L4850: A Dynamically Typed Language for Instruction

L4850 is a programming language designed for practice interpretation. L4850 is a simple dynamically typed language that contains elements of functional, object-oriented and imperative paradigms. L4850 is intended to be simple enough to interpret in a single semester by any student willing to put in some effort. Each feature included in the language was added specifically to help students learn about the meaning of programs written in different paradigms. The goal of this project is to help student understand what it means to express something in a language by writing an operational semantics (fancy words for interpreter). In the process, I hope that you will understand object-oriented design better and also learn about functional programming.

1 L4850 Syntax

1.1 Lexical Properties of L4850

- 1. In L4850, blanks are significant.
- 2. In L4850, all keywords are reserved; that is, the programmer cannot use an L4850 keyword as the name of a variable. The valid keywords are:

```
⟨ASSIGN⟩

ightarrow assign
\langle COND \rangle
                                  \rightarrow cond
(DEFCLASS)

ightarrow defclass
(DEFUNC)
                                  \rightarrow defunc
\langle \mathbf{ELSE} \rangle
                                         else
\langle \mathbf{FALSE} \rangle
                                         false
\langle \mathbf{FI} \rangle
                                         fi
\langle \mathbf{FUNC} \rangle
                                  \rightarrow func
\langle \mathbf{IF} \rangle

ightarrow if
\langle LOAD \rangle
                                  \rightarrow load
\langle \mathbf{METHOD} \rangle

ightarrow method

ightarrow new
\langle NEW \rangle
\langle THEN \rangle
                                       then
\langle \mathbf{OT} \rangle
                                         to
\langle \mathbf{TRUE} \rangle
                                         true
\langle VARS \rangle
                                        vars
\langle WITH \rangle

ightarrow with
```

(Note that L4850 is *case sensitive*, that is, the variable X differs from x. Thus, cond is a keyword, but COND can be a variable name.)

3. The following special characters have meanings in an L4850 program.

```
\langle AND \rangle
                                                       \rightarrow &&
\langle COMMA \rangle
\langle \mathbf{DIVIDE} \rangle
                                                       \rightarrow /
\langle \mathbf{DOT} \rangle
\langle \mathbf{EQUAL} \rangle
                                                       \rightarrow ==
\langle \mathbf{GREATER} \rangle
                                                       \rightarrow >
                                                       → >=
(GREATEREQUAL)
⟨INVOKE⟩
                                                       \rightarrow ->
\langle LB \rangle
\langle LBK \rangle
                                                       \rightarrow [
\langle LESS \rangle
                                                       \rightarrow <
\langle LESSEQUAL \rangle
                                                       \rightarrow <=
\langle \mathrm{LP} 
angle
                                                       \rightarrow (
(MINUS)
\langle \mathbf{MULTIPLY} \rangle
\langle NOT \rangle
                                                       \rightarrow !
\langle NOTEQUAL \rangle
                                                       \rightarrow !=
\langle \mathbf{OR} \rangle
                                                       \rightarrow \square
\langle PLUS \rangle
\langle RB \rangle
\langle RBK \rangle
                                                       \rightarrow ]
\langle \mathbf{RP} \rangle
                                                       \rightarrow )
```

- 4. Comments are delimited by a //. All characters following the //on the same line are part of the comment.
- 5. Identifiers are written with upper and lowercase letters and are defined as follows:

6. Constants are defined as follows:

1.2 Context-free Grammar

The following grammar describes the context-free syntax of L4850:

```
\mathbf{program} \qquad \quad \rightarrow \quad ( \ \mathbf{functionDef} \ \\
```

classDefexpressionloadFile)+

 $\mathbf{functionDef} \qquad \rightarrow \quad \langle \mathbf{DEFUNC} \rangle \ \langle \mathbf{ID} \rangle \ \langle \mathbf{LP} \rangle \ (\ \mathbf{idList} \)? \ \langle \mathbf{RP} \rangle \ \mathbf{expressionList}$

 $\mathbf{idList} \qquad \qquad \rightarrow \quad \langle \mathbf{ID} \rangle \; (\; \langle \mathbf{COMMA} \rangle \; \langle \mathbf{ID} \rangle \;)^*$

 $classDef \rightarrow \langle DEFCLASS \rangle \langle ID \rangle \langle LB \rangle (classVars)? (methods)? \langle RB \rangle$

 $\mathbf{classVars} \qquad \quad \rightarrow \quad \langle \mathbf{VARS} \rangle \ \mathbf{idList}$

methods \rightarrow ($\langle METHOD \rangle \langle ID \rangle \langle LP \rangle$ (idList)? $\langle RP \rangle$ expressionList)⁺

loadFile $\rightarrow \langle LOAD \rangle \langle STRING \rangle$

expressionList $\rightarrow \langle LB \rangle$ (expression)⁺ $\langle RB \rangle$

expression \rightarrow compExpr $(\log Op compExpr)^*$

 $|\langle NOT \rangle compExpr$

 $\log \mathrm{Op} \hspace{1cm} o \hspace{1cm} \langle \mathrm{OR} \rangle \hspace{1cm} | \hspace{1cm} \langle \mathrm{AND} \rangle$

 $compExpr \rightarrow addExpr (compOp addExpr)^*$

 $\mathbf{compOp} \qquad \qquad \rightarrow \quad \langle \mathbf{EQUAL} \rangle \ \mid \ \langle \mathbf{NOTEQUAL} \rangle \ \mid \ \langle \mathbf{LESS} \rangle \ \mid \ \langle \mathbf{LESSEQUAL} \rangle$

 $|\langle GREATER \rangle|\langle GREATEREQUAL \rangle|$

 $\mathbf{addExpr} \qquad \quad \rightarrow \quad \mathbf{mulExpr} \; (\; \mathbf{addOp} \; \mathbf{mulExpr} \;)^*$

 $\mathbf{addOp} \qquad \qquad \rightarrow \ \, \langle \mathbf{PLUS} \rangle \ \mid \ \langle \mathbf{MINUS} \rangle$

 $mulExpr \rightarrow factor (mulOp factor)^*$

 $\mathbf{mulOp} \qquad \qquad \rightarrow \quad \langle \mathbf{MULTIPLY} \rangle \; \mid \; \langle \mathbf{DIVIDE} \rangle$

```
factor \rightarrow operand (call)?
```

$$operand \rightarrow varRef \mid constant \mid newExpr \mid ifExpr$$

call $\rightarrow \langle INVOKE \rangle \langle LP \rangle$ (paramList)? $\langle RP \rangle$

paramList \rightarrow expression ($\langle COMMA \rangle$ expression)*

 $\mathbf{varRef} \qquad \rightarrow \langle \mathbf{ID} \rangle \; (\; \langle \mathbf{DOT} \rangle \; \langle \mathbf{ID} \rangle \;)?$

 $constant \qquad \rightarrow \ \, \langle INTUM \rangle \, \mid \, \langle FLOATNUM \rangle \, \mid \, list \, \mid \, \langle STRING \rangle$

 $|\langle TRUE \rangle | \langle FALSE \rangle$

 $list \hspace{1cm} \rightarrow \hspace{1cm} \langle LBK \rangle \hspace{1cm} constantList \hspace{1cm} \langle RBK \rangle$

 $constantList \rightarrow constant (\langle COMMA \rangle constant)^*$

 $\mathbf{newExpr} \qquad \rightarrow \ \, \langle \mathbf{NEW} \rangle \, \, \langle \mathbf{ID} \rangle$

 $if Expr \hspace{1cm} \rightarrow \hspace{1cm} \langle IF \rangle \hspace{1cm} expression \hspace{1cm} \langle THEN \rangle \hspace{1cm} expression \hspace{1cm} \langle ELSE \rangle \hspace{1cm} expression \hspace{1cm} \langle FI \rangle$

funcExpr $\rightarrow \langle FUNC \rangle \langle LP \rangle (idList)? \langle RP \rangle expressionList$

 $assignExpr \rightarrow \langle ASSIGN \rangle expression \langle TO \rangle \langle ID \rangle$

 $condExpr \rightarrow \langle COND \rangle condClauses$

condClauses \rightarrow ($\langle LB \rangle$ expression expression $\langle RB \rangle$)⁺

with Expr $\rightarrow \langle WITH \rangle \langle LP \rangle$ variable Defs $\langle RP \rangle$ expression List

variableDefs \rightarrow ($\langle LBK \rangle \langle ID \rangle$ expression $\langle RBK \rangle$)*

2 L4850 Notes

2.1 Function Declarations

The semantics of function definition cause the variable to be added to the environment with a closure as a value.

Example:

```
defunc even (n) {
    if (n == 0) then
        true
    else
        odd->(n - 1)
    fi
}
defunc odd(n) {
    if (n == 0) then
        false
    else
        even->(n - 1)
    fi
}
```

In the above program, the variables **even** and **odd** are added to the environment with closures for values. The value of a function definition is a closure.

2.2 Assignment Expressions

A variable is given a new value using an assign expression.

```
assign 2+3 to x
```

In the previous expression, x is given the value 5 and 5 is the return value of the entire assign expression.

2.3 If Expression

If the first expression in an if is true, then the expression in the then-part is evaluated and returned. Otherwise the expression in the else-part is evaluated and returned.

```
if x == 5 then 5 else 4 fi
```

The value returned by this expression is 5.

2.4 Cond Expression

The cond expression is a concise format for a sequence of if-then-else-if expressions. For example,

```
if x == 5 then
   5+2
else
   if x == 6 then
   5+3
   else
      if x == 7 then
      5+4
      else
      5+5
      fi
   fi
fi
```

can equivalently be expressed as

```
cond
  {x == 5 5+2}
  {x == 6 5+3}
  {x == 7 5+4}
  {true 5+5}
```

2.5 With Expressions

Variables are given scope in L4850 using a with expression. For example,

```
with ([x 5]) {
  x + 10
}
```

gives scope to the variable x in the body of the with expression and gives it an initial value of 5. The expression evaluation returns the value 10.

2.6 Function Invocation

A function invocation is done with the -> operator. Thus, the expression even->(n-1) in the previous declaration of odd calls the function even with the value of n-1 as an argument. L4850 also defines the language-supported functions in Table 1.

Function Prototype	Parameter Types		Description
first->(L)	L	a list	Return the first element of a
			L
rest->(L)	L	a list	Return a list contain all ele-
			ments of L except the first
insert->(e,L)	е	any value	Insert e onto the front of L
	L	a list	
list->(e,)	е,	one or more values	create a list containing all
			parameters
emptyp->(L)	L	a list	Return true iff L is an empty
			list
pairp->(L)	L	a list	true iff L is a non-empty list
listp->(L)	L	a list	Return true iff L is a list
equalp->(L1,L2)	L1,L2	lists	Return true iff L1 and L2 are
			equivalent lists
length->(L)	L	a list	Return the length of Lt
numberp->(e)	е	any value	Return true iff e is a number
exit->()			exit the interpreter

Table 1: L4850 Language-supported Functions

2.7 Function Expression

 ${\it L}4850$ supports first-class functions. Thus, functions are values just as numbers are values. The expression

func
$$(x) \{ x + 2 \}$$

declares a nameless function that adds 2 to its argument. A function expression can be put anywhere in a program that any other expression can be put. Thus, it can be assigned to a variable in an assignment expression or called in an invocation as in

$$(func (x) { x + 2 }) \rightarrow (3)$$

which returns the value 5.

2.8 Expressions

L4850 arithmetic expressions compute simple values of type integer or float For both integer and floating point numbers, arithmetic and comparison are defined.

Coercion: If an expression contains operands of only one type, evaluation is straight forward. When an operand contains mixed types, the situation is more complex. If an arithmetic expression has an integer operand and a float operand, the integer operand should be converted to a float before the operation is performed.

Relational and logic operators always produce a boolean. To perform a comparison between an integer and a float, the integer is converted to a float. Logic operators are only defined on boolean values. In addition, boolean values may not be used in relational and arithmetic operators.

Operator Precedence Operator precedences in L4850 are specified in the table below. Multiplication and division have the highest priority, && and | | have the lowest.

Operator	Precedence
*, /	5
+, -	4
<, <=, =, >=, >, !=	3
!	2
&&,	1

Precedence has already been encoded in the grammar in Section 1.2.

2.9 Classes and Objects

L4850 supports simple classes without inheritance. Classes may contain instance variables and methods. All variables have private access and all methods have public access. The code below defines a class with one instance variable and getter and setter methods for that variable.

```
defclass test {
    vars x
    method setX(n) { assign n to x }
    method getX() { x }
}
```

To instantiate an object, the new operator is used. In the expression

```
assign new test to t
```

a new instance of test is assigned to the variable t. Methods are accessed using the . operator and invoked using the -> operator. As an example, the following code returns the value 0.

```
defunc f (x) {
    assign new test to t
    t.setX->(x)
    t.getX->()
}
```