

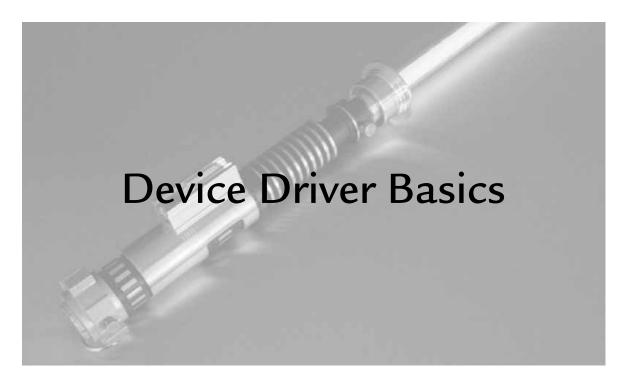
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Embedded I have

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Chapter 8



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One of the more challenging aspects of system design is partitioning functionality in a rational manner. The familiar device driver model found in UNIX and Linux provides a natural partitioning of functionality between your application code and hardware or kernel devices. In this chapter, we develop an understanding of this model and the basics of Linux device driver architecture. After reading this chapter, you will have a solid foundation for continuing your study of device drivers using one of the excellent texts listed at the end of this chapter.

This chapter begins by presenting Linux device driver concepts and the build system for drivers within the kernel source tree. We examine the Linux device driver architecture and present a simple working example driver. We introduce the user space utilities for loading and unloading kernel modules. We present a simple application to illustrate the interface between applications and device drivers. We conclude this chapter with a discussion of device drivers and the GNU Public License.

8.1 Device Driver Concepts

Many experienced embedded developers struggle at first with the concepts of device drivers in a virtual memory operating system. This is because many popular legacy real-time operating systems do not have a similar architecture. The introduction of virtual memory and kernel space versus user space frequently introduces complexity that is not familiar to experienced embedded developers.

One of the fundamental purposes of a device driver is to isolate the user's programs from ready access to critical kernel data structures and hardware devices. Furthermore, a well-written device driver hides the complexity and variability of the hardware device from the user. For example, a program that wants to write data to the hard disk need not care if the disk drive uses 512-byte or 1024-byte sectors. The user simply opens a file and issues a write command. The device driver handles the details and isolates the user from the complexities and perils of hardware device programming. The device driver provides a consistent user interface to a large variety of hardware devices. It provides the basis for the familiar UNIX/Linux convention that everything must be represented as a file.

¹ The terms *module* and *device driver* are used here interchangeably.

8.1.1 Loadable Modules

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Loadable mod le ha e pa ic la impo ance o embedded em . Loadable mod le enhance field pg ade capabili ie ; he mod le i elf can be pda ed in a li e em i ho he need fo a eboo . Mod le can be o ed on media o he han he oo (boo) de ice, hich can be pace con ained.

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Loadable mod le a e in alled af e he ke nel ha boo ed. S a p c ip can load de ice d i e mod le, and mod le can al o be demand loaded hen needed. The ke nel ha he capabili o e e a mod le hen a e ice i e e ed ha e i e a pa ic la mod le.

Te minolog ha ne e been anda di ed hen di c ing ke nel mod le . Man e m ha e been and con in e o be ed in e changeabl hen di c ing loadable ke nel mod le . Th o gho hi and la e chap e , he e m *device driver, loadable kernel module (LKM), loadable module,* and *module* a e all ed o de c ibe a loadable ke nel de ice d i e mod le.

8.1.2 Device Driver Architecture

The basic Linux device driver model is familiar to UNIX/Linux system developers. Although the device driver model continues to evolve, some fundamental constructs have remained nearly constant over the course of UNIX/Linux evolution. Device drivers are broadly classified into two basic categories: *character devices* and *block devices*. Character devices can be thought of as serial streams of sequential data. Examples of character devices include serial ports and keyboards. Block devices are characterized by the capability to read and write blocks of data to and from random locations on an addressable medium. Examples of block devices include hard drives and floppy disk drives.

8.1.3 Minimal Device Driver Example

Because Linux supports loadable device drivers, it is relatively easy to demonstrate a simple device driver skeleton. Listing 8-1 illustrates a loadable device driver module that contains the bare minimum structure to be loaded and unloaded by a running kernel.

Listing 8-1

Minimal Device Driver

```
/* Example Minimal Character Device Driver */
#include linux/module.h>

static int __init hello_init(void)
{
    printk("Hello Example Init\n");

    return 0;
}

static void __exit hello_exit(void)
{
    printk("Hello Example Exit\n");
}

module_init(hello_init);
module_exit(hello_exit);

MODULE_AUTHOR("Chris Hallinan");
MODULE_DESCRIPTION("Hello World Example");
MODULE_LICENSE("GPL");
```

The skeletal driver in Listing 8-1 contains enough structure for the kernel to load and unload the driver, and to invoke the initialization and exit routines. Let's look at how this is done because it illustrates some important high-level concepts that are useful for device driver development.

A device driver is a special kind of binary module. Unlike a stand-alone binary executable application, a device driver cannot be simply executed from a command prompt. The 2.6 kernel series requires that the binary be in a special "kernel object" format. When properly built, the device driver binary module contains a .ko suffix. The build steps and compiler options required to create the .ko module object can be quite complex. Here we outline a set of steps to harness the power of the Linux kernel build system without requiring you to become an expert in it, which is beyond the scope of this book.

8.1.4 Module Build Infrastructure

A device driver must be compiled against the kernel on which it will execute. Although it is possible to load and execute kernel modules built against a different kernel version, it is risky to do so unless you are certain that the module does not rely on any features of your new kernel. The easiest way to do this is to build the module within the kernel's own source tree. This ensures that as the developer changes the kernel configuration, his custom driver is automatically rebuilt with the correct kernel configuration. It is certainly possible to build your drivers outside of the kernel source tree. However, in this case, you are responsible for making sure that your device driver build configuration stays in sync with the kernel you want to run your driver on. This typically includes compiler switches, location of kernel header files, and kernel configuration options.

For the example driver introduced in Listing 8-1, the following changes were made to the stock Linux kernel source tree to enable building this example driver. We explain each step in detail.

- 1. Starting from the top-level Linux source directory, create a directory under .../drivers/char called examples.
- 2. Add a menu item to the kernel configuration to enable building examples and to specify built-in or loadable kernel module.
- 3. Add the new examples subdirectory to the .../drivers/char/Makefile conditional on the menu item created in step 2.

- 4. C ea e a makefile f he ne examples di ec , and add he hellolo m d le bjec be c m iled c ndi i nal n he men i em c ea ed in e 2.
- 5. Finall, c ea e he die hellol.c ce file f m Li ing 8.1.

Adding he examples diec nde he .../drivers/char bdiec i elf-e lana . Af e hi diec i cea ed, file a e cea ed in hi diec : he m d le ce file i elf f m Li ing 8-1 and he makefile f he examples diec . The makefile f examples i i e i ial. I ill c n ain hi ingle line:

Adding he men i em he ke nel c nfig a i n ili i a li le m e in l ed. Li ing 8-2 c n ain a a ch ha, hen a lied he .../drivers/char/Kconfig file f m a ecen Lin elea e, add he c nfig a i n men i em enable examples c nfig a i n i n. F h e eade n familia, i h he diff/patch f ma, each line in Li ing 8-1 eceded b a ingle l (+) cha ac e i in e ed in he file be, een he indica ed line (h e i h he leading + cha ac e).

Listing 8-2

Kconfig Patch for Examples

```
diff -u ~/base/linux-2.6.14/drivers/char/Kconfig
./drivers/char/Kconfig
--- ~/base/linux-2.6.14/drivers/char/Kconfig
+++ ./drivers/char/Kconfig
@@ -4,6 +4,12 @@

menu "Character devices"

+config EXAMPLES
+ tristate "Enable Examples"
+ default M
+ ---help---
+ Enable compilation option for driver examples
+
config VT
    bool "Virtual terminal" if EMBEDDED
    select INPUT
```

When a lied Koonfig in he .../drivers/char bdi ec f a ecen Lin ke nel, hi a ch e l in a ne ke nel c nfig a i n i n called CONFIG_EXAMPLES. As a reminder from our discussion on building the Linux kernel in Chapter 4, "The Linux Kernel—A Different Perspective," the configuration utility is invoked as follows (this example assumes the ARM architecture):

\$ make ARCH=ARM CROSS_COMPILE=xscale_be- gconfig

After the configuration utility is invoked using a command similar to the previous one, our new Enable Examples configuration option appears under the Character devices menu, as indicated in the patch. Because it is defined as type tristate, the kernel developer can choose from three choices:

- (N) No. Do not compile examples.
- (Y) Yes. Compile examples and link with final kernel image.
- (M) Module. Compile examples as dynamically loadable module.

Figure 8-1 shows the resulting <code>gconfig</code> screen with the new configuration option added. The dash (-) in the check box selects (M)odule, as indicated in the M column on the right. A check mark in the check box selects (Y)es, indicating that the driver module should be compiled as part of the kernel proper. An empty check box indicates that the option is not selected.

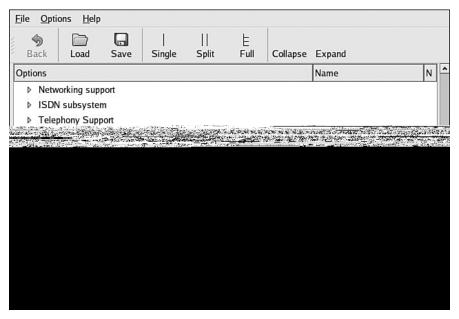


FIGURE 8-1
Kernel configuration with Examples module

Now that we have added the configuration option to enable compiling our examples device driver module, we need to modify the makefile in .../drivers/char to instruct the build system to descend into our new examples subdirectory if the configuration option CONFIG_EXAMPLES is present in our configuration. Listing 8-3 contains the patch for this against the makefile in a recent Linux release.

Listing 8-3

Makefile Patch for Examples

```
diff -u ~/base/linux-2.6.14/drivers/char/Makefile
./drivers/char/Makefile
--- ~/base/linux-2.6.14/drivers/char/Makefile
+++ ./drivers/char/Makefile
@@ -88,6 +88,7 @@
obj-$(CONFIG_DRM) += drm/
obj-$(CONFIG_PCMCIA) += pcmcia/
obj-$(CONFIG_IPMI_HANDLER) += ipmi/
+obj-$(CONFIG_EXAMPLES) += examples/
obj-$(CONFIG_HANGCHECK_TIMER) += hangcheck-timer.o
```

The patch in Listing 8-3 adds the single line (preceded by the + character) to the makefile found in .../drivers/char. The additional lines of context are there so that the patch utility can determine where to insert the new line. Our new examples directory was added to the end of the list of directories already being searched in this makefile, which seemed like a logical place to put it. Other than for consistency and readability, the location is irrelevant.

Having completed the steps in this section, the infrastructure is now in place to build the example device driver. The beauty of this approach is that the driver is built automatically whenever a kernel build is invoked. As long as the configuration option defined in Listing 8-3 is selected (either M or Y), the driver module is included in the build.

Building for an arbitrary ARM system, the command line for building modules might look like this:

```
$ make ARCH=arm CROSS_COMPILE=xscale_be- modules
```

Listing 8-4 shows the build output after a typical editing session on the module (all other modules have already been built in this kernel source tree.)

Listing 8-4

Module Build Output

```
$ make ARCH=arm CROSS_COMPILE=xscale_be- modules
   CHK    include/linux/version.h
make[1]: 'arch/arm/kernel/asm-offsets.s' is up to date.
make[1]: 'include/asm-arm/mach-types.h' is up to date.
   CC [M] drivers/char/examples/hello1.o
   Building modules, stage 2.
   MODPOST
   LD [M] drivers/char/examples/hello1.ko
```

8.1.5 Installing Your Device Driver

Now that this driver is built, we can load and unload it on a running kernel to observe its behavior. Before we can load the module, we need to copy it to an appropriate location on our target system. Although we could put it anywhere we want, a convention is in place for kernel modules and where they are populated on a running Linux system. As with module compilation, it is easiest to let the kernel build system do that for us. The makefile target modules_install automatically places modules in the system in a logical layout. You simply need to supply the desired location as a prefix to the default path.

In a standard Linux workstation installation, you might already know that the device driver modules live in /lib/modules/<kernel-version>/... ordered in a manner similar to the device driver directory hierarchy in the Linux kernel tree.² The <kernel-version> string is produced by executing the command uname -r on your target Linux system. If you do not provide an installation prefix to the kernel build system, by default, your modules are installed in your own workstation's /lib/modules/... directory. This is probably not what you had intended. You can point to a temporary location in your home directory and manually copy the modules to your target's file system. Alternatively, if your target embedded system uses NFS root mount to a directory on your local development

² This path is used by Red Hat and Fedora distributions, and is also required by the File System Hierarchy Standard referenced at the end of this chapter. Other distributions might use different locations in the file system for kernel modules.

workstation, you can install the modules directly to the target file system. The following example assumes the latter.

```
$ make ARCH=arm CROSS_COMPILE=xscale_be-
INSTALL_MOD_PATH=/home/chris/sandbox/coyote-target
modules_install
```

This places all your modules in the directory coyote-target, which on this example system is exported via NFS and mounted as root on the target system.³

8.1.6 Loading Your Module

Having completed all the steps necessary, we are now in a position to load and test the device driver module. Listing 8-5 shows the output resulting from loading and subsequently unloading the device driver on the embedded system.

Listing 8-5

Loading and Unloading a Module

You should be able to correlate the output with our device driver source code found in Listing 8-1. The module does no work other than printing messages to the kernel log system via printk(), which we see on our console. When the module is loaded, the module-initialization function is called. We specify the initialization function that will be executed on module insertion using the module_init() macro. We declared it as follows:

```
module_init(hello_init);
```

In our initialization function, we simply print the obligatory hello message and return. In a real device driver, this is where you would perform any initial resource allocation and hardware device initialization. In a similar fashion, when we unload the module (using the modprobe -r command), our module exit routine is called.

³ Hosting a target board and NFS root mount are covered in detail in Chapter 12, "Embedded Development Environment".

⁴ If you don't see the messages on the console, either disable your syslogd logger or lower the console loglevel. We describe how to do this in Chapter 14, "Kernel Debugging Techniques".

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As shown in Listing 8-1, the exit routine is specified using the module_exit() macro.

That's all there is to a skeletal device driver capable of live insertion in an actual kernel. In the sections to follow, we introduce additional functionality to our loadable device driver module that illustrates how a user space program would interact with a device driver module.

8.2 Module Utilities

We had a brief introduction to module utilities in Listing 8-5. There we used the module utility modprobe to insert and remove a device driver module from a Linux kernel. A number of small utilities are used to manage device driver modules. This section introduces them. You are encouraged to refer to the man page for each utility, for complete details. In fact, those interested in a greater knowledge of Linux loadable modules should consult the source code for these utilities. Section 8.6.1, "Suggestions for Additional Reading" at the end of this chapter contains a reference for where they can be found.

8.2.1 insmod

The insmod utility is the simplest way to insert a module into a running kernel. You supply a complete pathname, and insmod does the work. For example:

\$ insmod /lib/modules/2.6.14/kernel/drivers/char/examples/hello1.ko

This loads the module hellol.ko into the kernel. The output would be the same as shown in Listing 8-5—namely, the Hello message. The insmod utility is a simple program that does not require or accept any options. It requires a full pathname because it has no logic for searching for the module. Most often, you will use modprobe, described shortly, because it has many more features and capabilities.

8.2.2 Module Parameters

Many device driver modules can accept parameters to modify their behavior. Examples include enabling debug mode, setting verbose reporting, or specifying module-specific options. The insmod utility accepts parameters (also called *options* in some contexts) by specifying them after the module name. Listing 8-6 shows our

modified hellol.c example, adding a single module parameter to enable debug mode.

Listing 8-6

Example Driver with Parameter

```
/* Example Minimal Character Device Driver */
#include <linux/module.h>
static int debug_enable = 0;
                                    /* Added driver parameter */
module_param(debug_enable, int, 0); /* and these 2 lines */
MODULE_PARM_DESC(debug_enable, "Enable module debug mode.");
static int __init hello_init(void)
    /* Now print value of new module parameter */
    printk("Hello Example Init - debug mode is %s\n",
            debug_enable ? "enabled" : "disabled")
    return 0;
}
static void __exit hello_exit(void)
    printk("Hello Example Exit\n");
module_init(hello_init);
module_exit(hello_exit);
MODULE_AUTHOR("Chris Hallinan");
MODULE_DESCRIPTION("Hello World Example");
MODULE_LICENSE("GPL");
```

Three lines have been added to our example device driver module. The first declares a static integer to hold our debug flag. The second line is a macro defined in .../include/linux/moduleparam.h that registers the module parameter with the kernel module subsystem. The third new line is a macro that registers a string description associated with the parameter with the kernel module subsystem. The purpose of this will become clear when we examine the modinfo command later in this chapter.

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If we now use insmod to insert our example module, and add the debug_enable option, we should see the resulting output, based on our modified hellol.c module in Listing 8-6.

```
$ insmod /lib/modules/.../examples/hello1.ko debug_enable=1
Hello Example Init - debug mode is enabled
```

Or, if we omit the optional module parameter:

```
$ insmod /lib/modules/.../examples/hello1.ko
Hello Example Init - debug mode is disabled
```

8.2.3 Ismod

The Ismod utility is also quite trivial. It simply displays a formatted list of the modules that are inserted into the kernel. Recent versions take no parameters and simply format the output of /proc/modules.⁵ Listing 8-7 is an example of the output from Ismod.

Listing 8-7

Ismod Example Output Format

\$ lsmod		
Module	Size	Used by
ext3	121096	0
jbd	49656	1 ext3
loop	12712	0
hello1	1412	0
\$		

Notice the rightmost column labeled Used by. This column indicates that the device driver module is in use and shows the dependency chain. In this example, the jbd module (journaling routines for journaling file systems) is being used by the ext3 module, the default journaling file system for many popular Linux desktop distributions. This means that the ext3 device driver depends on the presence of jbd.

8.2.4 modprobe

This is where the cleverness of modprobe comes into play. In Listing 8-7, we see the relationship between the ext3 and jbd modules. The ext3 module depends on the

⁵ /proc/modules is part of the proc file system, which is introduced in Chapter 9, "File Systems".

jbd module. The modprobe utility can discover this relationship and load the dependent modules in the proper order. The following command loads both the jbd.ko and ext3.ko driver modules:

\$ modprobe ext3

The modprobe utility has several command line options that control its behavior. As we saw earlier, modprobe can be used to remove modules, including the modules upon which a given module depends. Here is an example of module removal that removes both jbd.ko and ext3.ko:

\$ modprobe -r ext3

The modprobe utility is driven by a configuration file called modprobe.conf. This enables a system developer to associate devices with device drivers. For a simple embedded system, modprobe.conf might be empty or might contain very few lines. The modprobe utility is compiled with a set of default rules that establish the defaults in the absence of a valid modprobe.conf. Invoking modprobe with only the -c option displays the set of default rules used by modprobe.

Listing 8-8 represents a typical modprobe.conf, which might be found on a system containing two Ethernet interfaces; one is a wireless adapter based on the Prism2 chipset, and the other is a typical PCI Ethernet card. This system also contains a sound subsystem based on an integrated Intel sound chipset.

Listing 8-8

Typical modprobe.conf File

```
$ cat /etc/modprobe.conf
alias eth1 orinoci_pci
options eth1 orinoco_debug=9
alias eth0 e100
alias snd-card-0 snd-intel8x0
options snd-card-0 index=0
$
```

When the kernel boots and discovers the wireless chipset, this configuration file instructs modprobe to load the <code>orinoco_pci</code> device driver, bound to kernel device <code>eth1</code>, and pass the optional module parameter <code>orinoco_debug=9</code> to the device driver. The same action is taken upon discovery of the sound card hardware. Notice the optional parameters associated with the sound driver <code>snd-intel8x0</code>.

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8.2.5 depmod

How does modprobe know about the dependencies of a given module? The depmod utility plays a key role in this process. When modprobe is executed, it searches for a file called modules.dep in the same location where the modules are installed. The depmod utility creates this module-dependency file.

This file contains a list of all the modules that the kernel build system is configured for, along with dependency information for each. It is a simple file format: Each device driver module occupies one line in the file. If the module has dependencies, they are listed in order following the module name. For example, from Listing 8-7, we saw that the ext3 module had a dependency on the jbd module. The dependency line in modules.dep would look like this:

```
ext3.ko: jbd.ko
```

In actual practice, each module name is preceded by its absolute path in the file system, to avoid ambiguity. We have omitted the path information for readability. A more complicated dependency chain, such as sound drivers, might look like this:

```
snd-intel8x0.ko: snd-ac97-codec.ko snd-pcm.ko snd-timer.ko \
    snd.ko soundcore.ko snd-page-alloc.ko
```

Again, we have removed the leading path components for readability. Each module filename in the modules.dep file is an absolute filename, with complete path information, and exists on a single line. The previous example has been truncated to two lines, to fit in the space on this page.

Normally, depmod is run automatically during a kernel build. However, in a cross-development environment, you must have a cross-version of depmod that knows how to read the modules that are compiled in the native format of your target architecture. Alternatively, most embedded distributions have a method and init script entries to run depmod on each boot, to guarantee that the module dependencies are kept up-to-date.

8.2.6 rmmod

This utility is also quite trivial. It simply removes a module from a running kernel. Pass it the module name as a parameter. There is no need to include a pathname or file extension. For example:

```
$ rmmod hello1
Hello Example Exit
```

The only interesting point to understand here is that when you use rmmod, it executes the module's *_exit() function, as shown in the previous example, from our hellol.c example of Listings 8-1 and 8-6.

It should be noted that, unlike modprobe, rmmod does not remove dependent modules. Use modprobe -r for this.

8.2.7 modinfo

You might have noticed the last three lines of the skeletal driver in Listing 8-1, and later in Listing 8-6. These macros are there to place tags in the binary module to facilitate their administration and management. Listing 8-9 is the result of modinfo executed on our hellol.ko module.

Listing 8-9

modinfo Output

The first field is obvious: It is the full filename of the device driver module. For readability in this listing, we have truncated the path again. The next lines are a direct result of the descriptive macros found at the end of Listing 8-6—namely, the filename, author, and license information. These are simply tags for use by the module utilities and do not affect the behavior of the device driver itself. You can learn more about modinfo from its man page and the modinfo source itself.

One very useful feature of modinfo is to learn what parameters the module supports. From Listing 8-9, you can see that this module supports just one parameter. This was the one we added in Listing 8-6, debug_enable. The listing gives the name, type (in this case, an int), and descriptive text field we entered with the MODULE_PARM_DESC() macro. This can be very handy, especially for modules in which you might not have easy access to the source code.

8.3 Driver Methods

8.3 Driver Methods

We've covered much ground in our short treatment of module utilities. In the remaining sections of this chapter, we describe the basic mechanism for communicating with a device driver from a user space program (your application code).

We have introduced the two fundamental methods responsible for one-time initialization and exit processing of the module. Recall from Listing 8-1 that these are module_init() and module_exit(). We discovered that these routines are invoked at the time the module is inserted into or removed from a running kernel. Now we need some methods to interface with our device driver from our application program. After all, two of the more important reasons we use device drivers are to isolate the user from the perils of writing code in kernel space and to present a unified method to communicate with hardware or kernel-level devices.

8.3.1 Driver File System Operations

After the device driver is loaded into a live kernel, the first action we must take is to prepare the driver for subsequent operations. The open() method is used for this purpose. After the driver has been opened, we need routines for reading and writing to the driver. A release() routine is provided to clean up after operations when complete (basically, a close call). Finally, a special system call is provided for nonstandard communication to the driver. This is called ioctl(). Listing 8-10 adds this infrastructure to our example device driver.

Listing 8-10

Adding File System Ops to Hello.c

```
#include <linux/module.h>
#include <linux/fs.h>

#define HELLO_MAJOR 234

static int debug_enable = 0;
module_param(debug_enable, int, 0);
MODULE_PARM_DESC(debug_enable, "Enable module debug mode.");

struct file_operations hello_fops;

static int hello_open(struct inode *inode, struct file *file)
{
    printk("hello_open: successful\n");
    return 0;
}
```

```
static int hello_release(struct inode *inode, struct file *file)
    printk("hello_release: successful\n");
    return 0;
}
static ssize_t hello_read(struct file *file, char *buf, size_t count,
                loff_t *ptr)
{
    printk("hello_read: returning zero bytes\n");
    return 0;
}
static ssize_t hello_write(struct file *file, const char *buf,
                size_t count, loff_t * ppos)
{
    printk("hello_read: accepting zero bytes\n");
    return 0;
static int hello_ioctl(struct inode *inode, struct file *file,
                unsigned int cmd, unsigned long arg)
{
    printk("hello_ioctl: cmd=%ld, arg=%ld\n", cmd, arg);
    return 0;
static int __init hello_init(void)
    int ret;
    printk("Hello Example Init - debug mode is %s\n",
            debug_enable ? "enabled" : "disabled");
    ret = register_chrdev(HELLO_MAJOR, "hello1", &hello_fops);
         if (ret < 0) {
             printk("Error registering hello device\n");
             goto hello_fail1;
    printk("Hello: registered module successfully!\n");
    /* Init processing here... */
    return 0;
hello_fail1:
    return ret;
static void __exit hello_exit(void)
    printk("Hello Example Exit\n");
```

8.3 Driver Methods

```
struct file_operations hello_fops = {
   owner: THIS_MODULE,
   read: hello_read,
   write: hello_write,
   ioctl: hello_ioctl,
   open: hello_open,
   release: hello_release,
};

module_init(hello_init);
module_exit(hello_exit);

MODULE_AUTHOR("Chris Hallinan");
MODULE_DESCRIPTION("Hello World Example");
MODULE_LICENSE("GPL");
```

This expanded device driver example includes many new lines. From the top, we've had to add a new kernel header file to get the definitions for the file system operations. We've also defined a major number for our device driver. (*Note to device driver authors:* This is <u>not</u> the proper way to allocate a device driver major number. Refer to the Linux kernel documentation (.../Documentation/devices.txt) or one of the excellent texts on device drivers for guidance on the allocation of major device numbers. For this simple example, we simply choose one that we know isn't in use on our system.)

Next we see definitions for four new functions, our open, close, read, and write methods. In keeping with good coding practices, we've adopted a consistent naming scheme that will not collide with any other subsystems in the kernel. Our new methods are called hello_open(), hello_release(), hello_read(), and hello_write(), respectively. For purposes of this simple exercise, they are donothing functions that simply print a message to the kernel log subsystem.

Notice that we've also added a new function call to our hello_init() routine. This line registers our device driver with the kernel. With that registration call, we pass a structure containing pointers to the required methods. The kernel uses this structure, of type struct file_operations, to bind our specific device functions with the appropriate requests from the file system. When an application opens a device represented by our device driver and requests a read() operation, the file system associates that generic read() request with our module's hello_read() function. The following sections examine this process in detail.

8.3.2 Device Nodes and mknod

To understand how an application binds its requests to a specific device represented by our device driver, we must understand the concept of a device node. A device node is a special file type in Linux that represents a device. Virtually all Linux distributions keep device nodes in a common location (specified by the Filesystem Hierarchy Standard⁶), in a directory called /dev. A dedicated utility is used to create a device node on a file system. This utility is called mknod.

An example of node creation is the best way to illustrate its functionality and the information it conveys. In keeping with our simple device driver example, let's create the proper device node to exercise it:

\$ mknod /dev/hello1 c 234 0

After executing this command on our target embedded system, we end up with a new file called /dev/hellol that represents our device driver module. If we list this file to the console, it looks like this:

```
$ 1s -1 /dev/hello1
crw-r--r- 1 root root 234, 0 Jul 14 2005 /dev/hello1
```

The parameters we passed to mknod include the name, type, and major and minor numbers for our device driver. The name we chose, of course, was hellol. Because we are demonstrating the use of a character driver, we use c to indicate that. The major number is 234, the number we chose for this example, and the minor number is 0.

By itself, the device node is just another file on our file system. However, because of its special status as a device node, we use it to bind to an installed device driver. If an application process issues an <code>open()</code> system call with our device node as the path parameter, the kernel searches for a valid device driver registered with a major number that matches the device node—in our case, 234. This is the mechanism by which the kernel associates our particular device to the device node.

As most C programmers know, the open() system call, or any of its variants, returns a reference (file descriptor) that our applications use to issue subsequent file system operations, such as read, write, and close. This reference is then passed to the various file system operations, such as read, write, or their variants.

⁶ Reference to this standard is found in the "Suggestions for Additional Reading," at the end of this chapter.

For those curious about the purpose of the minor number, it is a mechanism for handling multiple devices or subdevices with a single device driver. It is not used by the operating system; it is simply passed to the device driver. The device driver can use the minor number in any way it sees fit. As an example, with a multiport serial card, the major number would specify the driver. The minor number might specify one of the multiple ports handled by the same driver on the multiport card. Interested readers are encouraged to consult one of the excellent texts on device drivers for further details.

8.4 Bringing It All Together

Now that we have a skeletal device driver, we can load it and exercise it. Listing 8-11 is a simple user space application that exercises our device driver. We've already seen how to load the driver. Simply compile it and issue the make modules_install command to place it on your file system, as previously described.

Listing 8-11

Exercising Our Device Driver

```
#include <stdio.h>
#include <stdlib.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <unistd.h>
int main(int argc, char **argv)
    /* Our file descriptor */
    int fd;
    int rc = 0;
    char *rd_buf[16];
    printf("%s: entered\n", argv[0]);
    /* Open the device */
    fd = open("/dev/hello1", O_RDWR);
    if ( fd == -1 ) {
        perror("open failed");
        rc = fd;
        exit(-1);
    printf("%s: open: successful\n", argv[0]);
```

8.5 Device Drivers and the GPL

Much discussion and debate surrounds the issue of device drivers and how the terms of the GNU Public License apply to device drivers. The first test is well understood: If your device driver (or any software, for that matter) is based, even in part, on existing GPL software, it is called a *derived work*. For example, if you start with a current Linux device driver and modify it to suit your needs, this is certainly considered a derived work, and you are obligated to license this modified device driver under the terms of the GPL, observing all its requirements.

This is where the debate comes in. First, the disclaimer. This is not a legal opinion, and the author is not a lawyer. Some of these concepts have not been tested in court as of this writing. The prevailing opinion of the legal and open source communities is that if a work can be proven⁸ to be independently derived, and a given device driver does not assume "intimate knowledge" of the Linux kernel, the developers are free to license it in any way they see fit. If modifications are made to the kernel to accommodate a special need of the driver, it is considered a derived work and, therefore, is subject to the GPL.

A large and growing body of information exists in the open source community regarding these issues. It seems likely that, at some point in the future, these concepts will be tested in a court of law and precedent will be established. How long that might take is anyone's guess. If you are interested in gaining a better understanding of the legal issues surrounding Linux and open source, you might enjoy www.open-bar.org.

8.6 Chapter Summary

This chapter presented a high-level overview of device driver basics and how they fit into the architecture of a Linux system. Armed with the basics, readers new to device drivers can jump into one of the excellent texts devoted to device driver writers. Consult Section 8.6.1 for references.

• Device drivers enforce a rational separation between unprivileged user applications and critical kernel resources such as hardware and other devices, and present a well-known unified interface to applications.

⁸ This practice is not unique to open source. Copyright and patent infringement is an ongoing concern for all developers.

```
/* Issue a read */
rc = read(fd, rd_buf, 0);
if ( rc == -1 ) {
    perror("read failed");
    close(fd);
    exit(-1);
}
printf("%s: read: returning %d bytes!\n", argv[0], rc);
close(fd);
return 0;
```

This simple file, compiled on an ARM XScale system, demonstrates the binding of application to device driver, through the device node. Like the device driver, it doesn't do any useful work, but it does demonstrate the concepts as it exercises some of the methods we introduced in the device driver of Listing 8-10.

First we issue an open() system call⁷ on our device node created earlier. If the open succeeds, we indicate that with a message to the console. Next we issue a read() command and again print a message to the console on success. Notice that a read of 0 bytes is perfectly acceptable as far as the kernel is concerned and, in actual practice, indicates an end-of-file or out-of-data condition. Your device driver defines that special condition. When complete, we simply close the file and exit. Listing 8-12 captures the output of running this example application on an ARM XScale target:

Listing 8-12

Using the Example Driver

```
$ modprobe hello1
Hello Example Init - debug mode is disabled
Hello: registered module successfully!
$ ./use-hello
./use-hello: entered
./use-hello: open: successful
./use-hello: read: returning zero bytes!
$
```

⁷ Actually, the open() call is a C library wrapper function around the Linux sys_open() system call.

- The minimum infrastructure to load a device driver is only a few lines of code. We presented this minimum infrastructure and built on the concepts to a simple shell of a driver module.
- Device drivers configured as loadable modules can be inserted into and removed from a running kernel after kernel boot.
- Module utilities are used to manage the insertion, removal, and listing of device driver modules. We covered the details of the module utilities used for these functions.
- Device nodes on your file system provide the glue between your userspace application and the device driver.
- Driver methods implement the familiar open, read, write, and close functionality commonly found in UNIX/Linux device drivers. This mechanism was explained by example, including a simple user application to exercise these driver methods.
- We concluded this chapter with an introduction to the relationship between kernel device drivers and the Open Source GNU Public License.

8.6.1 Suggestions for Additional Reading

Linux Device Drivers, 3rd Edition Alessandro Rubini and Jonathan Corbet O'Reilly Publishing, 2005

Filesystem Hierarchy Standard Edited by Rusty Russel, Daniel Quinlan, and Christopher Yeoh The File Systems Hierarchy Standards Group www.pathname.com/fhs/

Rusty's Linux Kernel Page Module Utilities for 2.6 Rusty Russell http://kernel.org/pub/linux/kernel/people/rusty/