EL-GY 6483 Communications

•	You are encouraged to work in groups, and to use the Internet or any other tools available to learn the	ae
	material in order to answer these questions.	

1. The ATmega128 microcontroller includes a UART that can be used to provide a serial interface. The following code snippet is often seen in programs that use the UART interface:

```
while(!(UCSROA & 0x20));
UDRO = x;
```

where x is a previously declared and initalized uint8_t; UCSROA and UDRO are defined in header files to refer to memory locations corresponding to the USART Control and Status Register A and USART Data Register, respectively; and the UART interface has already been configured, i.e. is ready for use.

(a) Refer to page 188 of the manual for the ATmega128, (available online). What does each line of the code snippet above do, with respect to the peripheral registers? What does the code snippet as a whole do?

Solution:		

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- (e) On this microcontroller, the baud rate is set by writing the value $UBRR = \frac{f}{16} \frac{f}{B} \frac{osc}{des}$ 1 to a UBRR register, where f_{osc} is the oscillator frequency in Hz and B_{des} is the desired baud rate in bits per second. The achieved baud rate is then $B_{ach} = \frac{f_{osc}}{16(UBRR+1)}$ (See page 172-173 of the ATmega128 reference manual for more details.) Because we can only write integer values to the register, not all baud rates can be achieved exactly.
 - What is the closest we can get to 57600 baud (i.e., what is B_{ach}) if f_{osc} is 8 MHz? (Assume U2X is 0.)
 - What value should be written to the UBRR register to achieve this baud rate?
 - What is the percent error in this case, calculated as $\left(\frac{B_{ach}}{B_{des}} 1\right) \times 100\%$?

Solution:		

(f) Suppose the other communication partner is an ATmega128 using $f_{osc} = 2$ MHz. (Assume U2X is 0.) What will its B_{ach} be if it tries to operate at 57600 baud? What will be the total error between the pair, and is it less than the maximum error recommended in Table 75 of the ATmega128 reference manual (page 186)?

Solution:	6

Draw a connection diagram for this configuration. What is the total number of wires?
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