

TUTORIAL 2

Q1.) Find time complexity:

```
void fun(int n)
{
    int j=1, i=0;
    while (i<n)
    {
        i = i+j;
        j++;
    }
}
```

Soln. >

i	j
0	1
1	2
1+2	3
1+2+3	4

$\Rightarrow$  ~~k~~th Term is of form  $= \frac{k(k+1)}{2}$

$$\frac{k(k+1)}{2} > n \quad (\text{For loop end})$$

$$\Rightarrow k^2 > n$$

$$k > \sqrt{n} = O(\sqrt{n})$$

Q2.) Write recurrence relation & function to print Fibonacci Series.  
Solve & get complexity. Also find space complexity.

Soln

```
int fibo(int n)
{
    if (n <= 1) return n;
    return fibo(n-1) + fibo(n-2);
}
```

$$\Rightarrow T(n) = T(n-1) + T(n-2) \quad \text{2 ~~times~~ } , T(0) = T(1) = 1$$

let  $T(n-1) = T(n-2)$  lower bound

$$T(n) = 2 + 2(T(n-2))$$

$$= 2^2 (T(n-2 \times 2))$$

$$= 2^3 (T(n-2 \times 3))$$

$$= 2^k (T(n-2k))$$

$$n - 2k = 0 \Rightarrow k = \frac{n}{2}$$

$$= 2^{n/2} T(1) = 2^{n/2} = O(2^n)$$

$$M \quad T(n-2) \approx T(n-1) \quad (\text{Upper bound})$$

$$T(n) = 2T(n-1)$$

$$= 2(2T(n-2))$$

$$= 2^k T(n-k)$$

$$\text{Let } n-k=0 \Rightarrow k=n$$

$$\Rightarrow T(n) = 2^n T(0)$$

$$= 2^n = O(2^n)$$

Space complexity -  $O(n)$

Q3) Write programs with complexity:

(a)  $n \log n$

```
for (int i=0; i<n; i++)
```

```
for (int j=0; j<n; j*=2)
```

```
cout << "Hi";
```

(b)  $n^3$

```
for (int i=0; i<n; i++)
```

```
for (int j=0; j<n; j++)
```

```
for (int k=0; k<n; k++)
```

```
cout << "Hi";
```

(c)  $\log(\log n)$

```
for (int i=1; i<=n; i*=2)
```

```
for (int j=1; j<=n; j*=2)
```

```
cout << "Hi";
```



Q4 Solve:

$$T(n) = T\left(\frac{n}{4}\right) + T\left(\frac{n}{2}\right) + cn^2$$

$$\text{Let } T\left(\frac{n}{2}\right) \geq T\left(\frac{n}{4}\right)$$

$$\text{So, } T(n) = 2T\left(\frac{n}{2}\right) + cn^2$$

Applying master's theorem,

$$T(n) \quad a=2 \quad b=2 \quad f(n) = cn^2$$

$$n^{\log_b a} = n^{\log_2 2} = n$$

$$\infty \quad n < n^2$$

$$\Rightarrow T(n) = O(n^2)$$

Q5 Find time complexity:

int fun(int n)

{ for (i=1 to n, i++)

for (j=1 to n, j+=i)

{ }

}

$$i=1 \quad j=1, 2, 3, 4 \dots$$

$$i=2 \quad j=1, 3, 5 \dots$$

$$i=3 \quad j=1, 4, 7 \dots n$$

$$i=4$$

$$i=n$$

$$n = \frac{n(n+1)}{2}$$

$$n = \frac{n+1}{2}$$

$$= \frac{n+2}{3}$$

$$= \frac{n}{4} \text{ times}$$

$$= \frac{n}{n} \text{ times}$$

Hence total time complexity

$$= n \left( \frac{n}{1} + \frac{n}{2} + \frac{n}{3} + \dots + \frac{n}{n} \right)$$

$$= n \left( 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right)$$

$$= n \log n = O(n \log n)$$

Q6) for (int i=2; i <= n; i = pow(i, k))

{

}

Find T-Complexity.

$$\downarrow = 2, 2^k, (2^k)^k, ((2^k)^k)^k \dots$$

$$= 2, 2^k, 2^{k^2}, 2^{k^3} \dots$$

Loop will end if  $i > n$

$$\Rightarrow 2^{k^m} > n$$

$$k^m \log_2 2 > \log_2 n$$

$$k^m > \log_2 n$$

$$m \log k > \log(\log n)$$

$$m > \frac{\log(\log n)}{\log(k)} \rightarrow \text{const.}$$

$$T(n) = O(\log(\log n))$$

Q 7 → Write a recurrence relation when Quick sort repeatedly divides the array into 2 parts of 99% & 1%, Derive T.C., Show recursion tree & find the difference in heights of both extreme parts. What do you understand by this analysis?

Soln Dividing 99% to 1% is worst case of QuickSort

Recurrence relation →

$$T(n) = T(0) + T(n-1) + O(n)$$

$$T(n) = T(n-1) + O(n)$$

$$\text{let } O(n) = cn$$

Applying ~~Master~~ Backward substitution,

$$T(n) = T(n-1) + cn$$

$$T(n) = T(n-2) + c(n-1) + cn$$

$$T(n) = T(n-3) + c(n-2) + c(n-1) + cn$$

$$= T(n-3) + 3cn - c(2+1+0)$$

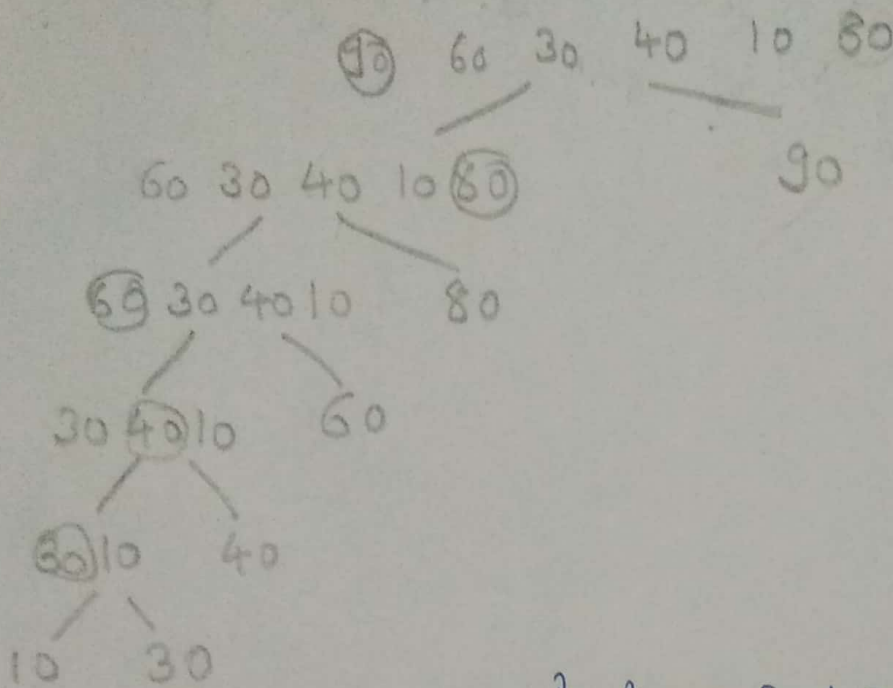
$$T(n) = T(n-k) + kn - c(k-1)$$



Put  $k=1$

$$T(n) = T(1) + k(n-1)n - c(n-2) = kn^2 - kn - cn + C_2$$

$$= O(n^2)$$



Difference in height =  $5 - 1 = 4$

We understand that for worst case QuickSort has complexity of  $O(n^2)$ .

Q8) Arrange in increasing order :

a)  $n, n!, \log n, \log(\log n), \sqrt{n}, \log(n!), n \log n, (\log n)^2, 2^n, 2^{2n}$   
 $4^n, n^{1/2}, 100, (\log n)^2$   
 $100 < \log(\log n) < \log n < (\log n)^2 < \sqrt{n} < n < n \log n < n^2 < 2^n$   
 $< n! < 4^n = 2^{2n}$

b)  $2(2^n), 4n, 2n, 1, \log n, \log(\log n), \sqrt{\log(n)}, \log 2n, 2 \log n, n$   
 $, \log(n!), n!, n^2, n \log n$   
 $1 < \log(\log n) < \sqrt{\log n} < \log n < \log 2n < 2 \log n < n \log n < n^2 < 2 \cdot 2^n < n!$

(C)  $8^{2n}$ ,  $\log_2 n$ ,  $n \log_8 n$ ,  $n \log_2 n$ ,  $\log n!$ ,  $n!$ ,  $\log_8(n)$ , 96  
 $8n^2$ ,  $7n^3$ ,  $5n$

Soln 96,  $\log_8 n$ ,  $\log_2 n$ ,  $n \log_6 n$ ,  $n \log_{10} n$ ,  ~~$n \log_2 n$~~ ,  $\log(n!)$   
 $5n$ ,  $8n^2$ ,  $7n^3$ ,  ~~$n!$~~ ,  $8^{2n}$ .