

9/29/20 Notes

Control structures and flow control

Boolean Algebra

- set mathematics on logical foundation

- Boolean logic has true and false

- 3 basic operators

- \wedge called AND (S&H in E)

- \vee OR which \rightarrow (1L in C) $\text{B&D} \rightarrow \text{O}$ (1L in C) \rightarrow 1

And (conjunction)

- associativity identity

- De Morgan's law $\neg(A \wedge B) \equiv \neg A \vee \neg B$ (distributive over AND)

- $A \wedge 1 = A$

- $A \wedge 0 = 0$

- $A \wedge A = A$ (idempotent law) \wedge is commutative and associative

Or (inclusion / disjunction)

- $A \vee (B \vee C) = (A \vee B) \vee C$ (associative over OR)

- De Morgan's law - distributive

- $A \vee 0 = A$ (absorption law - distributes over OR)

- $A \vee 1 = 1$

- $A \vee A = A$

Not (negation)

- De Morgan:

$\neg(A \wedge B) = \neg A \vee \neg B$

Exclusive-or (XOR)

first thing or second thing but not both

- $A \oplus B = (A \vee B) \wedge \neg(A \wedge B)$

- $A \oplus 0 = (A \wedge \neg B) \vee (\neg A \wedge B)$

Algebraic properties

- $A \oplus A = 0$

- $A \oplus 0 = A$

- $A \oplus 1 = \neg A$

- $A \oplus (B \oplus C) = (A \oplus B) \oplus C$

$\oplus(\text{XOR})$	T	F
T	F	T
F	T	F

Logical Operators

< less

L= less than equal

== equal

> greater

>= greater than equal

true and false in

- In C, zero (0) is false

- All that is not false is true

- Logical op. have type int

- You can have true and false if you

- # include <stdbool.h>

if()

- if execs next statement if Bool exp. is true

- Even though {} are not required for a single statement, always use them

- Why?

- it avoids errors when adding statements

if () { ... } else { ... }

Nested if ()

if() { ... } else if () { ... } ... else { ... }

Boolean Operators

&& = and

|| = or

~ = not

Short circuit evaluation

- false is anything is false regardless of what else is true
 - true is anything is true is either true or false now
 - stop evaluating as soon we know the result
 - suppose we eval exp:
we could be dividing by zero
 - we could follow a null path

Switch (~~at~~ esp) &

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$\mu \text{g} = " \text{Cug"}$

break), top of file: 3001 rev 339 Mar 30

Case 1, had strong physical signs, and many.

msg.

breakfast

Contra

default: ~~most part~~ ~~and~~ ~~not~~ ~~most~~ ~~not~~ ~~most~~ ~~not~~

~~longest~~ - The last word is derived from

3

switch ()

- allows you to select among a fixed set of alts.

equivalent goto code

What's wrong w/ goto?

- it's far about how you got there
- you could goto middle of if statement
- switch statements
- a loop

It is not even clear what that would mean.

What is a loop?

- repeats a seq. of code.
- progs. spend vast majority of their execution time in loops
- we will focus on loops: while, for, and do-while
- you can also create loops w/ do goto last day.

while()

- top - first loop
- first is evaluated before entering loop
- executes statements as long as 2nd cond. remains true

Equivalent goto code.

- you can implement it w/ goto statements
- not recommended

for()

• also top first loops

• 3 parts:

- initialization

- test, and

- increment all together

By convention they are related, but not required to be related.

equivalent while()

- This is the ~~that~~ equivalent while statement
- the while statement is complete
- which means you can implement any loop

do {} while

- bottom-test loop

- used when you want to perform the statement at least once
- continues to execute the enclosed statement over and over until the break condition is met

Infinite loops,

- executes forever
- the one you choose is a matter of style not of safety.
- How do you ever escape?
- break;

break

- immediately exits enclosing loop

design 2:

Factorial example

you don't
need to do
factorials if

when can I use goto?

- one place: non-local error handling
- this is when an exceptional condition that you cannot handle occurs
- It is not pretty.

Continue:

- You may have times when you want to stop remember of a loop
- for please use sparingly

Let's compute \sqrt{d} !

- \sqrt{d} is the same as: $d^2 - 2 = 0$
- we also know: $0 \leq x \leq d$,

Bisection Method

To look for a solution of $d^2 - 2 = 0$ we can use the bisection method.
We start with an interval $[a, b]$ where $a < b$.
Then we calculate the midpoint $m = \frac{a+b}{2}$.
If $f(m) > 0$, then the solution must be in $[a, m]$.
Otherwise, it must be in $[m, b]$.

After each step, we have a new interval $[a, b]$ which is half as long as the previous one.
This process continues until the width of the interval is small enough.
For example, if we start with $[0, 1]$ and we want the width to be less than 0.001 ,
then we need to do 10^3 steps.
This is because $10^{-3} = \frac{1}{1000}$ and $\frac{1}{2^{1000}} < 10^{-3}$.

Another question is: how many digits will we get?
The answer is: $\log_{10}(n)$ where n is the number of steps.

10/1/21 Notes
Word of Caution

- functions C are not the same in math function
 - they may / not return a value (will return, nothing)
 - may / may not have side-effects
- subroutines in other languages
- In Java, it's a method

How is it different?

- math has domain and range
- In C, we call function
 - function it returns values

functions in programming.

- block of code that performs a certain task
- defined exactly once
- must be declared before they are used
- can declare & call a function as many times as desired
- main()
 - is a special function
 - Is run when program starts
 - All other functions are subordinate main()

Why do we like functions?

- functions should
 - define abstraction
 - give names to seqs. of code
 - hide the implementation
- use them to:
 - refactor repeated seqs. of code
 - simplify code to aid understanding
- Functions should never be
 - Arbitrary seqs. of static methods

Function definition.

- function head (or function declaration) \rightarrow function name and parameters
 - return type \rightarrow function's output
 - defines type of function's return value
 - return type could be void or any obj. type (except. array)
- function-name
- parameters
 - contained in comma-separated list at def.
 - if function has no parameters \rightarrow then void.
- Function block / body
 - declarations
 - declared variables inside a func. bld. are local
- return new

Return values.

- function returning a value
- may be `int`
- return a struct (not recommended)
- return a pointer that points to the heap memory
- not an array

Function naming.

- same naming rule as variable
- can't start
 - start w/ a number or only punctuation other than - and _
- use snake case
- my_function_name

Parameters or Arguments

- if we have $f(x) = x \log(x+1)$ and we write $f(a)$ we substitute a for x and get $a \log(a+1)$.
- this is call-by-value and so may support this as default subs in macros and C Preprocessor.
- most prog. lang. use either call-by-name, call-by-reference or both.
- C uses call-by-value, except for arrays, and only b/c of their relation to pointers.

Parameters

- formal parameters
 - name of parameter w/ its value body
- Actual parameters
 - name of value that is passed to func
 - the value can be copied to the formal param.
 - or affect reference to the actual parameter
- Copy-in - Copy out means that if modified to base copied, value is copied back out
- C does not support this

Call-by-value

- all function use this.
- Arguments passed into a function are copied
 - any changes made to the parameters inside the function has no effect
- the called function copies the values in

(all - by - reference)

• ref. i

Ternary operator in switch statement from question

if h = (x < 1) ? 1 : x; // already returning 0 so no need

if (x < 1) {

return 1;

else { // already returning 0 so no need

return x; // already returning 0 so no need

swap() prob doesn't do what it's intended

(does not have i) which call-by-reference was wrong

points to a local variable that goes to zero

• Address initial & values are passed to arguments

void swap (int *a, *b) { // address of a and b

int temp = *a;

*a = *b; // swap the actual memory

*b = temp; // swapped the original memory back

return;

}

int void main (void) { // output passing to for n-

int x = 5;

int y = 7;

swap (&x, &y); // swap the actual memory

printf ("val x = %d\n",

printf ("val y = %d\n",

Function prototypes

- syntax:

~~function-name (parameters)~~

return-type function-name (parameters);

- prototype must be declared either at beginning of program or in included header files.

#include

- a preprocessor directive,
- Before compilation, C source files are processed by preprocessor
- preprocessing is a macro processor to transform programs before compilation
- works in C or C++ through text replacement
- works on .c or .cpp file, .h or .hpp file
- used to include functions defined in other libraries

#define

- A preprocessor directive that defines a macro for the program
- C preprocessor performs all text replacement for defined macros prior to compilation

Conditional Directives

- a set of preprocessor directives that vary conditionally to include code selectively.
- #if, #else, #endif statements

#if, #else, #endif statements will execute when macro is defined.

Header file

should only contain things that are shared b/t source files

Header is. # program only.

#ifndef checks macro will always work (use it more portably)

Effect

- extends visibility of var and function such that they can be called by any program file.

static

• can be declared in and outside functions.

18/4/21 Note

On the nature of Numbers

- We all think that we know numbers
- 1, 2, 3, 4 ... are all numbers right!
- they are representation of numbers
- you cannot hold a number in hand
- Numbers exist independently of their representation
- Our method for writing numbers comes from Hindu-Arabic numerical system.

Kinds of numbers

- $N = \{1, 2, 3, \dots\}$
 - $Z = \{..., -3, -2, -1, 0, 1, 2, 3\}$
 - $Q = \{q = \frac{p}{n} : p, n \in Z\}$
 - $R = Q \cup$ the irrational numbers.
 - $C = R \cup$ the imaginary numbers
- $\pi, e, \sqrt{2}, \phi$ are all irrational

Integers and Natural Numbers

- $Z = \{..., -3, -2, -1, 0, 1, 2, 3\}$ denotes integers
- $Z^+ = \{1, 2, 3, \dots\}$ + integers
- $Z^- = \{..., -3, -2, -1\}$ set of negative integers
- 0 is special if it is neither + or -
- $N \subseteq O \subseteq Z^+$

Computers can represent none of these sets.

- Computers do arithmetic in finite field
- which means they have one large limit of size
- the memory can handle up to 10 digits
- digit is usually called a bit
- a bit is either 0 or 1

Positional Number System

every member of N can be written as

$$\cdot q_k b^k + q_{k-1} b^{k-1} + \dots + q_1 b^1 + q_0 b^0$$

$$146_2 = 1 \times 10^3 + 4 \times 10^2 + 6 \times 10^1 + 2 \times 10^0$$
$$= 74A_{16} = 2652_8 = 200000_3$$
$$= 110101010_2$$

All represent the same number.

Licence plates

numbers are uniform in base 36

$A = 10, B = 11, C = 12, \dots$

$A24 BC = 10 \times 36^5 + 2 \times 36^4 + \dots$

Specifying an integer in C

undefined } short
long } int
long long }

unsigned } char

C Integer data types

char \rightarrow 8 bits

short \rightarrow 16 bits

int \rightarrow probably 16 bits

long \rightarrow probably 32 bits

long long \rightarrow probably 64 bits

#include <stdint.h>

signed	unsigned	size
int8_t	uint8_t	8 bits
int16_t	uint16_t	16 bits
int32_t	uint32_t	32 bits
		64 bits

Binary arithmetic

Addition works as expected

- $0+0=0$
- $1+0=1$
- $1+1=10$ (the 0 carry for 1)

Multiplication is even simpler:

- $0 \times 0 = 0$
- $1 \times 0 = 0$
- $1 \times 1 = 1$

$$\cdot 101 + 11 = 1000$$

$$\cdot 101 \times 101 = 11001$$

Arithmetic in a Finite field.

• in a finite field, we can define for every number its additive inverse.

if i is a number in a finite field and i is its additive inverse then $i + i = 0$.

Suppose we take k bits then we can have any positive integer from 0 to $2^k - 1$

Two's Complement Arithmetic

To get the additive inverse of a number m ,

- Flip the bits in m , $0 \rightarrow 1_x$ and $1 \rightarrow 0$, and then
- add one to result

Real Numbers

- Real numbers (\mathbb{R}) are:
- Continuous
- Uncountably infinite

There are just as many numbers between any two any as there are all of \mathbb{R} .

- Include all of the:
 - integers (\mathbb{Z})
 - Rational numbers (\mathbb{Q})
 - Irrational numbers ($\mathbb{R} - \mathbb{Q}$)

Floating point numbers

- Are a proper subset of the real numbers.
- $\mathbb{F} \subset \mathbb{R}$
- are a proper subset of rational numbers.
 - $\mathbb{F} \subset \mathbb{Q}$
- are a proper subset of the integers.
 - $\mathbb{F} \subset \mathbb{Z}$
- mistakes to think that a real, rational, floating point, are approximately

Floating point number

float → 32

double → 64

long double → 128

single precision

0 0 ... 0 0 1 ... 0
↑ ↑
sign exponent fraction
(8 bits) (23 bits)

Intel Extended Precision (64 bits)
sign (1 bit) (150 bits)

Convert to floating point precision

Big endian, little endian

Little endian = low address byte

Big endian = high address byte

Big endian vs. Little endian

Big

12 34 56 78

Little

78 56 34 12

Random numbers

- True random numbers cannot be created by computers.
- Why?
 - Programs are inherently deterministic.
 - It has the advantage of repeatability.
 - It has the disadvantage of predictability.

Mersenne Twister

- A state-of-the-art pseudo number generator.

Arithmetic operators

- These operators follow the precedence and associativity rules that you learned in high school algebra.
- Modulo is the remainder when you divide 2 integers.

Type promotion

With mix of express types C will promote the low type to a higher type.

Operator precedence

* / %

+ -

<< >>

< <= > >=

== !=

&

^

!

&&

||

?:

= += -= += /= % = ...

10/6/21 Notesⁿ

Numerical Computation

Basic Arithmetic operation

- computer can only do basic operation with e.g. add, sub, mult, div
 - mult. \rightarrow shift + add
 - Div \rightarrow shift + sub
 - add \rightarrow is little more than exclusive-or
- trig functions?
 - processor do them internally, they are still only basic operations

Absolute Value

- $\text{abs}(x) = \text{if } x < 0 \text{ then } -x \text{ else } x$

(can I just use the library?)

- yes, it's best to use the library
- the library should be designed to be:
 - careful w/ numerical precision
 - fast + portable.

Taylor Series

- if $f(x)$ is a real or complex valued function, and it is infinitely differentiable then
- $$f(x) = \sum_{n=0}^{\infty} \frac{f^{(n)}(a)}{n!} (x-a)^n$$

• for example consider e^x centered @ $a=0$!

$$e^x = 1 + \frac{x}{1!} + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \quad -\infty < x < \infty$$

Increasingly better approximation,
- each deriving up add a term

e^x
but, the simplest Taylor series
Computing e^x
 $|k| > x^k$ for all x , eventually (when k gets large enough),
so,

x^k
gets smaller w/ each step

we can stop when it is "small enough" for our computation

Computing $\sin(x)$

If \sin ,

- periodic on $[-\pi, \pi]$, meaning $\sin(9\pi) = \sin(\pi)$
- infinitely differentiable, making a good candidate for a Taylor series
- where $\sin(0) = 0$

Computing the Terms.

- we skip even numbered terms
- the signs alternate b/t terms
- remove extra factors of 2π

"Padé" approximants

Padé approximants \rightarrow ratio of 2 polynomials that correspond to a series but p. well to compute and fit better around the center

What happens if the series converge too slowly?

- Series for $\log(1+x)$:

$$x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots + O(x^n)$$

- Converges very slowly unless x is close to 0.

- We can do little algebra, remember!

- $\log(x) = -\log(\frac{1}{x})$; for $x > 0$

(recall)

- $x = \log(e^x) = e^{\log x}$

(you can invert this)

Inverting a function

'less simple' functions that give a stable'

• \sqrt{x}

- we can't really expand a Taylor series for $x^{1/2}$

\sqrt{x} w/ binary & search

- goes that it is in the middle

- is the given too small? Look in the larger half.

- is the given too big? Look in the smaller half.

\sqrt{b} using a Newton Iteration

- Isaac Newton coming to the rescue again

- we use this iterate:

$$x_{k+1} \rightarrow \frac{1}{2}(x_k + \frac{b}{x_k})$$

- As before, we iterate until the approximation is good enough

$\log(x)$ using Newton's Method.

Newton's formula provides a linear approximation of the function

Solve for x_{n+1}

$$x_{n+1} = x_n + \frac{f(x_n)}{f'(x_n)}$$

Find the root of $f(x) = x - c$

we iterate this equation until answer is close enough.

$$y_{n+1} = y_n + \frac{x - c}{c}$$

Floating point arithmetic

normal mathematics

real numbers are exact

Computer arithmetic

approximate numbers

Floating points have round off error

Round-off errors

errors occur when doing any calculation w/ floating-point
rounding errors occur to fit values

IEEE 754 defines standard rounding modes

Some dangers of floating point arithmetic

add + sub

Compare floating-point numbers

- What do we know?
- floating point numbers are not real numbers
 - we need to take care when working with floating-point numbers

What should we do?

- use a good library function for floating-point comparison
- understand the difference between floating-point and real numbers
 - $\text{real} = \text{all}$ the ~~real~~ numbers
 - $\text{floating} =$ all the floating-point numbers