ($s[n][sum]$) Print "Yes"

else Print "No"

> Stop