CHAPTER 1

- Components of a computer: ALU and Control Unit (CPU), Memory, Input, and Output.
- 2. Functions of various components:

.CPU: It processes and stores binary data, transfers data from and to memory and I/O devices, and provides timing to all the operations. It includes ALU, register arrays, and control unit. The ALU performs the arithmetic and logic operations, and the control unit provides timing.

.Input - provides binary data as an input to the CPU.

.Output - accepts binary data from the CPU.

- 3. A microprocessor functions as the CPU of a microcomputer, and includes the ALU, register arrays, and the control unit on one chip; it is manufactured using the LSI technology. On the other hand, the CPU is designed with various discrete boards. Functionally, both are similar; however, technology and processes used for designing is different.
- A microprocessor is one component of a microcomputer, and the microcomputer is a complete computer consists of a microprocessor, memory, input, and output.
- 5, 6. See Summary: Scale of Integration
- Four bytes.
- 8. The machine language of the 8085 are the commands to the microprocessor given in binary. These are the binary instructions the processor can understand and execute. The assembly language comprise of mnemonics (group of letters to represent commands) assigned by the manufacturer for the convenience of the users.
- 9, 10, 11. See Summary: Computer Languages
- 12. The assembly language mnemonics represent instructions to the microprocessor; therefore, when they are translated into machine language, there is one-to-one correspondence between the mnemonics and the machine code. The assembly language programs are compact, require less memory space, and are efficient. The high level languages are written in English-like statements, and when these statements are translated in machine language, the object code tends to be large, and requires large memory. The execution of the programs written in high level languages is less efficient than that of assembly language programs. http://jntu.blog.com

- 13, 14. See Summary: Computer Languages
- http://jntu.blog.com 15. ASCII codes in Hex: A = 41, Z = 5A, and m = 6D
- 16. See Summary: Computer Languages

CHAPTER 2

- 1. Memory Read, Memory Write, I/O Read, and I/O Write.
- 2. A bus is group of lines (wires or conductors) which carry digital information.
- The function of the address bus is to carry a binary address of a memory location or an I/O
 device. The address bus is unidirectional, and the information flows from the MPU to
 peripherals and memory.
- A microprocessor with 14 address lines is capable of addressing 16 K (2) memory locations.
- 21 address lines.
- Data bytes are transferred in both directions between the MPU and memory/peripherals.
- IOR (I/O Read), IOW (I/O Write), MEMR (Memory Read), and MEMW (Memory Write).
- In memory write operation, the control signal required is MEMW, and the direction of the data flow is from the MPU to memory.
- 9. The accumulator is an 8-bit register and it is a part of the ALU. All 8-bit arithmetic and logic operations are performed in relation to the accumulator content, and the result is stored in the accumulator (with a few exceptions).
- 10. A flag is the output of a given flip-flop to indicate certain data conditions.
- 11. The program counter and the stack pointer store memory addresses of 16 bits.
- The program counter always points to the next memory location; therefore, the content of the program counter will be 2058H.
- 13. 128 registers and 128 X 4 = 512 memory cells.
- 14. 1024 bits are can be stored by this chip; however, it can not be specified as a 128-byte memory chip because the byte indicates 8-bit memory registers; this chip has 4-bit registers.

http://jntu.blog.com 15. 8-bit word size. 16. 8 chips. 17. 4 chips. 18. 32 chips. The WR signal enables the input buffer of a memory chip so that information can be stored (written) in the selected memory register. 20. 11 address lines. 21. 16 pages and the last location is 2FFFH. 22. The starting address is F800H, and the memory map is F800H to FBFFH. 23. The starting address is: E000H. 24. The address ranges from FF00H to FFFFH. 25. The address of the selected register: $1000\ 0000\ 0100\ 0111 = 8047H$ 26. The memory map ranges from 2000H to 23FFH. The address of the selected register: 0010 0000 1111 1000 = 20F8H 27. 28. 8 address lines are required for a peripheral I/O port, and 16 address lines are required for a memory-mapped I/O port. 29. Tri-state devices are logic devices with three states; the third state is high impedance. In a bus-oriented system, devices are connected in parallel, and the buses are capable of driving one TTL logic device. The MPU communicates with one peripheral at a time, and other peripherals are placed in high impedance to avoid bus loading. 30. High impedance state. 31. From B to A. 32. None. The decoder is not enabled; all output lines will be high. 33. The line $6(O_6)$

0 0 1 (Complement of 1 1 0) http://jntu.blog.com

34.

- A transparent latch is a flip-flop; its output changes according to input when the clock signal is high, and it latches the input when the clock goes low. The latch is necessary for output devices to retain the result; otherwise, the result will disappear.
- The high-order address lines: A12-A15, the low-order address lines: A0-A10, and the don't care line: A11.
 - 37. This answer assumes the memory chips are 2048 X 4:

38. The memory occupies the memory space from F000H to FFFFH. The don't care line A11 generates additional address range. This is a 2K memory chip that occupies 4K of memory space in the map; thus wasting 2K of memory space. If A11 is assumed to be at logic 0 as in Q. 37, the address range is: F000H to F7FFH and if it is assumed to be at logic 1, the address range (also called foldback memory space) is: F800H to FFFFH.

CHAPTER 3

- The ALE signal goes high at the beginning of each machine cycle indicating the availability
 of an address on the address bus, and the signal is used to latch the low-order address bus.
 The IO/M signal is a status signal indicating whether the machine cycle is I/O or memory
 operation. The IO/M signal is combined with the RD and WR control signals to generate
 IOR, IOW, MEMR and MEMW control signals.
- 2. The low-order bus AD7-AD0 is used for two purposes. In the earlier part of a machine cycle, the bus is used for the low- order address of a memory location the 8085 is accessing, and in the latter part of the cycle the bus is used for data. By demultiplexing the bus, the address and the data are kept separate.
- In Fig. 3.22, the input signal RD and WR cannot be low at the same time. Therefore, the valid combinations of the input signals are:

IO/M	RD	WR	Out	tp://jntu:blog.co	m
0	0	0	Oo	Invalid	RD and WR cannot be active simultaneously
0	0	1	O_1	MEMR	M and RD active
0	1	0	O_2	MEMW	M and WR active
0	1	1	O_3	Irrelevant	Both RD and WR are inactive
1	0	0	O_4	Invalid	RD and WR cannot be active simultaneously
1	0	1	O ₅	IOR	IO and RD active
1	1	0	O ₆	IOW	IO and WR active
1	1	1	O ₇	Irrelevant	Both RD and WR are inactive

- See the answer of Q.3.
- In Fig. 3.23, the 74LS139 is enabled when IO/M is low. Therefore, the following memory control signals can be generated.

RD WR Decoder Output

0 0 O₀ - Invalid 0 1 O₁ - MEMR 1 0 O₂ - MEMW 1 O₃ - No operation

- 6. The output of the latch will be 05H; however, it will be not be latched until the ALE goes low.
- The output of the latch is 05H. At T2, the ALE is low; therefore, the latch will not be enabled, and it will continue to hold the previously latched byte (05H).
- 8. The crystal frequency should be = 2.2 MHz because the oscillator logic divides the input frequency by two.
 - 9. See the steps on page 66/67, Example 3.1.
 - 10. The sum of 87H + 79H = 100H. Therefore, the accumulator will have 00H, and the flags will be S = 0, CY = 1, Z = 1.
 - 11. 2060H. The program counter always points to the next machine code to be fetched.
 - 12. 18T X .2 micro-sec = 3.6 micro-sec.
 - 13. (A15-A8) = 20H, (AD7-AD0) = 47H, (PC) = 2076H
 - 14. RD and IO/M are asserted low.

- The second machine cycle is Memory Read; the processor reads the contents of memory in register B, and the control signal is RD http://intu.blog.com
- The fourth machine cycle is Memory Read; the processor reads the contents of memory in the accumulator.
- 17. (A15-A0) = 2050H(AD7-AD0) as data bus = Contents of location 2050H
- 18. (Refer to Instruction Set on pages 696-699)
 SUB B = OF (Opcode Fetch)
 ADI 47H = OF, MR (Memory Read)
 STA 2050H = OF, MR, MR, MW (Memory Write)
 PUSH B = OF, MW, MW
- Memory map: 6000H to 6FFFH
- Memory map: 8000H to 8FFFH
- OR gate
- 22. Connect RD to OE of the memory chip and IO/M to E2 of the decoder.
- - 24. Total range = 16K. Map = 8000H to BFFFH
- 25. A data byte entered at location 2100H will be accepted and stored at location 2000H. The address lines A10, A9, and A8 are not being used for memory addressing; therefore, they can assume 0 or 1 (don't care) logic state which results into multiple addresses for the same memory locations.
- 26. Memory address: 0800H-08FFH, and the foldback memory ranges from 0900 to 0FFFH.
- 27. Memory map: 3800H 3FFFH.
- 23. In Figure 3.19, three lines are don't care which can have (23) eight combinations. Thus, the memory chip will occupy the memory space equal to eight times its size.
- 29. ROM1: 0000H 1FFFH, ROM2: E000H FFFFH, R/WM1: 8000H 83FFH

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Memory map: 8000H to of FHH (Assume all don't care lines at 0)

Foldback Memory: 8400H to 9FFFH

The address range: 0000H to 3FFFH

32. The address range: 4000H to 7FFFH

33. The primary address range: 0000H to 1FFFH (Assumes A13 = 0)

The foldback or the mirror address range: 2000H to 3FFFH

- The mirror address range: 8000H to 9FFFH
- 35. The address range when Y1 is asserted: 4000 to 7FFFH
- 36. The total address range is: 4000H to BFFFFH. For a 16K memory chip, when A14 = 1, the address ranges from 4000H to 7FFFH as in Q. 35. When A14 = 0, the address ranges from 8000H to BFFFH. For a 32K memory chip, it is accessed either by Y1 or Y2; therefore, the address ranges from 4000H to BFFFH.
- 37. The opcode fetch cycle begins immediately after MEMW signal.
 1st MEMR ----> opcode fetch of the JMP instruction.
 4th MEMR ----> opcode fetch of the MVI instruction.
 6th MEMR ----> opcode fetch of the STA instruction.
- The last MEMR is the third byte of the STA instruction. It reads FFH.

CHAPTER 4

- The number of output ports in the pripheral I/O is restricted to 256 ports because the operand of the OUT instruction is 8 bit; it can have only 256 combinations.
- 2. Yes.
- The 8085 differentiate between the input and the output ports of the same address by the control signal. The input port requires the RD and the output ports requires the WR control signals.
- 4. WR (low) and IO/M (high). 4
 - Pulse going from high to low.
 - Trailing edge.

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7.

- Each LED requires 10 to 19 mA current for proper illumination. The latch cannot supply the necessary current when the output is logic high, but it can sink the necessary current when the logic level is low.
- RD (low) and IO/M (high).

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- A latch is necessary to hold the output data for display; however, the input data byte is obtained by enabling a tri-state buffer and placed in the accumulator.
- 10. RD, WR, and IO/M (low).
- 11. No.
- 12. 78H.

W

No. An output byte will be displayed temporarily until the WR signal is active, and then, it will disappear.

Memory-mapped I/O. LE is enabled when IO/M is low.

15. 8000H.



- Assuming A3 = 0, port address = F1H.
- 17. If A7 = 0, port address = 75H, and if A7 = 1, address = F5H.
- 18. If IO/M is connected to /E1 (active low), it will be a memory-mapped I/O. The port address = 00F5H.

Replace OUT F5H by STA 00F5H

- 19. MVI A, C0H ;Code for '0'
 OUT F5H
 HLT
- 20. In Q. 19, replace the code C0H by the code for letter 'H'. Code for H = 89H.
- If A7 is replaced by IO/M signal, the circuit will have three don't care address lines: A7, A4
 and A3 resulting in eight different addresses.

If A7 = 0, the addresses are: 04H, 0CH, 14H and 1CH.

If A7 = 1, the addresses are: 84H, 8CH, 94H, and 9CH (as shown in Section 4.34).

22. The port will be a memory-mapped I/O with an address = 00F8H.

- 23. Port A = Memory-mappedattp://jiphut.blog.com
 Port B = Memory-mapped Input Port
- Both are memory-mapped I/O ports. Assuming the address lines A15-A8 are at logic 0, Port A and Port B will be 0085H.
- In Figure 4.10, the output O5 is enabled by the address which is active for three T-states. On the other hand, the IOW signal requires WR signal which is active for approximately one and half T-states.
- 26. a. Machine Cycles: M1 M2 M3

IN 84H OF MR IORD JMP START OF MR MR

- b. $20T \times 0.5 = 10 \text{ micro-sec.}$
- c. Six times.

. 10 micro-sec (from beginning to th next beginning)

- e. There is no WR pulse in the routine. IO/M high or IORD can be used to sync the scope.
- 27. a. Machine Cycles: M1 M2 M3 M4

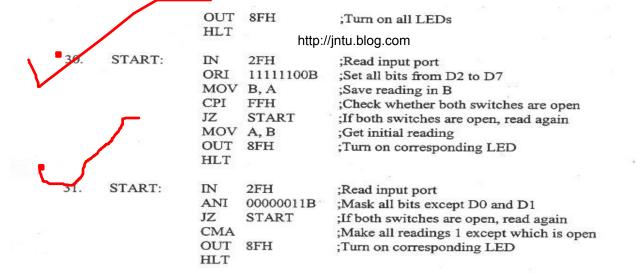
LDA FFF9H OF MR MR MR STA FFF8H OF MR MR MW MOV B,A OF JMP START OF MR MR

- b. FFF9H
- c. RD = 11 times and WR = 1 time.
- d. 40T X 0.5 micro-sec = 20 micro-sec.
- 28. In Figure 4.18, the address line A4 is don't care.

 Assuming A4 = 0: Input Port = 2FH and Output Port = 8FH.

 Assuming A4 = 1: Input Port = 3FH and Output Port = 9FH.
- 28. START:

IN 2FH ;Read input port
ANI 00000011B ;Mask all bits except D1 and D0
JNZ START ;If a switch is open, read again
MVI A, 00 ;This instruction is unnecessary
;Used here for clarity
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32.. The address of the latch enabled by Y3 = F5H and the address of the latch enabled by Y2 = F4H.

MVI A, 98H ;Common anode code for '9'
OUT F5H
MVI A, F8H ;Common anode code for '7'
OUT F4H
HLT

http://inturbleg.comuTIONS

CHAPTER 5

- The four categories of instructions that manipulate data are: data transfer (copy), arithmetic, logic, and branch.
- The task to be performed is called the opcode (operation code), and the data to be operated on is called the operand which may be specified as data, register or address.

Opcode: MOV and Operand: H,L

- 3. The machine code: 01 100 111 = 67H
- 4. (a) 2647H OPCODE = MVI OPERANDS = H, 47H
 - (b) C6F5H OPCODE = ADI OPERANDS = A (IMPLIED), F5H
 - (c) 91H OPCODE = SUB OPERANDS = A (IMPLIED), C
- 5. (a) HEX = 325020H OPCODE = STA OPERANDS = 2050H
 - (b) HEX = C27020H OPCODE = JNZ OPERANDS = 2070H
- 6. The SUB A instruction clears the accumulator. Z = 1, CY = 0
- 7. INSTRUCTION ADDRESS HEX

MVI B,4FH	2000	064F
MVI C,78H	2002	0E78
MOV A,C	2004	79
ADD B	2005	80
OUT 07H	2006	D307
HLT	2008	76

8. INSTRUCTION ADDRESS HEX

MVI A,8FH	2020	3E8F
MVI B,68H	2022	0668
SUB B	2024	90
ANI OFH	2025	E60F
STA 2070H	2027	327020
HLT	202A	76

INSTRUCTION ADDRESS HEX

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START: IN F2H http://jntu.blog.com

CMA 2002 2F ORA A 2003 B7 JZ START 2004 CA0020

10. Logical steps to add two Hex numbers:

Load A2H in one register.

Load 18H in second register.

Copy A2H in the accumulator.

Add the contents of the second register to the contents of the accumulator.

End of the program.

MVI B, A2H

MVI C, 18H

MOV A, B

ADD C

HLT

Register contents:

Initial: B=28H, A=97H

After the execution: A=28H, B=28H, C=28H

- 13. In Q. 6, if the code 07H (port address) is omitted, the processor assumes the opcode of the next instruction 76H (HLT) as the address of the output port, outputs the contents of the accumulator to the address 76H, and continues to the next code. After the next code, results are indeterminate.
- 14. In Q. 8, if the byte 0FH is omitted, the processor assumes the opcode 32H of the next instruction (STA) as the second byte of the ANI instruction. The processor is a sequential machine; it assumes the next code 20H (the low-order address of 2070H) as the opcode of the next instruction and continues.

CHAPTER 6

Section 6.1: Data Transfer (Copy) Operations

A B C D S Z CY

MVI A,00 00 NA NA NA

MVI B,F8 00 F8 MOV C,A 00 F8 00

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MOV D,B 00 F8 00 F8 HLT

- 3. MVI C, 65H MVI A, 92H OUT PORT1 ;Display 92H MOV A, C OUT PORT0 ;Display 65H HLT
- IN 07H
 OUT 00H ;Display data from input port 07H
 IN 08H
 MOV B, A ;Store data from port 08H
 HLT
- 5. 82H
- 6. 80H
- Both will be 80H

Section 6.2: Arithmetic Operations

8. A B S Z CY

MVI A,F2H F2 FF NA NA NA MVI B,7AH F2 7A NA NA NA ADD B 6C 7A 0 0 1
OUT PORTO 6C 7A NA NA NA HLT

- The instruction ADD A will add the content of the accumulator to itself; this is equivalent to multiplying by 2.
- 10. The instruction SUB A will clear the accumulator. The flag status will be: CY = 0, Z = 1.
- 11. A C S Z CY XX XX 0 0 0 0

MVI A,5EH 5E XX NA NA NA ADI A2H 00 XX http://jntū.bldg.com

13. MVI A, 3AH ADI 48H OUT PORT# HLT

The instruction DCR does not set the CY flag.



OUT PORT1

The S flag has no significance when subtracting unsigned numbers. If the the CY flag is set, it indicates a negative result.

1 0001 1001 = 19H Borrow flag is deleted by the CY of the result.

(PORT0) = B5H and (PORT1) = 19H

19.

If a number is added before clearing the accumulator, the result will include the residual contents of the accumulator.

Section 6.3: Logic Operations

- 20. The instruction XRA A clears the accumulator, and the flag status is: CY = 0, Z = 1.
- 21. The instruction ADD B sets the CY flag, but the instruction ORA A resets the CY flag.

A = 00 S = 0 Z = 1 CY = 0



The instruction ORA A will set the flag without affecting the content of the accumulator.

24. A B S Z CY

XRA A 00 0 1 0 MVI B,4AH 00 4A NA NA NA SUI 4FH B1 4A 1 0 D ANA B 00 4A 0 1 1 HLT

- MVI C, A8H
 MOV A, C
 ANI 0FH ;Masking byte to mask D7-D4
 OUT PORT0
 HLT
- 27. MVI D, 8EH
 ANI 0FH ;Mask D7-D4
 MOV D, A ;Save in D
 MVI E, F7H
 ANI 0FH ;Mask D7-D4 of second byte
 XRA D ;Exclusive OR masked bytes
 OUT PORT0
 HLT

MVI B, 91H
MVI C, 87H
MOV A, B
ANI 01H
MOV B, A
Save D0 from first byte
MOV A, C
ANI 01H
ANA B
AND bitshttp://intel.biog.com/H

http://jntu.blog.com OUT PORT1 ;Turn on/off light connected to D0 HLT

29. IN 07H
CMA ;Complement data from port 07H
ORA A ;Set Z flag if all switches are open
;Continue

Section 6.4: Branch Operations

- 30. 00
- 31. 28H
- 32. In the following program, explain the range of bytes that will be displayed at PORT2 for various values of BYTE1.

MVI A, BYTE1 MOV B, A SUI 50H JC DELETE MOV A, B SUI 80H JC DSPLAY

DELETE: XRA A

OUT PORT1

HLT

DSPLAY: MOV A, B

OUT PORT2

HLT

In this problem, all bytes from 50H to 7FH will be displayed at POPRT2.

- 33. The address of the output port = F2H. All positive signed numbers and zero will be displayed at port F2H.
- 34. : 00
- This routine displays the absolute value (magnitude) of BYTE1.
- 36. 59H

37. MVI D,9BH START: MVI E,A7H MOV A,D ADD E JC DSPLAY OUT OOH HLT DSPLAY: MVI A,01H OUT OOH HLT 38. XRA A ;Clear CY MVI B, FFH INR B MOV A, B JNC DSPLAY MVI A, 01H DSPLAY: OUT PORT# ;The output = 00H because INR does not HLT ;set CY flag. To clear the CY flag, the instructions such as ANA A, SUB A, ORA A can be used instead of the instruction XRA A. 39. ORA A ;Clear CY MVI C, FFH MOV A, C ADI 01H JNC DSPLAY MVI A, 01H DSPLAY: OUT PORT# ;The output = 01H because ADI sets CY HLT ;flag. 40. MVI B, BYTE1 MVI C, BYTE2 MOV A, B SUB C JNC DSPLAY ;Jump if result is positive CMA ;Take one's complement ADI 01H ;Find two's complement DSPLAY: OUT PORT1

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HLT

Section 6.6: Debugging a Program

41. Reference: Section 6.53

2000		ORG 2000H
2000 DBF1	START:	IN 0F1H
2002 47		MOV B, A
2003 DBF2		IN 0F2H
2005 E680		ANI 80H
2007 4F		MOV C, A
2008 78		MOV A, B
2009 E680		ANI 80H
200B A1		ANA C
200C C21720		JNZ SHTDWN
200F 78		MOV A, B
2010 E61F		ANI 1FH
2012 D3F3		OUT 0F3H
2014 C30020		JMP START
2017 3E40	SHTDWN:	MVI A, 40H
2019 D3F3		OUT 0F3H
201B 76		HLT
201C		END

42. In the following program, the instructions IN F1 and IN F2 are replaced by loading two data bytes 97H and 85H in registers D and E respectively.

2000	ORG 2000H	
2000 1697 STAR	T: MVI D,BYTE1 ;Si	imulate data from port F1
2002 1E85	MVI E,BYTE2 ;Si	mulate data from port F2
2004 7B	MOV A,E	1.70
2005 E680	ANI 80H ;Mask	S6'-S0'
2007 5F	MOV E,A ;Save	S7'
2008 7A	MOV A,D .	
2009 E680	ANI 80H	;Mask S6-S0
200B A3	ANA E	;Check S7 & S7'
200C C21720	JNZ SHTDWN	;If S7 & S7' are on, jump to ; intiate shut down procedure
200F 7A	MOV A,D	;If not, get data from port F1
2010 E61F	ANI 1FH	;Mask bits D7-D5
2012 D3F3	OUT PORT	;Turn conveyer belts
2014 C30020	JMP START	
2017 3E40 SHTDWN	MVI A,40H	;Set bit D6=1
2019 D3F3	OUT PORT	;Turn on emergency
201B 76	HLT	
0097 = BYTE	1 EQU 97H	
	h-44 //:4 h-	lan anna

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0085 =	BYTE2	EQU	85H
00F3 =	PORT	EQU	0F3H
201C		END	

43. This program turns on the LED indicator when the switch S7 is on.

2000		ORG 2000H	
2000 DBF1	START:	IN 0F1H	;Comments are same as illustrative
2002 47		MOV B,A	;program omitted here
2003 DBF2		IN 0F2H	
2005 E680		ANI 80H	
2007 4F		MOV C,A	;Save S7'
2008 78		MOV A,B	;Get data from port F1
2009 E680		ANI 80H	
200B CA1420		JZ TURNON	;If S7 =0, turn on belts
200E D3F3		OUT 0F3H	;Turn on LED to indicate S7 is on
2010 A1		ANA C	;Check S7 and S7'
2011 C21C20		JNZ SHTDW	N
2014 78	TURNON:	MOV A,B	
2015 E61F		ANI 1FH	
2017 D3F3		OUT 0F3H	
2019 C30020		JMP START	
201C 3E40	SHTDWN:	MVI A,40H	;Load byte to turn off belts and turn on emergency
201E D3F3		OUT 0F3H	(2) (2) (2) (2) (3) (3) (4) (4) (4) (4) (4) (4) (4) (4) (4) (4
2020 76		HLT	
		END	

CHAPTER 7

The following programs assume the systems R/W memory begins at location 2000H. The symbols XX in the assignments are assumed as memory page 20H.

Section 7.1

See Figures 7.1, 7.2, 7.3, & 7.4: pg. 81

Section 7.2

5. Location 2075H will contain F7H

```
Bhttp://jntu.blog.com
6.
                                                H
                                                     L
     MVI C,FF
                                  FF
     LXI H,2070H
                                  FF
                                                20
                                                     70
      LXI D,2070H
                                 FF
                                      20
                                           70
                                                20
                                                     70
      MOV M,C
                                 FF
                                      20
                                           70
                                                20
                                                     70
      LDAX D
                                 FF
                                      20
                                           70
                                                20
                                                     70
                        FF
      HLT
                        FF
                                  FF
                                      20
                                           70
                                                20
                                                     70
7.
      A = FFH (2070H) = FFH
     2075H and 2076H
8.
```

- 9. A = 00H D = 00H HL = 209FH
- 10. Clears locations 2090H to 209FH
- 11. LXI B,2090H
 SUB A
 MVI D,0FH
 LOOP: STAX B
 INX B
 DCR D
 JNZ LOOP
 HLT
- 12. Infinite loop. DCX instruction does not affect Z flag.
- 13. 7 times.
- DCX instruction does not affect Z flag.

15.	START:	LXI H, 2055H	;Index for data source
		LXI D, 2085H	;Index for data destination, starting at last location
		MVI B, 06H	;Byte counter
	NEXT:	MOV A, M	;Get data byte
		STAX D	;Store data byte
		INX H	;Next location
		DCX D	
	(1 ¹⁰⁾	DCR B	;Next count
		JNZ NEXT	;If counter is not 0, go back to transfer next byte
		HLT	

16. START: LXI H, 205FH ;Index pointing to last source byte LXI D, 2064H ;Index for destination

	` .	MVI B, 10H http://jnt	u.blogtcomunter
	NEXT:	MOV A, M	Get data byte
		STAX D	;Store data byte
		DCX H	;Next location
		DCX D	
		DCR B	;Next count
		JNZ NEXT	
	80	HLT	
17.	START:	LXI H, 2061H	;Index pointing to low-order reading
		LXI D, 2080H	;Index for storing low-order reading
		MVI B, 05H	;Counter for temp. readings
	NEXT:	MOV A, M	Get reading
		STAXD	;Store low-order reading
		INX H	, store ion order reading
		INX H	;Point to next low-order reading
		INX D	,
		DCR B	;Next count
		JNZ NEXT	
		HLT	
18.	START:	MVI B,6	;BYTE COUNT
		LXI H,2050H	;SOURCE
		LXI D,2050H	;DESTINATION
	LOOP:	MOV A,M	GET BYTE
		ORA A	;TEST IT FOR ZERO
		JNZ SKIP	
		STAX D	;NOT ZERO, SO STORE IT
		INX D	**
	SKIP:	INX H	GO ON TO NEXT
		DCR B	
		JNZ LOOP	
		HLT	
10	CT A DT	1377 D 2000011	
19.	START:	LXI D, 2060H	;Index for data
		MVI C, 05H	;Counter for data
	NEXT	MVI B, 00H	;Clear B to store partial sum
	NEXT:	LDAX D	
		ADD B	;Add data byte
10		MOV B, A	;Save partial sum
17		INX D	
		DCR C	;Next count
		JNZ NEXT	D: 1
		OUT PORT1	;Display sum
		HLT	

http://jntu.blog.com

For the given five bytes, the supplied B7H blog.com

20. START: LXI H, 2060H

LXI H, 2060H ;Index for data MVI C, 05H ;Counter for data

MVI B, 00H

;Clear B to store partial sum

NEXT: MOV A, M

ADD B

;If sum > FF, display 01

JC CARRY MOV B, A

;Save partial sum

DCR C ;Next c

JNZ NEXT

;Next count

OUT PORT1

;Display sum

HLT

CARRY: MVI A, 01H

OUT PORT1

HLT

First set display = 01H Second set display = AFH

 Clear register D to count the number of bytes added and insert the instruction INR D after the instruction JC CARRY. Display the contents of register D at PORT2 at the end.

First set display = 3 and Second set display = 5

Section 7.3:

- 22. Locations 2070H to 2074H will contain 01H, 02H, 03H, 04H, AND 05H
- 23. A H L Z CY M 2055H 20 55 NA NA NA LXI H,2055H 20 55 8A MVI M,8AH MVI A,76H 20 55 8A 76 00 20 55 8A ADD M 0 STA 2055H 00 20 55 NA NA NA 00 HLT
- 25. S = 0 Z = 1 CY is not affected (NA).

26. START: LXI D, 2050H

;Set up index pointer

LXI H, 2040H

;Point to counter location in memory

MVI M, 06H XRA A ;Load counter with the count ;Clear accumulator

MOV C, A ;Cle

;Clear C to save partial sum

		MOV B, A	http://jntu;blogacomto save carrys
	NEXT:	LDAX D	;Get data
		ADD C	;Add previous partial sum
		JNC SKIP	;If no carry, do not increment CY register
		INR B	;Increment carry register
	SKIP:	INX D	;Point to next byte
		INX H	SACRET CONTRACTOR TO CONTRACTOR TO CONTRACTOR CONTRACTO
		DCR M	;Next count
		JNZNEXT	;If all numbers are not yet added, get the next byte
		OUT PORT1	;Display sum
		MOV A, B	
		OUT PORT2	;Display CY register
		HLT	
27.	START:	LXI D, 2050F	H ;Set pointer for readings of furnace 1
		LXI H, 2060F	Fig. 3. 3. 3. 3. 3. 3. 3. 3. 3. 3. 3. 3. 3.
		MVI C, 05H	;Set up counter to count five readings
	NEXT:	LDAX D	;Place temp. reading from furnace 1 in A
		MOV B, M	;Place temp. reading from furnace 2 in B
		SUB B	
		JC SHTDWN	
		INX D	;Point to next memory location
		INX H	
		DCR C	;Decrement reading count
		JNZ NEXT	;Get next temp. reading
		MVI A, 01H	;Set bit $D0 = 1$
		OUT PORT1	;Turn on bit D0
		HLT	25 22 621 28 525 88
	SHTDWN	: MVI A, 0FFH	
		OUT PORT1 HLT	;Display FF for emergency
28.	Add ofter	the statement J	C SUTDWAI
20.	Aud, arter	JNZ NOTEQ	
		MVI A,80H	
		OUT PORTI	
	NOTEQ:	continue	
29.	START:	LXI H,2070H	H ;DATA LOCATION
		MVI B,4	:NUMBER OF PAIRS
	LOOP:	MOV A,M	GET FIRST BYTE OF PAIR
		INX H	POINT TO SECOND BYTE
		ADD M	;ADD 2ND TO 1ST
		DCX H	POINT TO LOCATION TO STORE SUM
			;NOTE: DCX AND INX DO NOT AFFECT CARRY
			http://jntu.blog.com

		MOV M,A	;STORE SUM
		INX H	HEPPINT FOR CORRY POSITION
		MVI A,0	GET CARRY INTO A
		ADC A,A	CTOP =
		MOV M,A	STORE IT TO 2ND LOCATION
		INX H	POINT TO NEXT PAIR
		DCR B	;NEXT COUNT
		JNZ LOOP HLT	;IF COUNT IS NOT ZERO, GET NEXT PAIR
30.	START:	LXI H,2070H	;Data location
	,	LXI D,2070H	;Storage location
		MVI B,4	;Number of pairs
	LOOP:	MOV A,M	;Get first byte of pair
		INX H	;Point to second byte
		SUB M	;Subtract 2nd from 1st
		STAX D	;Store sum
		INX D	;Bump sum pointer
		INX H	;Point to next pair
		DCR B	;Decrement count
		JNZ LOOP	;Go back for next operation
		HLT	
31.	START:	LXI H,2070H	;Data location
		LXI D,2070H	;Storage location
		MVI B,4	;Number of pairs
	LOOP:	MOV A,M	;Get first byte of pair
		INX H	;Point to second byte
		SUB M	;Subtract 2nd from 1st
		JC SKIP	;Skip if < 0 (negative) unsigned
		STAX D	;Store sum
		INX D	;Bump sum pointer
	SKIP:	INX H	Point to next pair
		DCR B	;Decrement count
		JNZ LOOP	;Go back for next operation
		HLT	
Section	on 7.4:		
Secul	JII /		
32. (a) $A = 6FH$,	CY = 1 (b) $A = 6$	EH, $CY = 1$
33.		A S	Z CY

33.		A	S	·Z	CY
	MVI A,80H	80			
	ORA A	80	1	0	0
*	RAR	00	0	1	1

34.	(a)	A	CYhttp://hdt	u.blog.com	A	CY
	MVI A,C5H	C5		MVI A, A7H	A7	
	ORA A	C5	0	ORA A	A7	0
	RAL	8A	1	RAR	4E	1
	RRC	45	0	RAL	A7	0

- 35. These instructions will move the MSD of a BCD number (7 in this case) from 10th to the 1's place. A = 0.7
- 36. (a) Multiply by 2 (b) Divide by 4
- 37. Multiply by 10.

38.	START:	MVI D, 0AH	;Set up counter to count ten readings
		LXI H, 2060H	;Set pointer to data location
		MVI C, 00H	;Clear register C to save partial sum
		MOV B, C	;Clear register B for carry
	LOOP:	MOV A, M	;Get current reading
		ORA A	;Set flags
		JM NEXT	;If D7 = 1, reject the data byte
		ADD C	;Add the current reading
		MOV C, A	;Save partial sum of the readings
		JNC NEXT	;If no carry, do not increment (B)
		INR B	;Add 1 to previous carry count
	NEXT:	INX H	;Go to next data location
		DCR D	;Next count
		JNZ LOOP	;Go back to gat next reading
		MOV A, C	;Get the final sum
		OUT PORT1	;Display low-order byte of the sum
		MOV A, B	;Get the carry count
		OUT PORT2	;Display carry count
		HLT	

For the given set of readings in assignment 38, the sum = 2A0H

39.	START:	MVI D, 0AH	;Set up counter to count ten readings
		LXI H, 2060H	;Set pointer to data location
		MVI C, 00H	;Clear register C to save partial sum
		MOV B, C	;Clear register B for carry
		MOV E, C	;Clear regi. E to count positive readings
	LOOP:	MOV A, M	;Get current reading
		ORA A	;Set flags
		JM NEXT	;If D7 = 1, reject the data byte
		INR E	;Positive reading found; count it
		ADD C	;Add the current reading
		la 44 a . /	first a laboratory

;Save partial sum of the readings MOV C, A http://jhfunblogreymdo not increment (B) JNC NEXT INR B ;Add 1 to previous carry count NEXT: INX H ;Go to next data location :Next count DCR D ;Go back to get next reading JNZ LOOP MOV A, C :Get the final sum ;Display low-order byte of the sum OUT PORT1 ;Get the carry count MOV A, B OUT PORT2 ;Display carry count ;Get number for positive readings MOV A, E ;Display number for positive reading **OUT PORT3** HLT ;Set pointer to data storage location LXI D, 2060H 40. START: LXI H, 2050H ;Set pointer to data location MVI C, 08H ;Set counter to count eight bytes MOV A, M ;Get the byte LOOP: ;Set flag if D7 = 1ORA A ;Jump to increment data pointer JM NEXT ;Place D0 in CY position RRC JC NEXT ;Jump to increment data pointer :Restore original byte RLC ;Save the byte at storage location STAX D ;Next storage location INX D :Next data location INX H NEXT: DCR C :Next count Jump back to get next byte JNZ LOOP HLT The following memory locations should have the data bytes as shown. Answer: 2060 = 52H2061 = 78H2062 = 62HSection 7.5: 41. A S Z CY MVI A,7FH 7F 7F ORA A 0 0 CPI A2H

```
42.
```

A S http://jntu.blog.com

LXI H,2070H MVI M,64H MVI A,8FH 8F CMP M 8F 0 0 0

- 43. 00, 00, 7A, 87, 00, 00
- 44. 100 (64H) < BYTE <= 200 (C8H)
- 45. 20, 64, 8F
- 46. If PORT1 < 32 (20H), CY flag is set.</p>
 or if PORT1 = > 160 (A0H), Minus flag is set.

47.	START:	LXI H, 2050H MVI C, 00H MOV B, C MOV D, C	;Set index to point to data ;Clear register C to save sum ;Clear (B) to save carry
	NXTBYT	: MOV A, M CMP D JZ DSPLAY ADD C JNC SAVE INR B	;Set register D to compare bytes ;Get data ;Check whether this the last byte ;If 'Yes', go to display section ;Add previous sum ;If there is no carry, go to save the sum ;Update carry register
	SAVE:	MOV C,A INX H JMP NXTBYT	;Save sum ;Point to next reading ;Go back to get next reading
	DSPLAY:	MOV A, C OUT PORT1 MOV A, B OUT PORT2 HLT	;Display low-order byte of the sum ;Copy the carry count to accumulator ;Display high-order byte of sum

48. START:

LXI H, 2050H MVI D, 40H ;Set index to point to data location ;Byte to be found in data string

MVI C, 08H

;Set up counter ;Get data byte

NXTBYT: MOV A, M CMP D

; Is the byte = 40H?

JZ DSPLAY ;

If yes, go to display its location

INX H

;Point to next data byte

DCR C

Next count

JNZ NXTBYT

;Jump to get next byte

	DSPLAY	MVI A, FFH OUT PORT1 HLT : MOV A, H OUT PORT1 MOV A, L OUT PORT2 HLT	http://jntublog.tom/0H is not in the set ;Load high-order memory address ;Display memory page number ;Load low-order memory address ;Display memory low-order address
49.	START:	LXI H, 2050H MVI C, 08H MVI B, 00H	;Set index to point to data location ;Set up counter
	NXTBY	CMP B JNC NEXT MOV B, A	;Clear (B) to save the highest reading ;Get data byte ;Is (B) > (A)? ;If yes, replace (B) with (A) ;Save the larger number
	NEXT:	INX H DCR C JNZ NXTBYT	;Point to next data byte ;Next count ;Jump to get next byte
		MOV A, B OUT PORT1 HLT	;Load the largest byte ;Display the largest byte in the string
50.	START:	LXI H, 2050H MVI C, 08H MVI B, 00H	;Set index to point to data location ;Set up counter ;Clear (B) to save the highest reading
	NXTBYT:	MOV A, M CMP B JC NEXT	;Get data byte ;Is (B) < (A)? ;If yes, replace (B) with (A)
	NEXT:	MOV B, A INX H DCR C JNZ NXTBYT MOV A, B OUT PORT1 HLT	;Save the larger number ;Point to next data byte ;Next count ;Jump to get next byte ;Load the largest byte ;Display the largest byte in the string
51.	START:	LXI H,2050H LXI D,2050H	;SOURCE POINTER ;SAVE POINTER
	LOOP:	MVI B,10 MOV A,M CPI 60	;BYTE COUNT
		JC REJECT CPI 101 JNC REJECT	;REJECT IF < 60 ;NOTE: IF SUBTRACT 100, WOULD REJECT 100 ;REJECT IF > 100

	10	STAX D INX D	http://jntu.blog.com SAVE IT
	REJECT	DCR B JNZ LOOP	;LOOP FOR NEXT
		HLT	14 H
52.	START:	LXI H,2050 LXI D,2050 MVI B,10 MVI C,0	
	LOOP:	MOV A,M CPI 60	
		JC REJECT	
		CPI 101	;NOTE: IF SUBTRACT 100, WOULD REJECT 100
		JNC REJECT	
		INX D	;OK, SO SAVE IT
		INR C	;AND COUNT IT
	REJECT:		;LOOP FOR NEXT
		DCR B	
		JNZ LOOP MOV A,C	DIGDI AM GOVERN
		OUT PORTI	;DISPLAY COUNT
53.	START:	LXI H,2070H	;SOURCE POINTER
		LXI D,2090H	SAVE POINTER
	LOOP	MOV A,M	E1
		CPI 0DH JZ ENDS	;CHECK FOR END OF STRING
		CPI 30H	;IF =0DH, THEN END OF SRTING
		JC REJECT	;REJECT IF < 30H
		CPI 3AH	;NOTE: IF SUBTRACT 39H, WOULD REJECT 39H
		JNC REJECT	;REJECT IF > 39H
		STAXD	;OK, SO SAVE IT
	REJECT:	INX D INX H	;LOOP FOR NEXT
	100001.	JMP LOOP	,LOOF FOR NEXT
	ENDS:	HLT	
54.	START:	LXI H,2070H LXI D,2090H	;SOURCE POINTER ;SAVE POINTER
	LOOP:	MVI C,0 MOV A,M	
			http://intu.blog.com

http://jntu.blog.com

CPI ODH CHECK FOR END OF STRING JZ ENDS http://intu.blog.comHEN END OF SRTING CPI 30H JC REJECT ;REJECT IF < 30H CPI 3AH ;NOTE: IF SUBTRACT 39H, WOULD REJECT 39H JNC REJECT ;REJECT IF > 39H STAXD ;OK, SO SAVE IT INX D INR C ;AND COUNT IT REJECT: INX H ;LOOP FOR NEXT JMP LOOP ENDS: MOV A,C ;DISPLAY COUNT OUT PORT1 HLT 55. START: LXI H,XXXXH SOURCE POINTER FOR SCANNER READINGS MVI C.0 COUNTER TO COUNT SETS LOOP: MOV A,M GET SCANNER READING ORA A ;SET ZERO FLAG JZ ENDS ;IF BYTE =0, THEN END OF STRING CPI A3H ;IS IT A 19" SET JNZ SKIP;IF NO, DO NOT COUNT IT INC C;COUNT THE SET SKIP: INX H JMP LOOP ENDS: HLT

56. This assignment is similar to the Illustrative Program 7.53 except that the order is descending and the number of bytes = 10. Therefore, the counter register C should be set for 9 and the instruction JC NXTBYT should be replaced by JNC NXTBYT.

Section 7.6:

- 57. (a) 256 (b) 2090H-2091H
 - (c) 1) A was not initialized to 00 2) Missing INX H in loop
- 58. SET 1 SET 2
 - (a) E7H 9BH
 - (d) DCR B hould be DCR C and add the instruction MOV A,B just before (STA 2070H) storing the answer

(a) 039CH

- (a) 016FH
- (d) Change JNZ 2008H to JNZ 2007H
- (d) ADC M should be ADD M

CHAPTER 8

- 1. 1168 uSec
- 234.67 mSec
- 468.584675 mSec
- Count 12FFH = 1X16 + 2X16 + 15X16 + 15X16

=4863

The delay in the loop:

- = T-states X Clock Period X Count
- = 64 X (.33X10) X 4863
- = 102 ms

The above calculations are based on the assumption that the JMP instruction takes 10 T-states in the last iteration, and the initial instruction LXIB is not part of the delay loop. To account for the 7 T-states in the last iteration, the delay of 0.99 micro-sec (.33x10 X 3 = .99 X 10) should be subtracted from the above delay.

- 234.1313 mSec
- 6. (a) 4
 - (b) infinite
 - (c) infinite or just once if the flag is set initially.
- 7. (a) infinite (b) infinite (c) 1
- 8. 1.465 mSec
- If the system frequency is 3.072 MHz, the clock period will be 325 ns. This will reduce the delay to .325 s.
- 10. COUNT = 0EH

- 11. BC = 34965₁₀ Insignificant difference when the delay is calculated with JNZ = 7
 T-states in the last iteration.
 http://jntu.blog.com
- 12. 12799 (31FFH)

Note: In the following assignments, the delay calculations are based on the SDK-85 system with the frequency of 3.072 MHz (T = 325.5 ns).

		Mnemonics	Comments	T-states
13.	START:	MVI B, 00H	;Initialize counter	
	DSPLAY:	OUT PORT1	;Display count	10
		LXI D, COUNT	;Load delay count	10
	LOOP:	DCX D		6
		MOV A, E		4
		ORA D	;Set Z flag if (D) & (E) = 0	4
		JNZ LOOP	;If $Z = 1$, repeat the loop	10/7
	- 2	INR B	:Next count	4
		MOV A, B	, a tone obtain	4
		CPI 21H	:Is count = 21H?	7
		JNZ DSPLAY	;If not, go to display count	10/7
		JMP START	;Reset counter, start again	

Calculations for 100 ms delay:

Delay outside the loop:

 $To = 45 \text{ T-states } \times 325.5 \text{ ns} = 14648 \text{ ns}$

Delay within the loop:

T = 24 T-states x 325.5 ns x COUNT

Total Delay T = To + T

100 ms = 14648 ns + (24 x 325.5 ns x COUNT)

COUNT =
$$\frac{100 \times 10 - 14648 \times 10}{7812 \times 10} = 12,798 = 31\text{FEH}$$

14.	START:	MVI B, 00H	http://jntu.blog.comounter
	DSPLY1:		,—————————————————————————————————————
		MVI C, COUNT	1 :COUNT1 = 10
	LOOP1:	LXI D, COUNT	
	LOOP2:	DCX D	,
		MOV A, E	
		ORA D	
		JNZ LOOP2	
		DCR C	
		JNZ LOOP1	
		INR B	
		MOV A, B	190 m
		CPI 0AH	;Check whether 9 is displayed
		JNZ DSPLY1	,
		DCR B	
	DSPLY2:	OUT PORT1	
		MVI C, COUNT	1 ;COUNT1 = 10
	LOOP3:	LXI D, COUNT?	;COUNT2 should provide 100 ms delay
	LOOP4:	DCX D	
		MOV A, E	
		ORA D	
		JNZ LOOP4	
		DCR C	
		JNZ LOOP3	
		MOV A, B	
		ORA A	;Set Z flag if $(A) = 0$
		JNZ DSPLY2	
		JMP START	
	For delay o	alculations and A	sei anno ent #7

For delay calculations, see Assignment #7.

15.	START:	MVI D, 00H	;Load bit pattern	T-States
	ROTATE:	MOV A, D		4
		CMA	;Complement bit pattern	4
		MOV D,A		4
		ANI 80H	;Mask D6-D0	7
		OUT PORT1		10
		MVI E, COUNT1		7
	LOOP1:	LXI B, COUNT2		10
	LOOP2:	DCX B		6
		MOV A, C		4
		ORA B		4
		JNZ LOOP2	•	10/7
		DCR E		4
		JNZ LOOP1		10
		JMP ROTATE ht	tp://jntu.blog.com	10

To design 5-second delay using the loop within the loop technique, it is assumed that the LOOP2 will provide 100 ms delay, and it will be repeated 50 times by using the COUNT1 = 32H in the outer loop.

LOOP1 Delay: =
$$24T \times (325.5 \text{ ns}) \times 50$$

= 0.39 ms

Delay Outside the loop: $= 39T \times (325.5 \text{ ns}) = 0.0127 \text{ ms}$

Total Delay = LOOP1 + LOOP2 x 50 + Outside Delay

5 s = .39 ms + (24 x 325.5 ns x COUNT2) 50 + .0128 ms

COUNT2 = 12799.78

The COUNT2 can be calculated by ignoring all the delays except in the LOOP2 as follows:

$$5 s = LOOP2 \times 50$$

 $= (24 \times 325.5 \text{ ns } \times \text{COUNT2}) \times 50$

COUNT2 = 12800

The difference between the two calculations is 0.22; it is insignificant.

16.	START:	MVI D, 00H	;Load bit pattern	
	ROTATE:	MOV A, D	Supplied the state of the state	4
		CMA	;Complement bit pattern	4
		MOV D,A		4
		ANI 01H	;Mask D7-D1	7
		OUT PORT1		10
		MVI B, COUNT		7
	LOOP:	DCR B		4
		JNZ LOOP		10/7
	88	JMP ROTATE		10

Total Delay T = To + T

200 s = $46 \times 325.5 \text{ ns} + 14 \times 325.5 \text{ ns} \times \text{COUNT}$

COUNT = 42.5 ≈ 42

17. TURNON: MVI A, 01H ;Bit pattern to turn on D0 7
OUT PORT1 ;On-period begins 10

	7.53	MVIB COINHIP:	/intublog.com s pulse	7
	LOOP1:	DCR B	Count for 200 s puise	4
		JNZ LOOP1		10/7
			pattern to turn off D0	7
		OUT PORT1 ;O:	ff-period begins	10
			Count for 400 s delay	7
	LOOP2:	DCR B	Count for 400 3 delay	4
		JNZ LOOP2		10/7
		JMP TURNON		10//
		Jiii Tolatoit		10
	On-period	Delay:		
	T = To +	T		
	200 = (24)	4 T x 325.5 ns) + (14	T x 325.5 ns x COUNT1)	
	COLDINA	(200 - 7.81) x 10		
	COUNTI	== = 4	2	
		(4.557 x 10)		
	Off-period	Delay:		100
	On penoe	Delay.		
	T = To +	T		
	400 = (34)	$T \times 325.5 \text{ ns}$) + (14 %	Γ x 325.5 ns x COUNT2)	
		00 - 11.06)		
C		= 85		
	(4.	.557 x 10)		
	CTART	3.00 T 101010100		
	START:	MVI L,10101010B	;ALTERNATING LIGHT PATTERN	V
	LIGHTS:	MOV A,L		
		RRC		
		OUT PORT		
		MOV L,A		
	OLITED		$;20 \times 50 \text{ mSec DELAY} = 1 \text{ Sec}$	
	OUTER:	LXI D,2559	;20 mSec DELAY	
	INNER:	DCX D		
		MOV A,D	50	
		ORA E		
		JNZ INNER	*** S	
		DCR B	672	
		JNZ OUTER		
		JMP LIGHTS		

18.

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