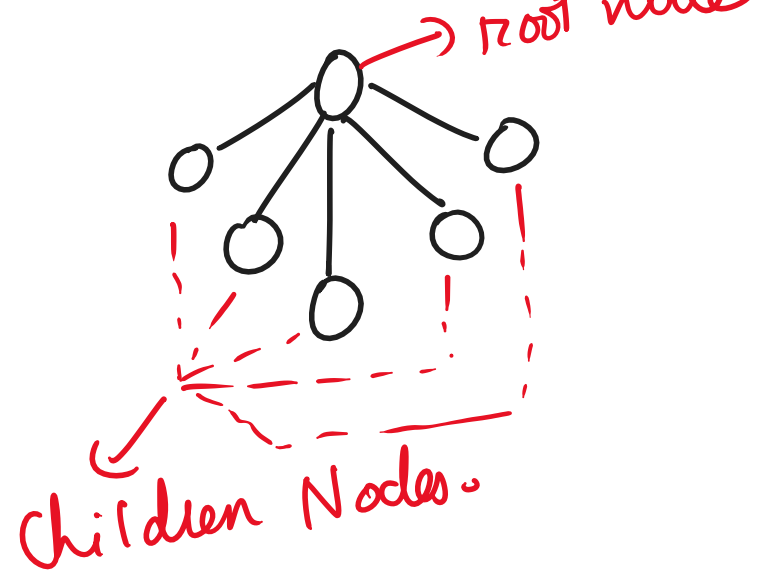


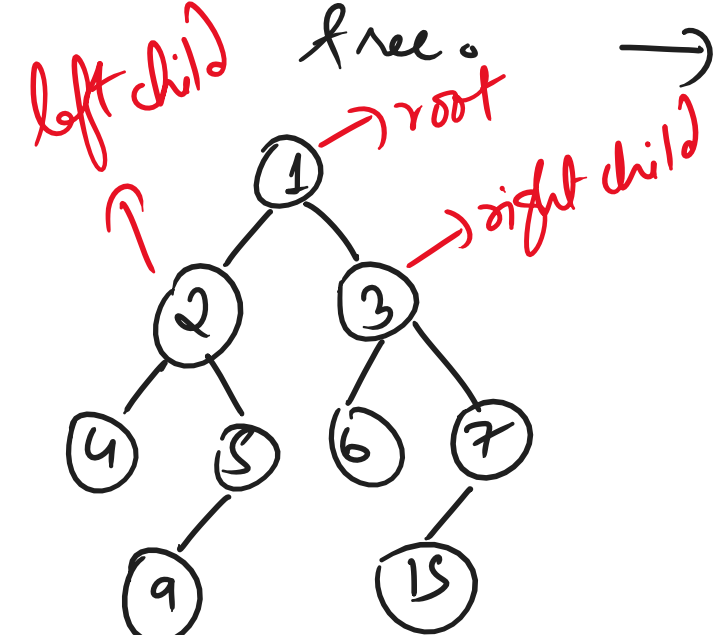
Introduction to Trees : \rightarrow

* Normal Tree : \rightarrow an entity containing nodes and children.



* The nodes are pointers having an address & some data.
* The starting node is called the "root" node.

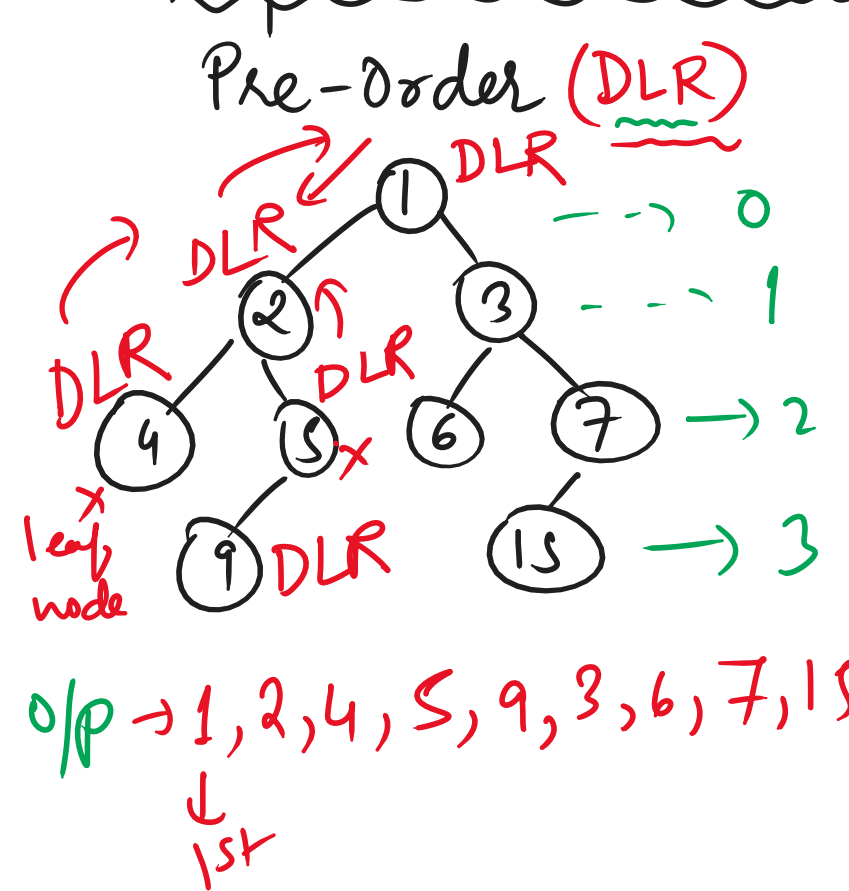
* Binary Tree : \rightarrow A tree having at most two children nodes (left & right) is called a binary tree.



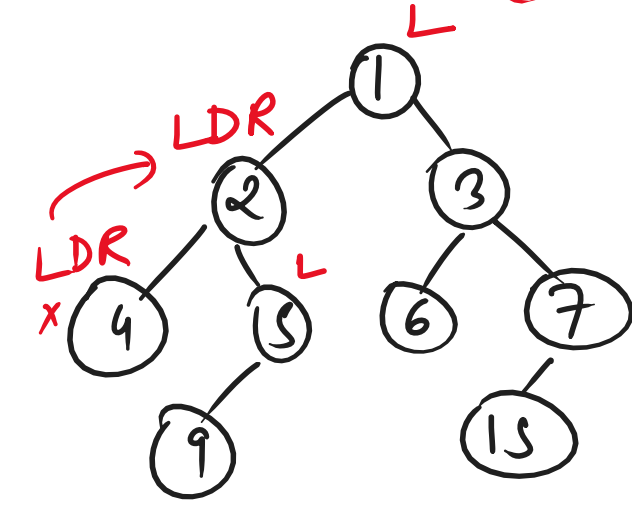
* Traversals : \rightarrow

- \rightarrow Depth First Traversal
- \rightarrow Breadth First Traversal (Level Order Traversal)

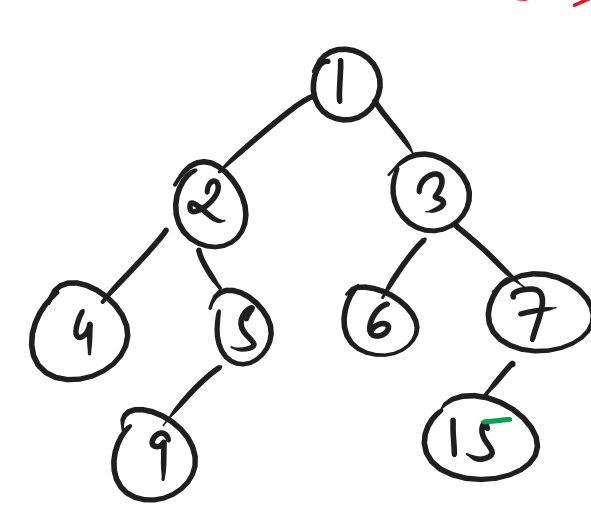
Depth First Traversal : \rightarrow (Recursion) \leftarrow (DFS)



* In-Order (LDR)

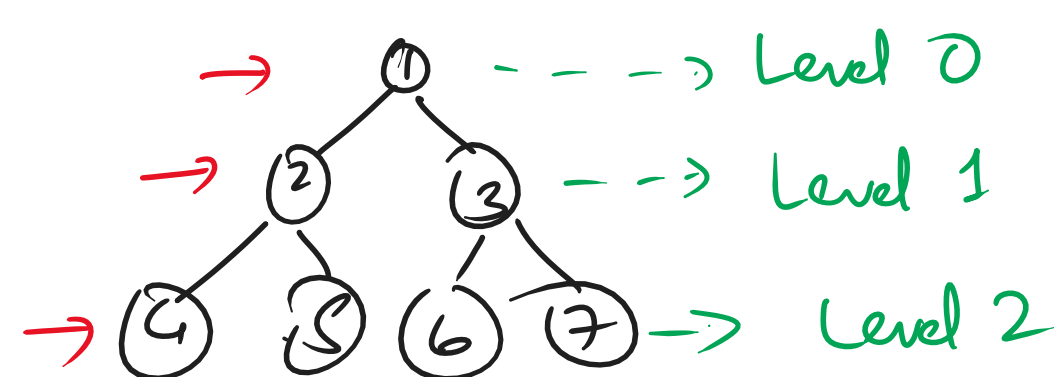


Post-Order (LRD)



The tree should be constructed Top to Bottom & left to right (Level By Level)

* Level Order Traversal : \rightarrow BFS (Queue)

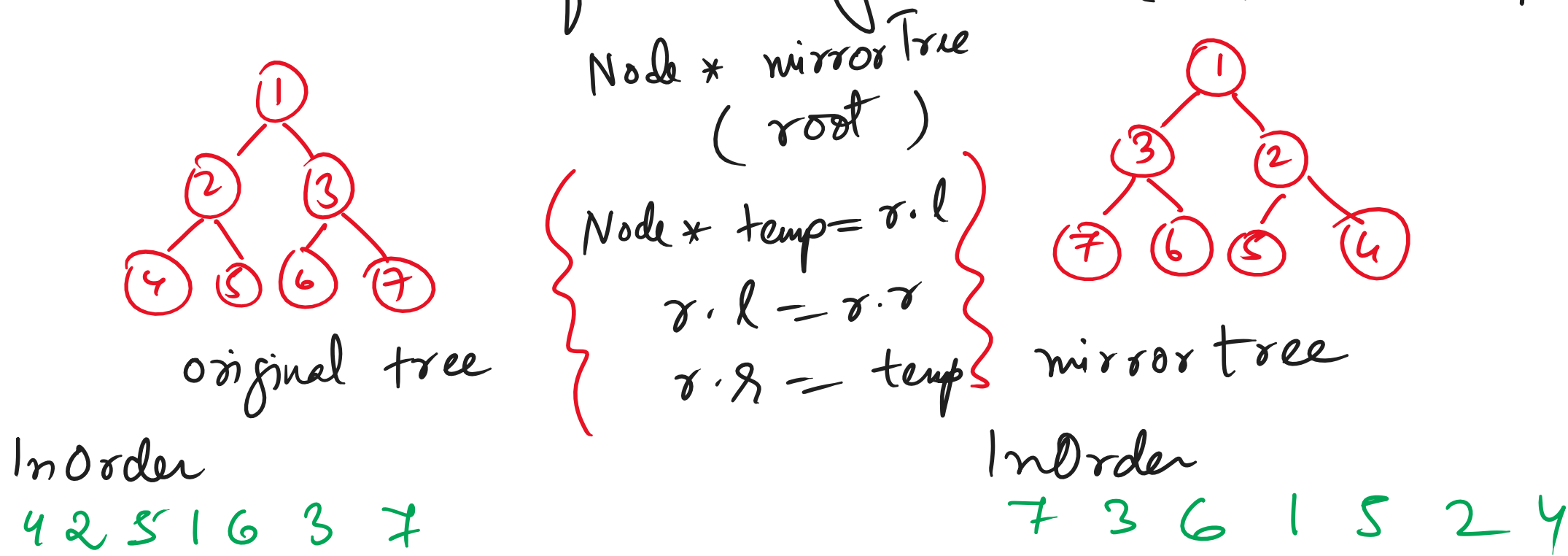


O/p is always top to bottom & left to right :

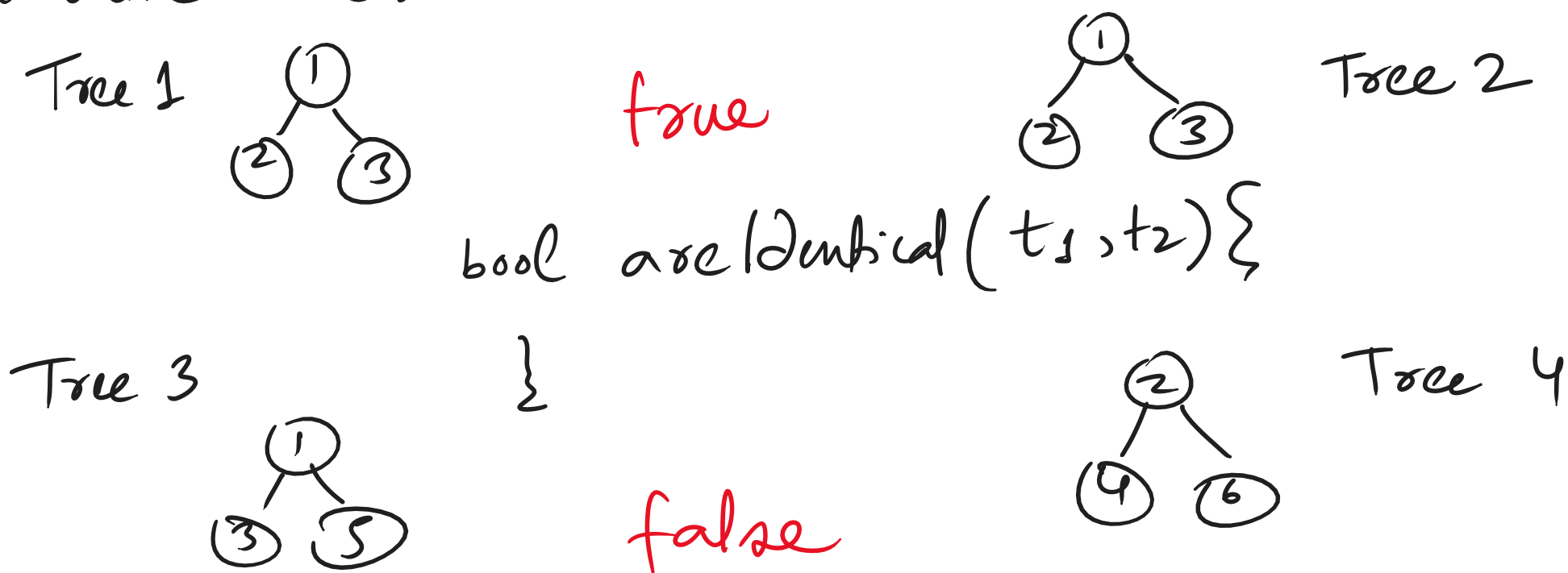
\rightarrow 1, 2, 3, 4, 5, 6, 7

Binary Tree Important Interview Questions : \rightarrow

* Mirror of a Binary Tree (TCS/Accenture/Infosys)

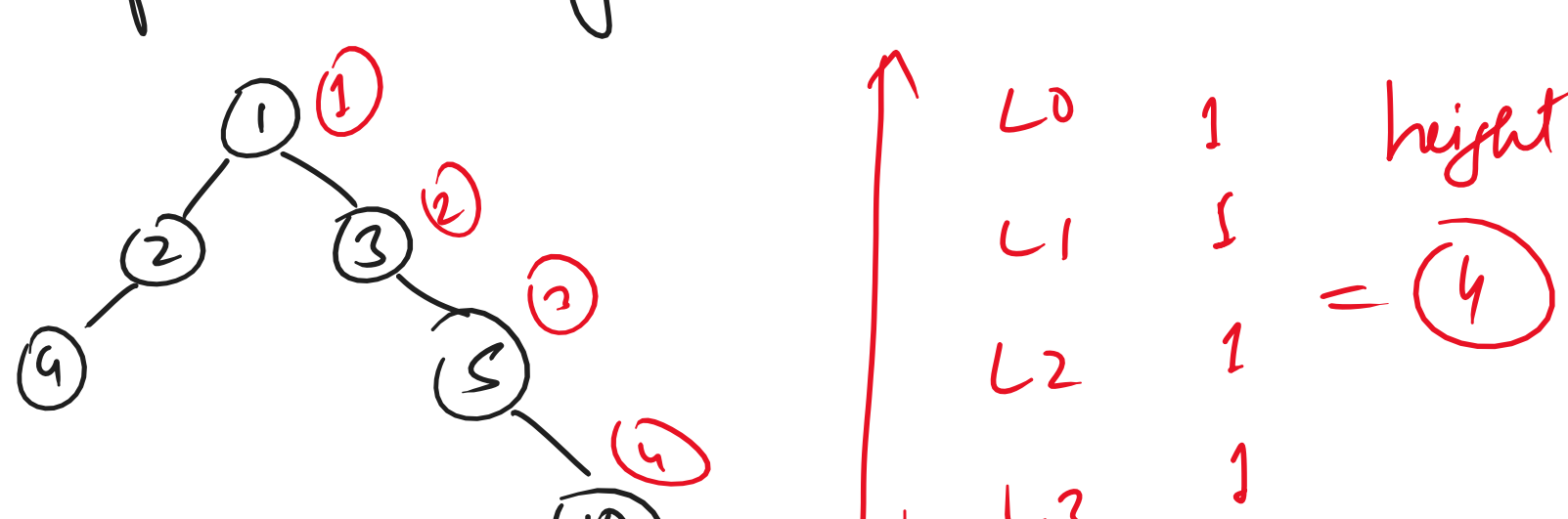


* Given two binary trees T1 & T2 write a function to determine whether they are identical or not.



$(t1 == null \&\& t2 == null) \rightarrow true$ because empty trees
 $(t1 == null \parallel t2 == null) \rightarrow false$ one is empty
 $(t1.data != t2.data) \rightarrow false$
 recursion

* Height of a Binary Tree : \rightarrow

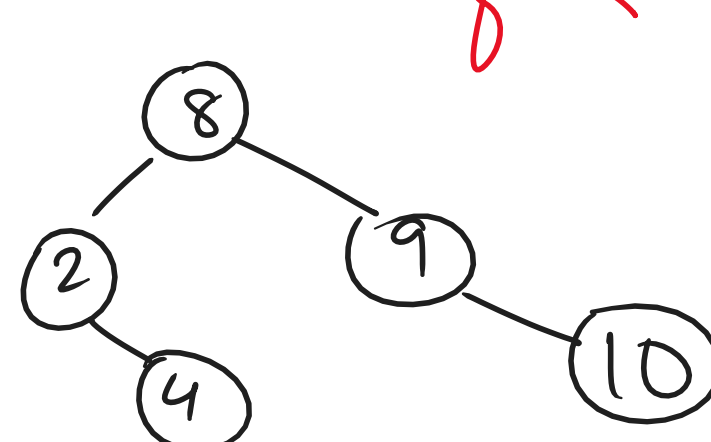


The maximum distance from the root node to any of its leaf nodes.

Binary Search Tree :

Every Node has a very important property.

Left < Node < Right (L < N < R) (\sqrt{n})

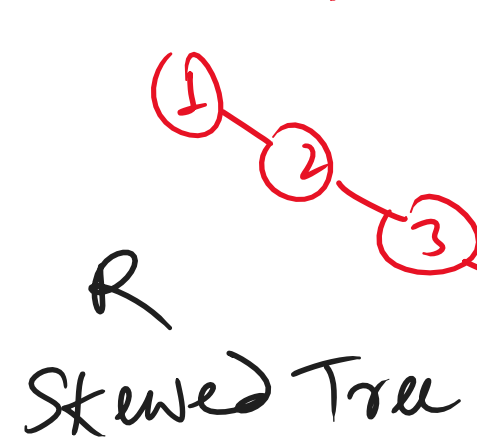


Time Complexity of search & insert = $\log(N)$

1, 2, 3, 4, 5

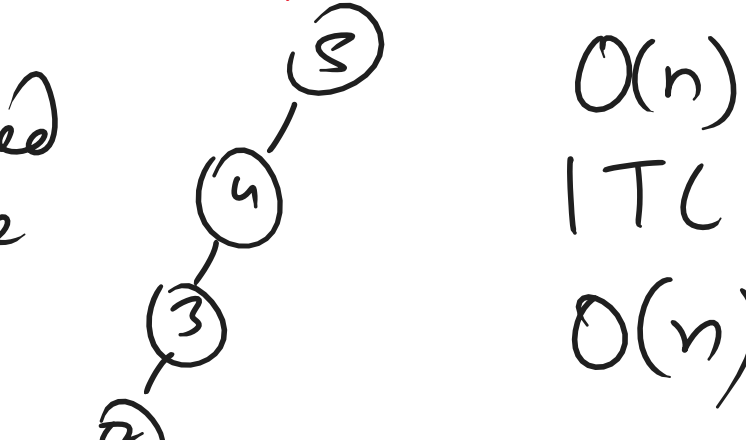
h = 5 = n

5, 4, 3, 2, 1 STC



1
2
3
4
5

L Skewed Tree



$O(n)$

ITC

$O(n)$