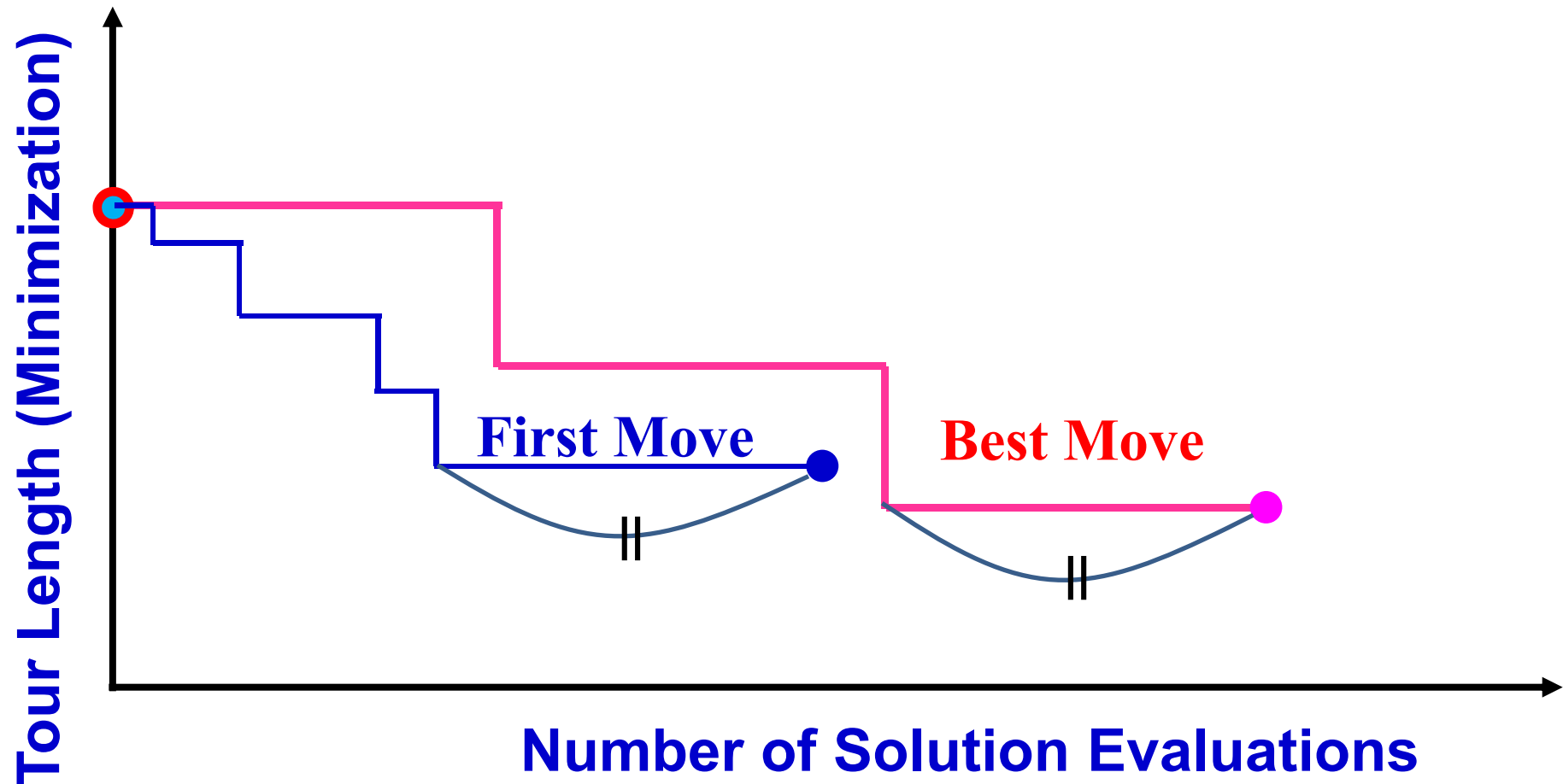


Optimization Methods

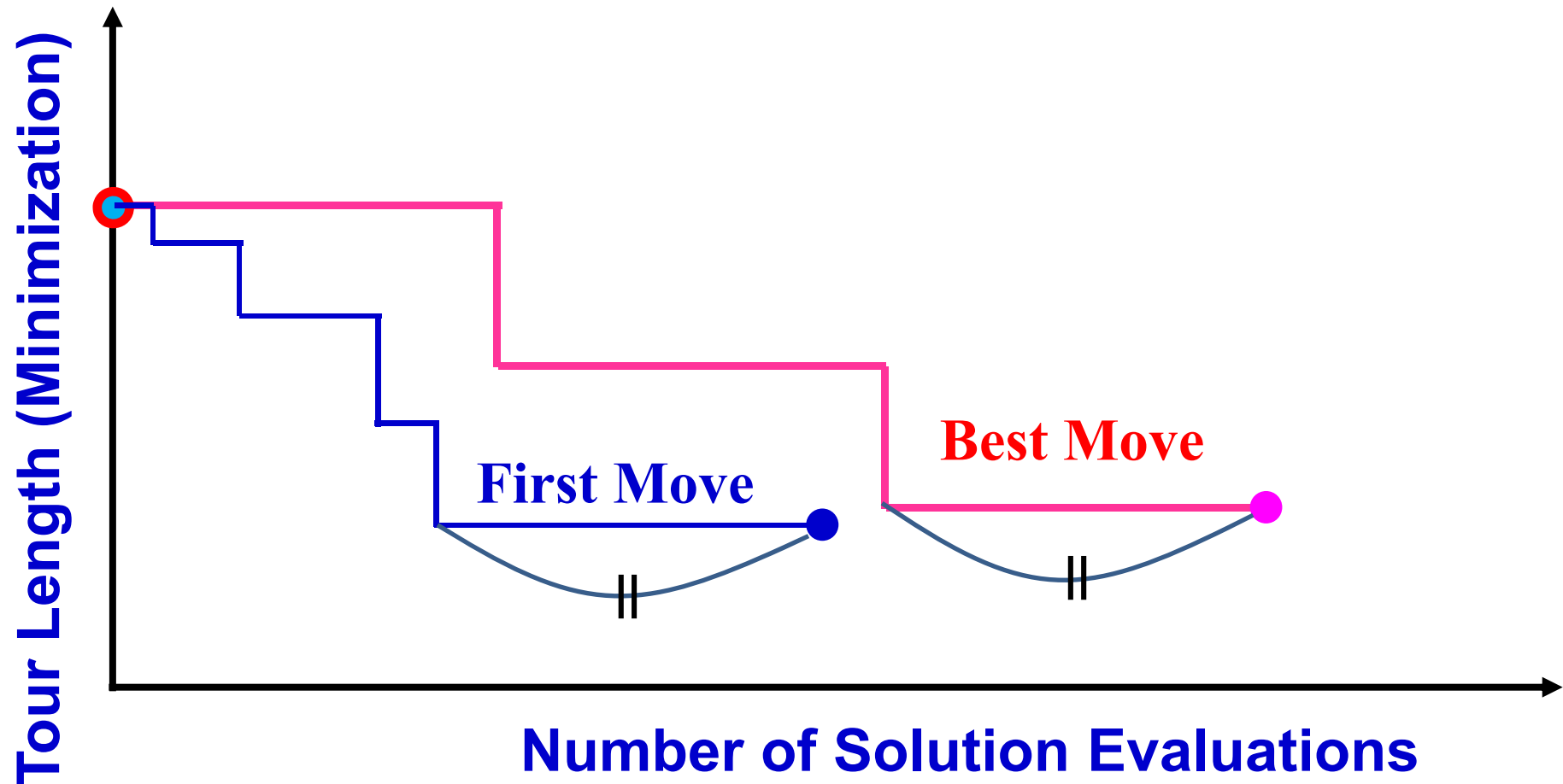
1. Introduction.
2. Greedy algorithms for combinatorial optimization.
3. LS and neighborhood structures for combinatorial optimization.
- 4. Variable neighborhood search, neighborhood descent, SA, TS.**
5. Branch and bound algorithms, and subset selection algorithms.
6. Linear programming problem formulations and applications.
7. Linear programming algorithms.
8. Integer linear programming algorithms.
9. Unconstrained nonlinear optimization and gradient descent.
10. Newton's methods and Levenberg-Marquardt modification.
11. Quasi-Newton methods and conjugate direction methods.
12. Nonlinear optimization with equality constraints.
13. Nonlinear optimization with inequality constraints.
14. Problem formulation and concepts in multi-objective optimization.
15. Search for single final solution in multi-objective optimization.
- 16: Search for multiple solutions in multi-objective optimization.

Best Move and First Move: Which is better ?



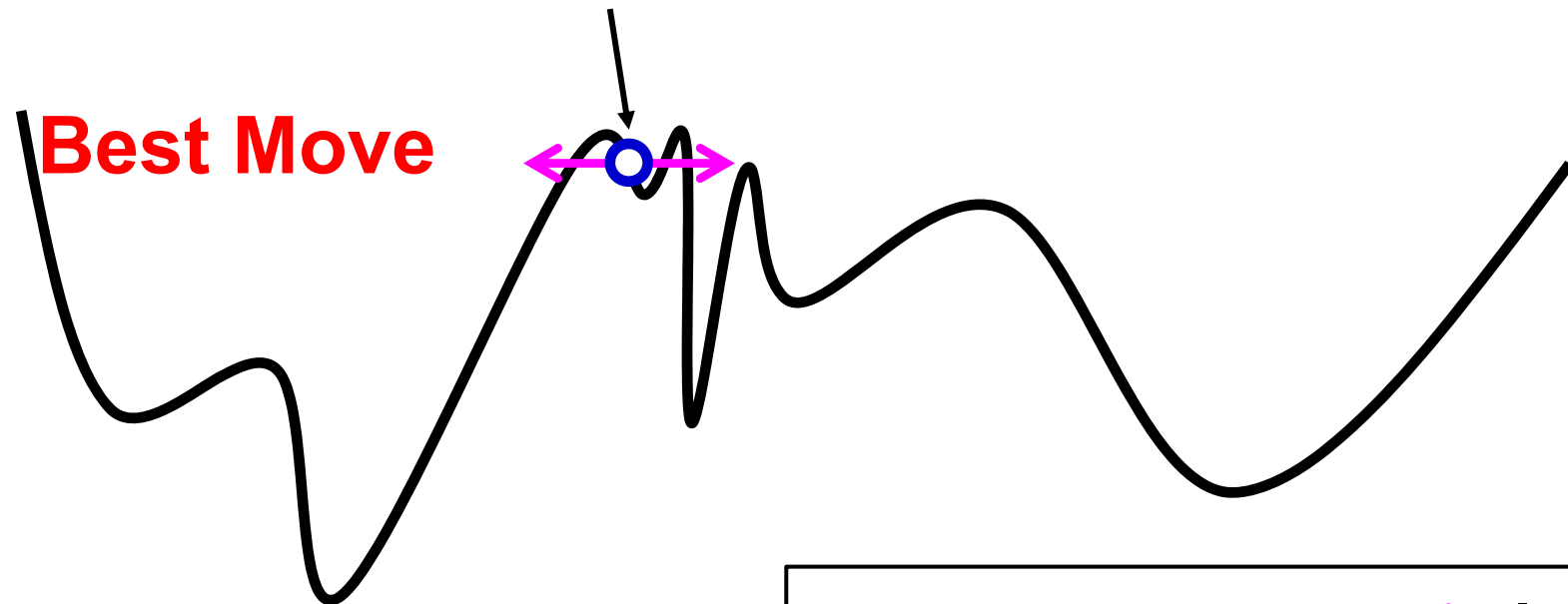
- (1) In general, the first move needs less computation load.
- (2) With respect to the final solution quality, we cannot say which is better between these two strategies.

Best Move and First Move: Which is better ?

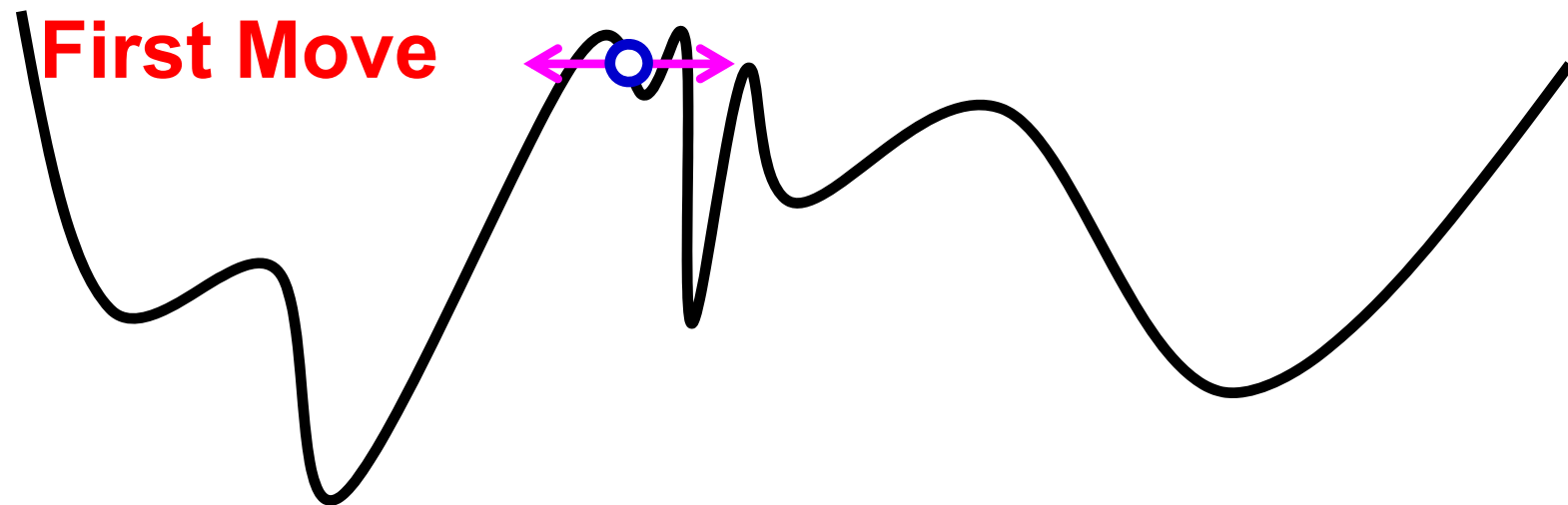


- (1) In general, the first move needs less computation load.
- (2) With respect to the final solution quality, we cannot say which is better between these two strategies.

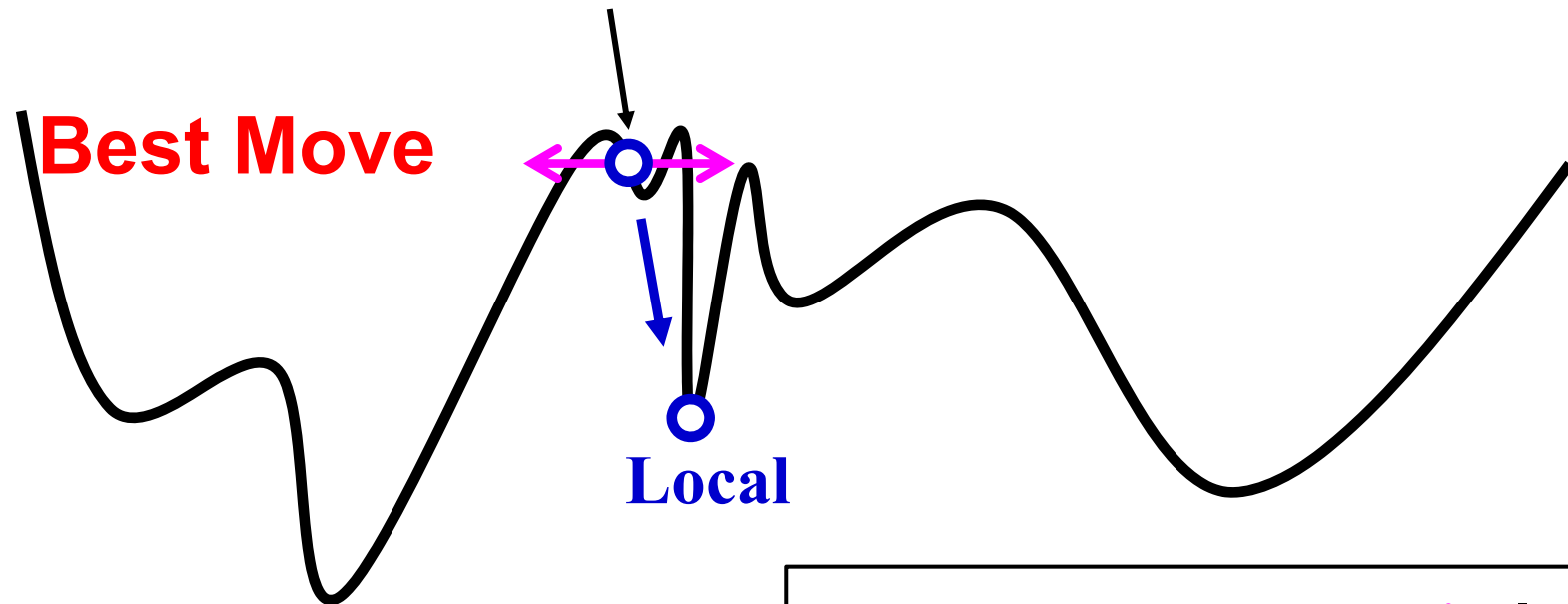
Initial Solution



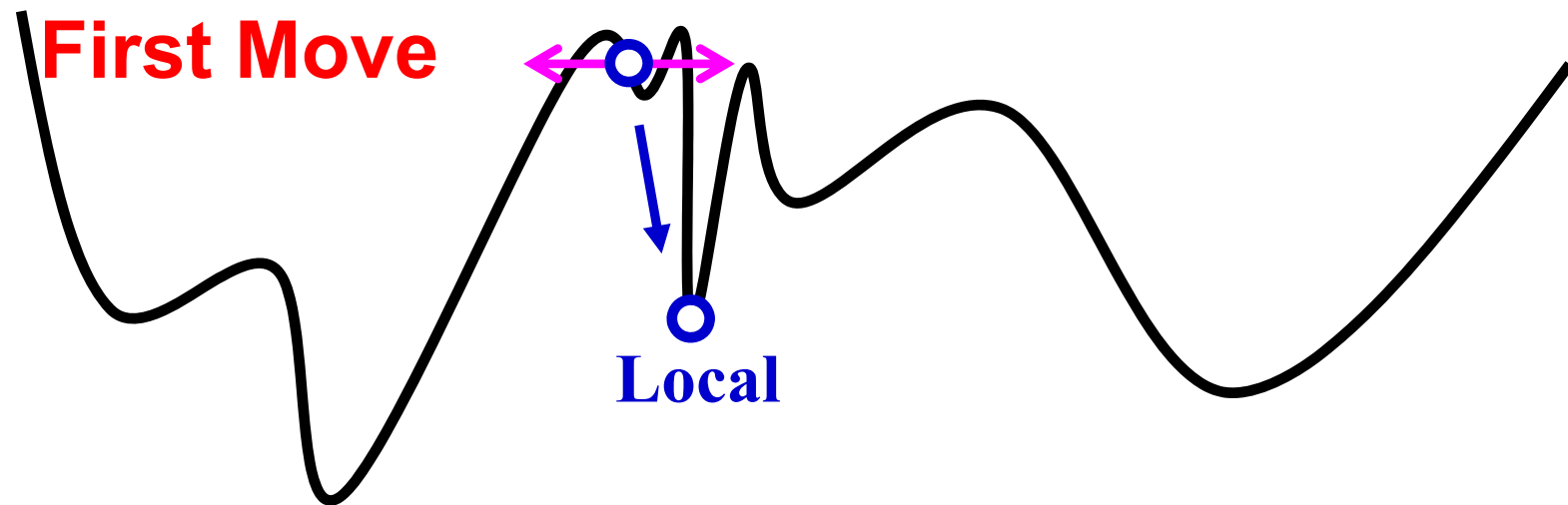
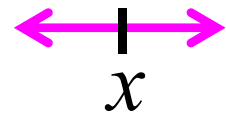
Neighborhood size 



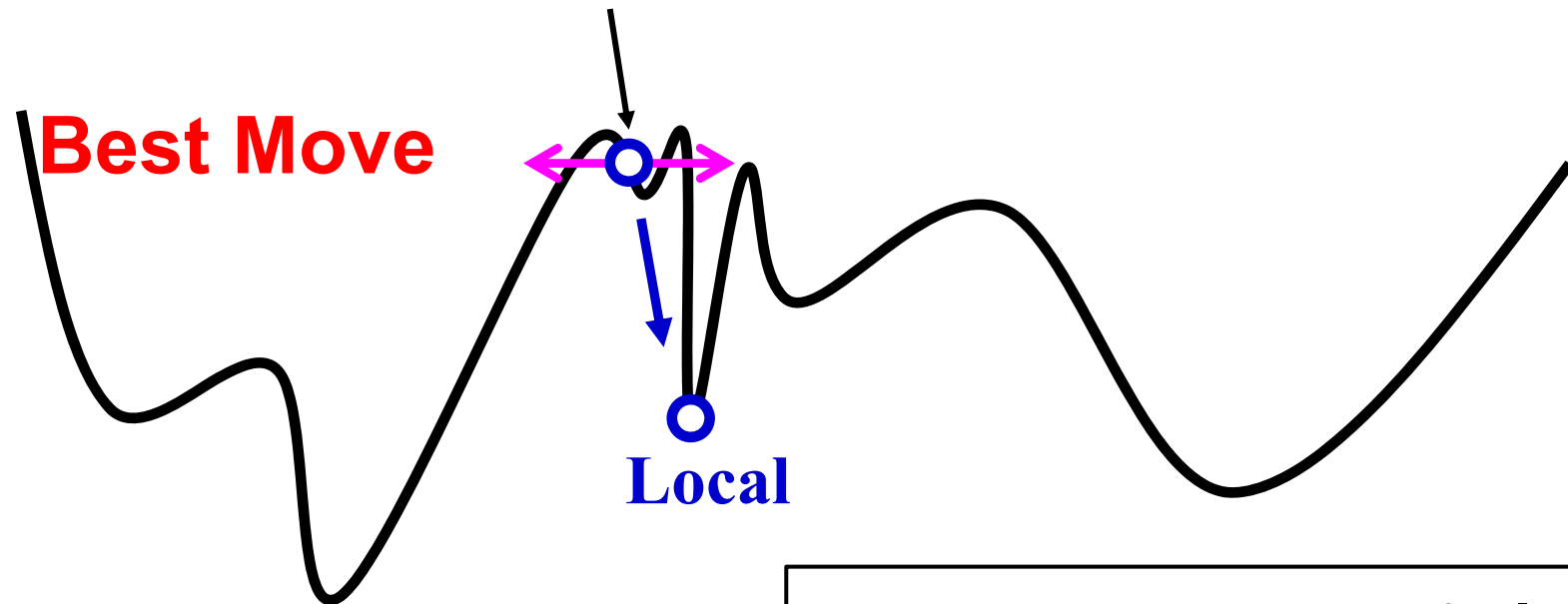
Initial Solution



Neighborhood size

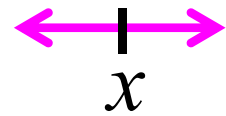


Initial Solution

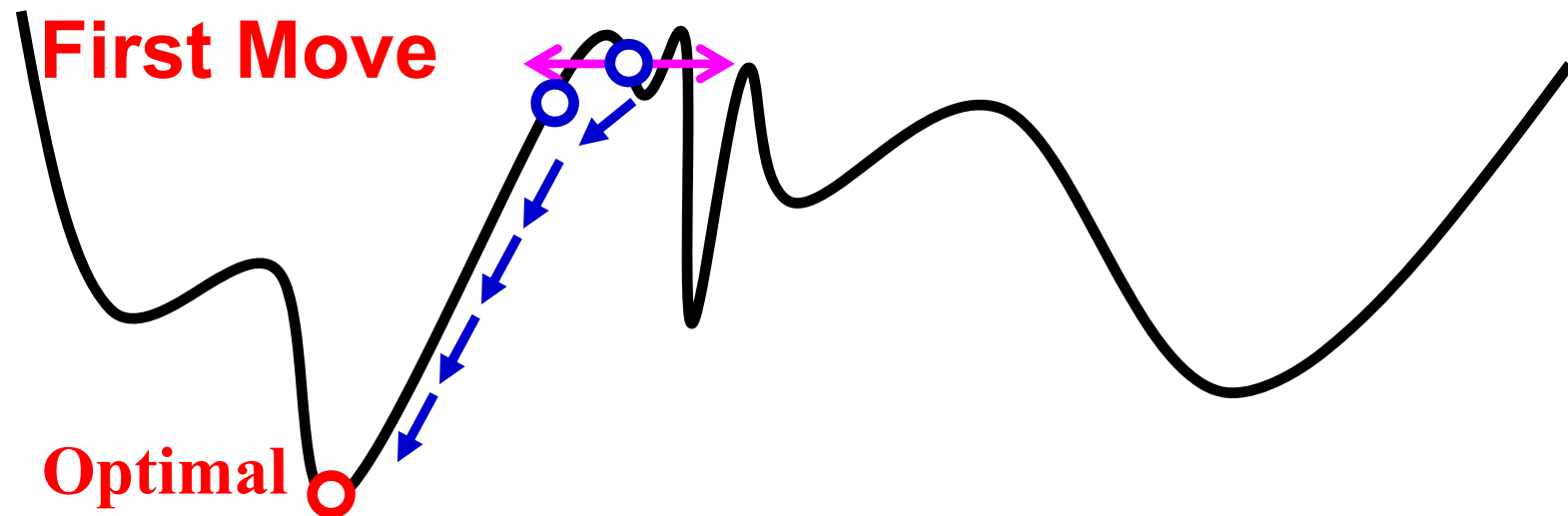


Local

Neighborhood size



x

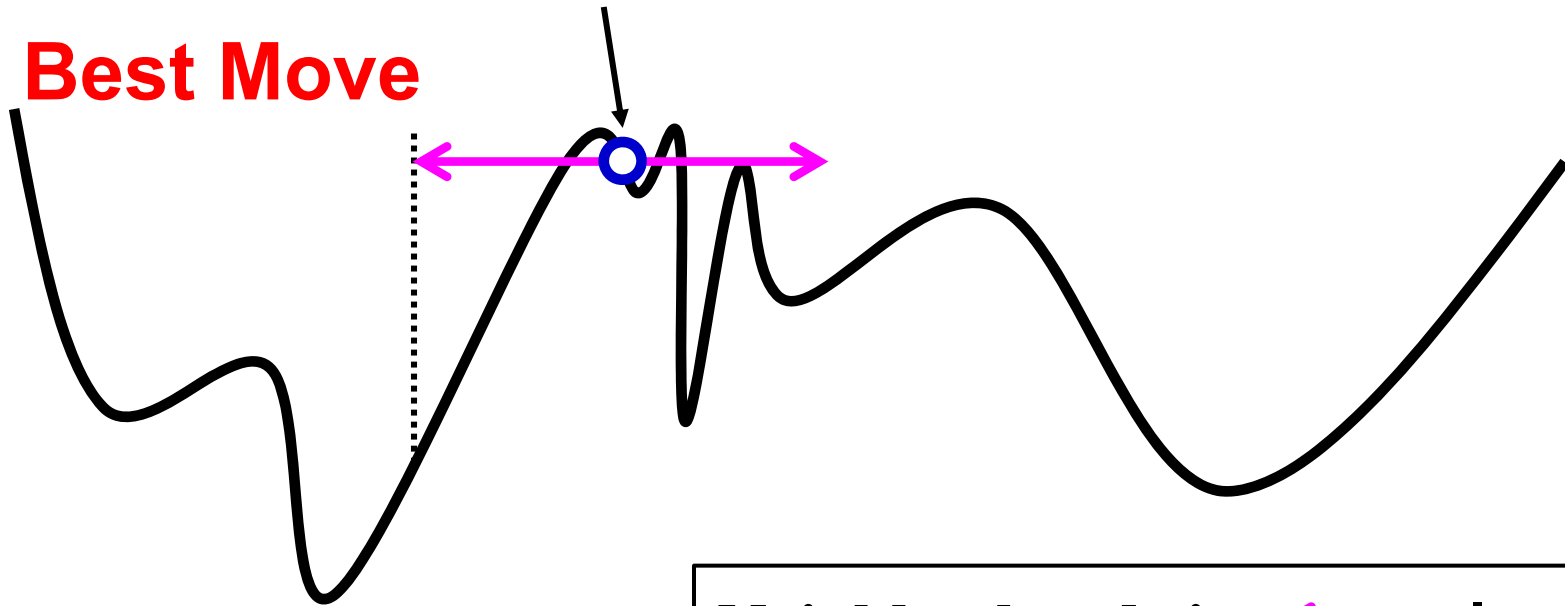


First Move

Optimal

Initial Solution

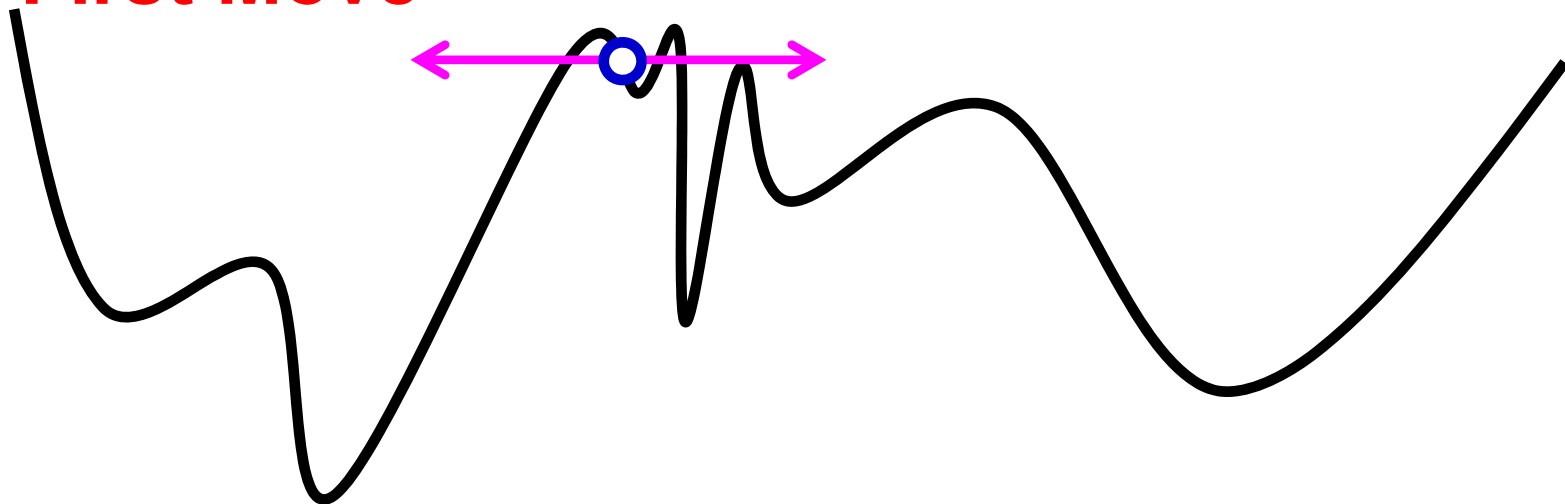
Best Move



Neighborhood size

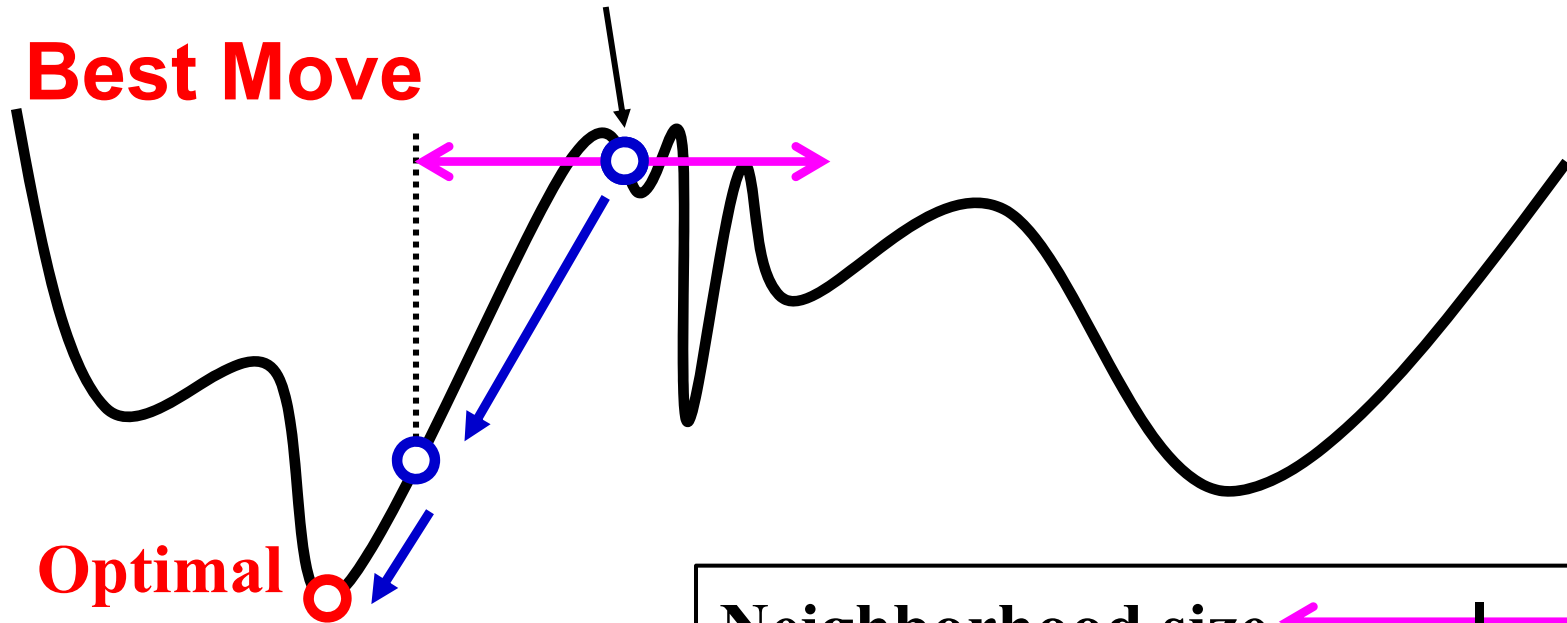
x

First Move



Initial Solution

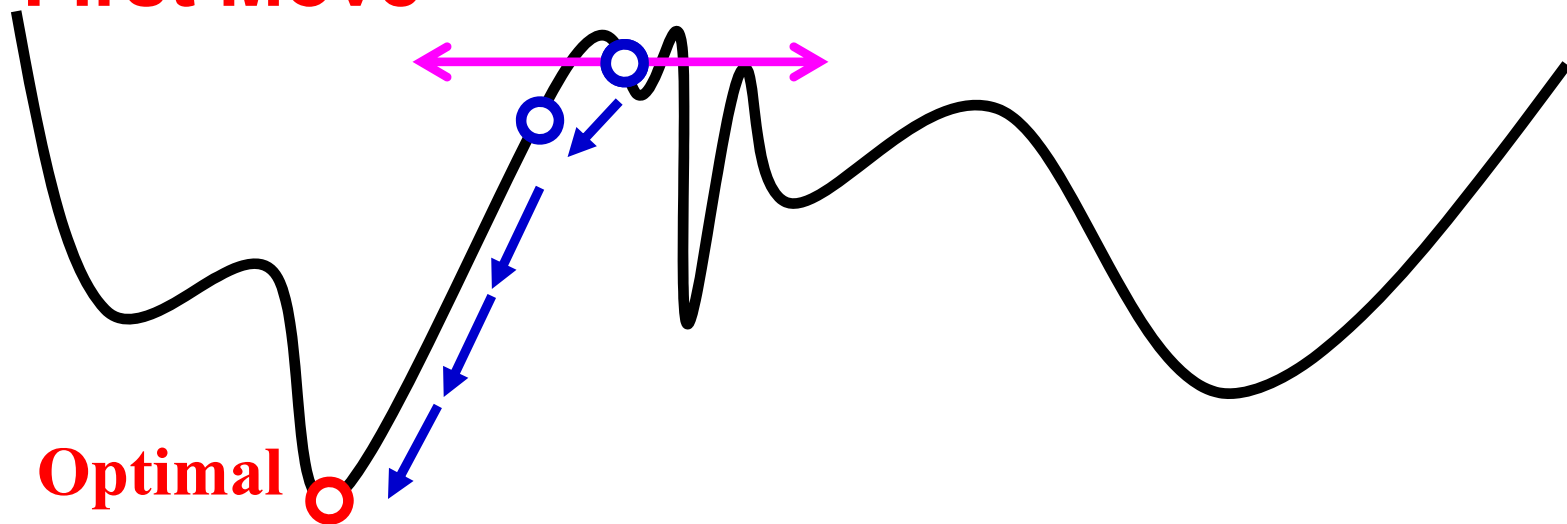
Best Move



Neighborhood size

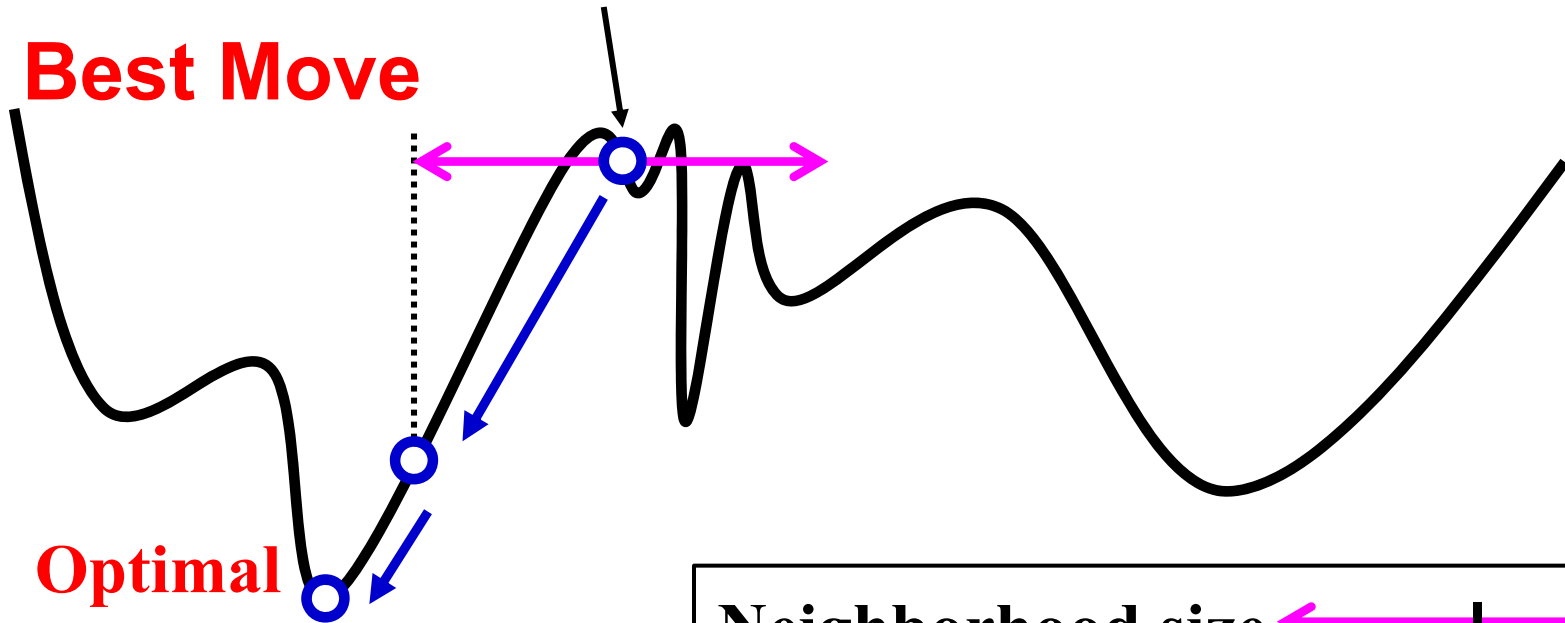
x

First Move



Initial Solution

Best Move

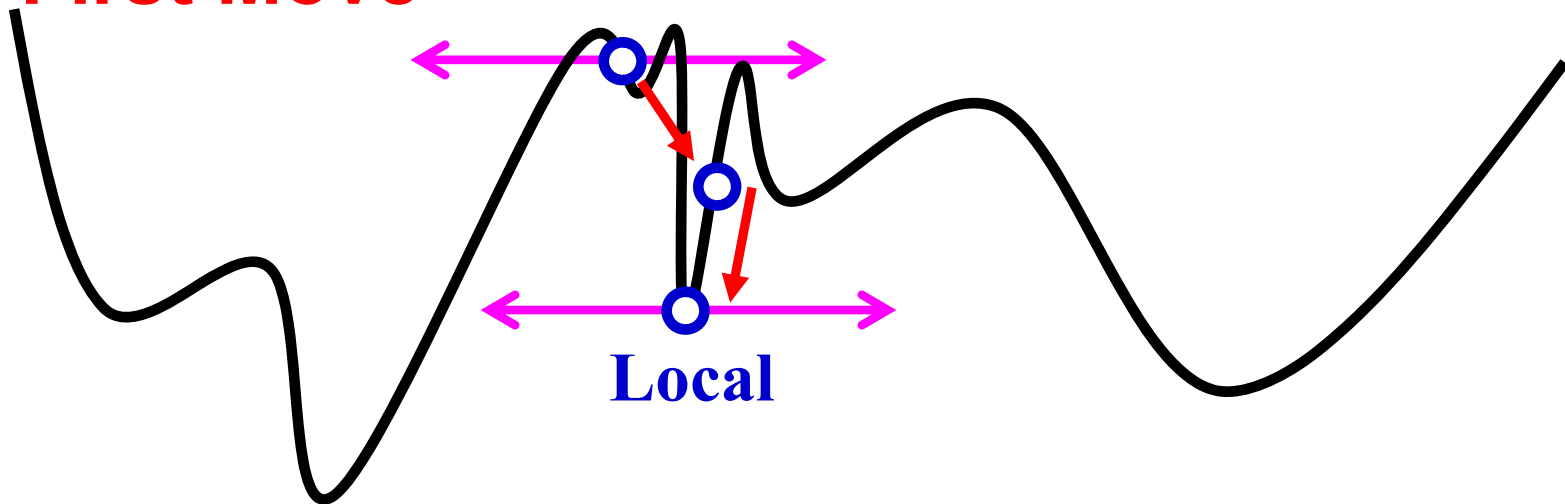


Optimal

Neighborhood size

x

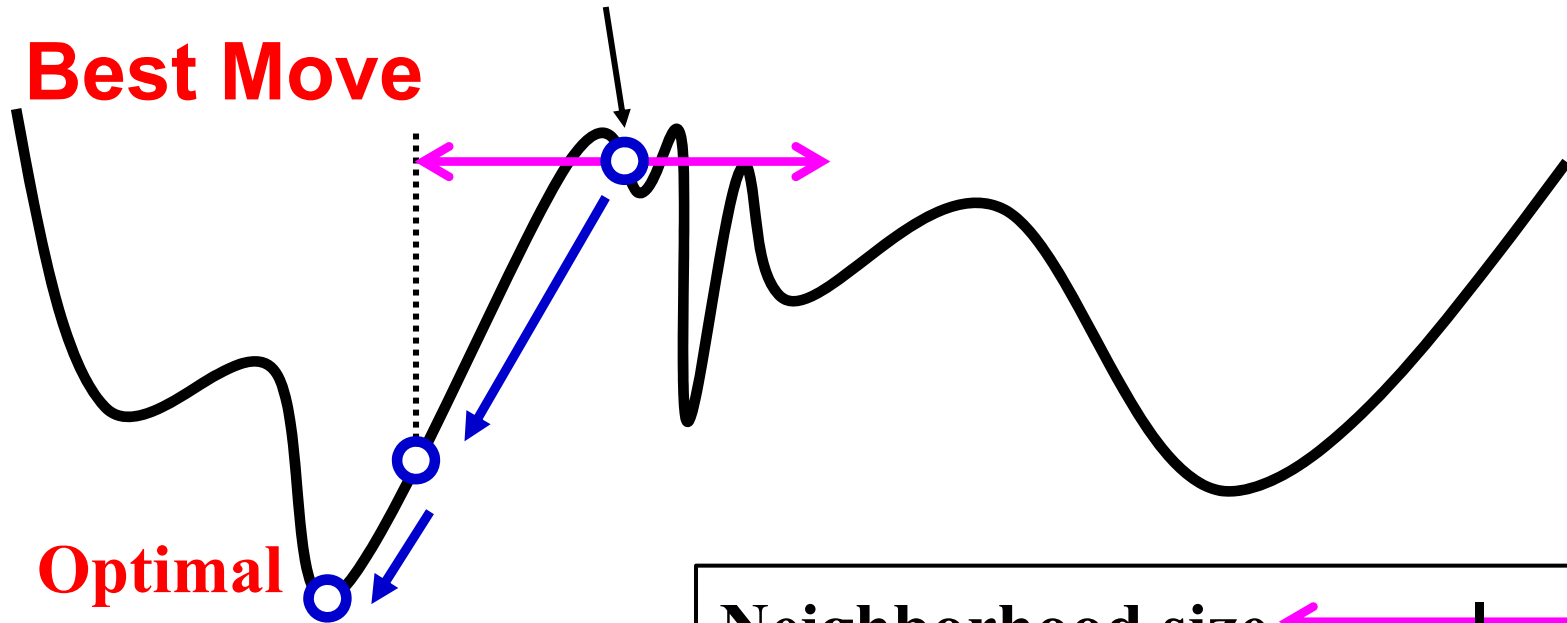
First Move



Local

Initial Solution

Best Move

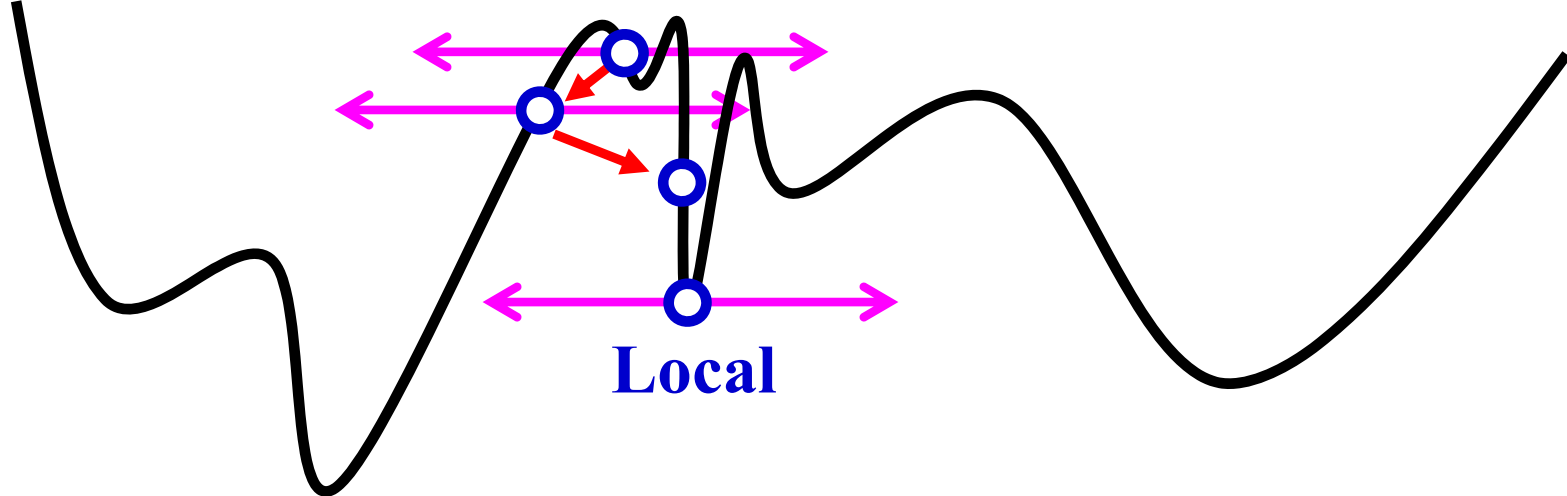


Optimal

Neighborhood size

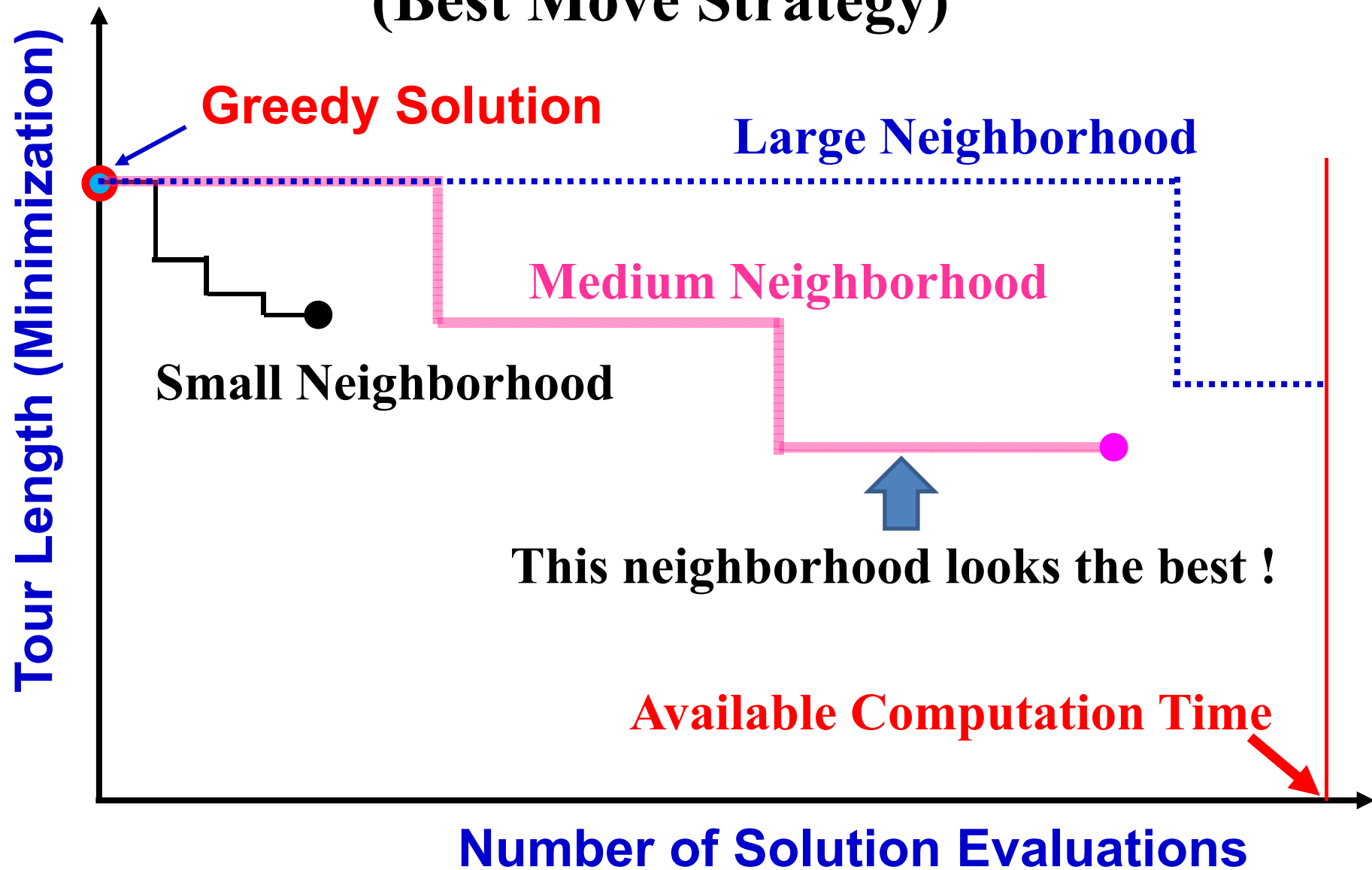
x

First Move

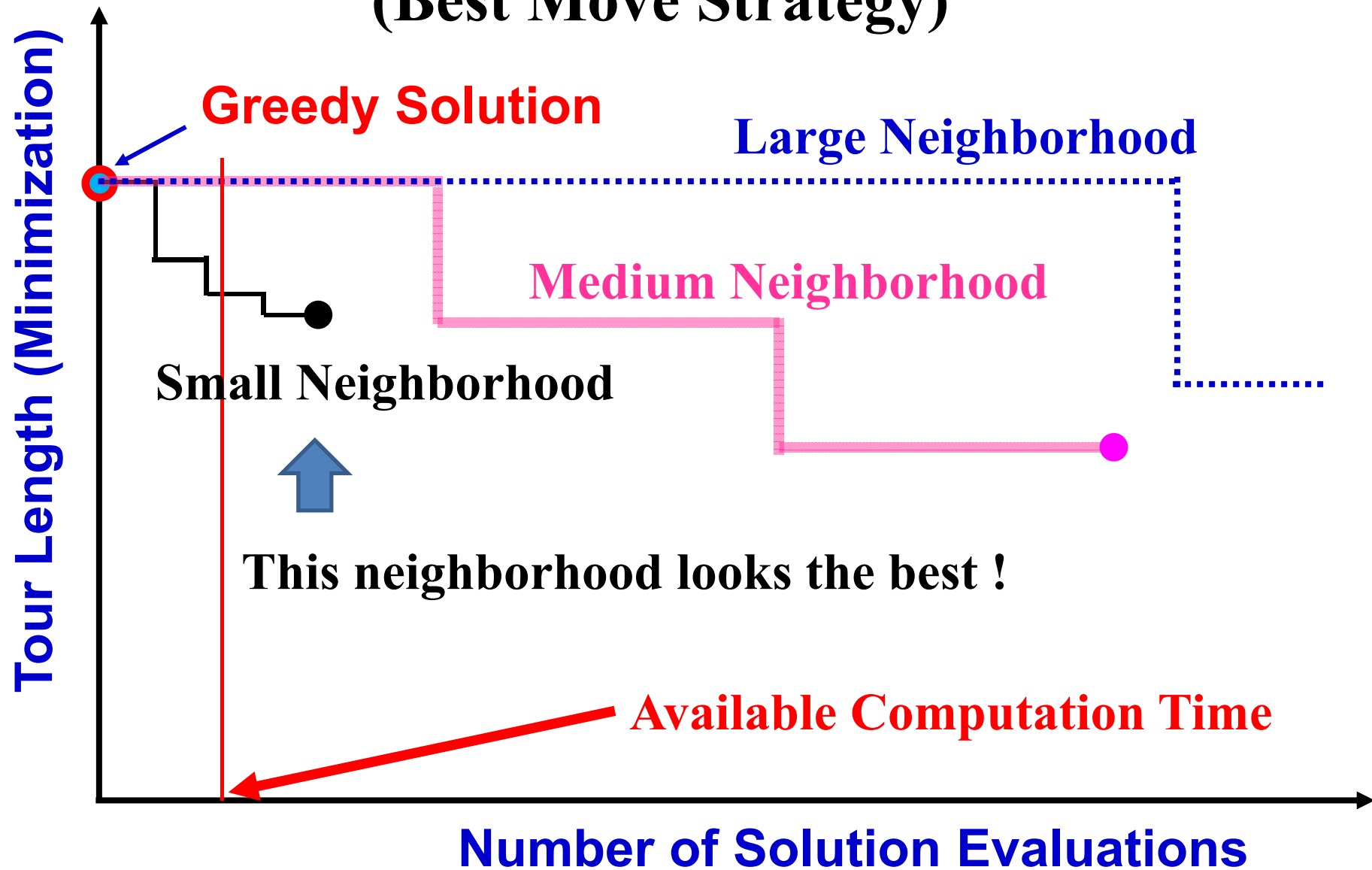


Local

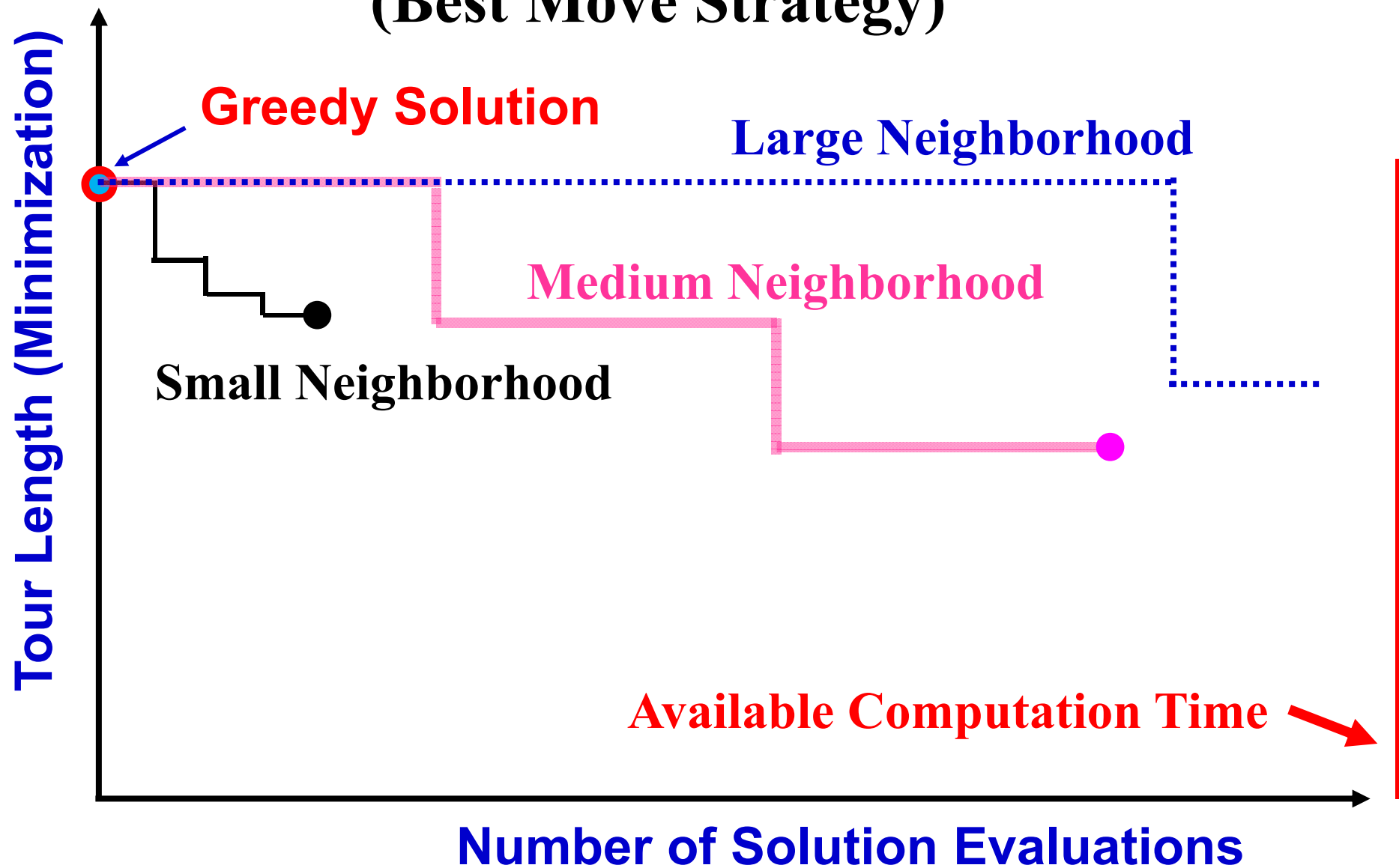
Behavior of Local Search (Best Move Strategy)



Behavior of Local Search (Best Move Strategy)



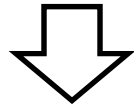
Behavior of Local Search (Best Move Strategy)



Local Search

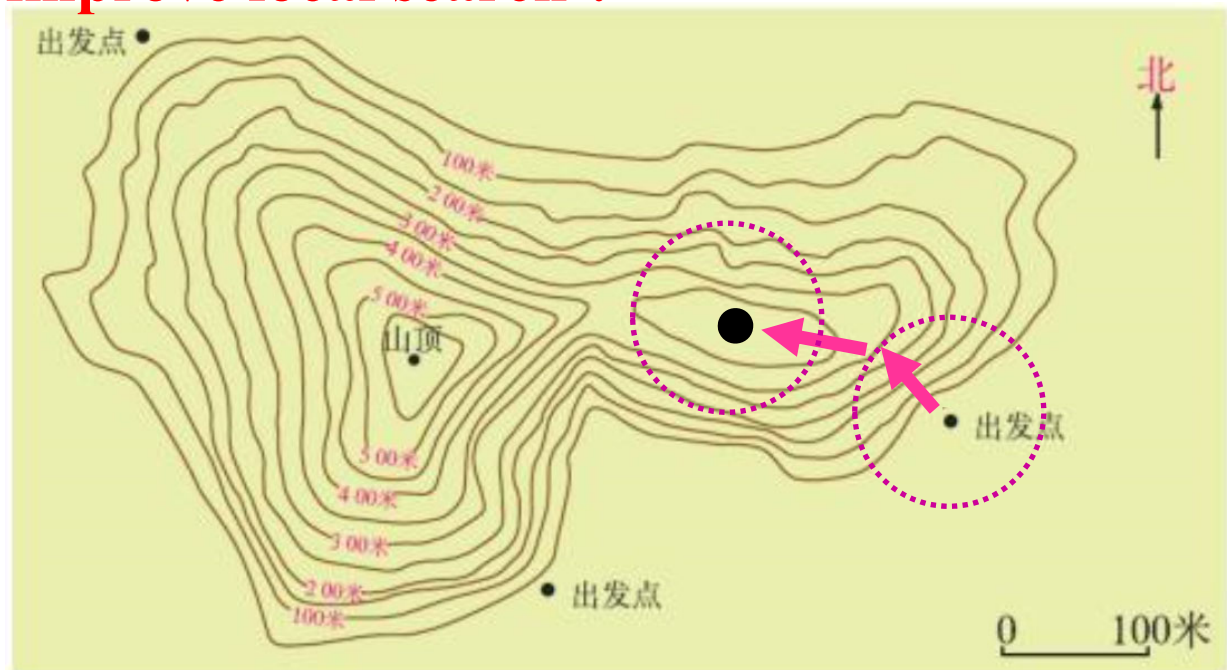
Implementation Issues

- (1) Specification of an initial solution
- (2) Specification of a neighborhood structure
- (3) Choice between the first move and the best move



Local Solution

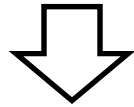
Question: How can we improve local search ?



Local Search

Implementation Issues

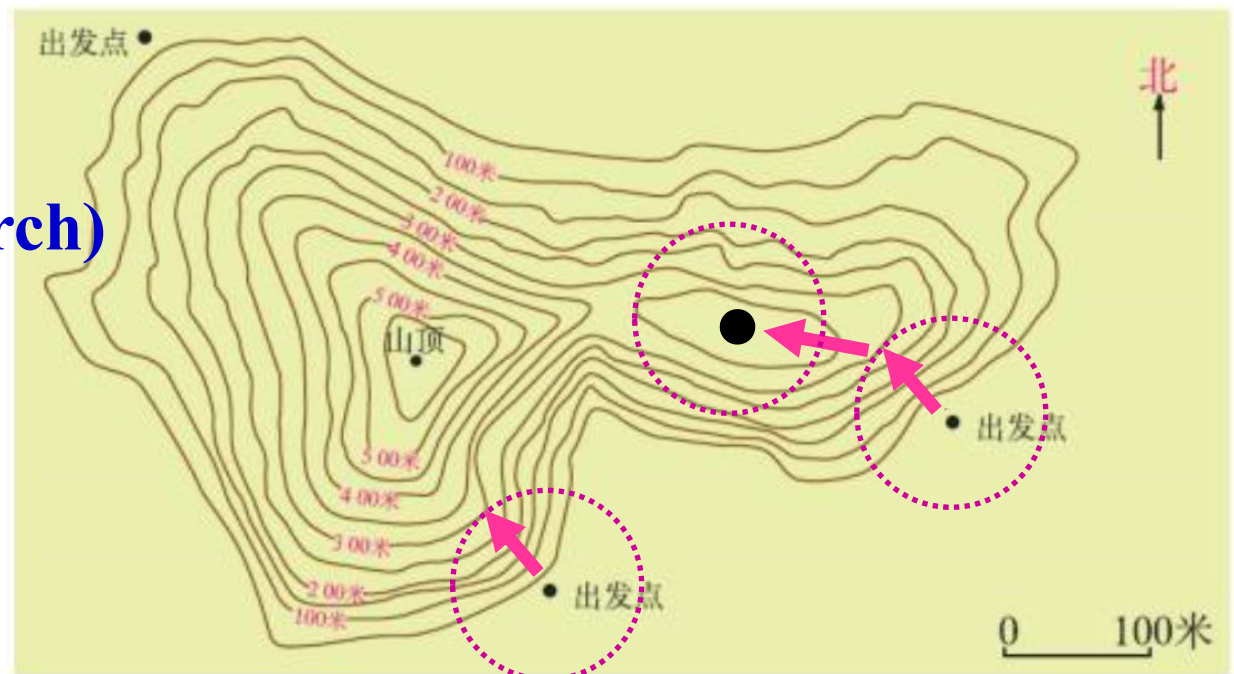
- (1) Specification of an initial solution
- (2) Specification of a neighborhood structure
- (3) Choice between the first move and the best move



Local Solution

Next Step

- Restart
(Iterated Local Search)

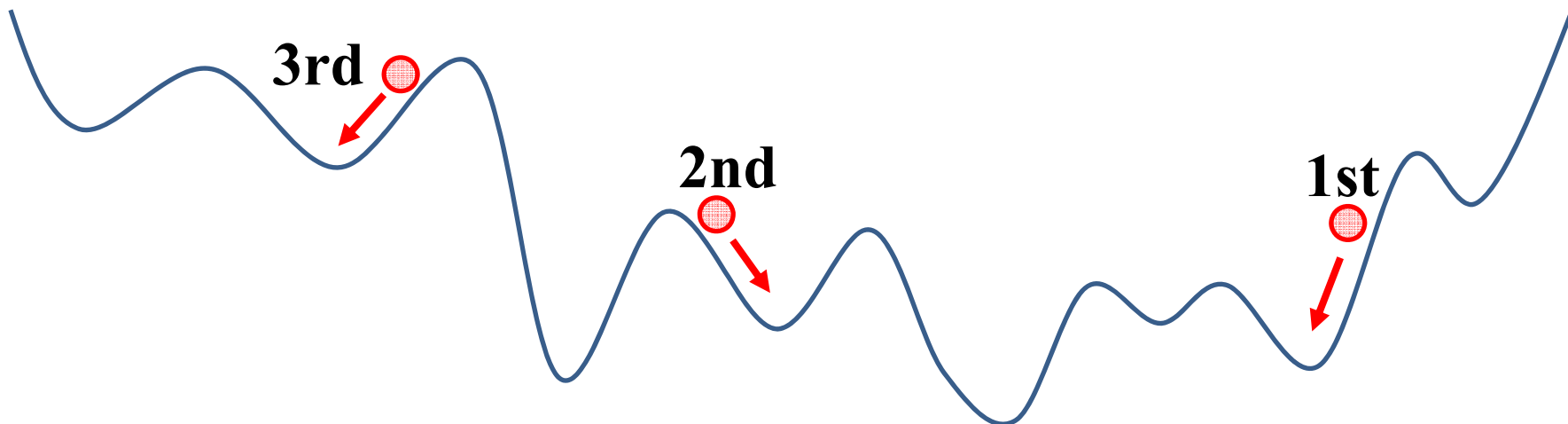


Iterated Local Search

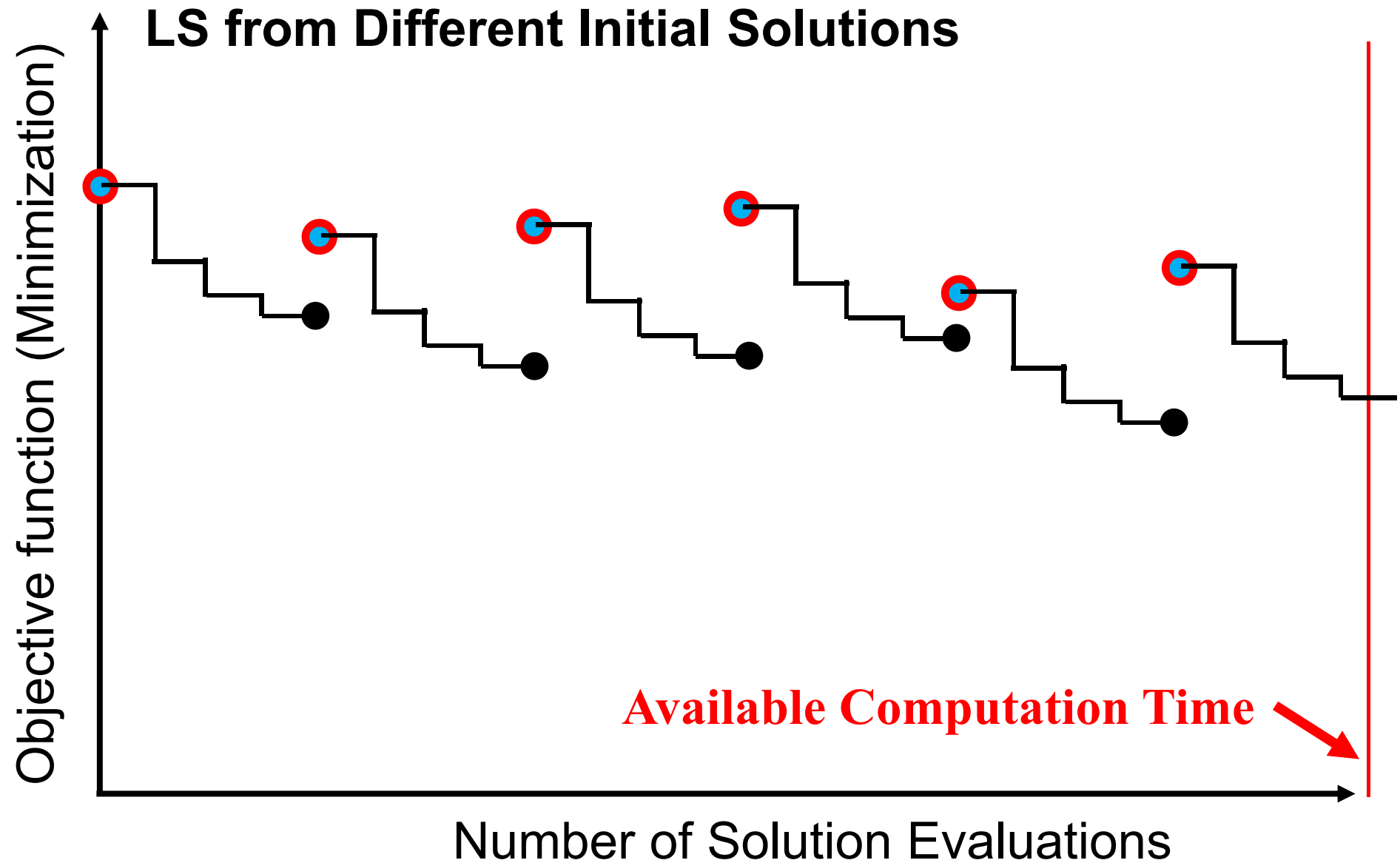
Iterate LS from Different Initial Solutions

Iterate the following steps:

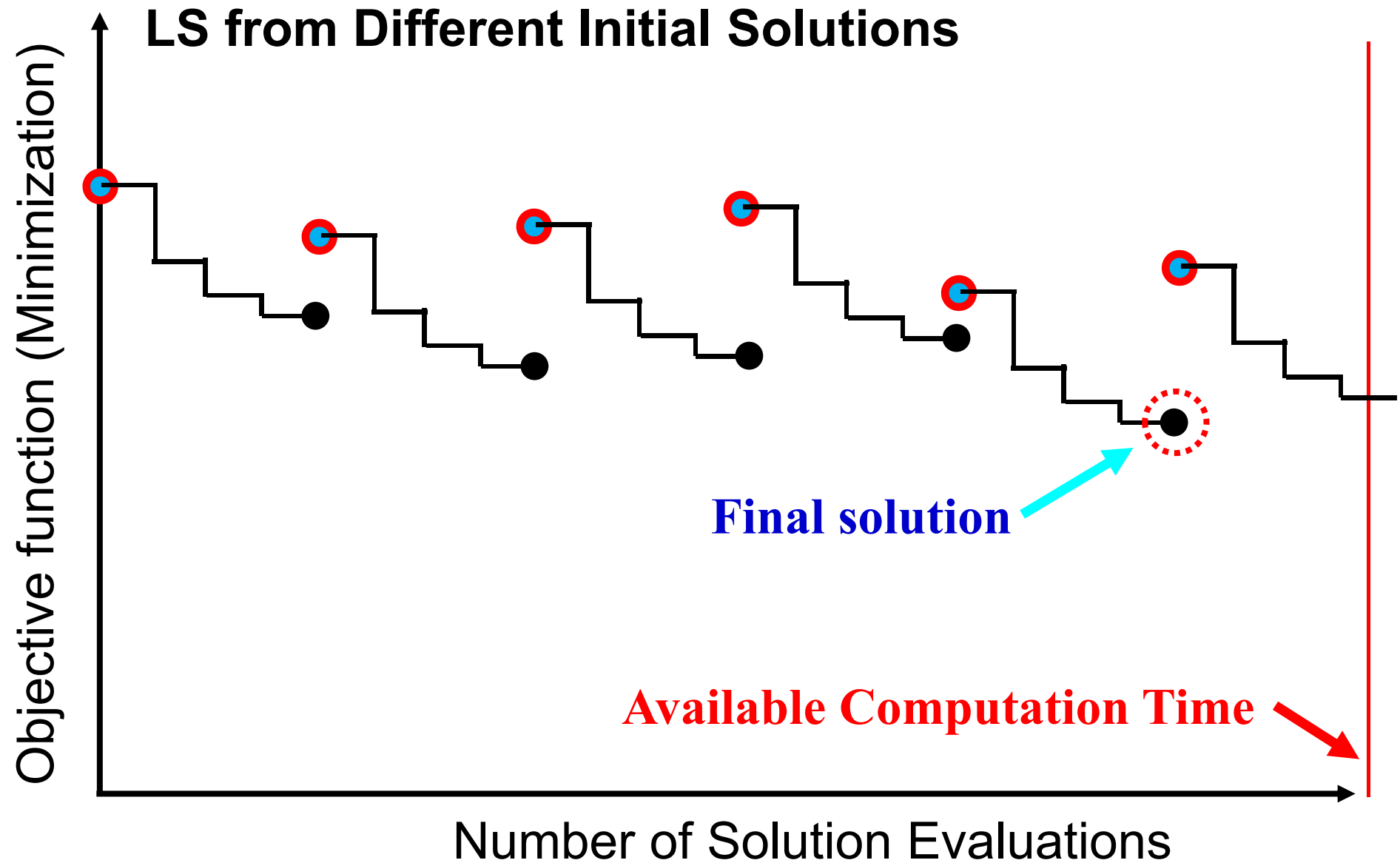
- (1) Generate an initial solution \mathbf{x} .
- (2) Iterate the following steps:
 - (i) The first move or the best move (if a neighbor \mathbf{y} of the current solution \mathbf{x} is better than \mathbf{x} , replace \mathbf{x} with \mathbf{y}).
 - (ii) If no better solution exists in the neighborhood of \mathbf{x} , terminate the current execution of LS (i.e., **restart with a new initial solution \mathbf{x} in (1)**).



Iterated Local Search



Iterated Local Search



Iterated Local Search

Iterate LS from Different Initial Solutions

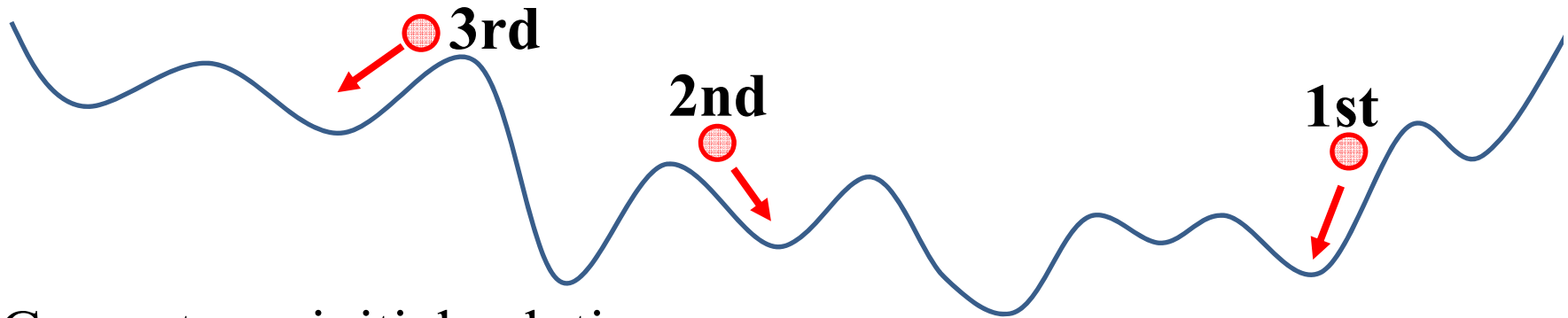
Iterate the following steps:

- (1) Generate an initial solution \mathbf{x} .
- (2) Iterate the following steps:
 - (i) The first move or the best move (if a neighbor \mathbf{y} of the current solution \mathbf{x} is better than \mathbf{x} , replace \mathbf{x} with \mathbf{y}).
 - (ii) If no better solution exists in the neighborhood of \mathbf{x} , terminate the current execution of LS (i.e., **restart with a new initial solution \mathbf{x} in (1)**).

Discussions:

How to specify an initial solution for each LS run in (1)?

- In the first LS run: **Greedy solution.**
- In the other LS runs: ????????????.



- (1) Generate an initial solution \mathbf{x} .
- (2) Iterate the following steps:
 - (i) The first move or the best move (if a neighbor \mathbf{y} of the current solution \mathbf{x} is better than \mathbf{x} , replace \mathbf{x} with \mathbf{y}).
 - (ii) If no better solution exists in the neighborhood of \mathbf{x} , terminate the current execution of LS (i.e., **restart with a new initial solution \mathbf{x} in (1)**).

Discussions:

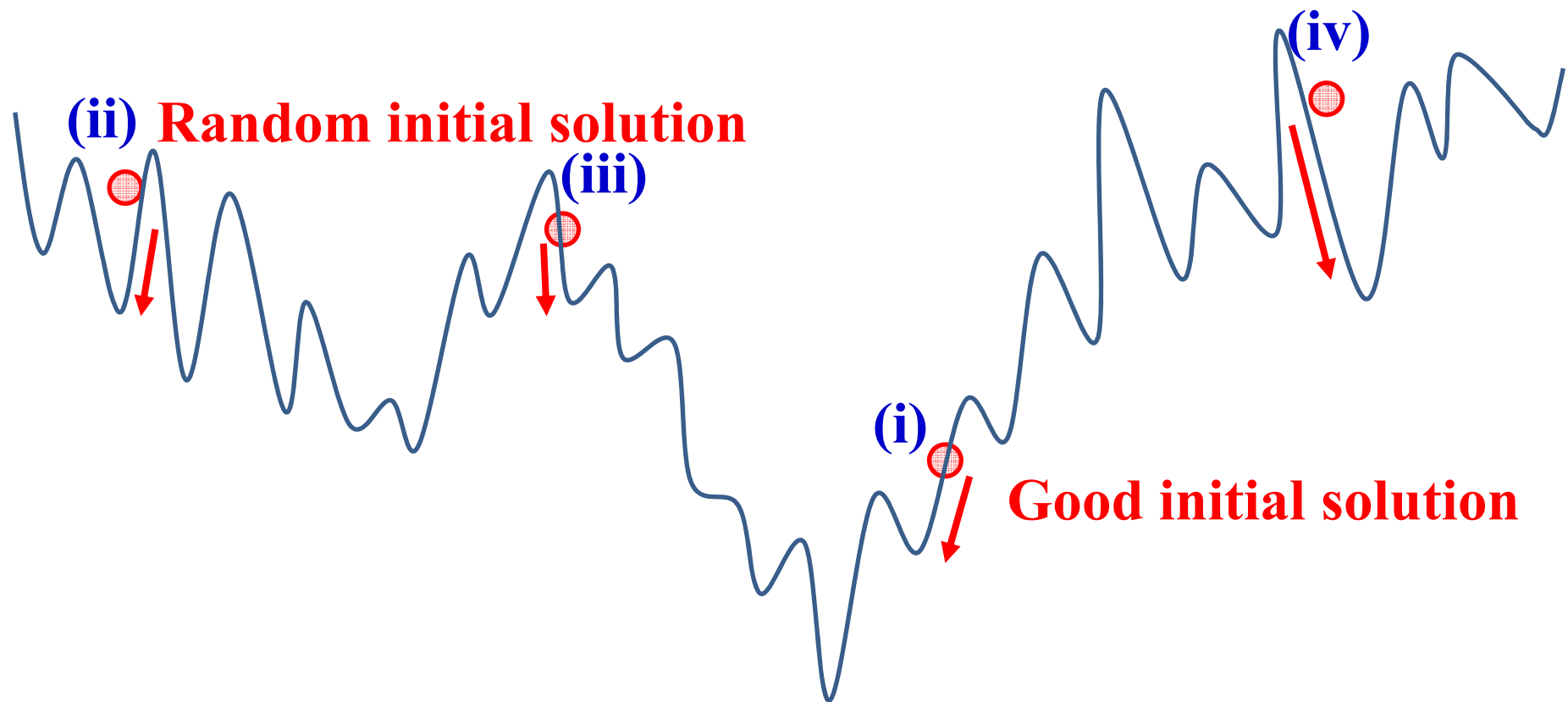
How to specify an initial solution for each LS run in (1)?

- In the first LS run: **Greedy solution.**
- In the other LS runs: ????????????.

Main Issue

Choice of Initial Solutions:

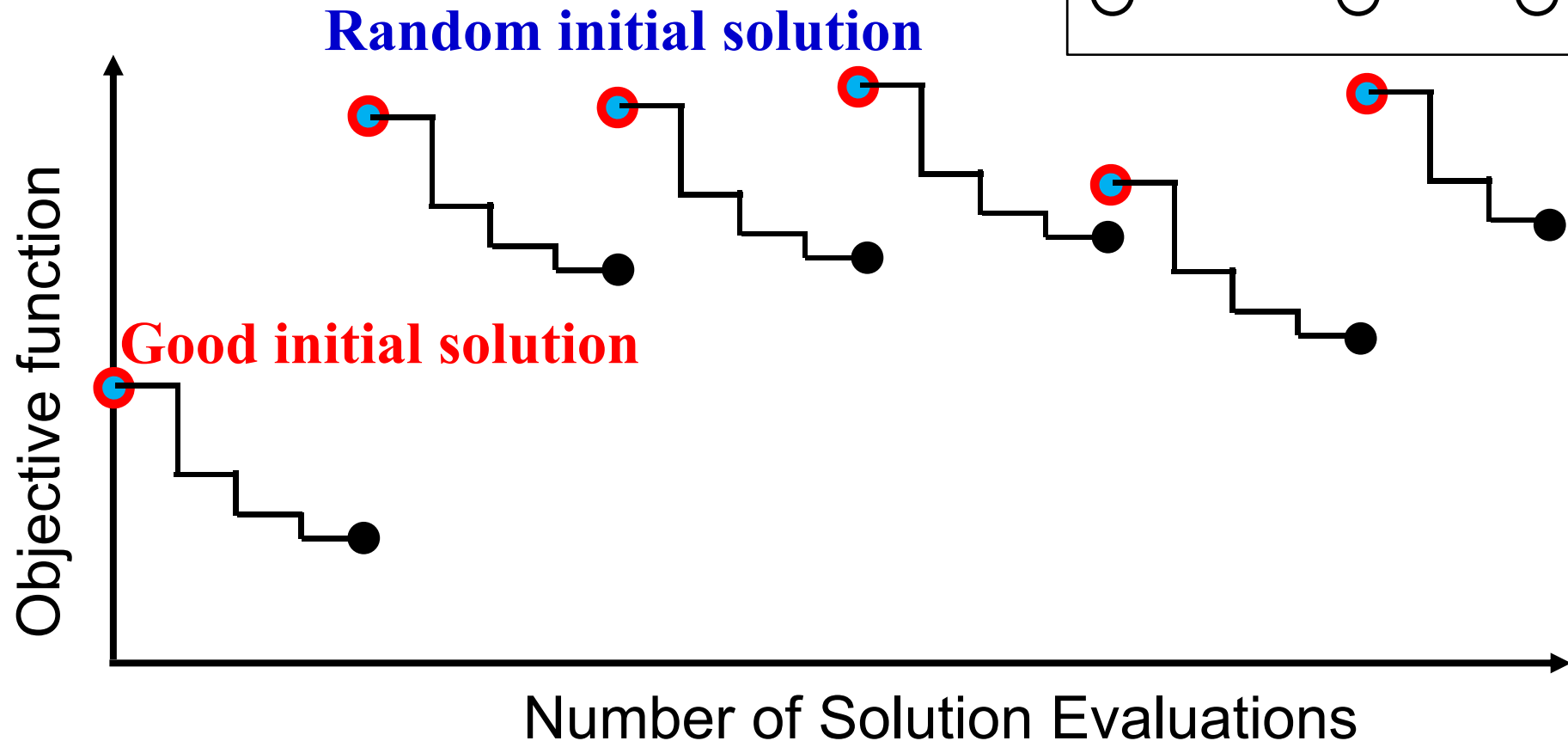
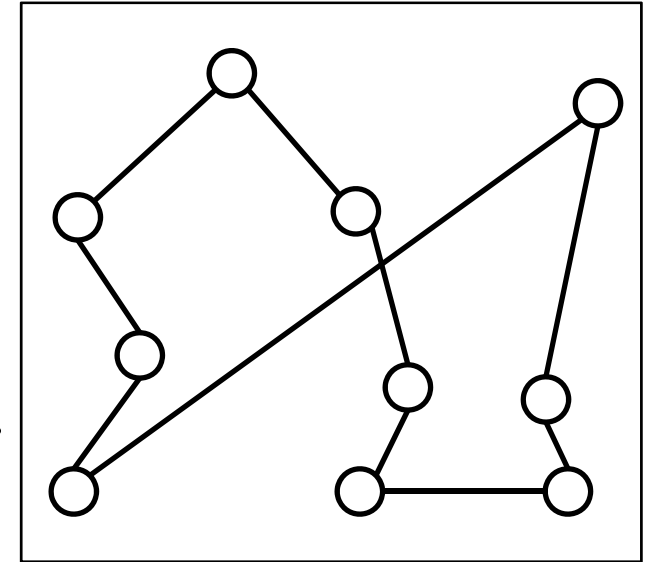
- First initial solution: **Greedy solution.**
- Other initial solutions: **Random solutions.**



Main Issue

Choice of Initial Solutions:

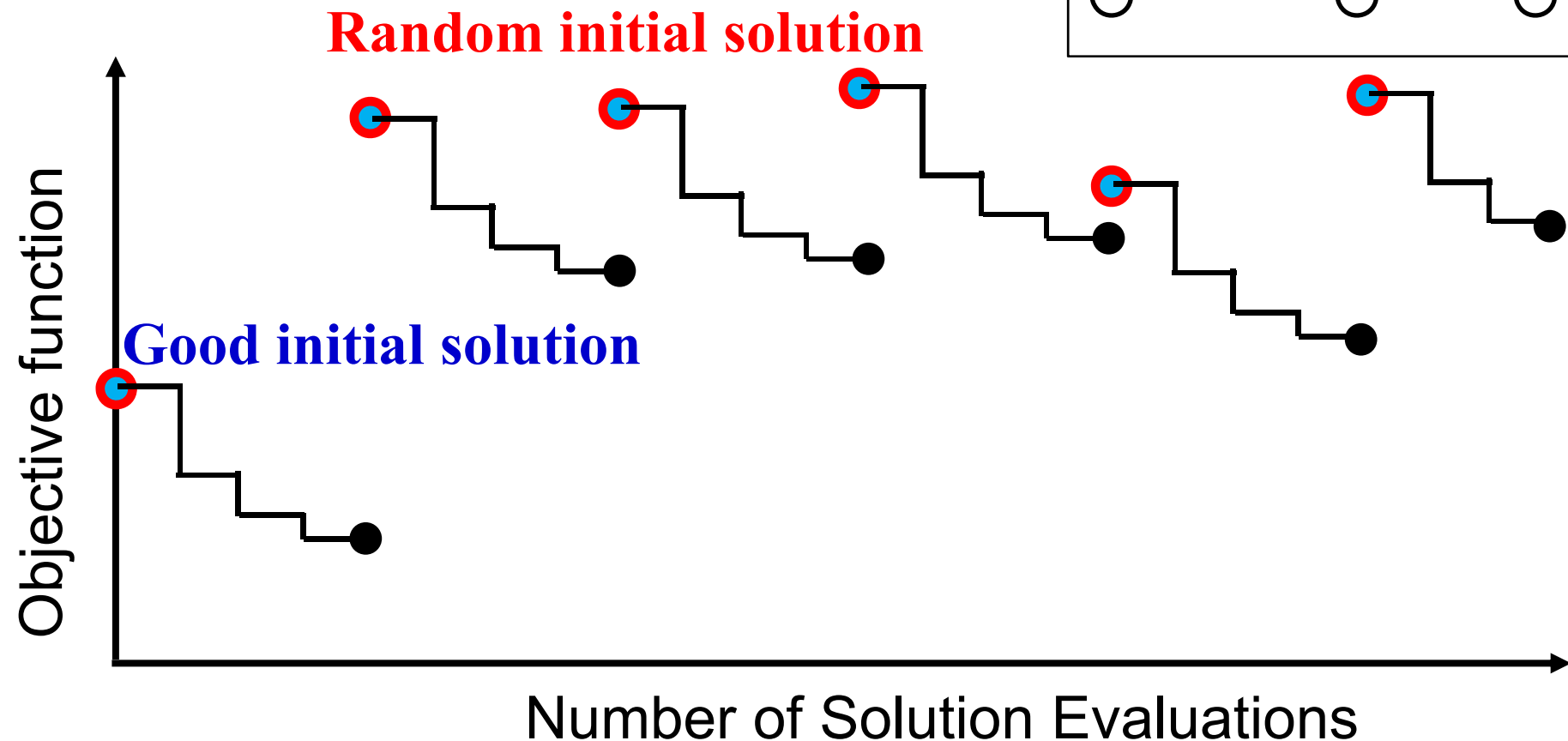
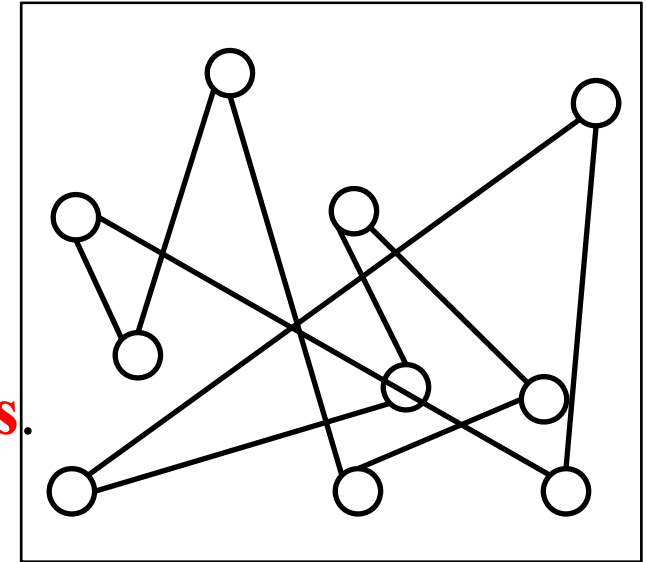
- First initial solution: **Greedy solution.**
- Other initial solutions: Random solutions.



Main Issue

Choice of Initial Solutions:

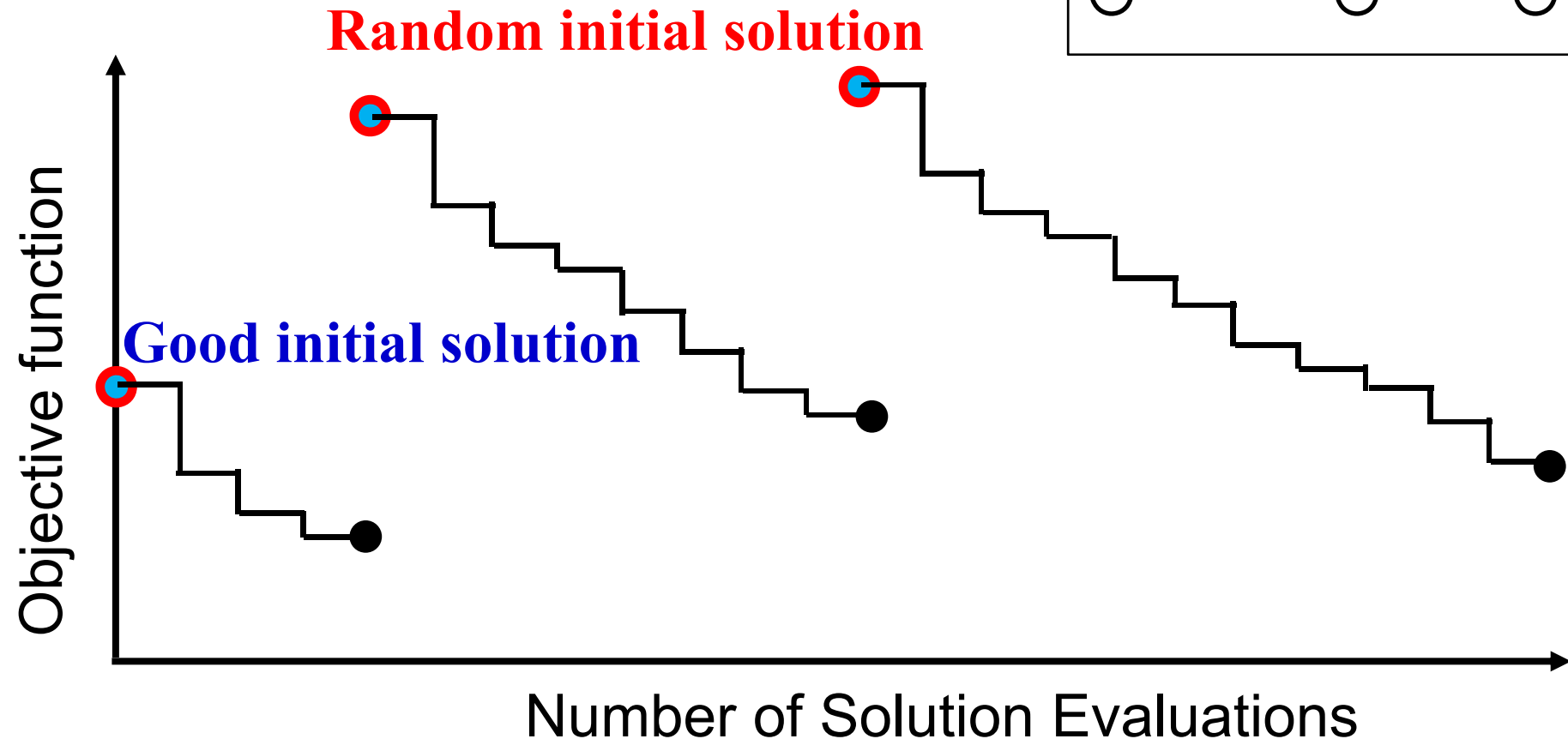
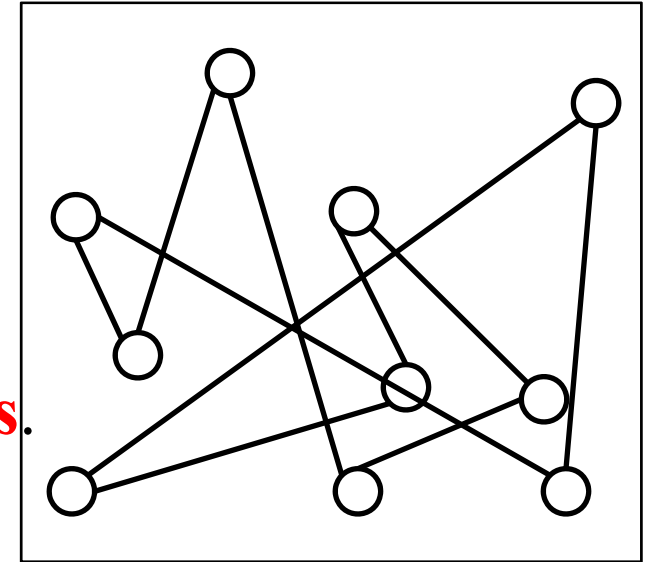
- First initial solution: Greedy solution.
- Other initial solutions: **Random solutions**.



Main Issue

Choice of Initial Solutions:

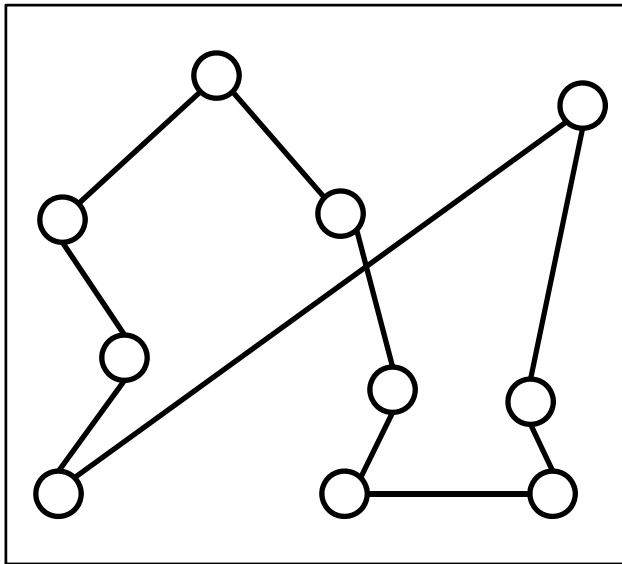
- First initial solution: Greedy solution.
- Other initial solutions: **Random solutions**.



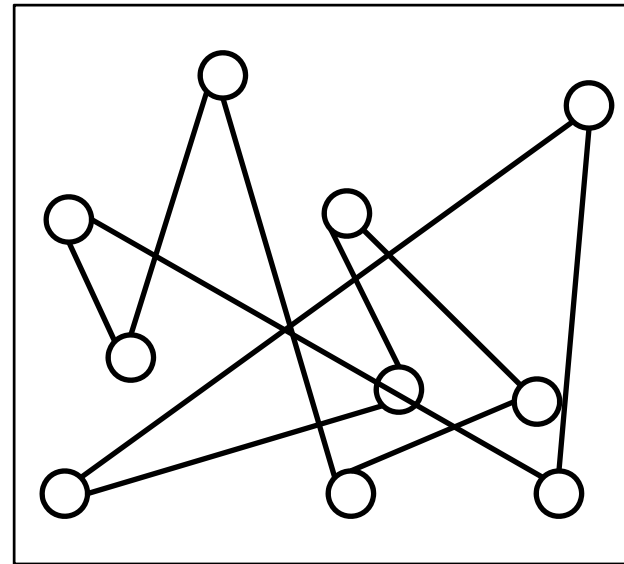
Main Issue

Choice of Initial Solutions:

- First initial solution: Greedy solution.
- Other initial solutions: Random solutions.



Good greedy solution

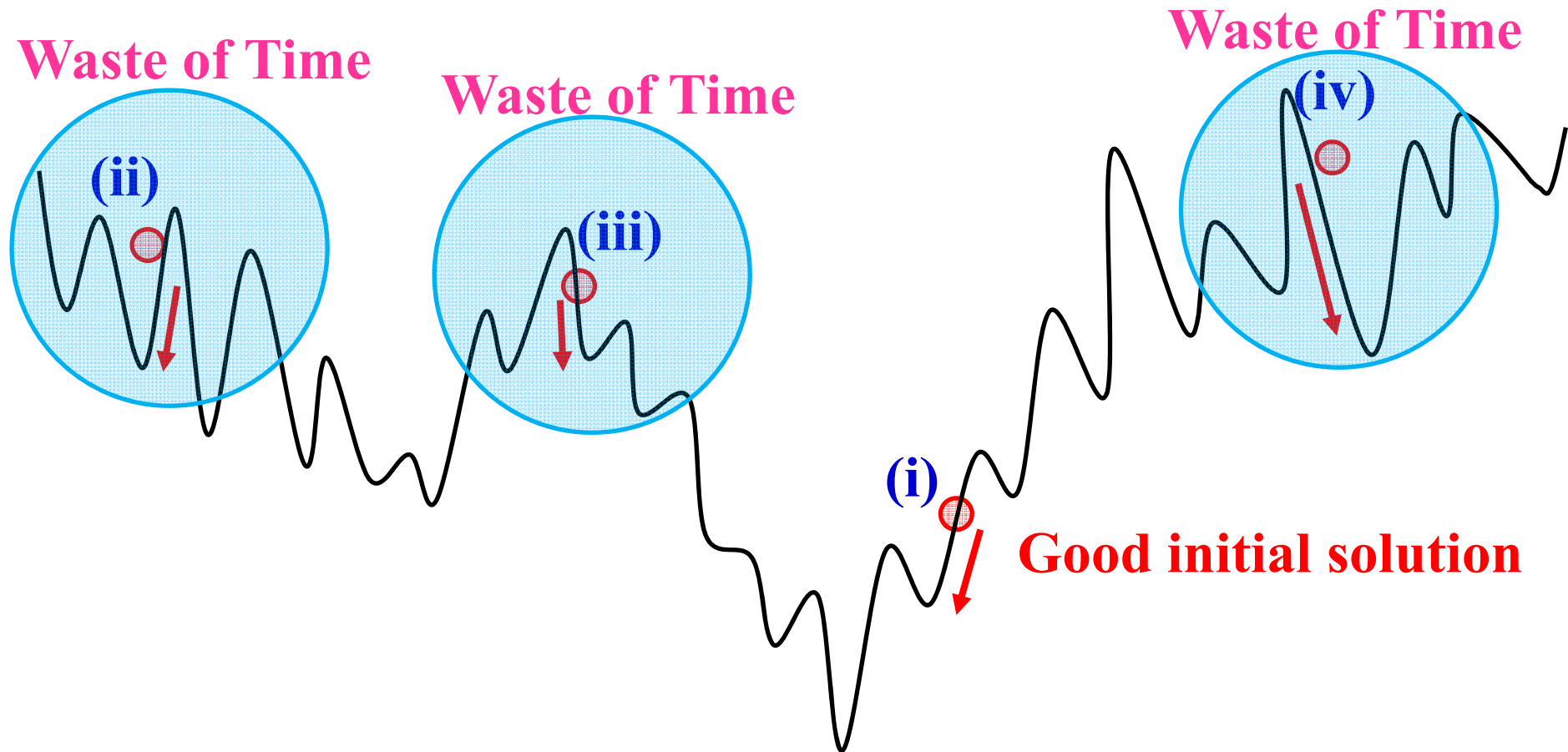


Poor random solution

Main Issue

Choice of Initial Solutions:

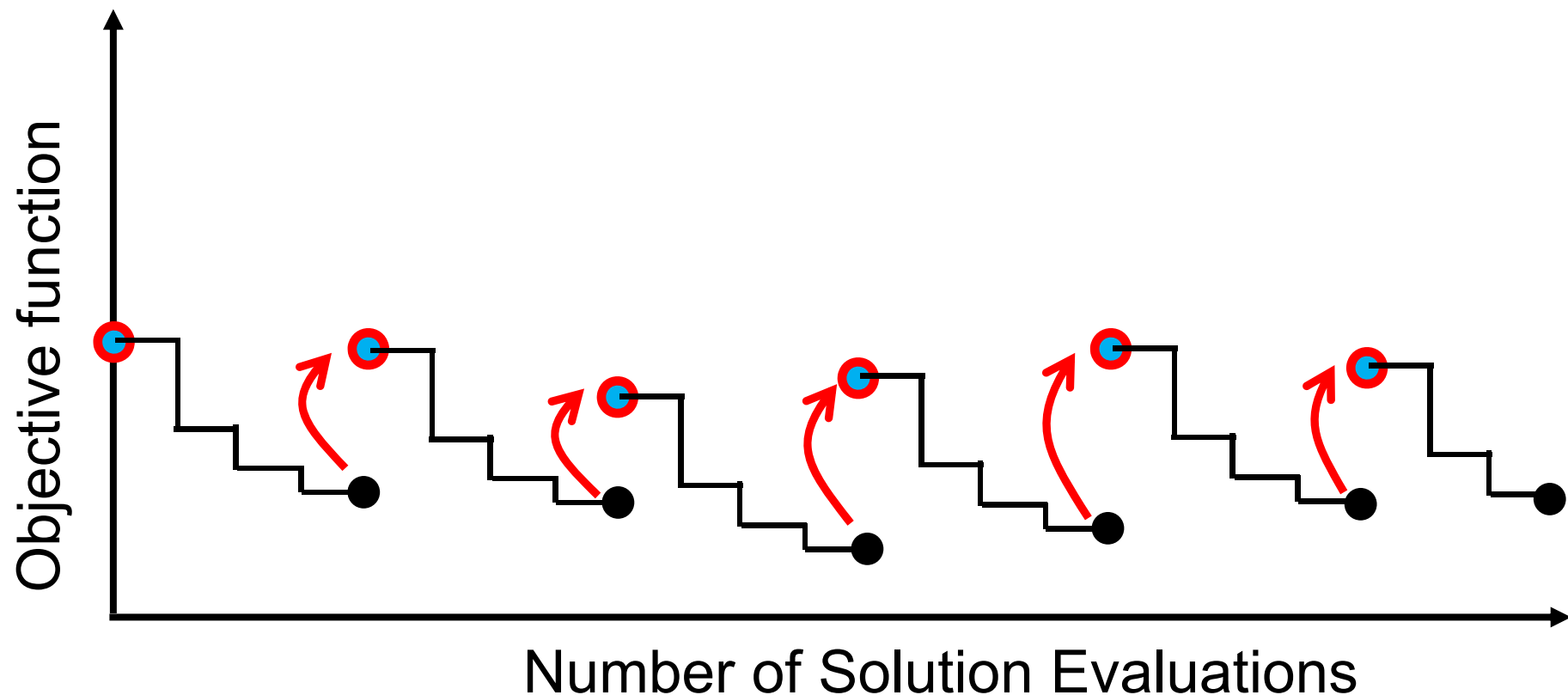
- First initial solution: Greedy solution.
- Other initial solutions: Random solutions.



One Idea

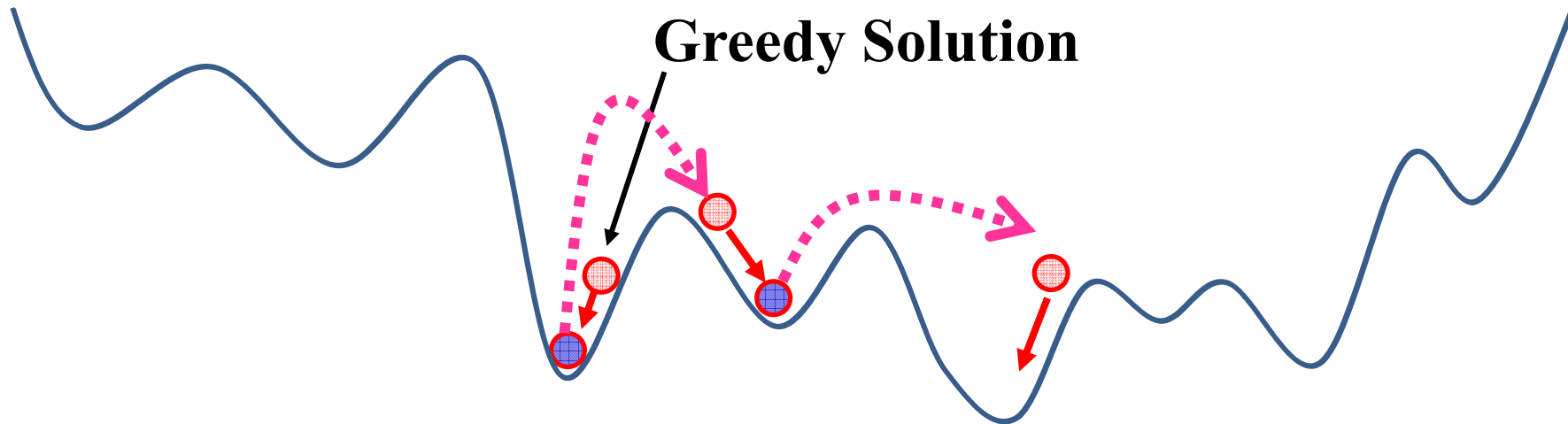
Use of the Current Solution

- First initial solution: Greedy solution.
- Other initial solutions: Generate from the final solution of the previous LS run

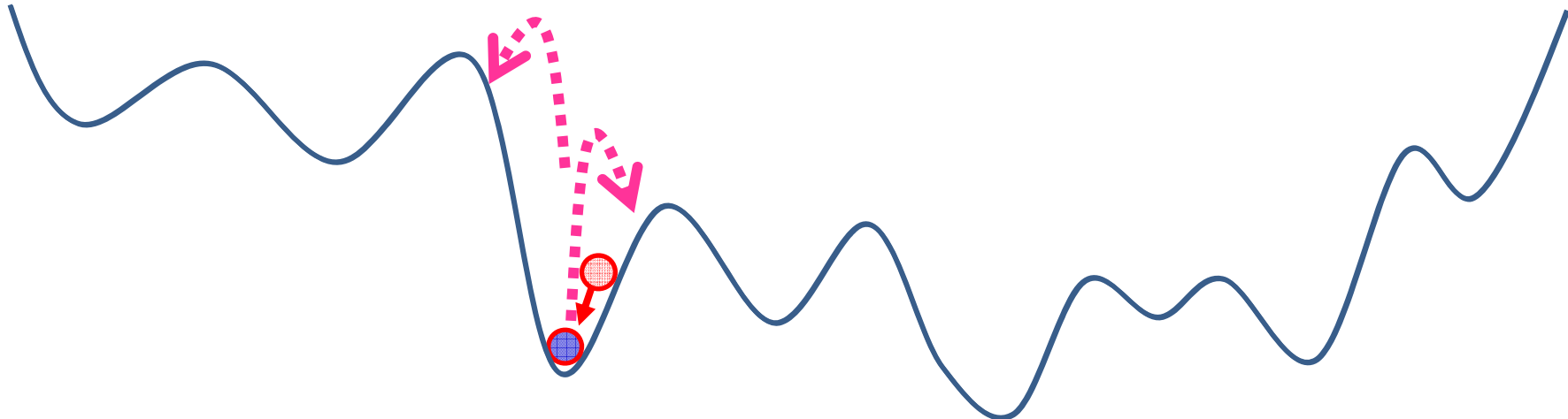


How to specify initial solutions

A new solution form the final solution

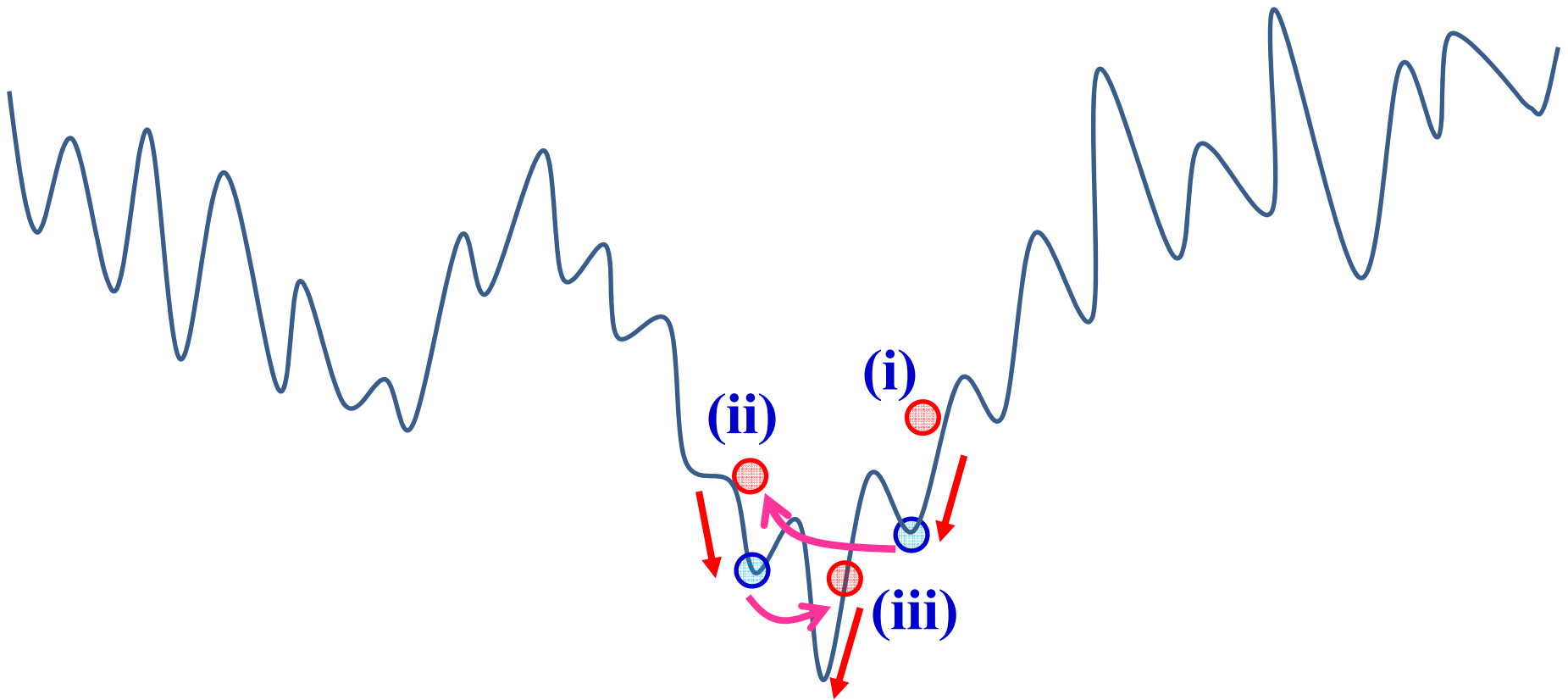


The following undesirable situation can happen.



Use of the final Solution

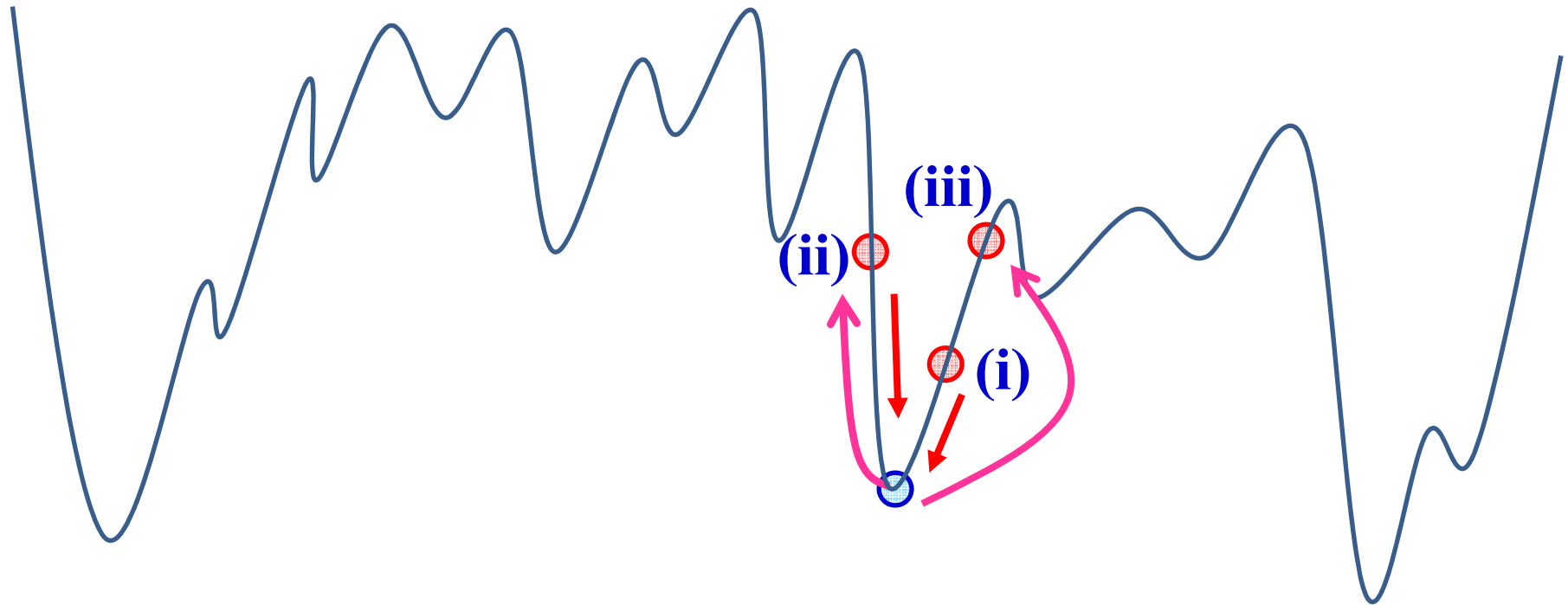
- First initial solution: Greedy solution.
- Other initial solutions: Generate from the final solution



The optimal solution can be efficiently searched by local search..

Use of the final Solution

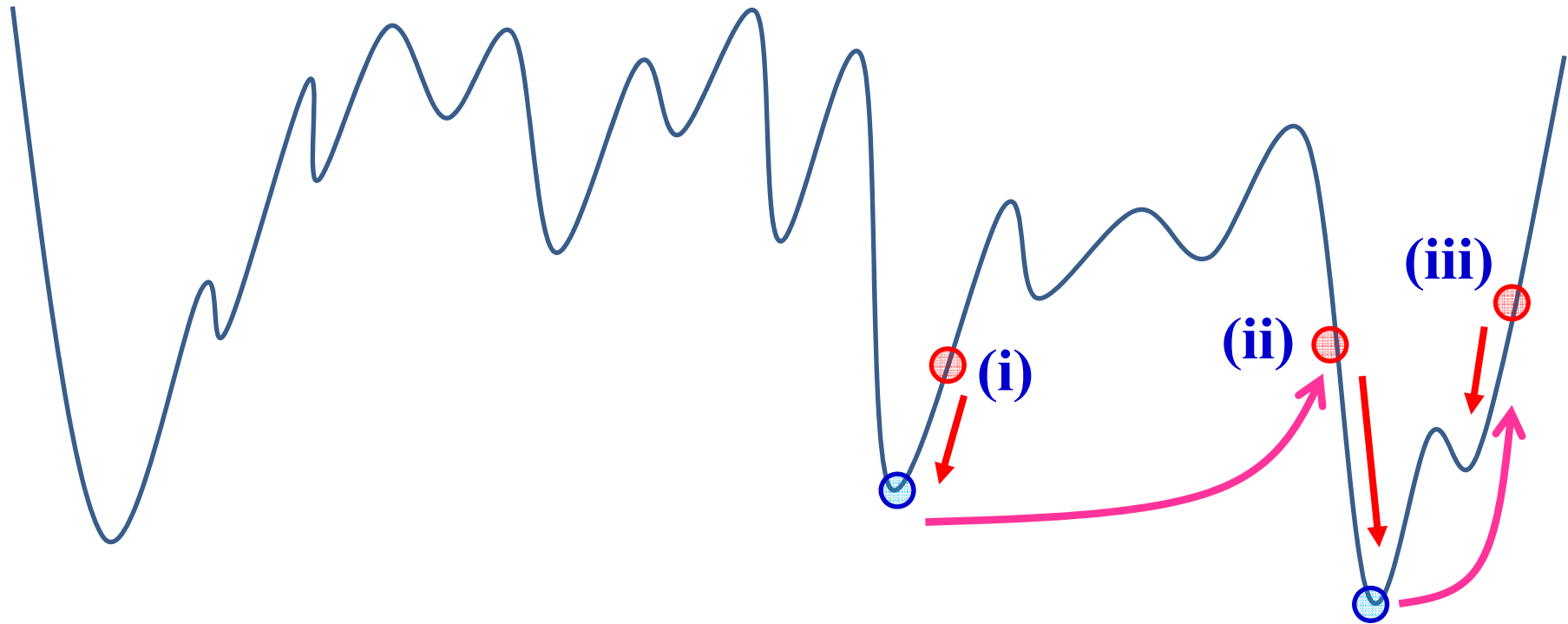
- First initial solution: Greedy solution.
- Other initial solutions: Generate from the final solution



The same solution can be obtained from different initial solutions.

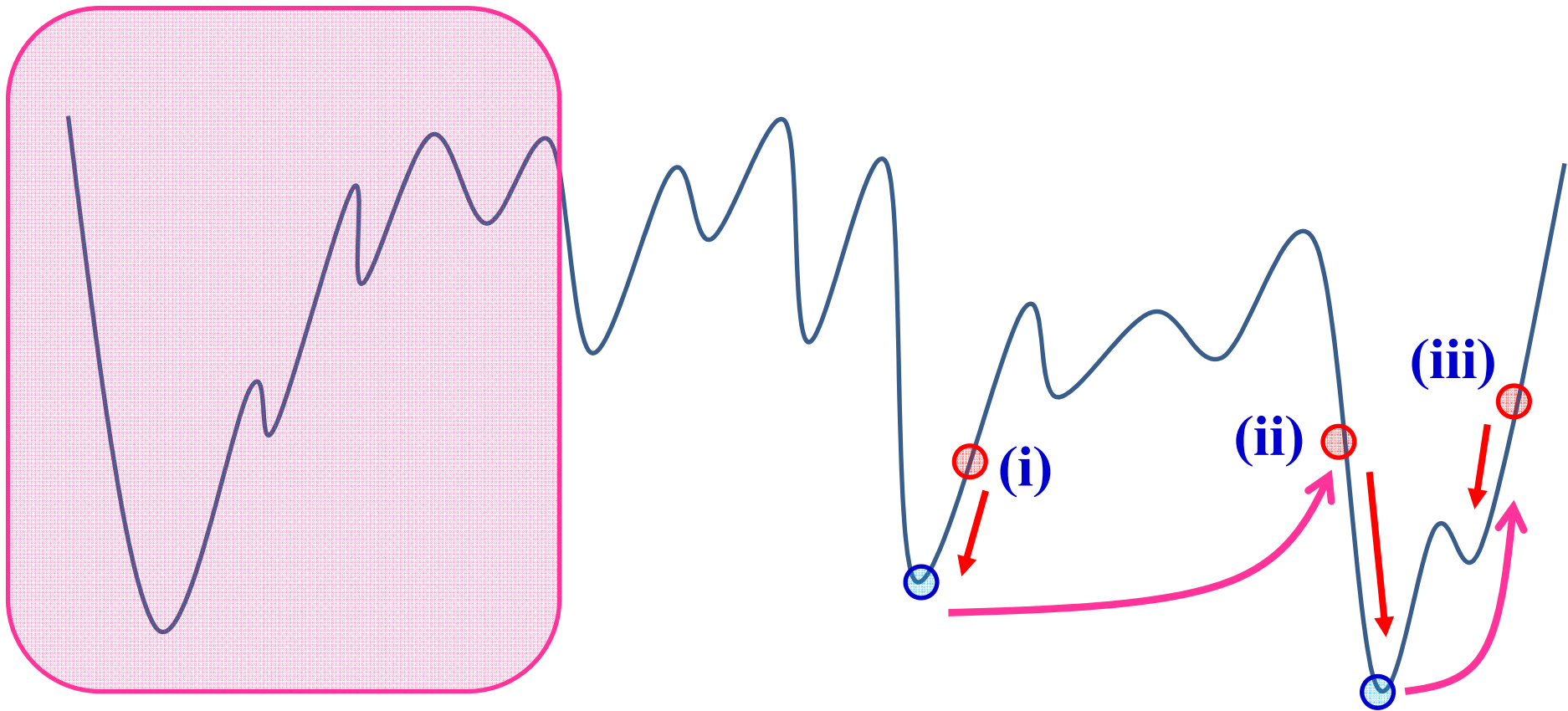
Good Initial Solutions

- Close to a good local solution
- Search for a different local solution



Good Initial Solutions

- Close to a good local solution
- Search for a different local solution

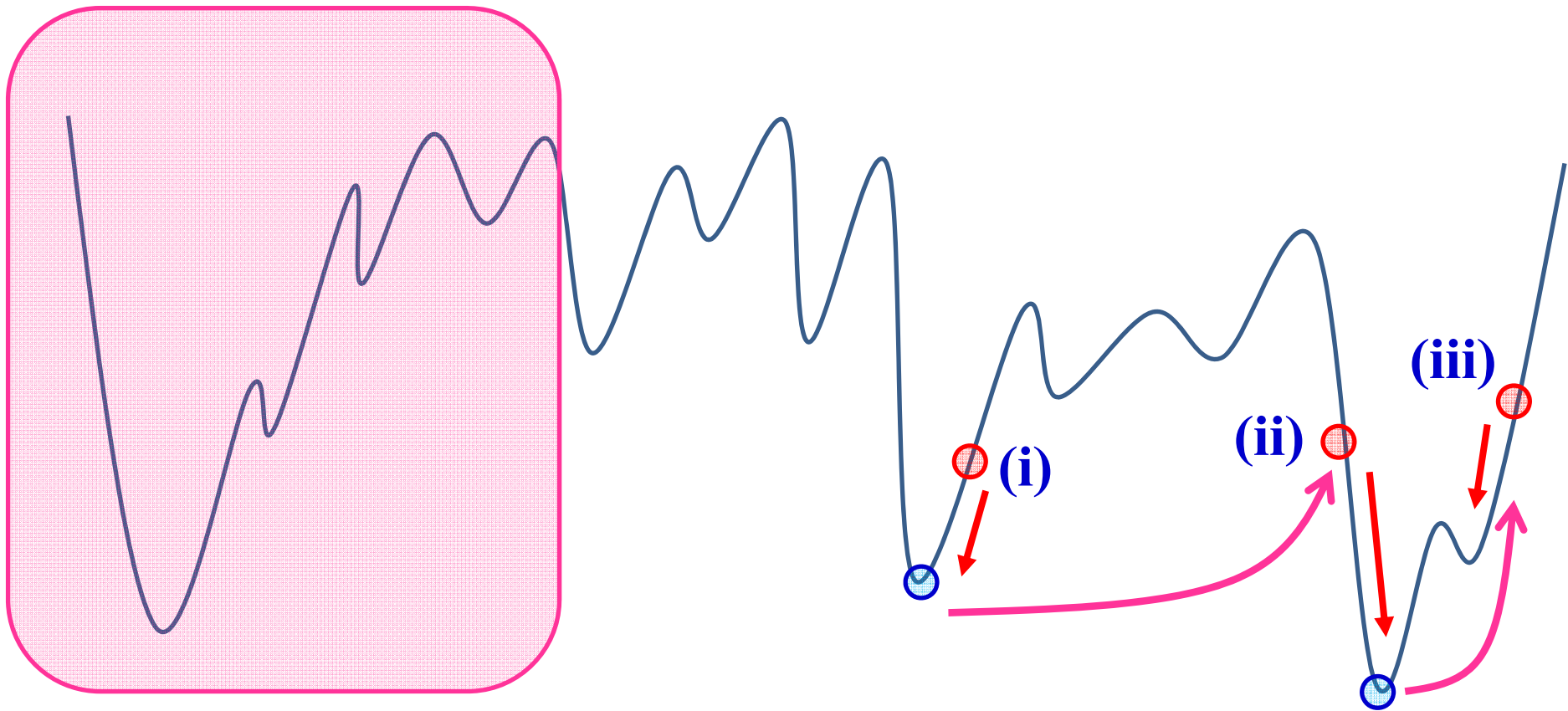


The search of this region is difficult.

Initial Solutions

Random: Wide search region & Inefficient search

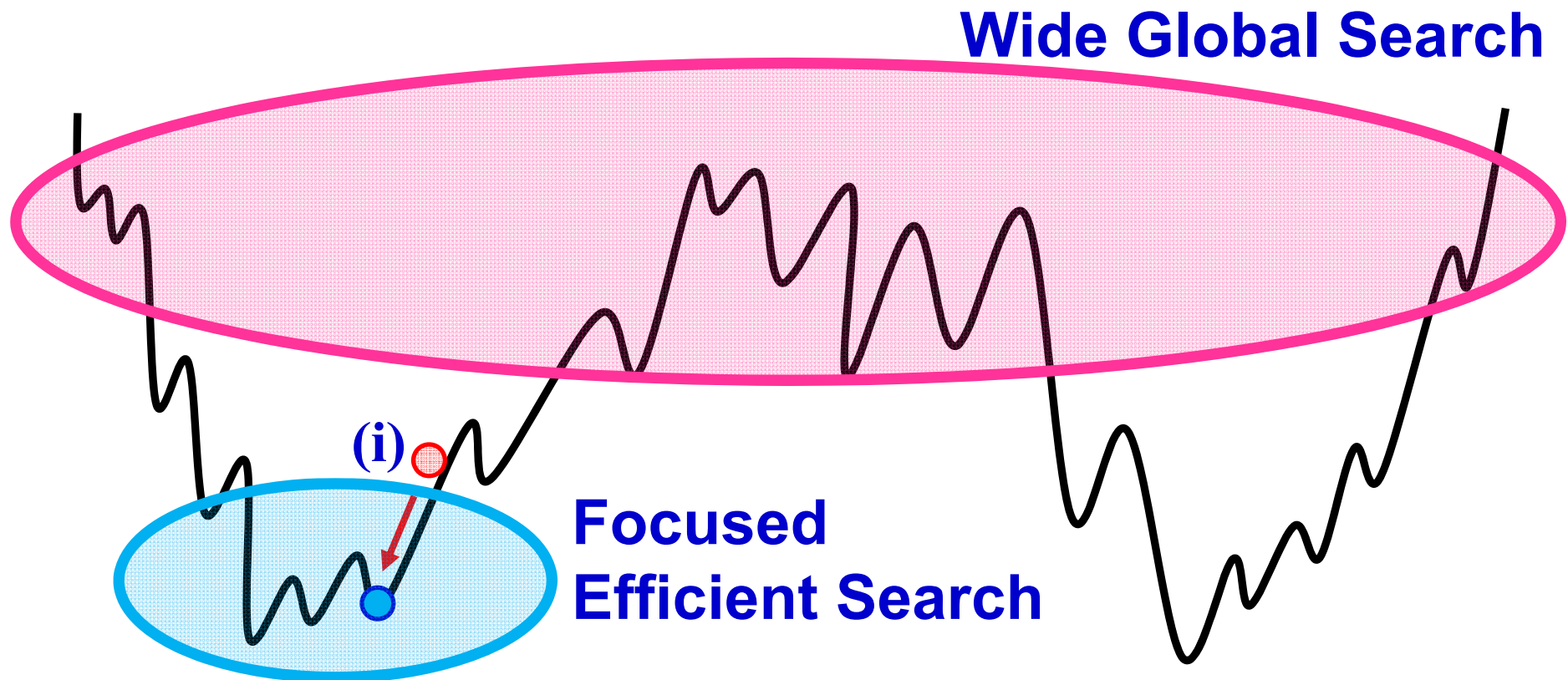
Focused: Narrow search region & Efficient search



The search of this region is difficult.

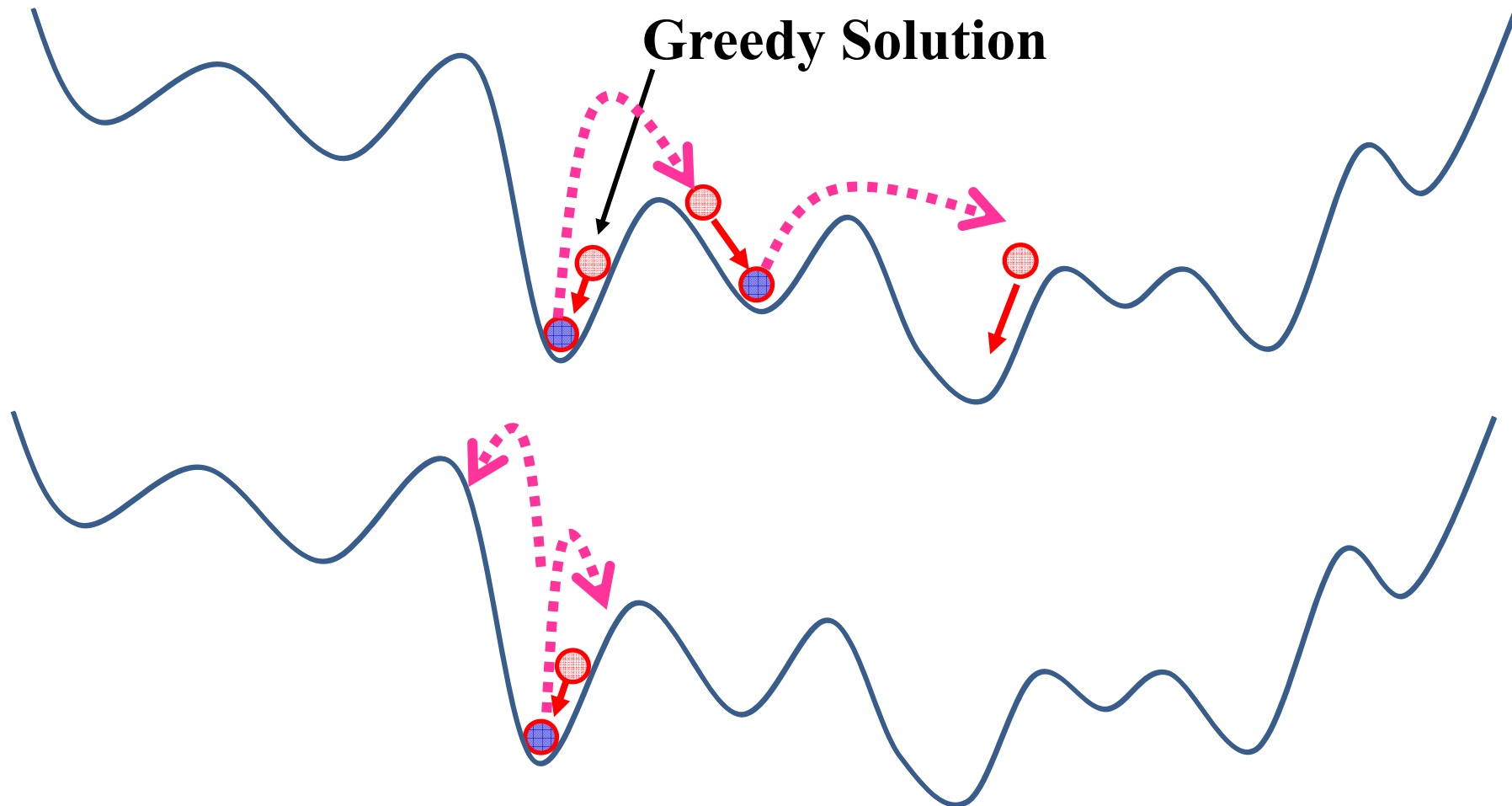
Optimization Algorithm Design:

Find a good balance between the wide global search and the focused efficient search (the good balance depends on the problem size and the available computation time)



General Guideline:

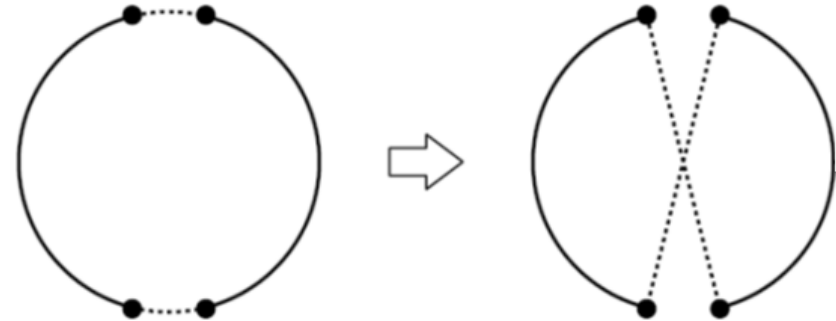
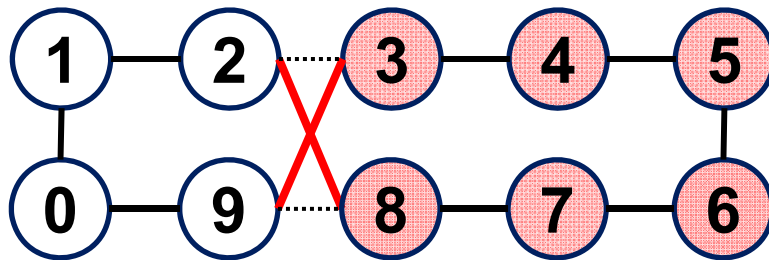
Generate good initial solutions which are not likely to go back to the same final solution.



Lab Session Task 1

Local search

- Neighborhood structure: Inversion (two-edge change)



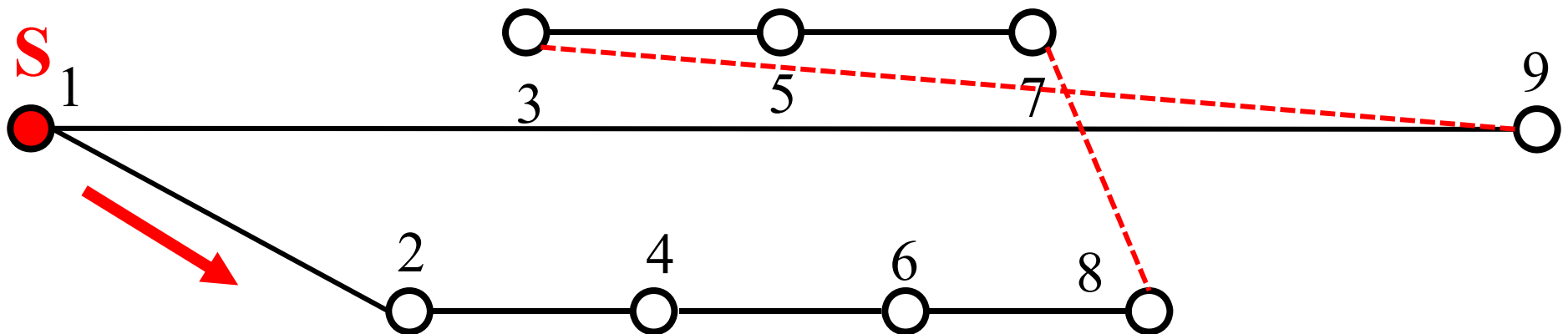
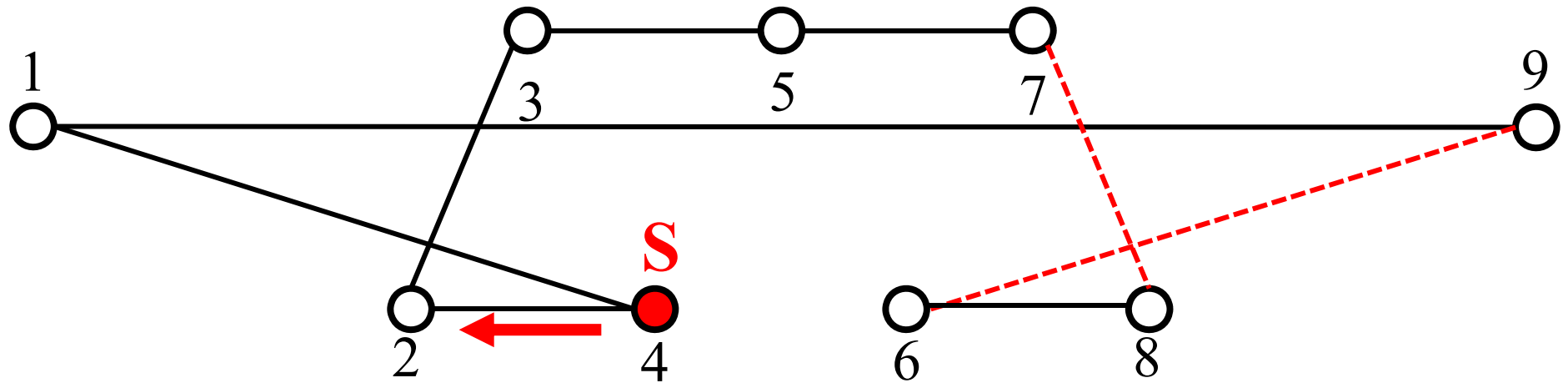
Creation of a new initial solution:

- 1st initial solution: Greedy solution

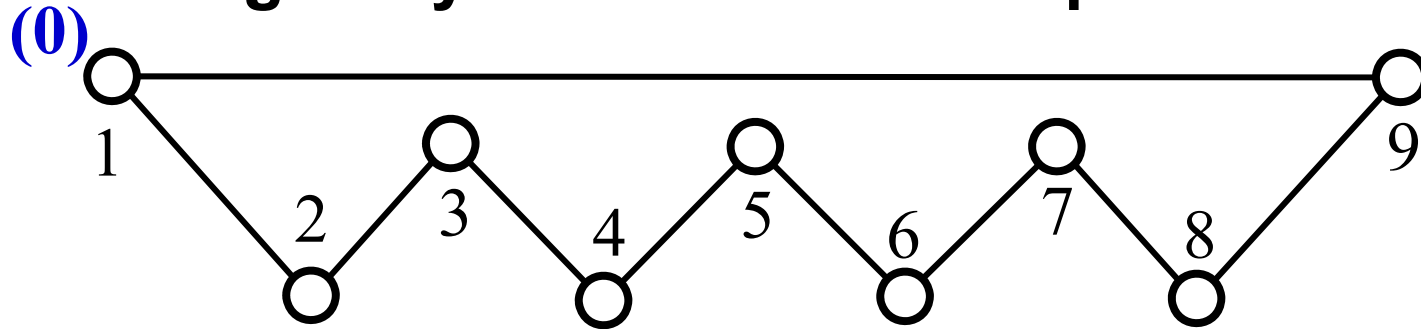
Task 1

Create a TSP problem where a non-optimal greedy solution is obtained from an initial city and the obtained greedy solution cannot be improved by inversion-based local search. It is enough to show that one greedy solution cannot be improved. You do not have to show that any greedy solution from an arbitrary initial city cannot be improved.

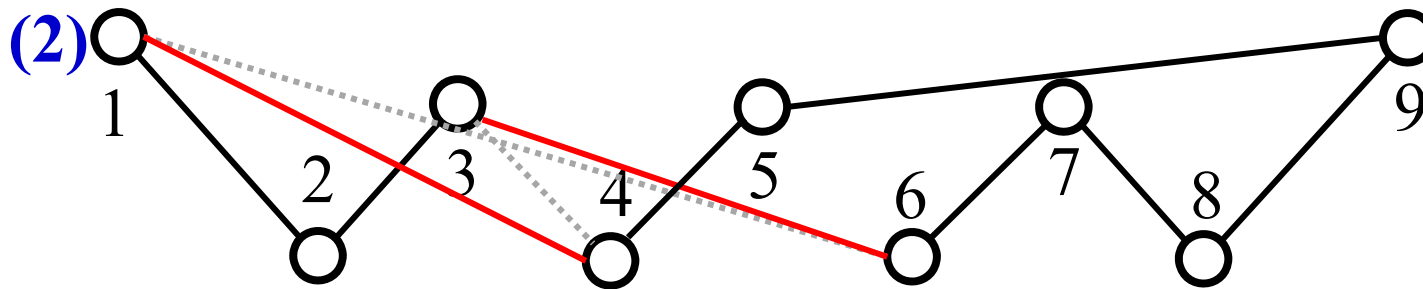
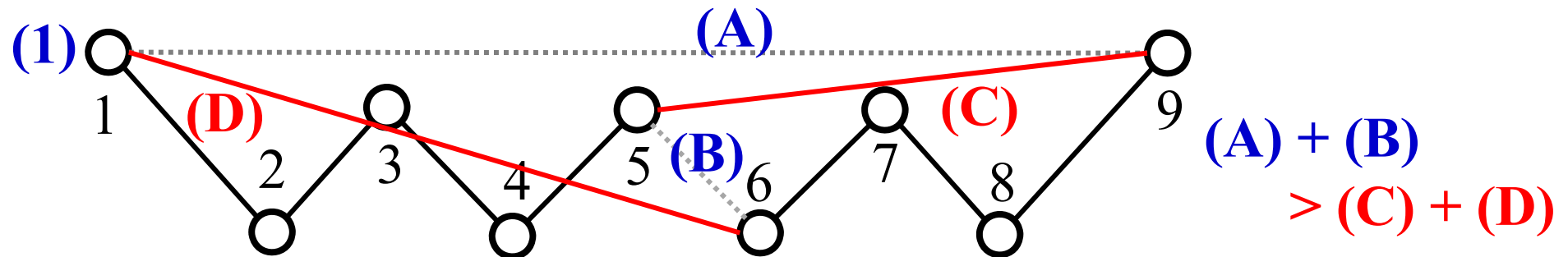
TSP Problem: In many cases, greedy solutions can be improved by inversion-based local search.

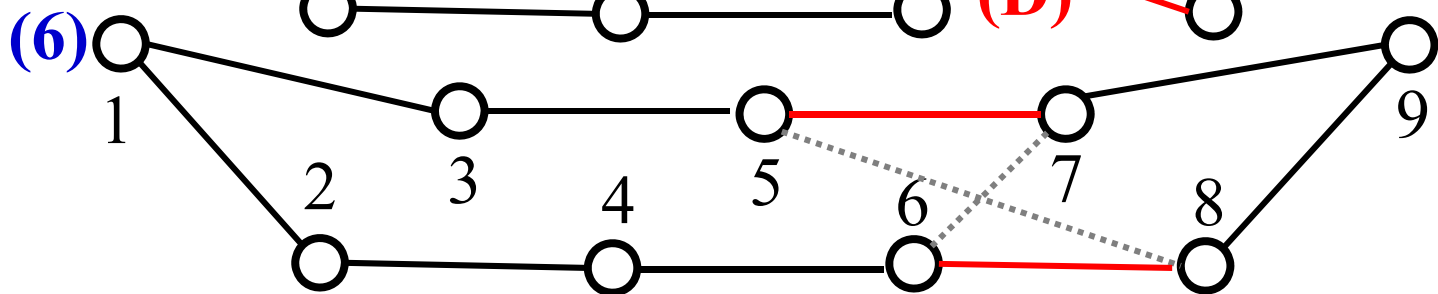
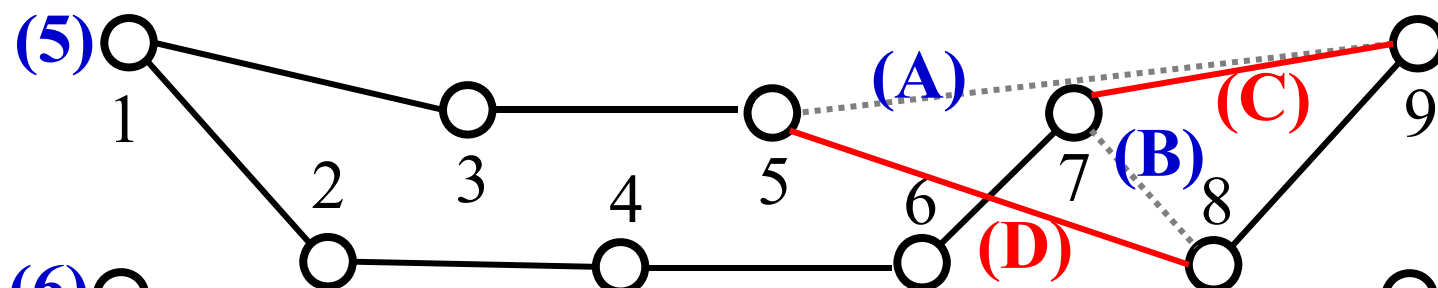
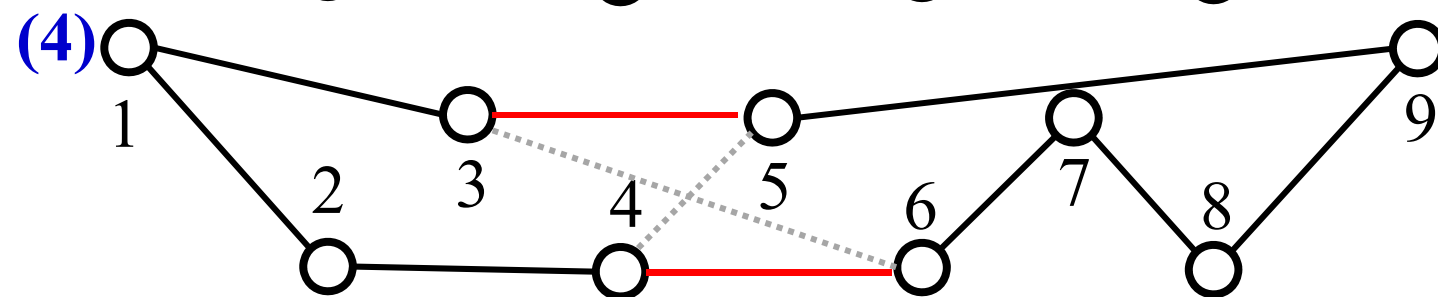
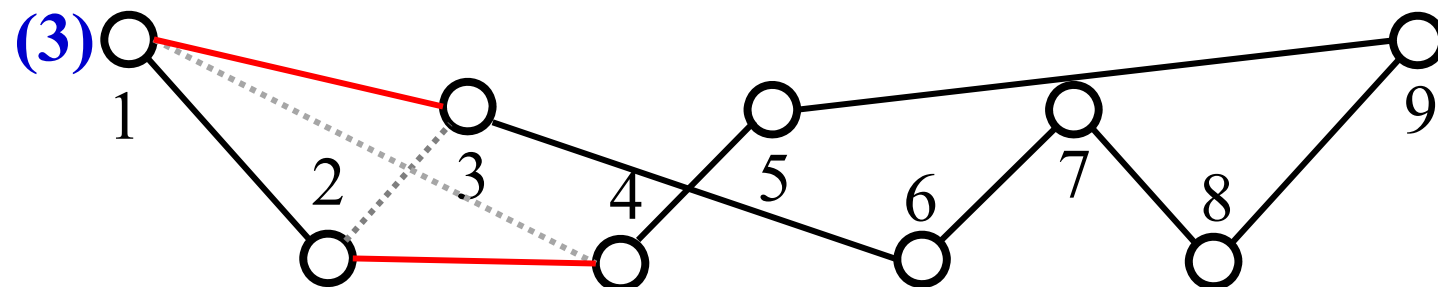
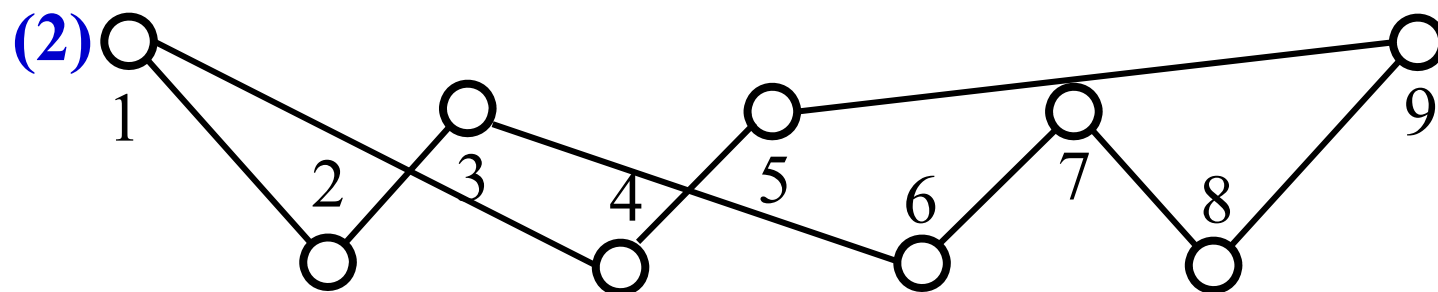


This greedy tour looks local optimal



However, this can be improved by inversion-based LS.



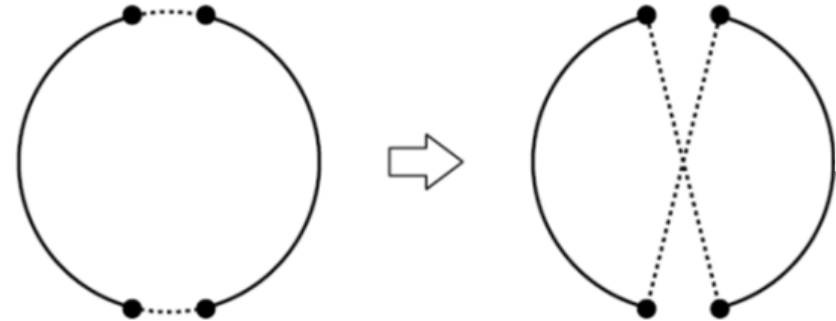
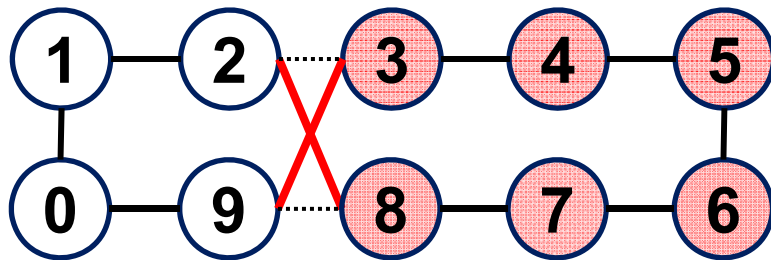


$$\begin{aligned} & \text{(A)} + \text{(B)} \\ & > \text{(C)} + \text{(D)} \end{aligned}$$

Lab Session Task 2

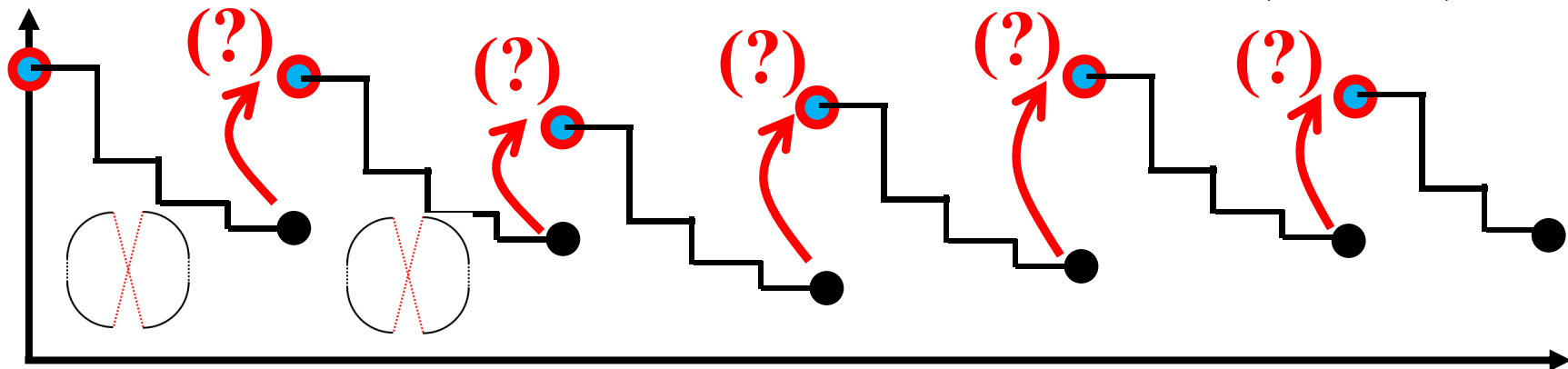
Local search

- Neighborhood structure: Inversion (two-edge change)



Creation of a new initial solution:

- 1st initial solution: Greedy solution
- Other initial solutions: Explain your own idea about how to create an initial solution for each local search run (Task 2).



The graph illustrates the performance of two algorithms in minimizing tour length. The y-axis represents the 'Tour Length (Minimization)' and the x-axis represents the 'Number of Solution Evaluations'.

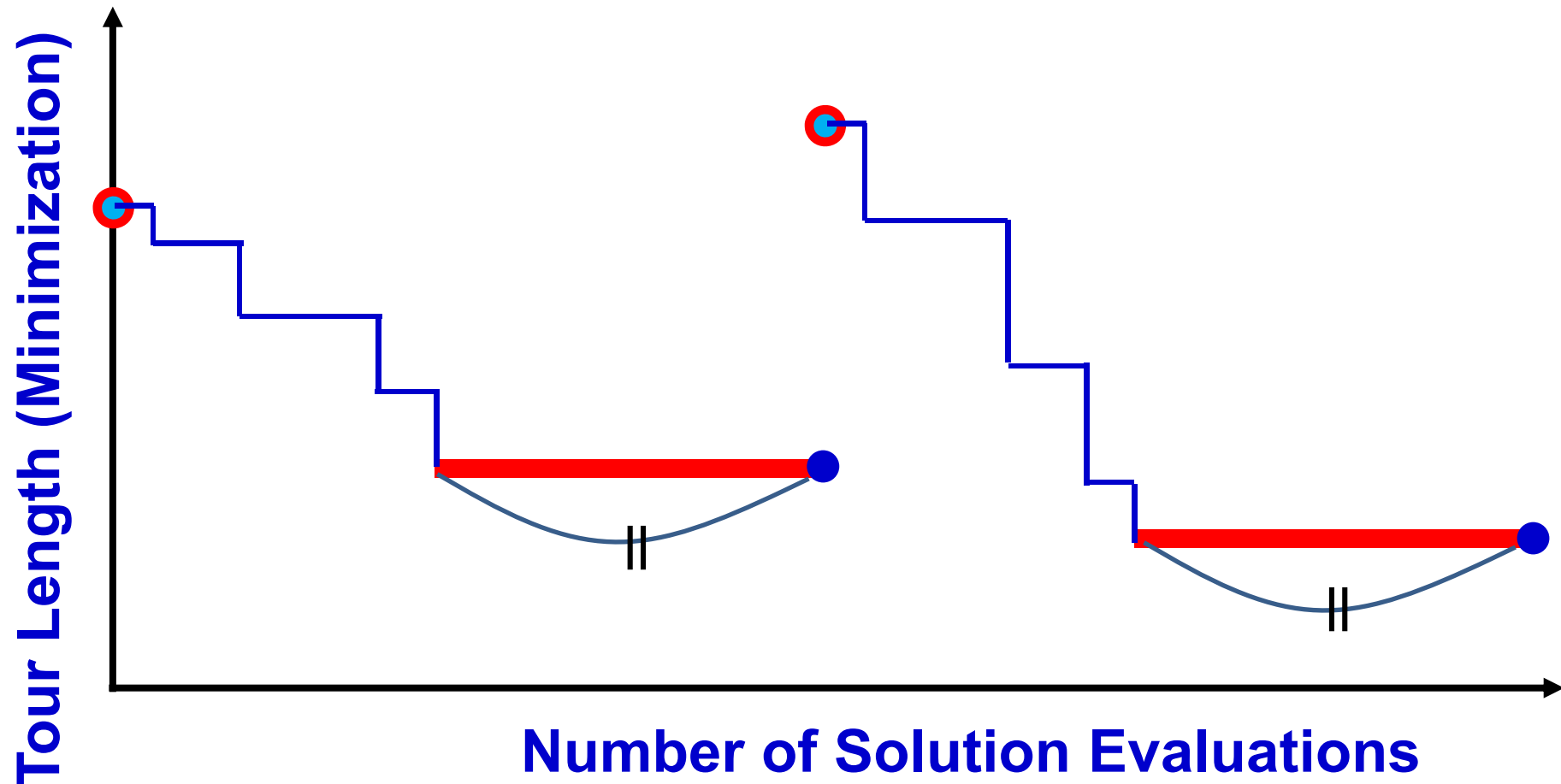
The blue line represents the greedy algorithm, which shows a steady decrease in tour length as the number of evaluations increases. The pink line represents the hill-climbing algorithm, which shows a more rapid initial decrease in tour length, followed by a plateau, and then a further decrease after a 'Best Move' is found. The hill-climbing algorithm eventually reaches a lower tour length than the greedy algorithm.

Key features of the graph include:

- Best Move:** A red circle on the pink line indicating the point where the hill-climbing algorithm found the best move.
- First Move:** A blue circle on the blue line indicating the point where the greedy algorithm found the first move.

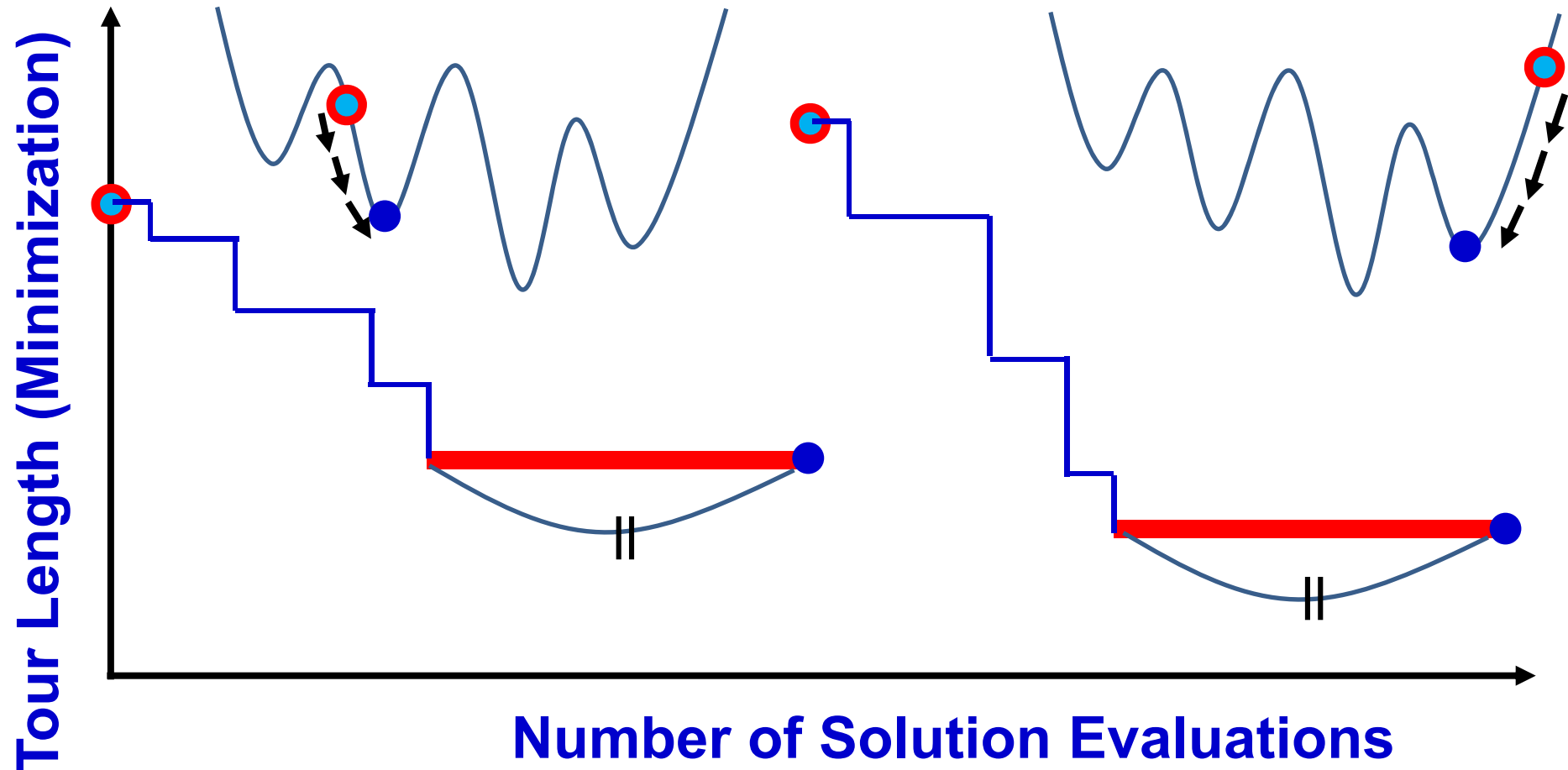
It may be a good idea to use the first move for examining more initial solutions (i.e., for finding more local solutions)

Early Termination



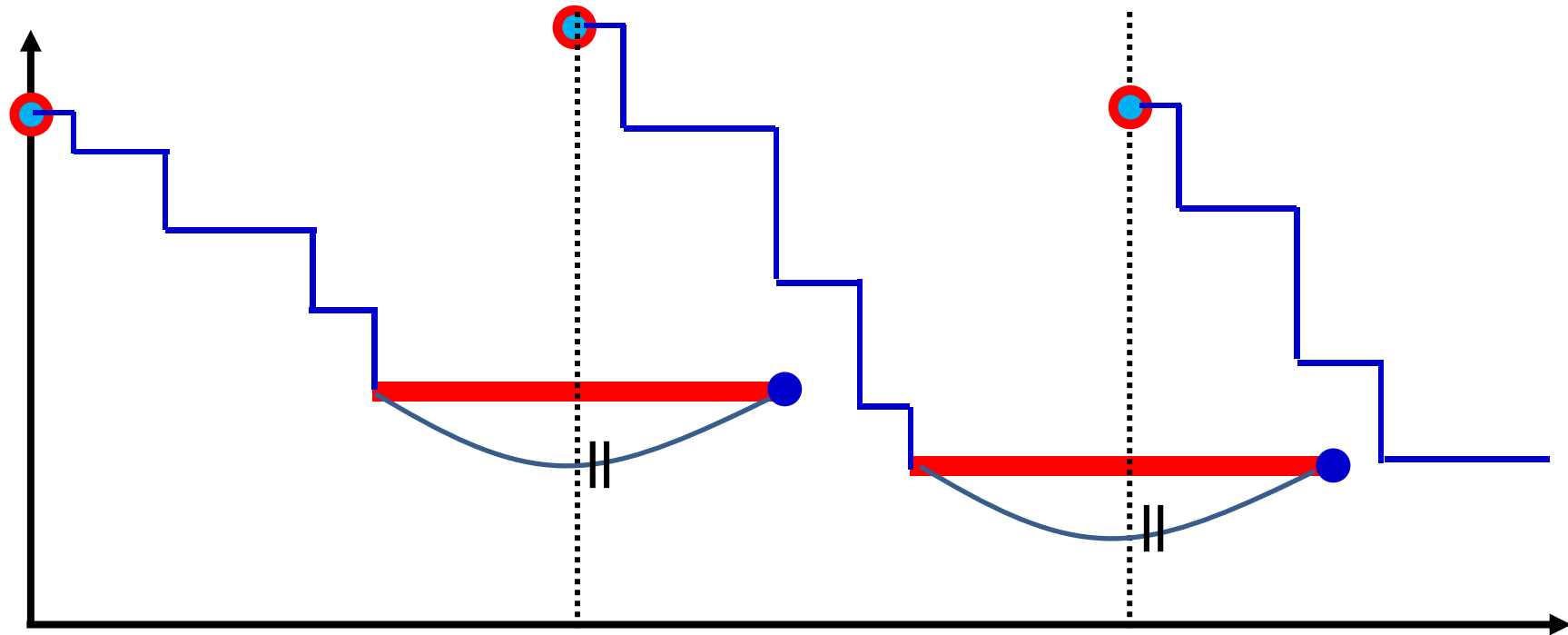
Even in the first move local search, all neighbors are examined before the termination of local search. This usually spends a long computation time with no performance improvement, which looks waste of time.

Early Termination



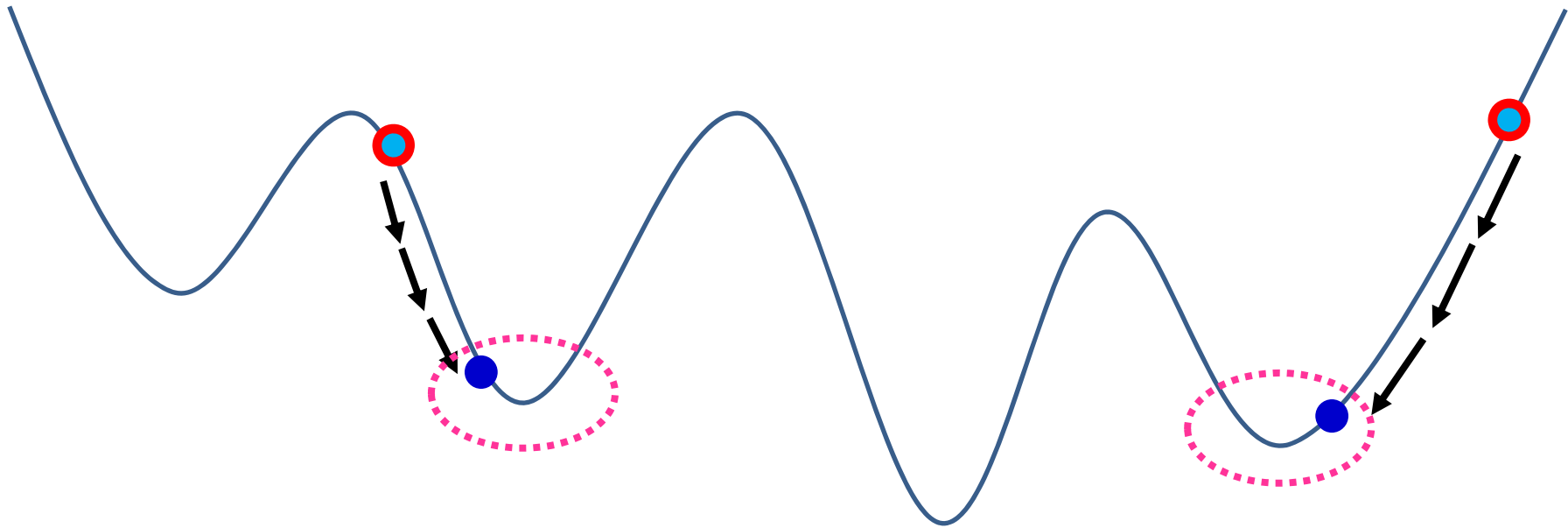
Even in the first move local search, all neighbors are examined before the termination of local search. This usually spends a long computation time with no performance improvement, which looks waste of time.

Early Termination

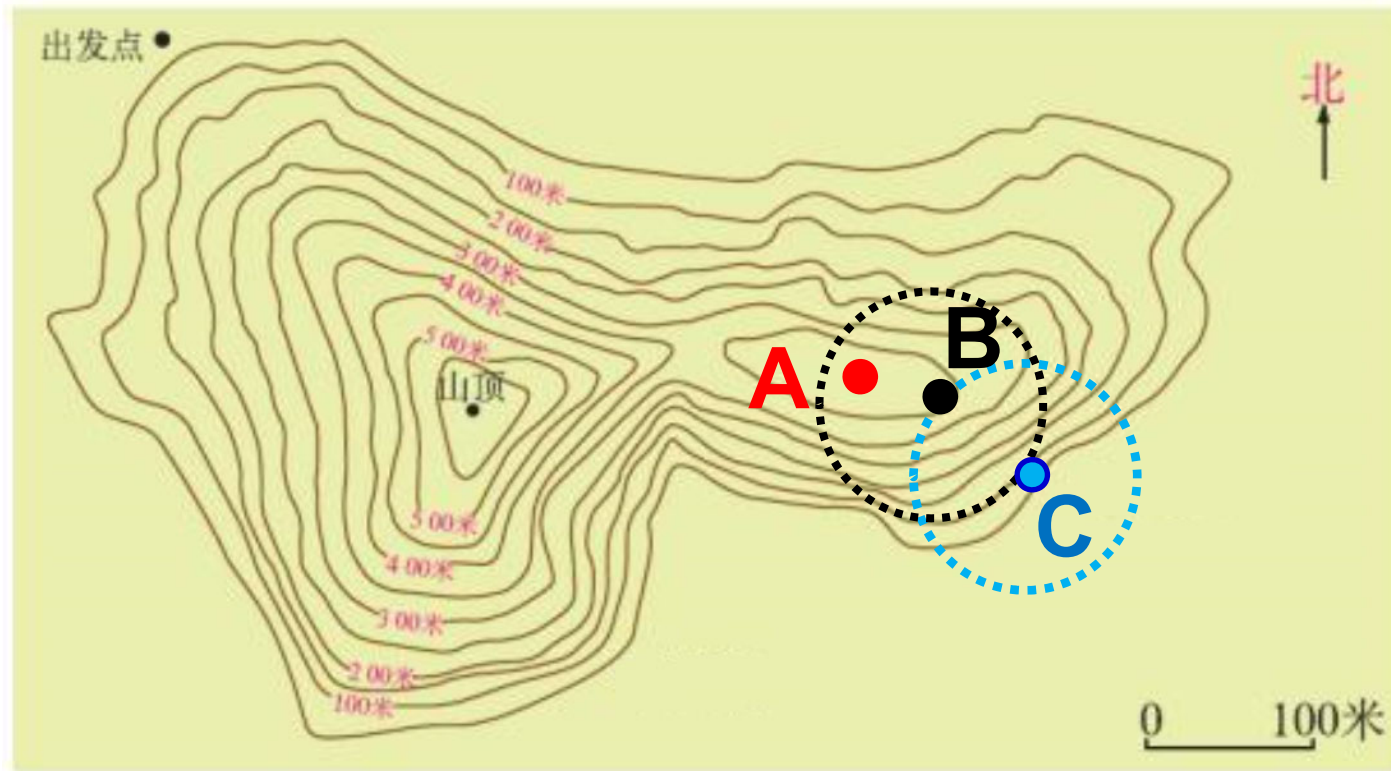


It may be a good idea for efficient search to terminate the local search before examining all neighbors (e.g., restart from a new initial solution after examining 100 neighbors or 50% of neighbors).

Early Termination



One negative effect is that the local search will terminate before finding a local solution. The question is which is a better search strategy between (i) to carefully search for a local solution around the current solution, and (ii) to quickly search for better solutions from many initial solutions. Usually, the improvement by local search around a local solution is not large. Thus, (ii) is more efficient in many cases.



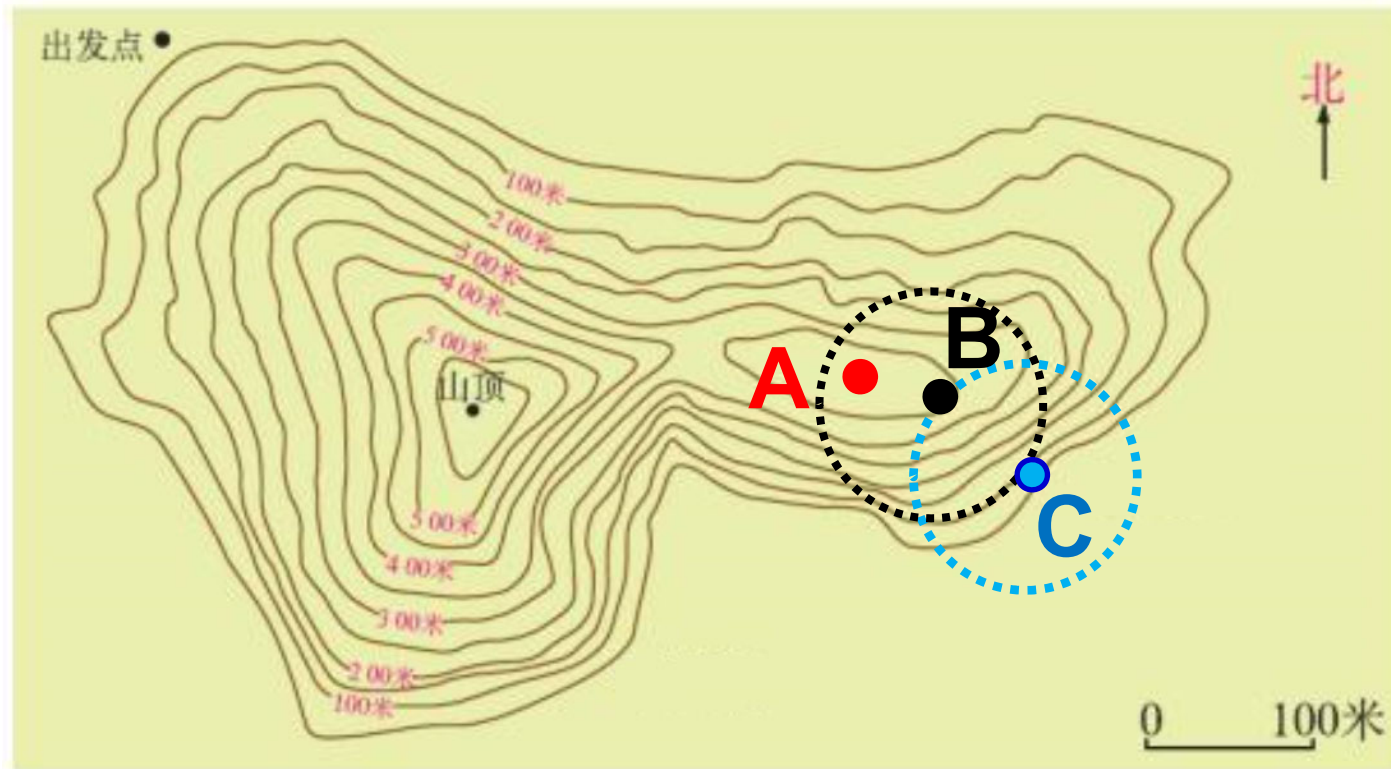
Around the local solution A: Current Solution B

(i) It is not easy to find a better solution

==> Early termination before finding the local solution.

(ii) The improvement by local search is small.

==> Early termination is not a big problem.



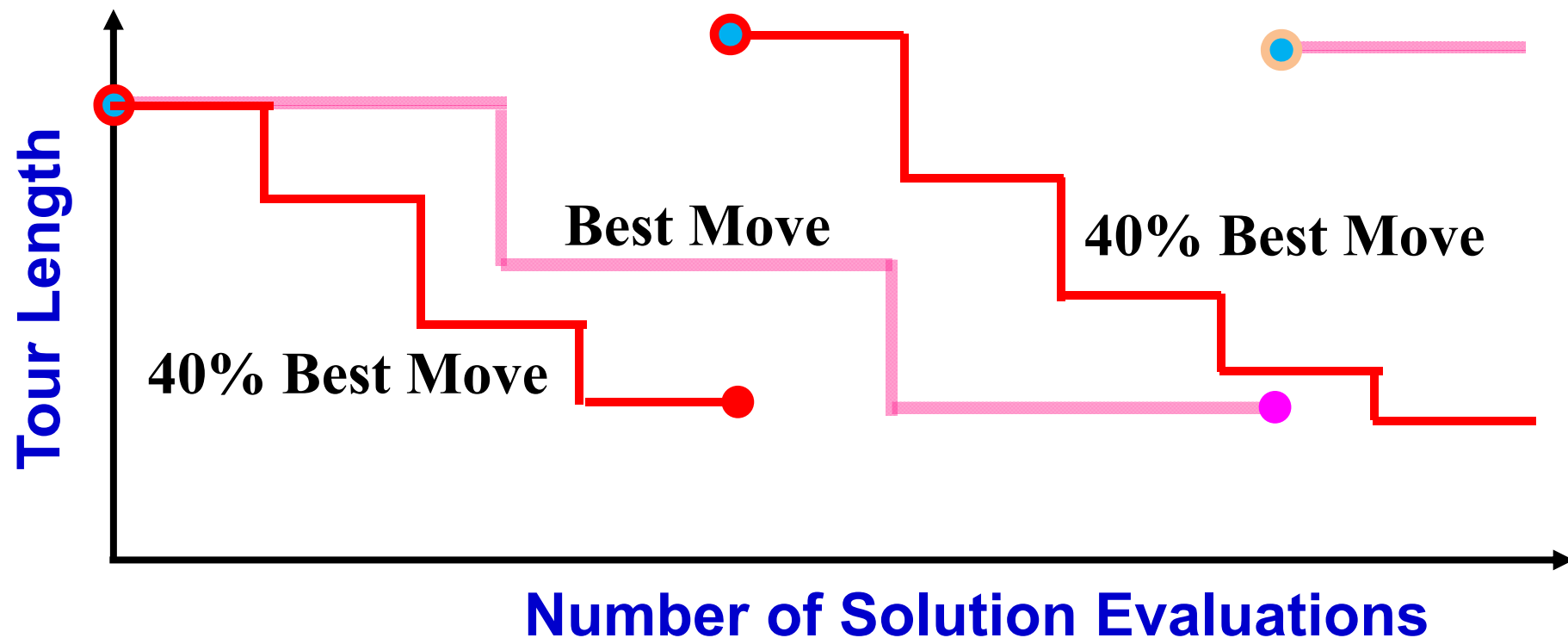
Around the local solution A: Current Solution B

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==> Early termination before finding the local solution.
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At solution C, it is easy to find a better solution.

Best Move with Early Termination

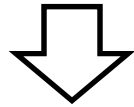
- (i) Examine $\alpha\%$ of neighbors.
- (ii) Choose the best neighbor among them.
- (iii) If the best neighbor is better than the current solution, move to the best neighbor. Otherwise, restart the local search from a new initial solution.



Local Search

Implementation Issues

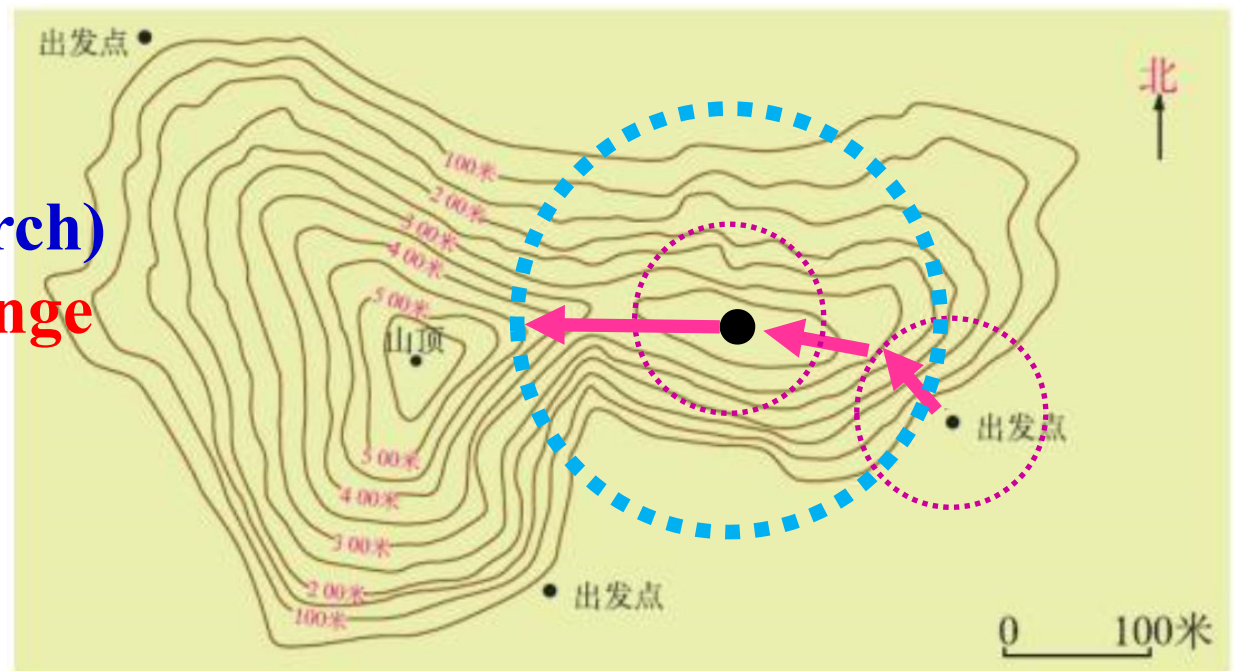
- (1) Specification of an initial solution
- (2) Specification of a neighborhood structure
- (3) Choice between the first move and the best move



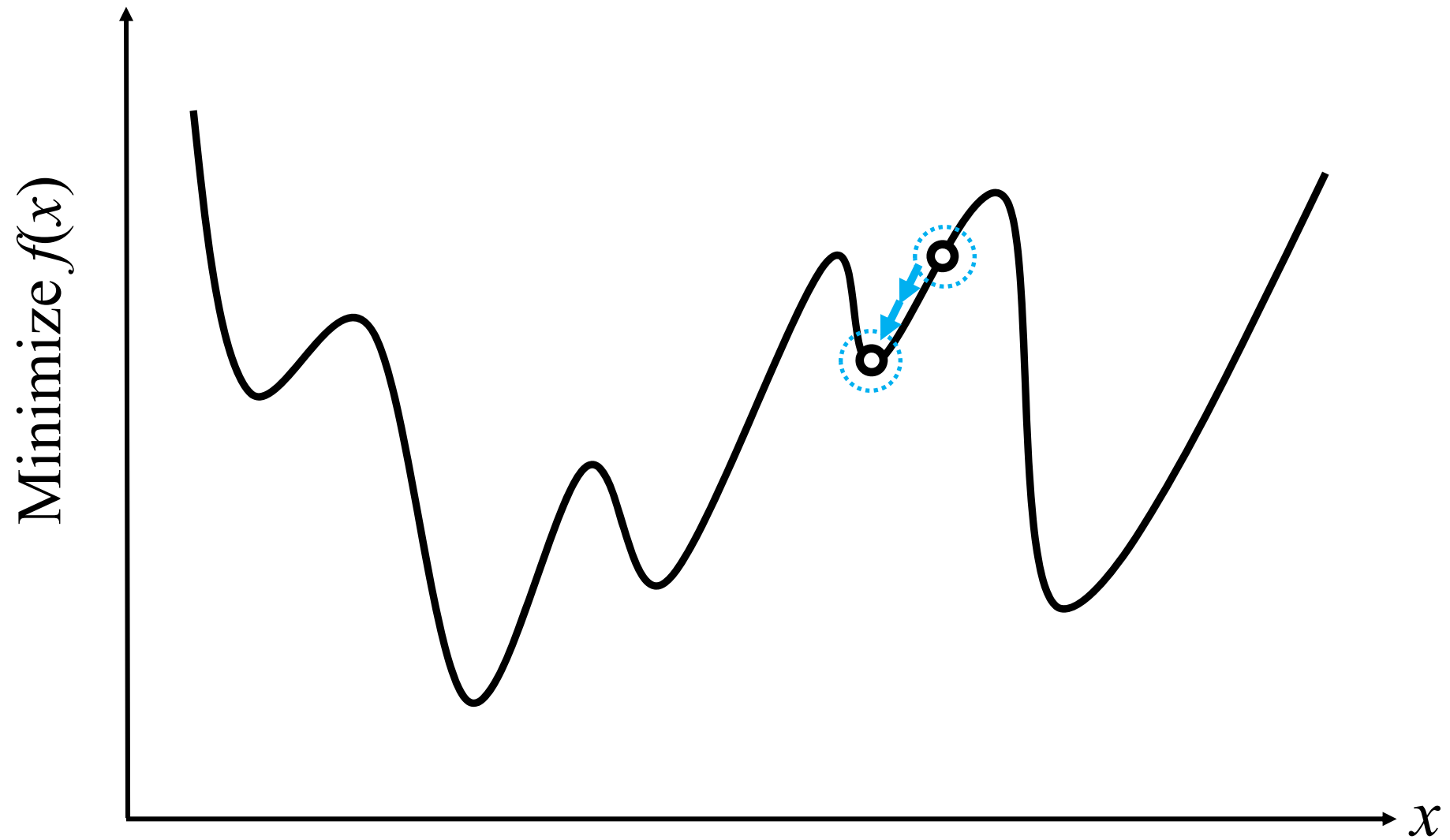
Local Solution

Next Step

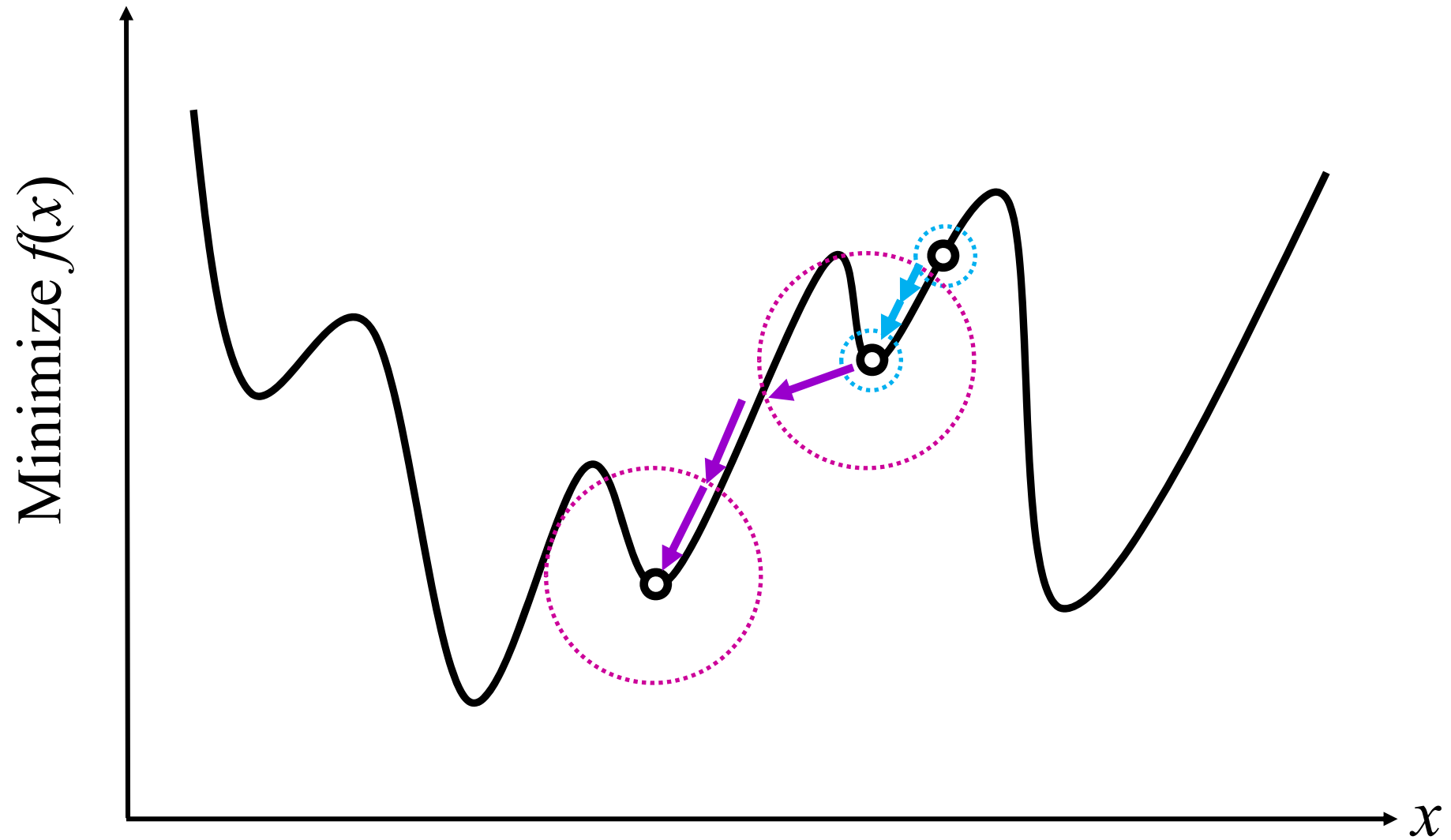
- Restart
(Iterated Local Search)
- Neighborhood Change



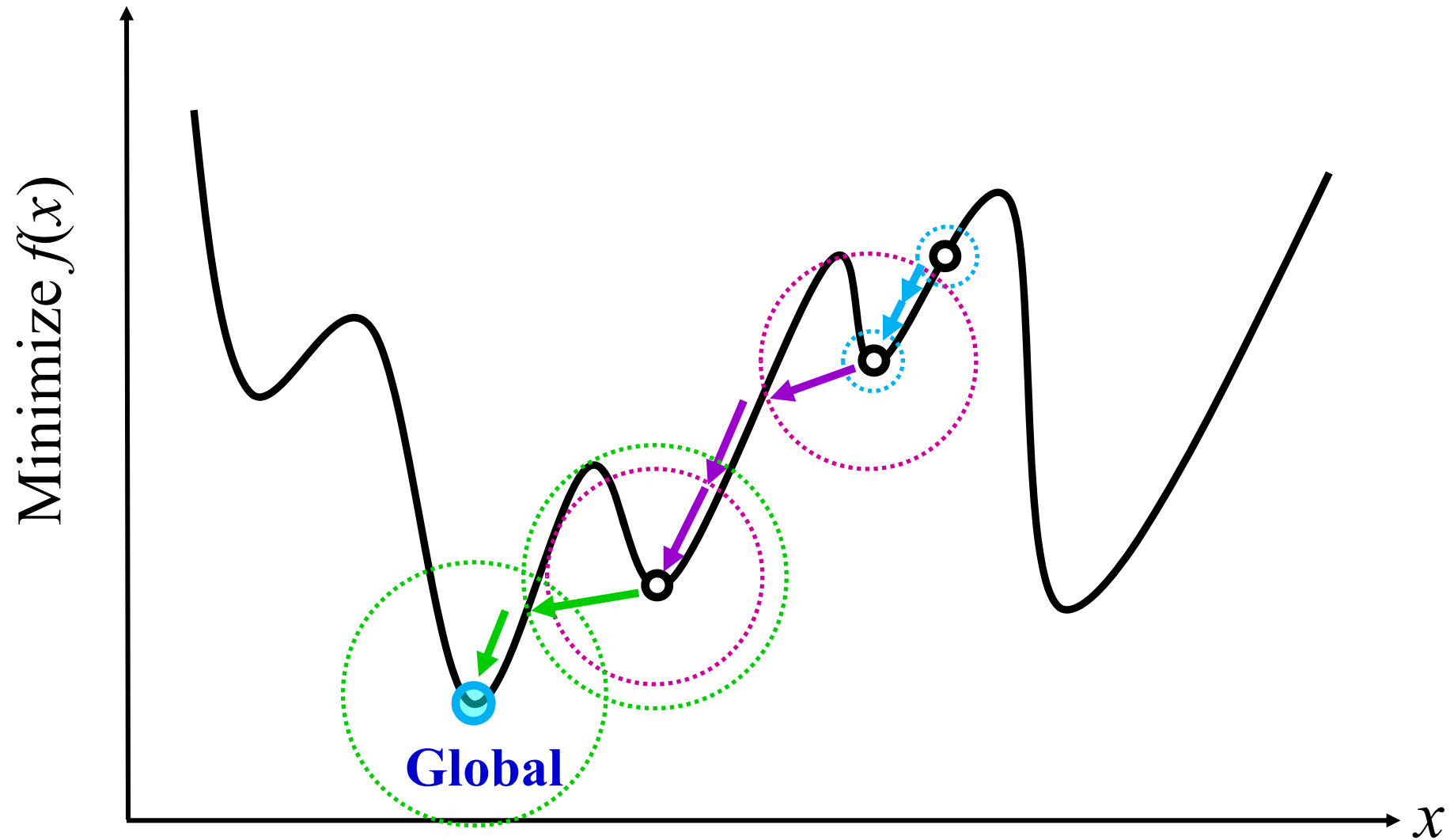
Local Search with Neighborhood Change

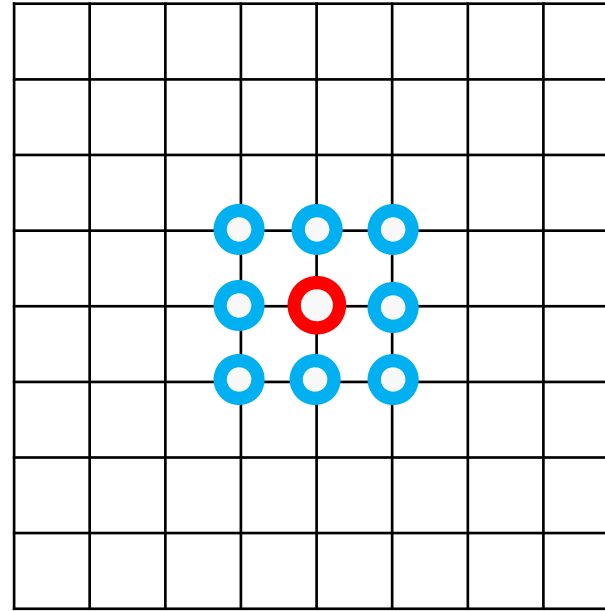
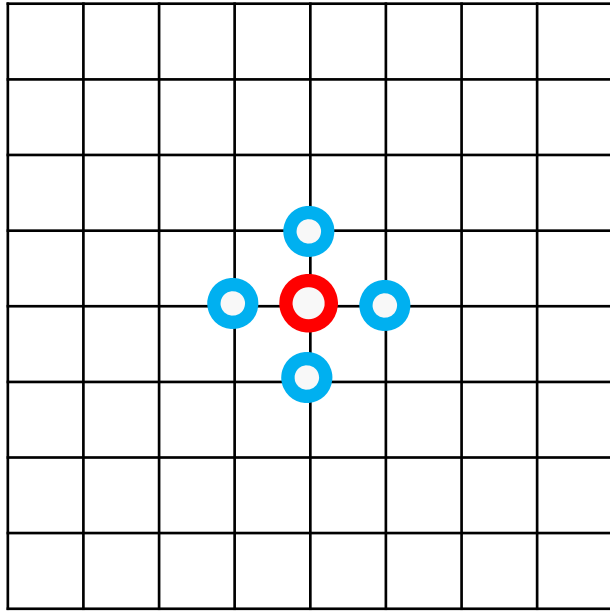


Local Search with Neighborhood Change

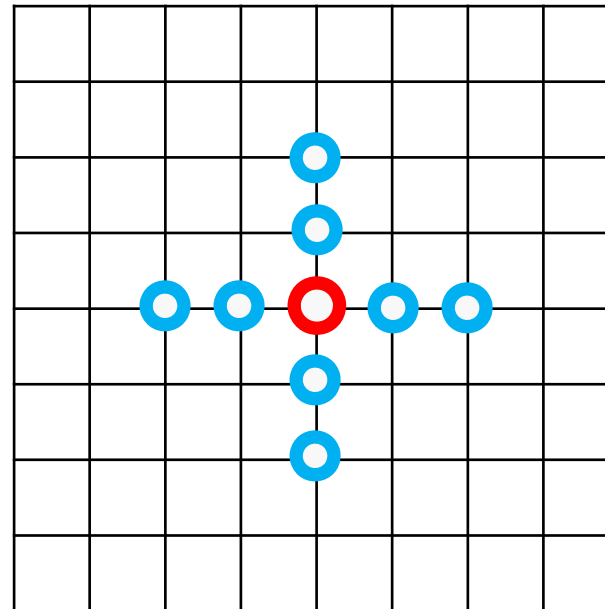
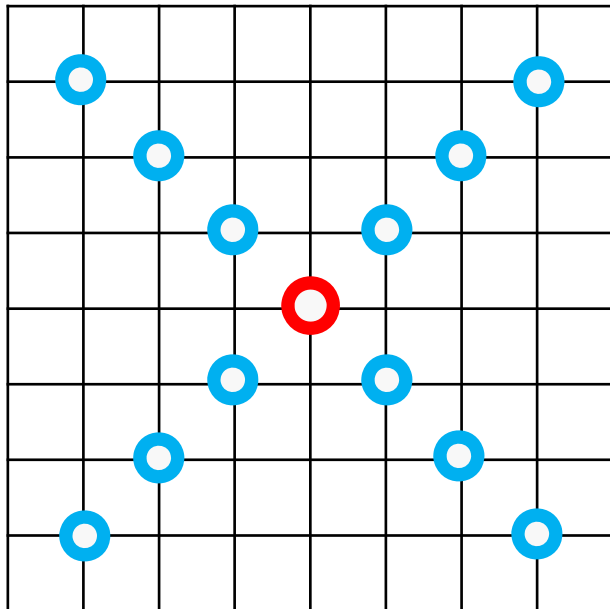


Local Search with Neighborhood Change





Local Search with Neighborhood Change



Adjacent two-city change

Size: n

0	1	2	3	4	5	7	6	8	9
---	---	---	---	---	---	---	---	---	---

Arbitrary two-city change

Size: $n(n-1)/2$

0	1	2	3	4	5	9	7	8	6
---	---	---	---	---	---	---	---	---	---

Insertion (Shift)

Size: $n(n-3)$

0	1	2	3	4	5	7	8	6	9
---	---	---	---	---	---	---	---	---	---

Inversion (Two-edge change)

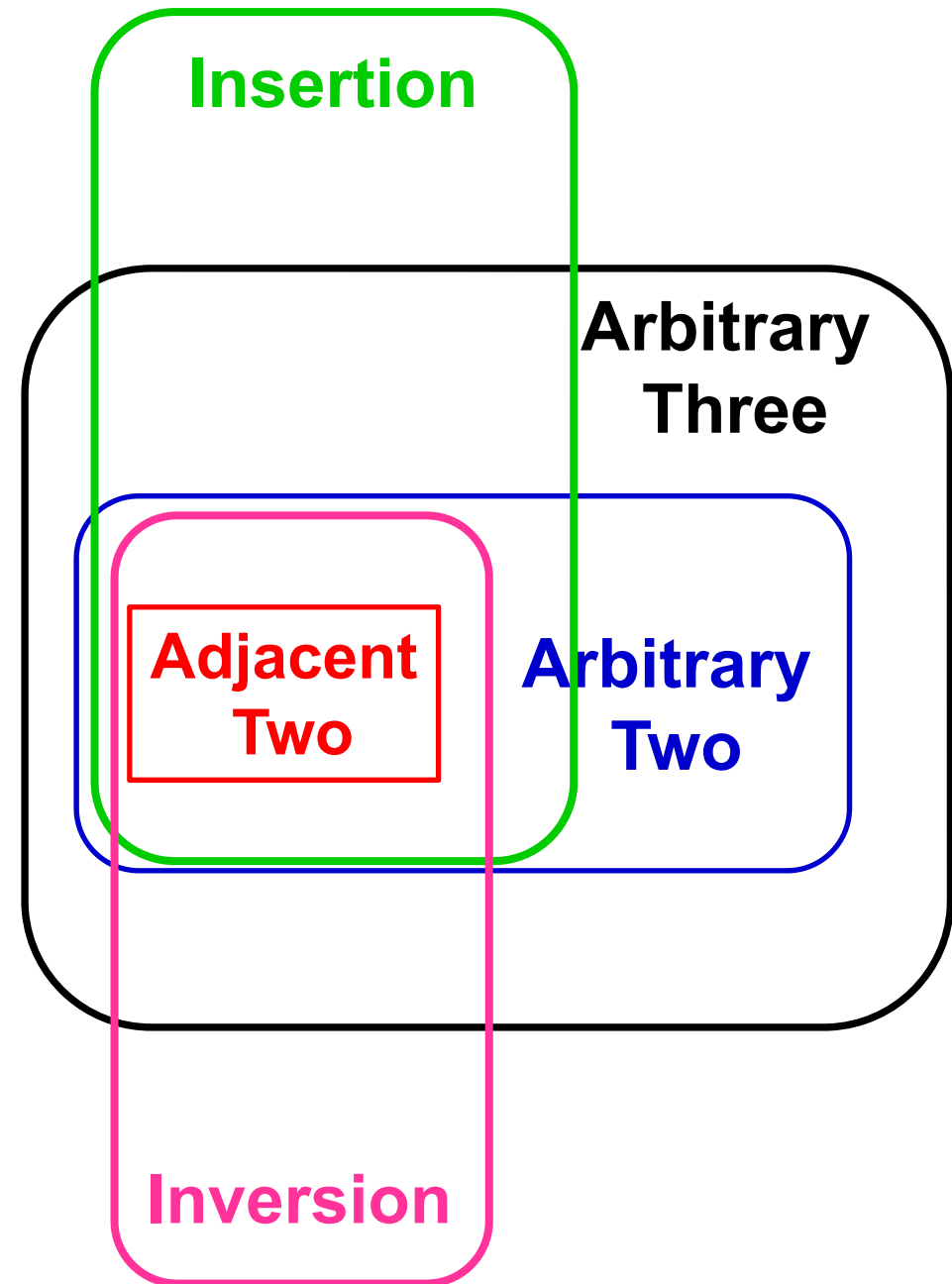
Size: $n(n-3)/2$

0	1	2	8	7	6	5	4	3	9
---	---	---	---	---	---	---	---	---	---

Arbitrary three-city change

Size: $n(n-1)(2n-1)/6$

0	1	2	6	4	5	9	7	8	3
---	---	---	---	---	---	---	---	---	---



Variable Neighborhood Search



Pierre Hansen

GERAD, HEC Canada

Verified email at hec.ca - [Homepage](#)

[Graph Theory](#) [VNS](#)

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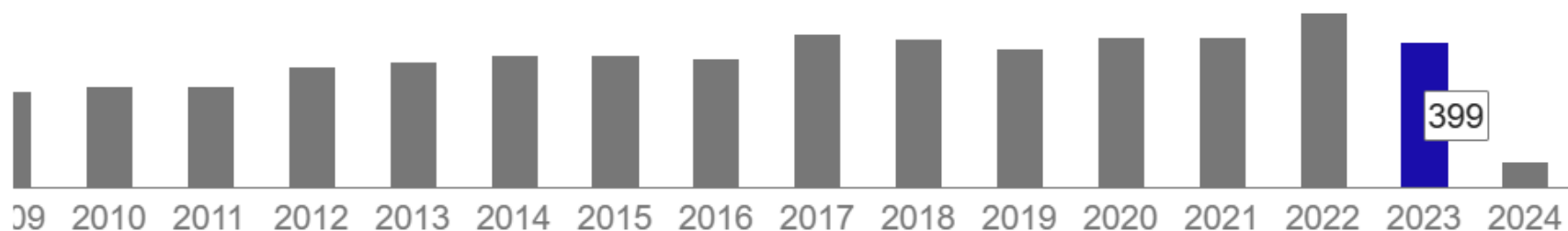
Variable neighborhood search

N Mladenović, P Hansen

Computers & operations research 24 (11), 1097-1100

6547 * 1997

Cited by 6547



Variable Neighborhood Search Algorithm

European Journal of Operational Research 130 (2001) 449–467

Invited Review

Variable neighborhood search: Principles and applications

Pierre Hansen *, Nenad Mladenović

Variable Neighborhood Search Algorithm

P. Hansen, N. Mladenović / European Journal of Operational Research 130 (2001) 449–467

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Fig. 1. Steps of the basic VNS.

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P. Hansen, N. Mladenović / European Journal of Operational Research 130 (2001) 449–467

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Local search run. Local solution: x''

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(c) *Move or not.* If this local optimum is better than **the incumbent**, move there ($x \leftarrow x''$), and continue the search with \mathcal{N}_1 ($k \leftarrow 1$); otherwise, set $k \leftarrow k + 1$;

If the current solution is improved, $\mathcal{N}_k \Rightarrow \mathcal{N}_1$. If not, $\mathcal{N}_k \Rightarrow \mathcal{N}_{k+1}$.

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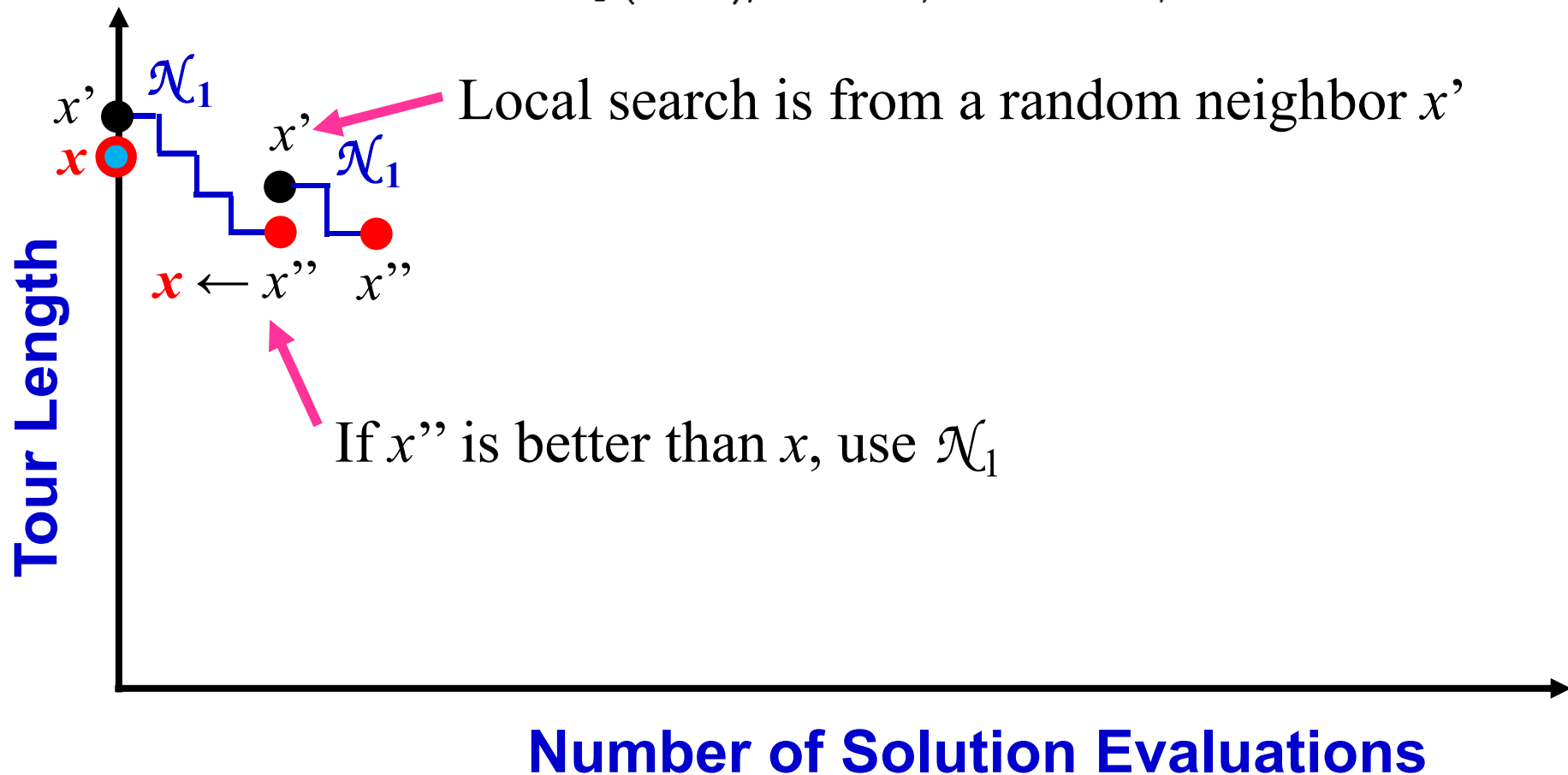
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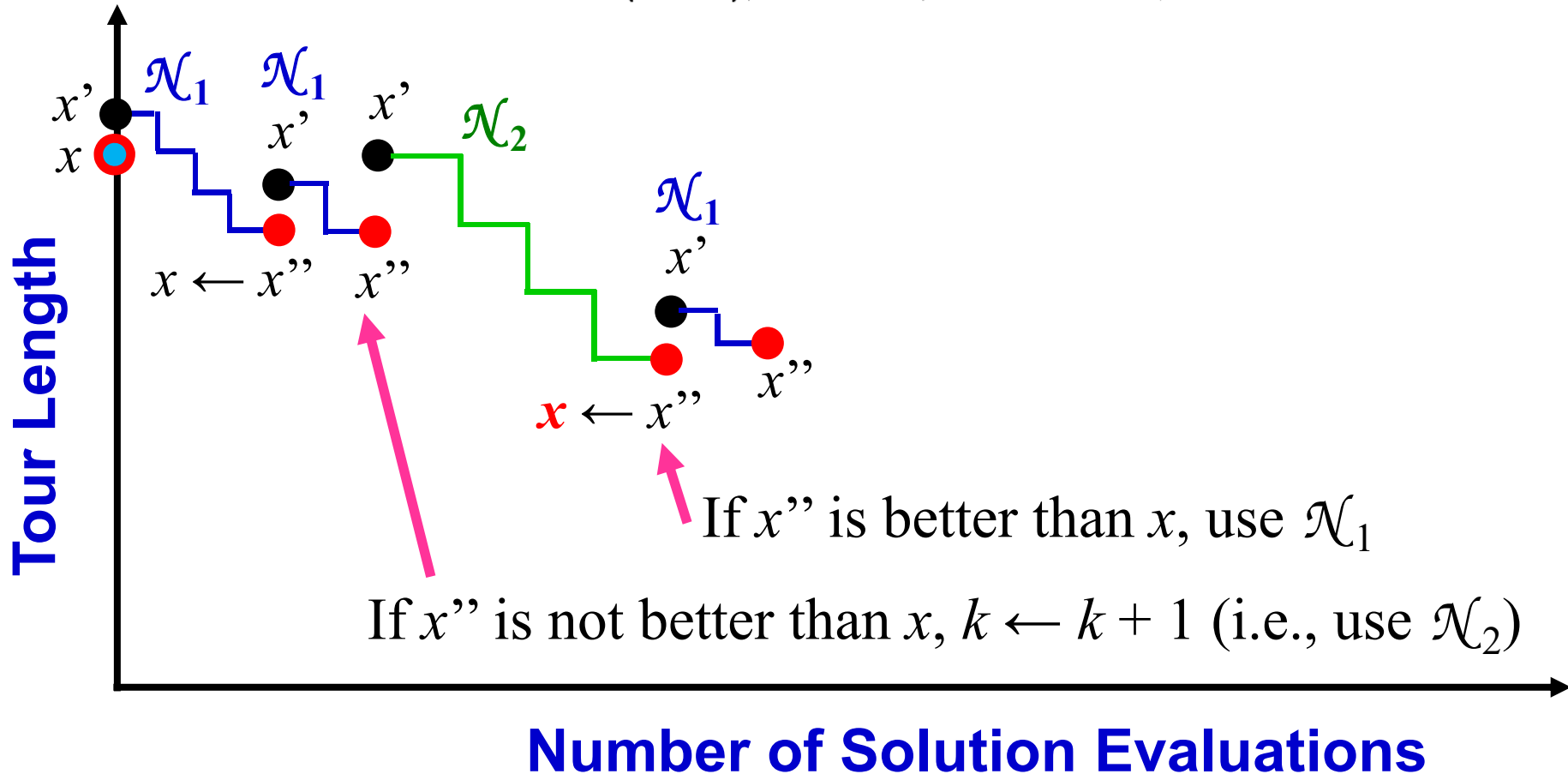
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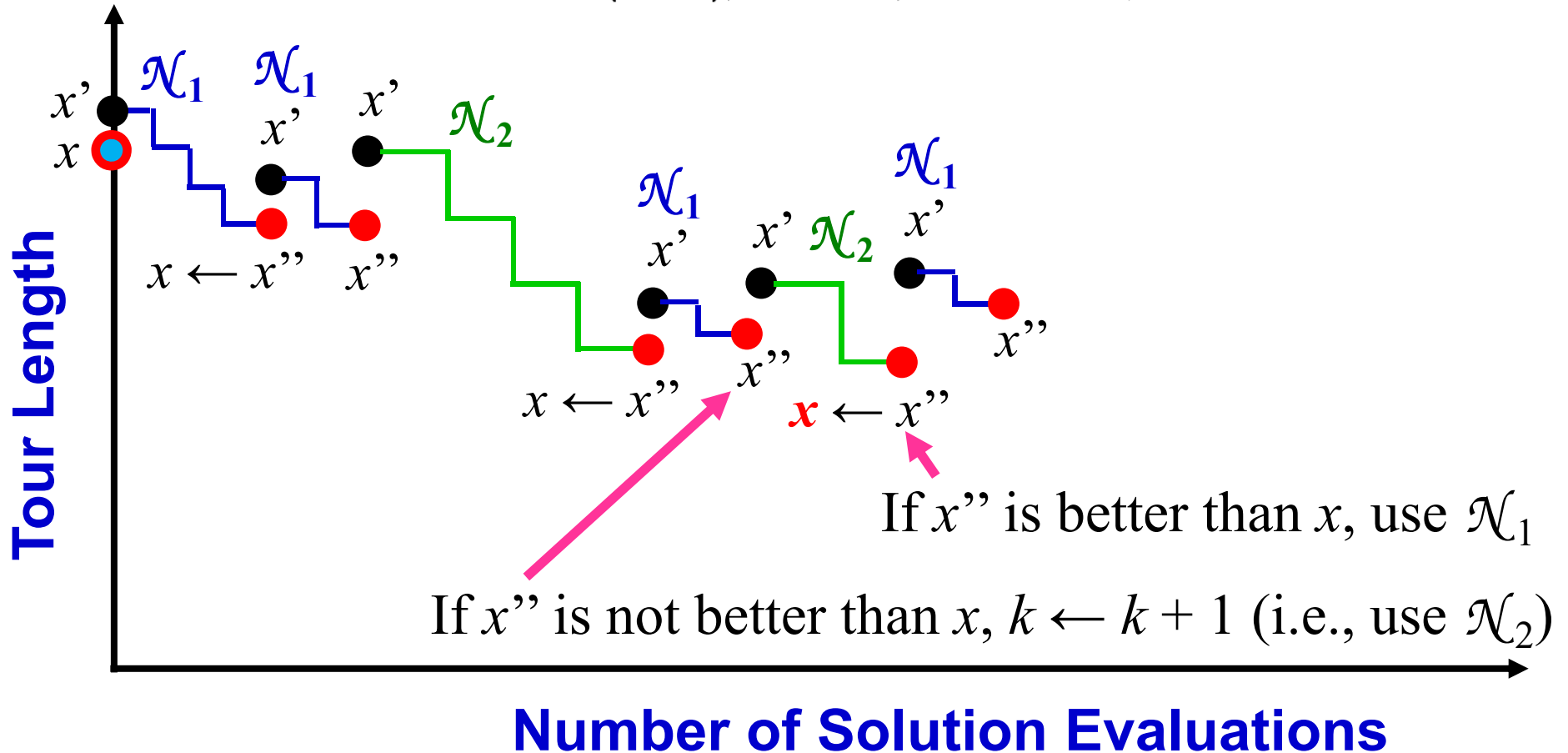
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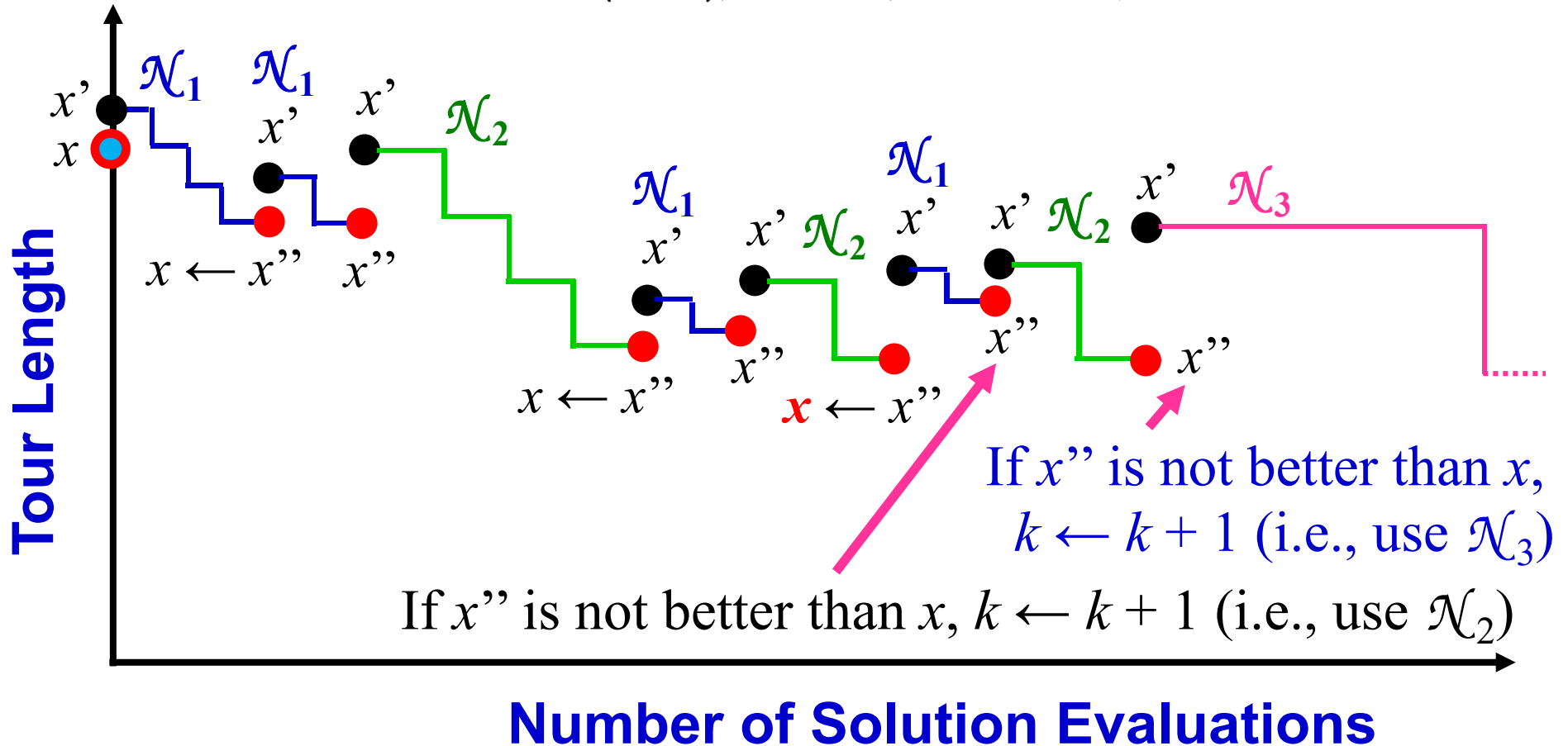
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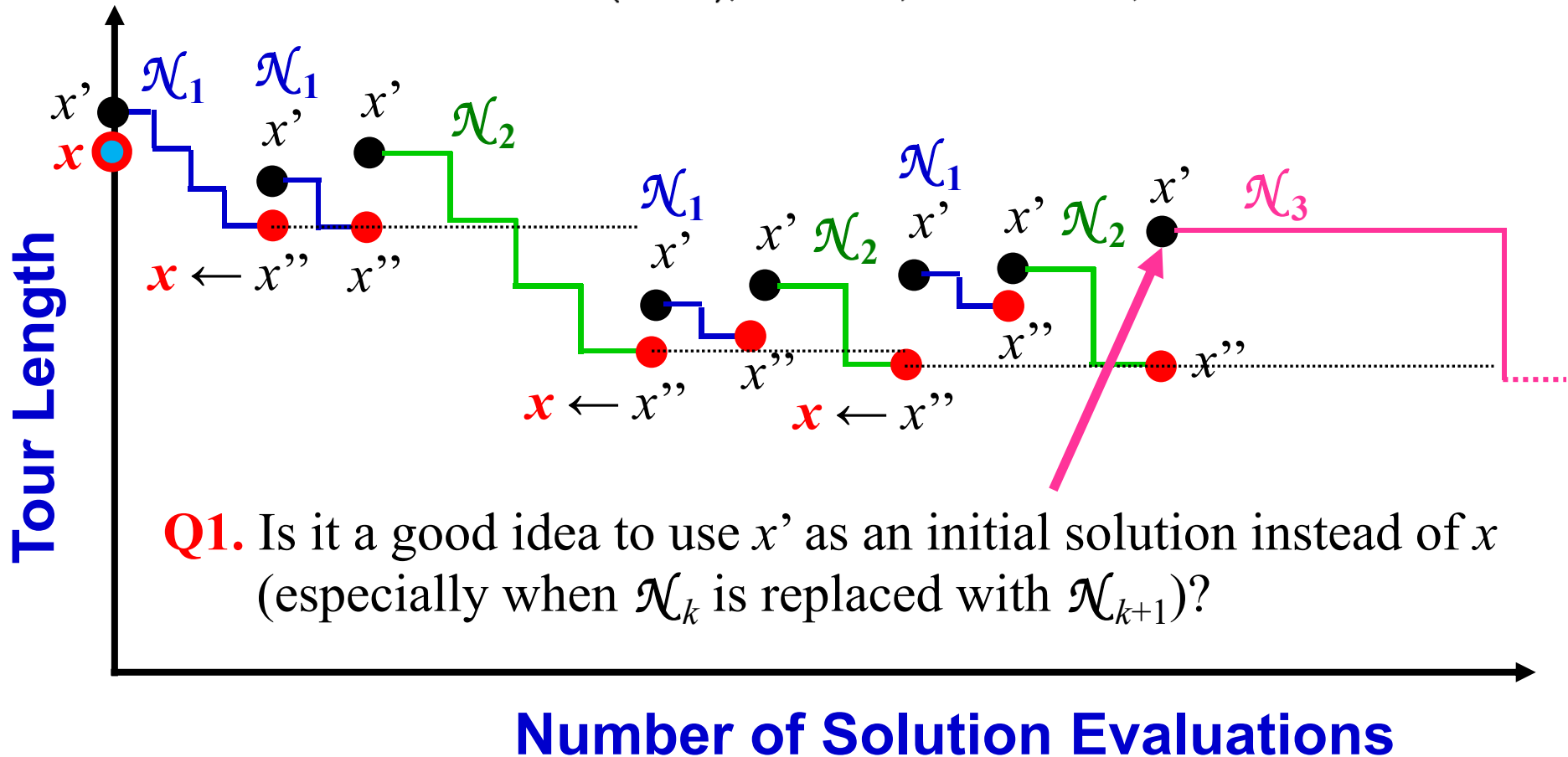
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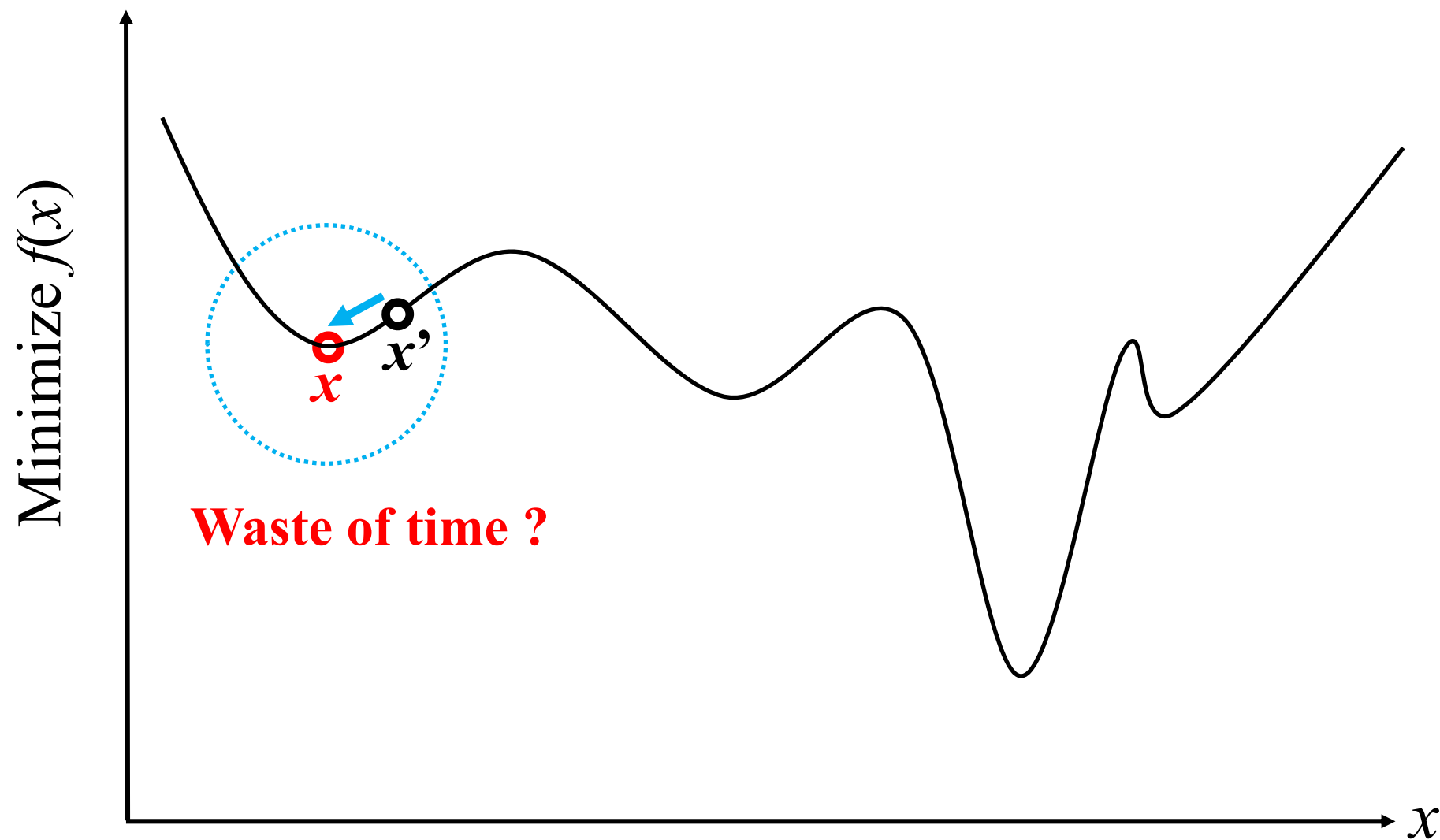
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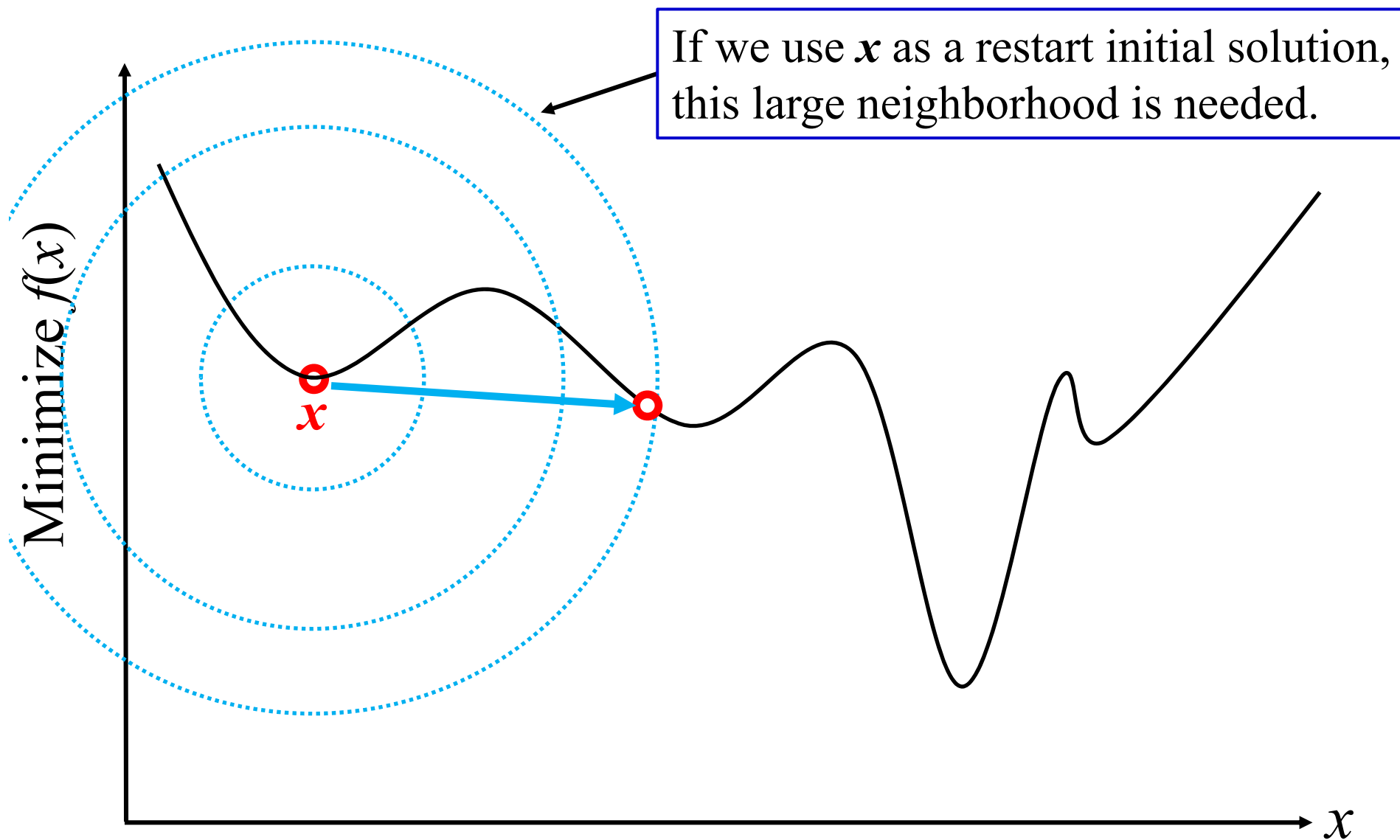
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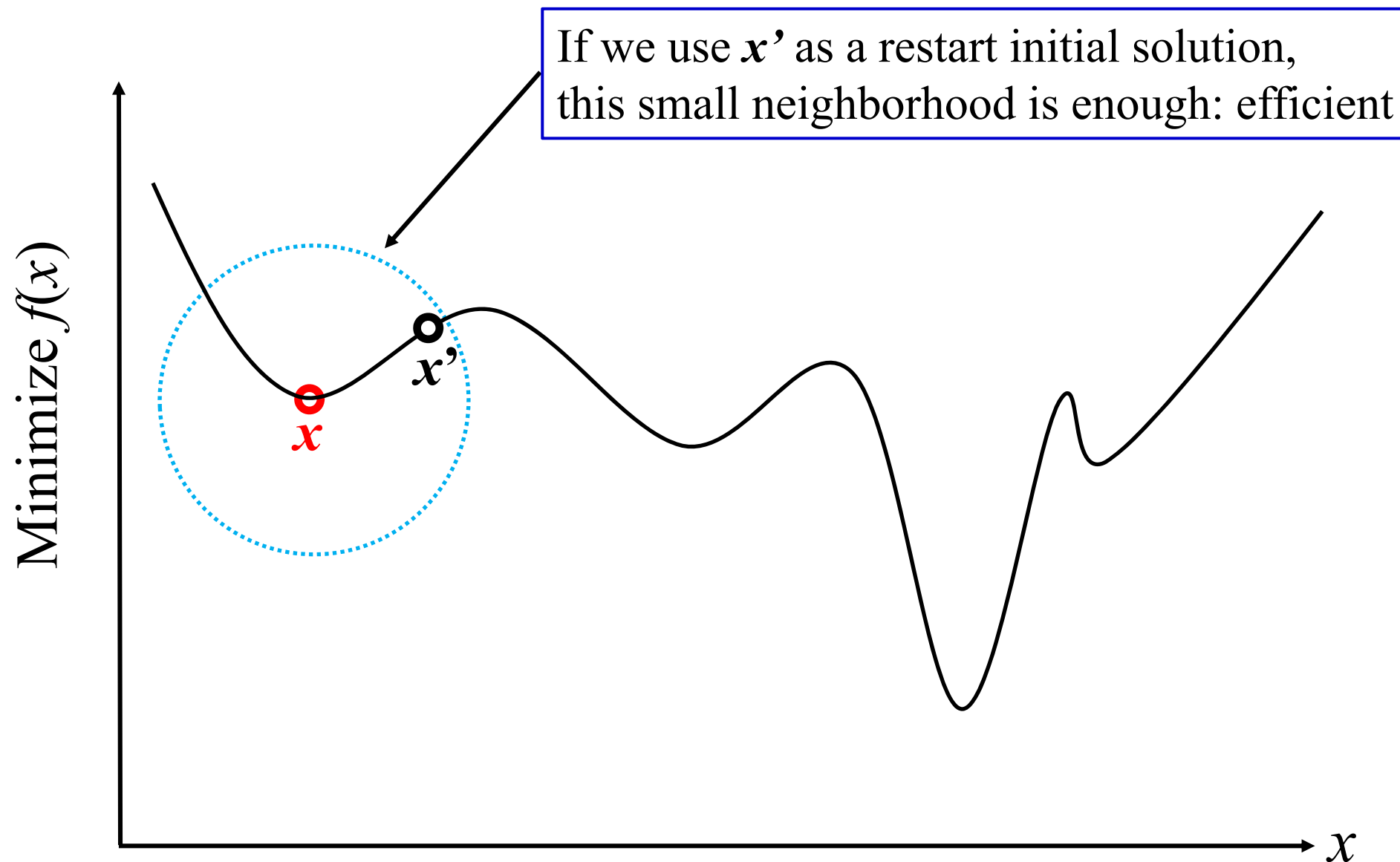
Q1. Is it a good idea to use x' instead of x ?



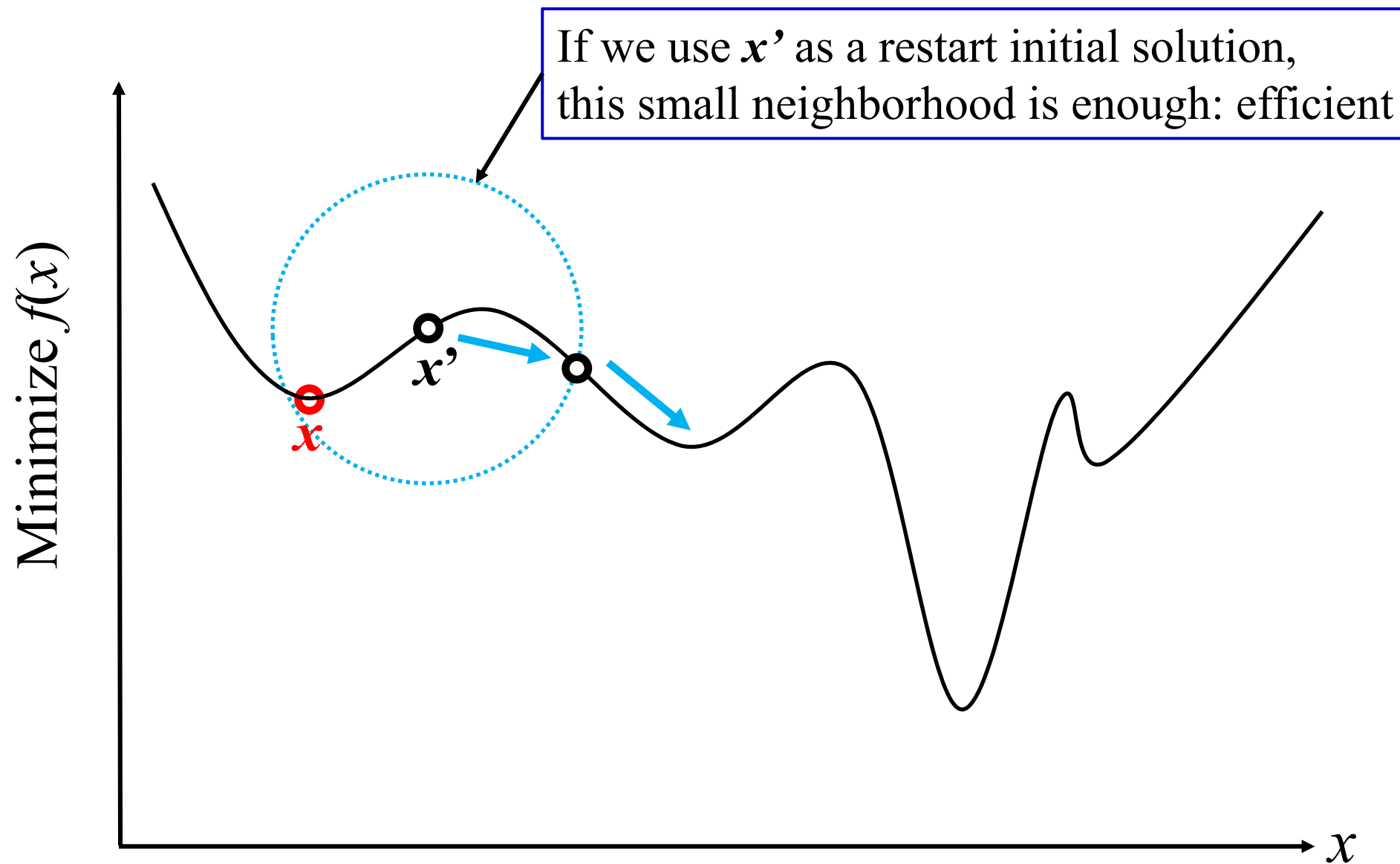
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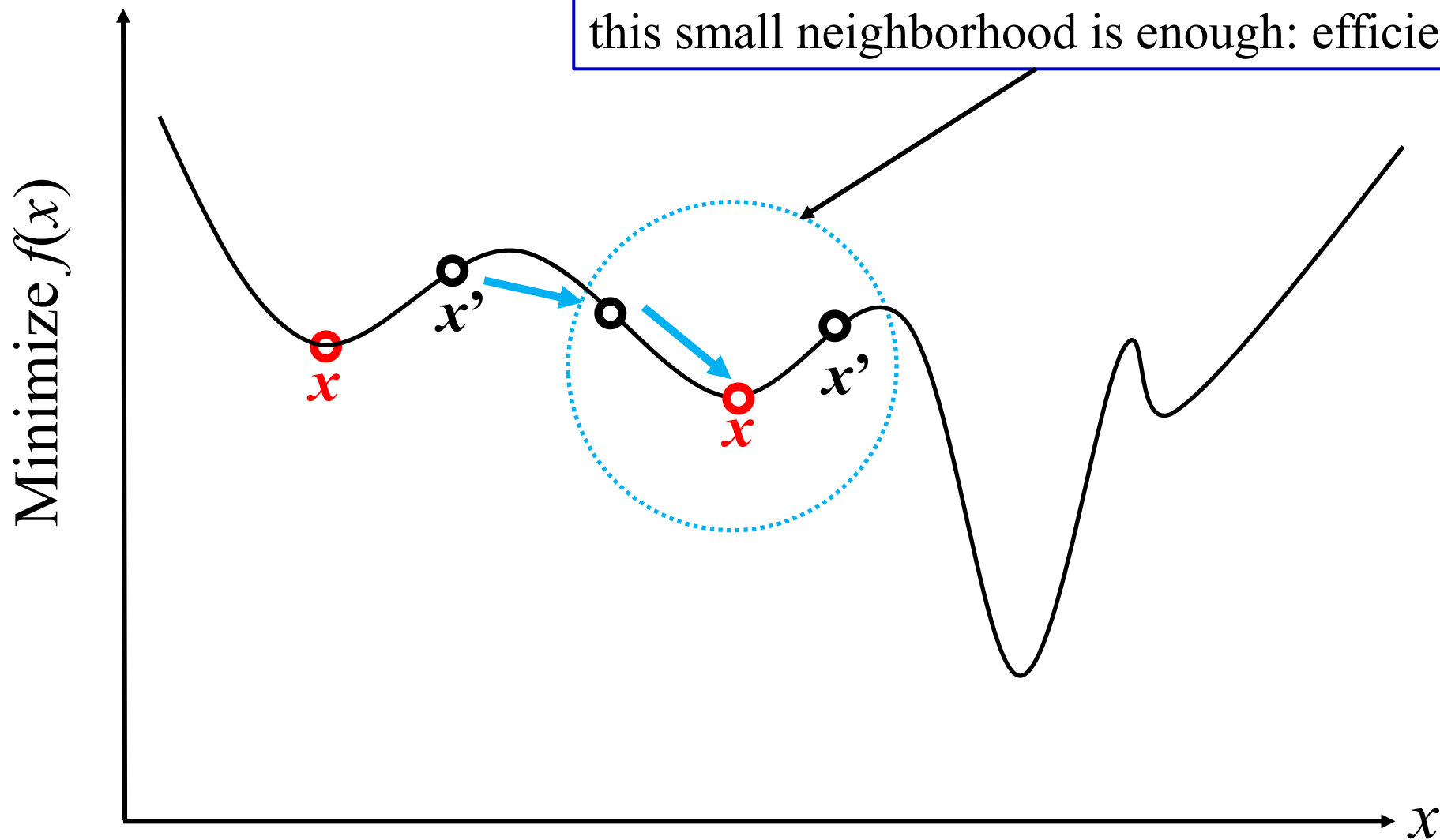


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If we use x' as a restart initial solution, this small neighborhood is enough: efficient



Initialization. Select the set of neighborhood structures \mathcal{N}_k , $k = 1, \dots, k_{max}$, that will be used in the search; find an initial solution x ; choose a stopping condition;

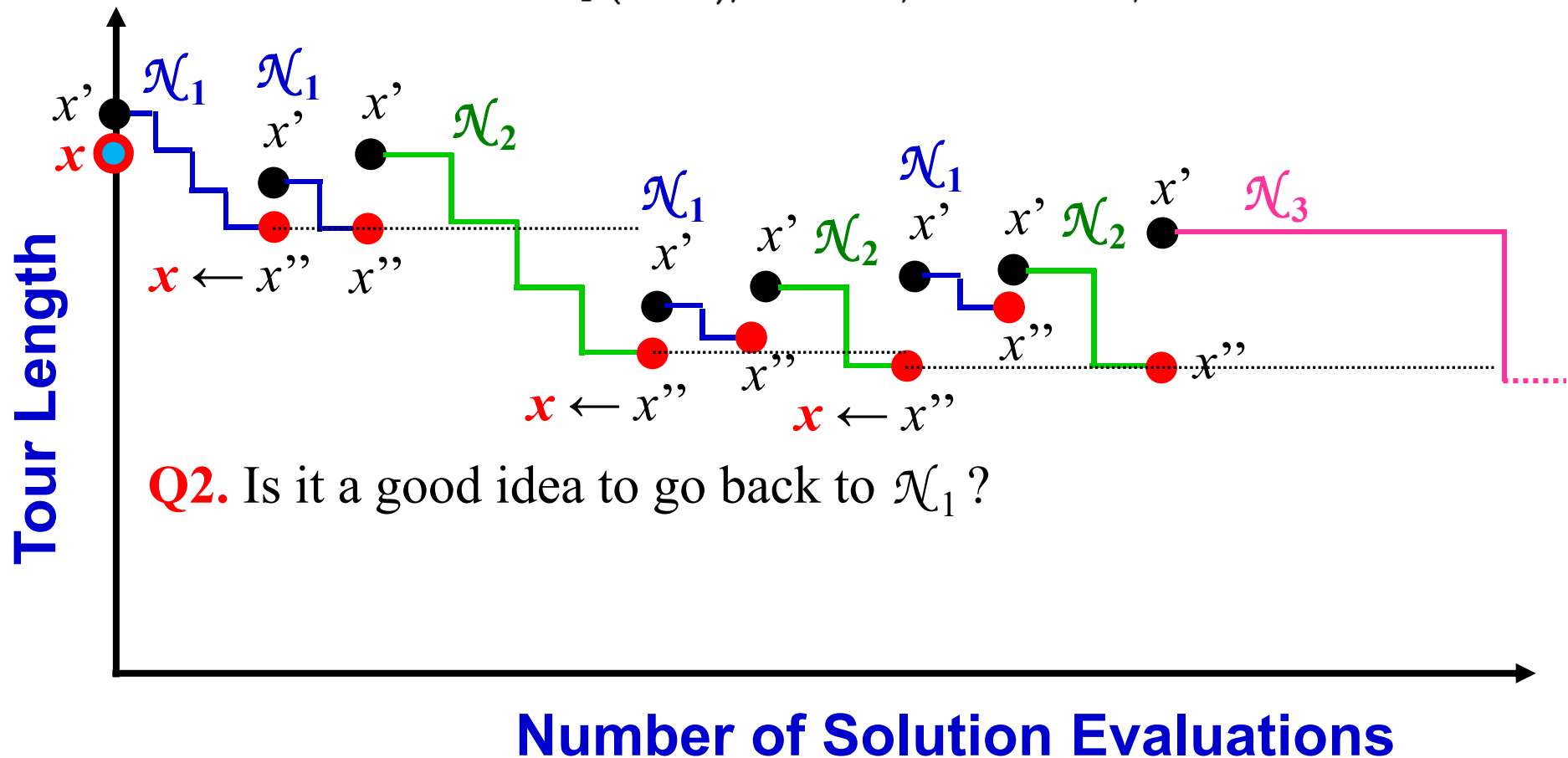
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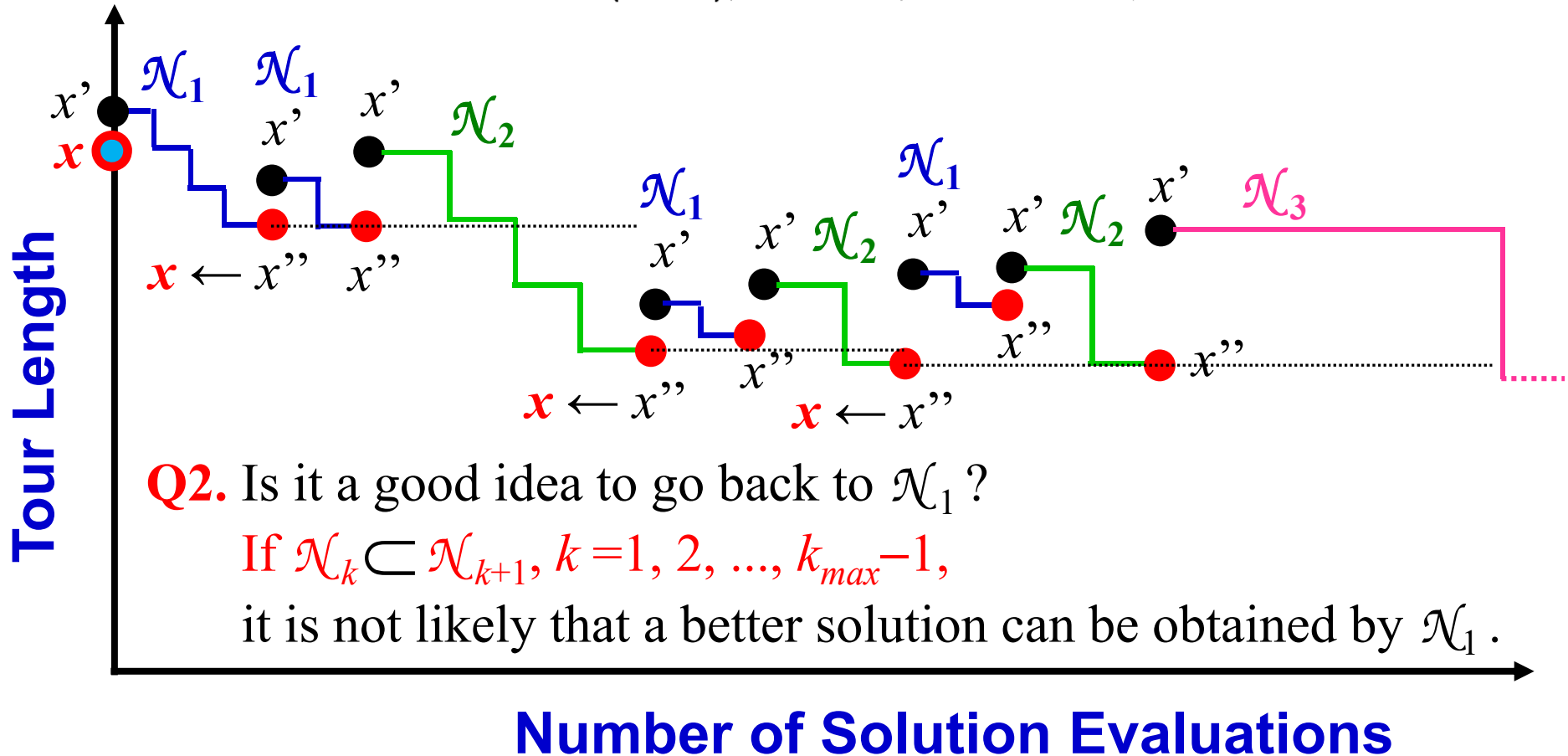
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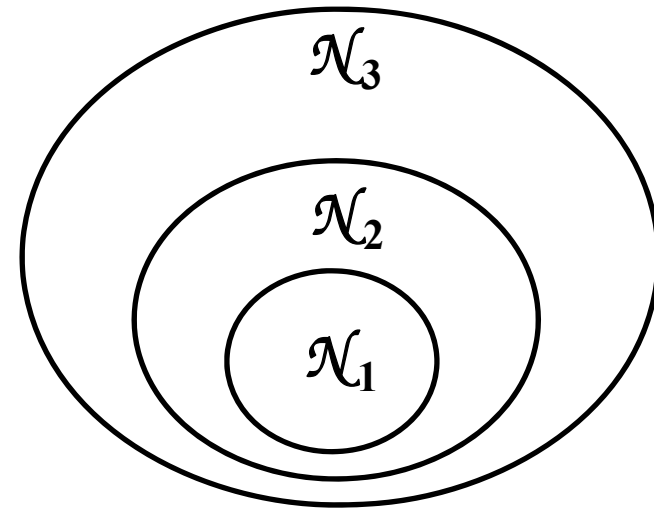
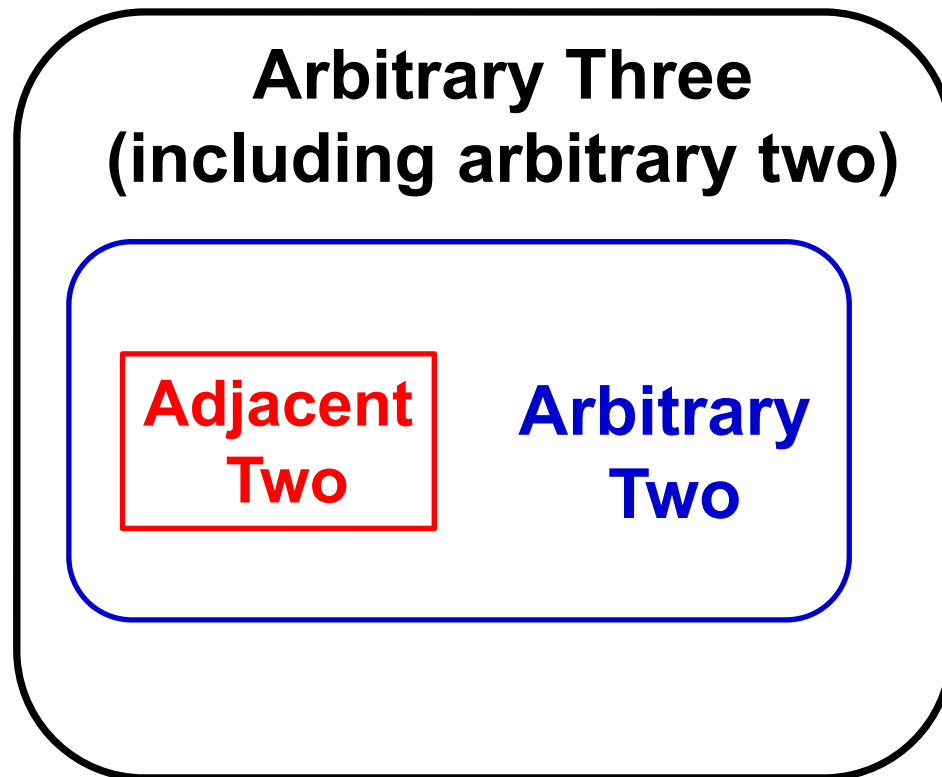
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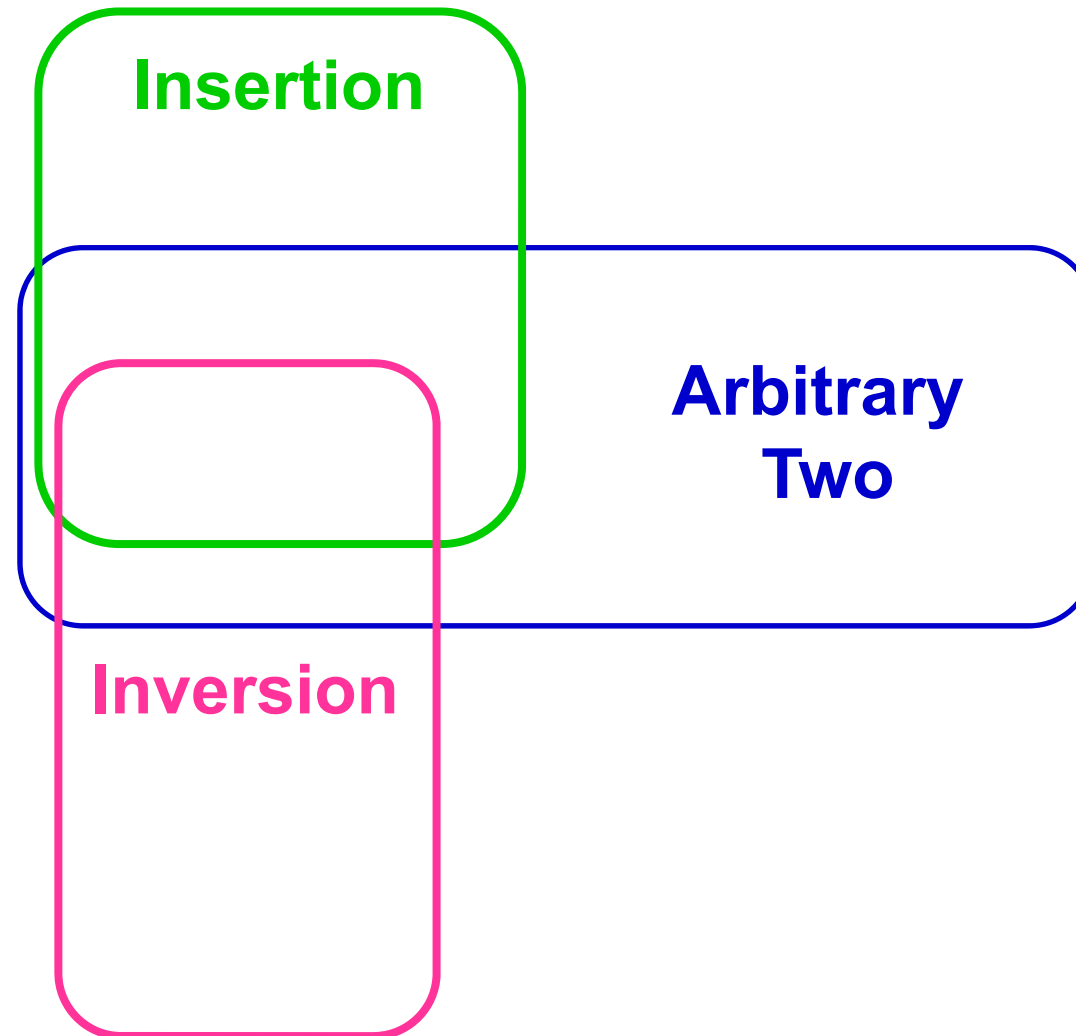
Q2. Is it a good idea to go back to \mathcal{N}_1 ?

If $\mathcal{N}_k \subset \mathcal{N}_{k+1}$, $k=1, 2, \dots, k_{max}-1$,
it is not likely that a better solution can be obtained by \mathcal{N}_1 .



Q2. Is it a good idea to go back to \mathcal{N}_1 ?

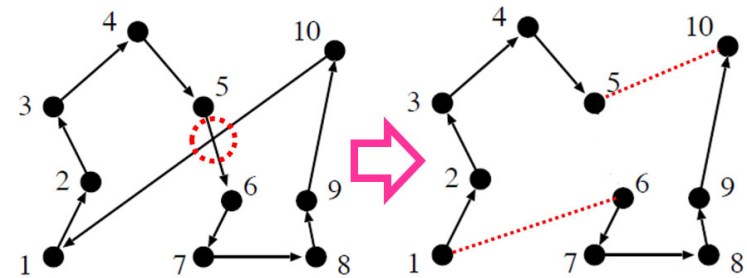
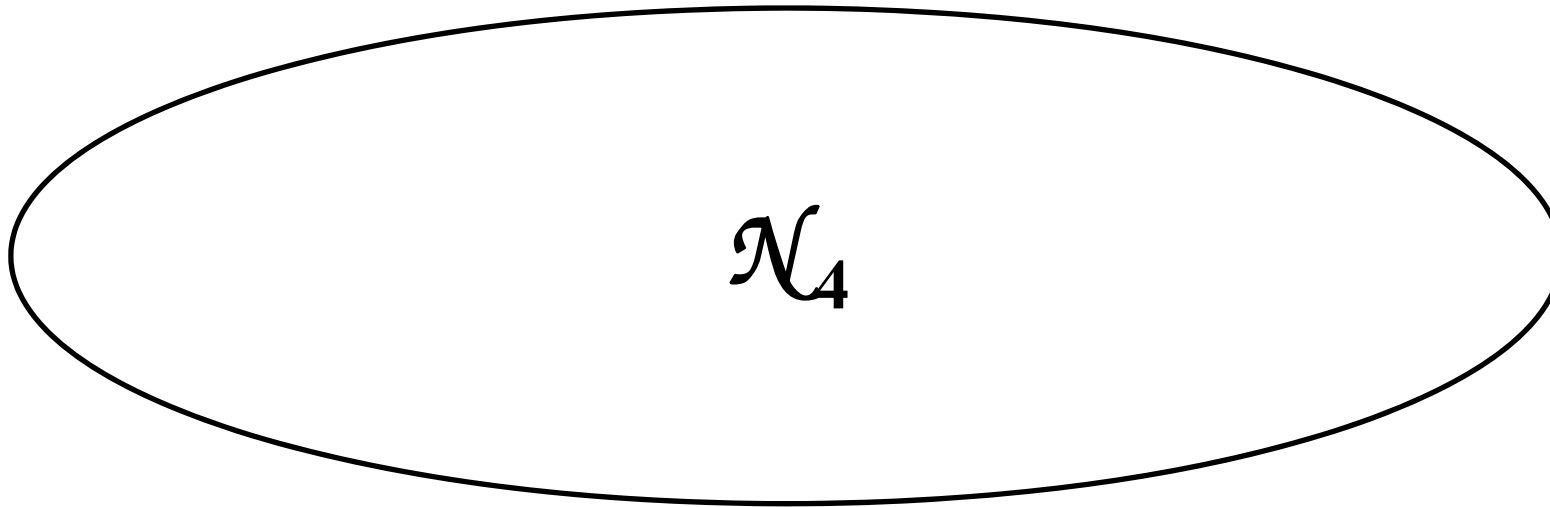
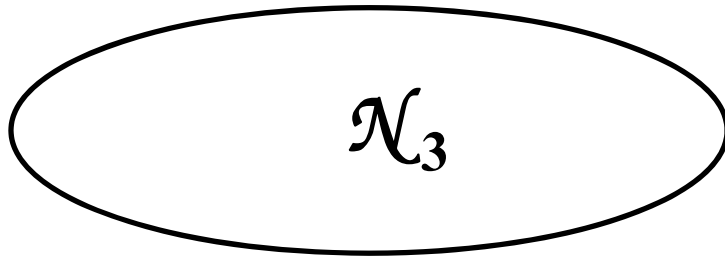
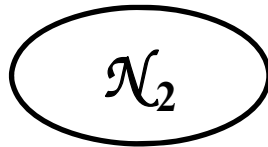
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Example of Neighborhood Structures of TSP (n -city problem)

\mathcal{N}_k : k -edge change neighborhood ($k = 2, 3, \dots, n$)

$$\mathcal{N}_1 = \emptyset$$

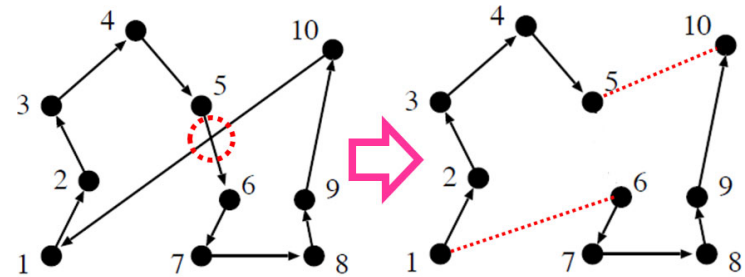
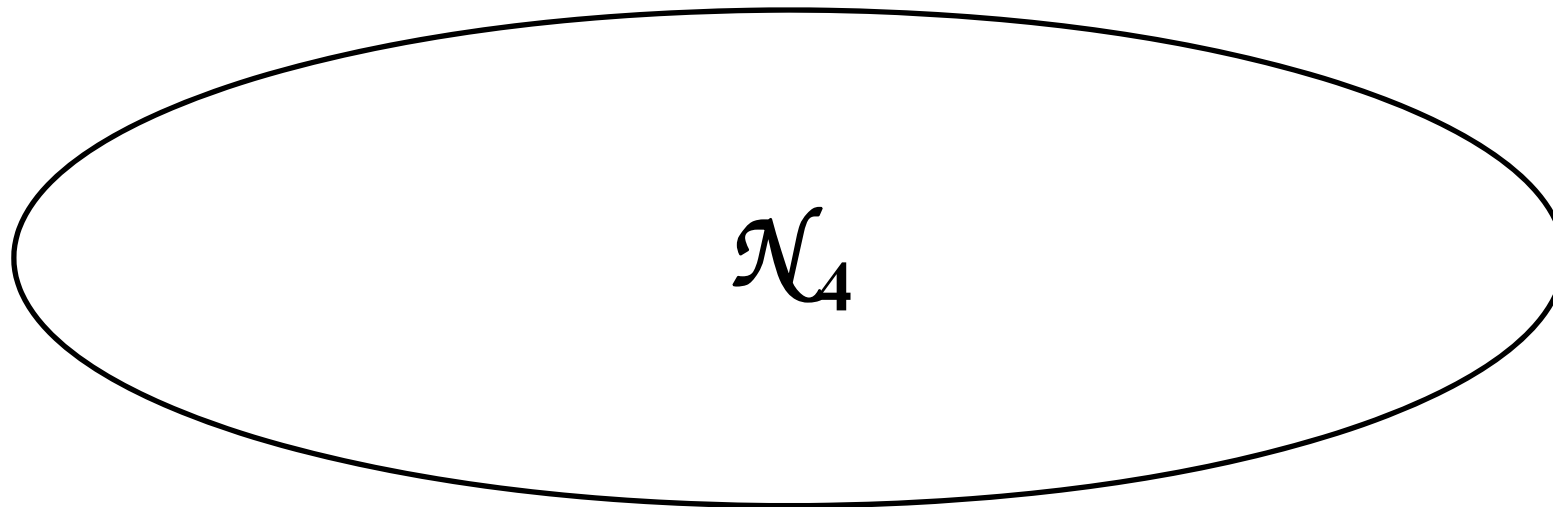
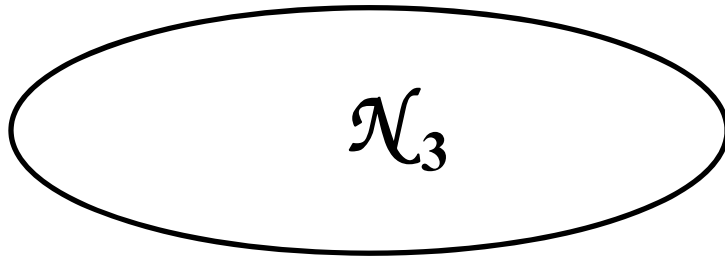
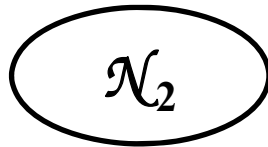


\mathcal{N}_k : 2-edge change neighborhood

Q2. Is it a good idea to go back to \mathcal{N}_1 ? **Yes ?** (Depends on the problem.)

\mathcal{N}_k : k -edge change neighborhood ($k = 2, 3, \dots, n$)

$$\mathcal{N}_1 = \emptyset$$



\mathcal{N}_k : 2-edge change neighborhood

Variant of Variable Neighborhood Search (VNS)

Variable Neighborhood Descent (VND)

European Journal of Operational Research 130 (2001) 449–467

Invited Review

Variable neighborhood search: Principles and applications

Pierre Hansen *, Nenad Mladenović

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Variable Neighborhood Descent (VND)

P. Hansen, N. Mladenović / European Journal of Operational Research 130 (2001) 449–467

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Neighborhood Structures: \mathcal{N}_k , $k = 1, 2, \dots, k_{max}$

\mathcal{N}_{k+1} is larger than \mathcal{N}_k , $k = 1, 2, \dots, k_{max}-1$.

In some cases, $\mathcal{N}_k \subset \mathcal{N}_{k+1}$, $k = 1, 2, \dots, k_{max}-1$.

\mathcal{N}_1 is used first, then \mathcal{N}_2 , \mathcal{N}_3 , ...

~~Local search is from a randomly selected neighbor in the current \mathcal{N}_k .~~

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If no improvement, use the next neighborhood ($\mathcal{N}_k \Rightarrow \mathcal{N}_{k+1}$).

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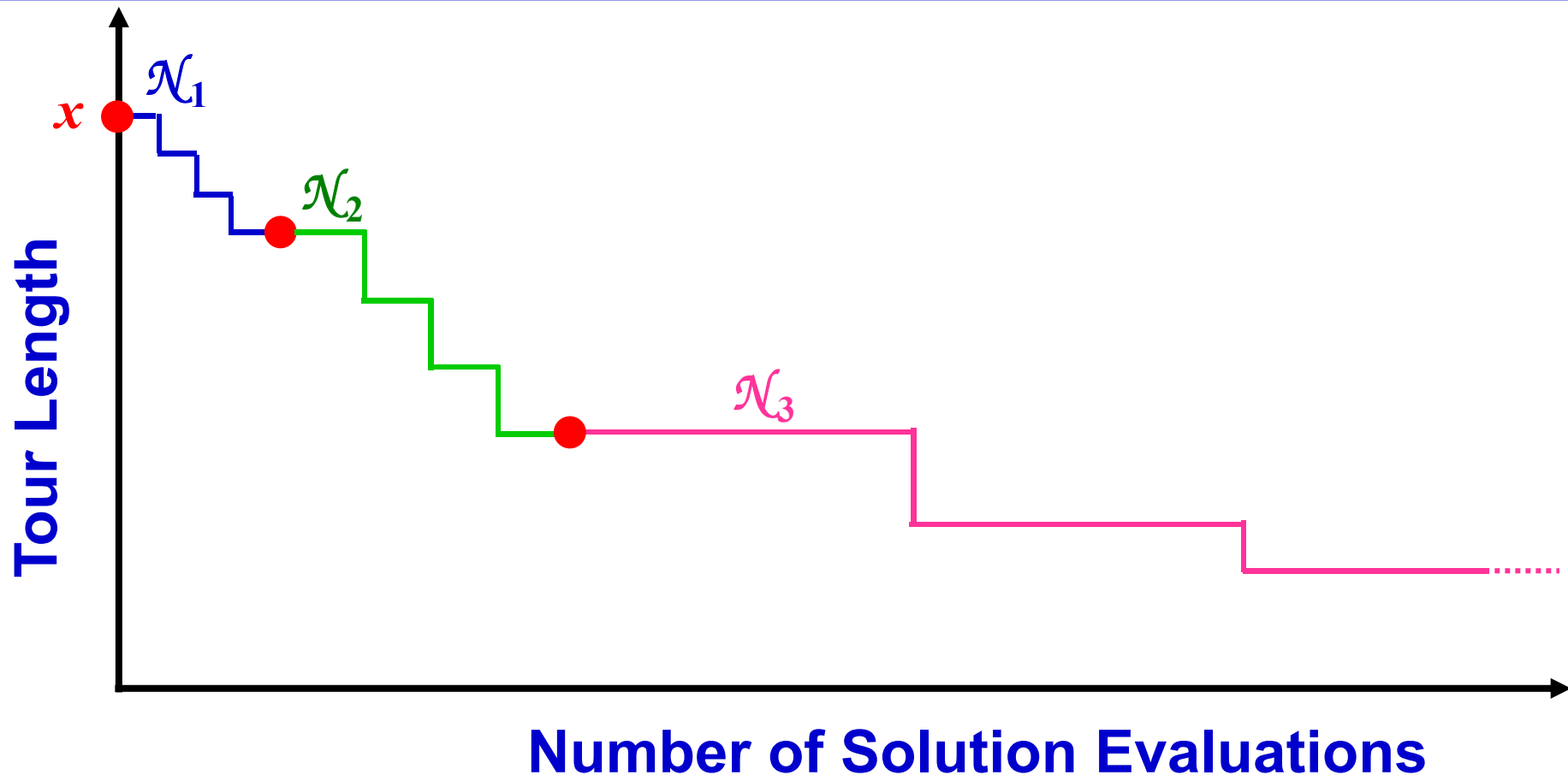
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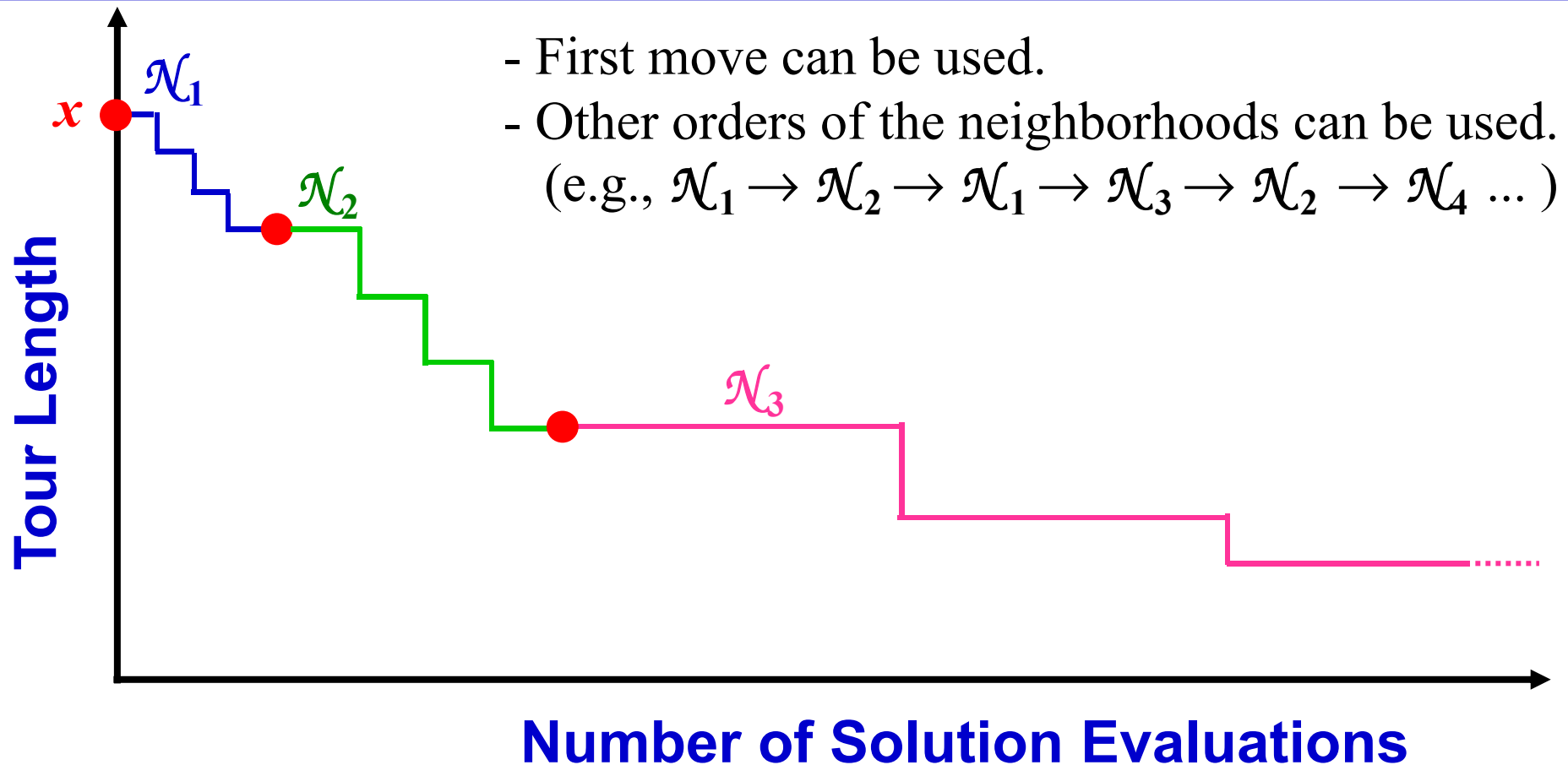
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A Tutorial on Variable Neighborhood Search

Pierre Hansen

GERAD and HEC Montreal

Nenad Mladenović

GERAD and Mathematical Institute, SANU, Belgrade

Basic Ideas behind Variable Neighborhood Search

Fact 1 *A local minimum with respect to one neighborhood structure is not necessary so for another;*

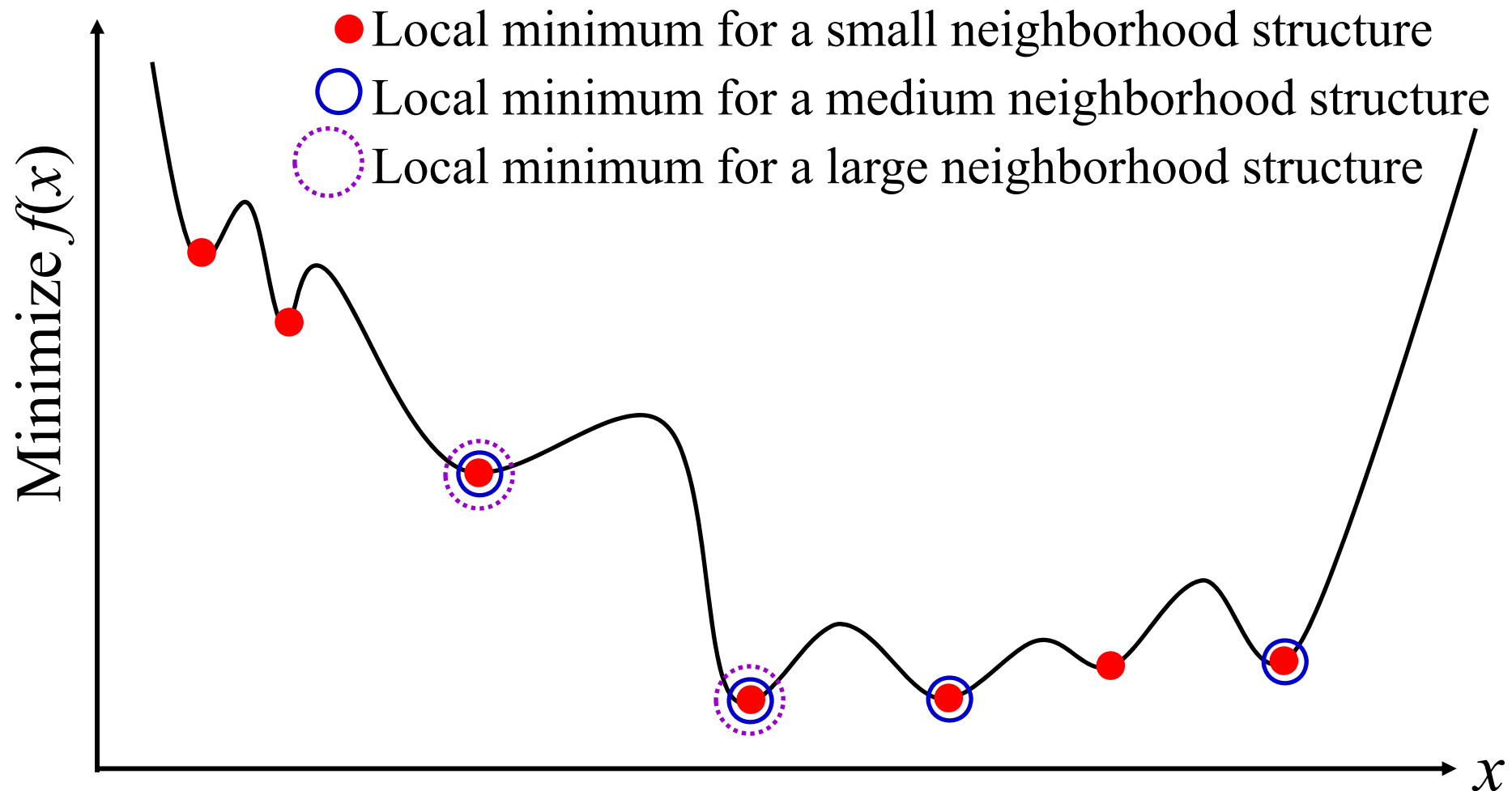
Fact 2 *A global minimum is a local minimum with respect to all possible neighborhood structures.*

Fact 3 *For many problems local minima with respect to one or several neighborhoods are relatively close to each other.*

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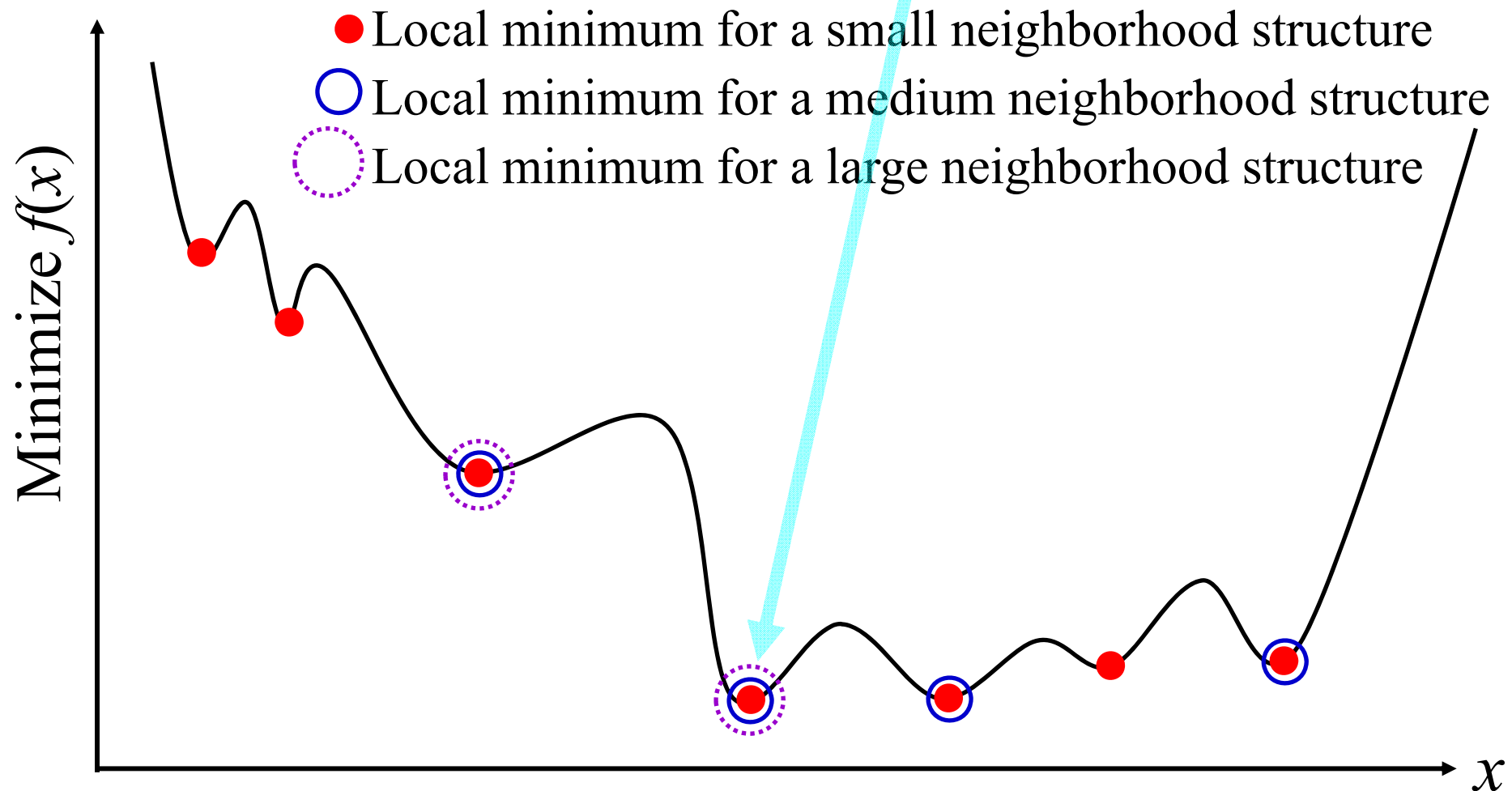
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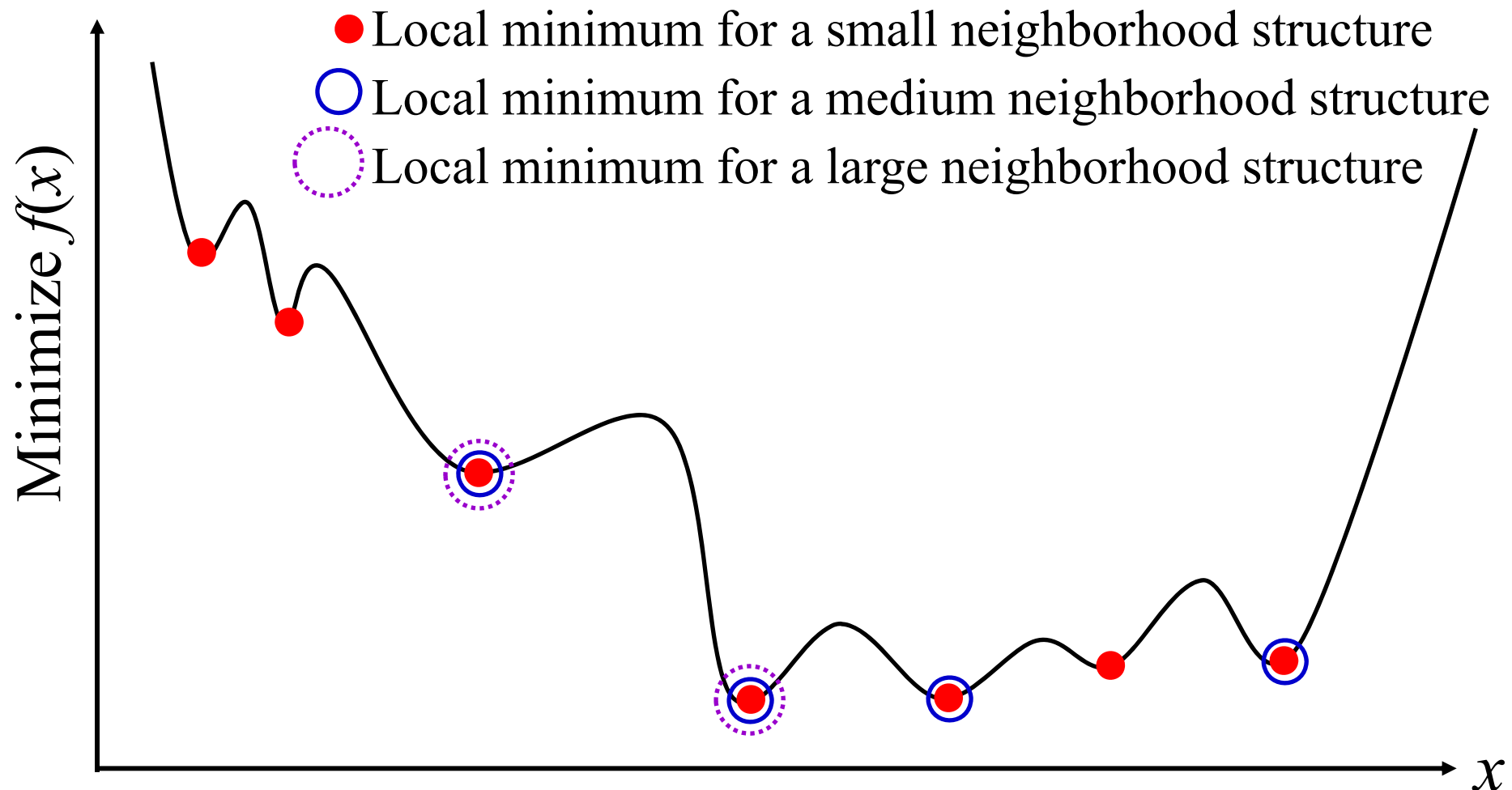
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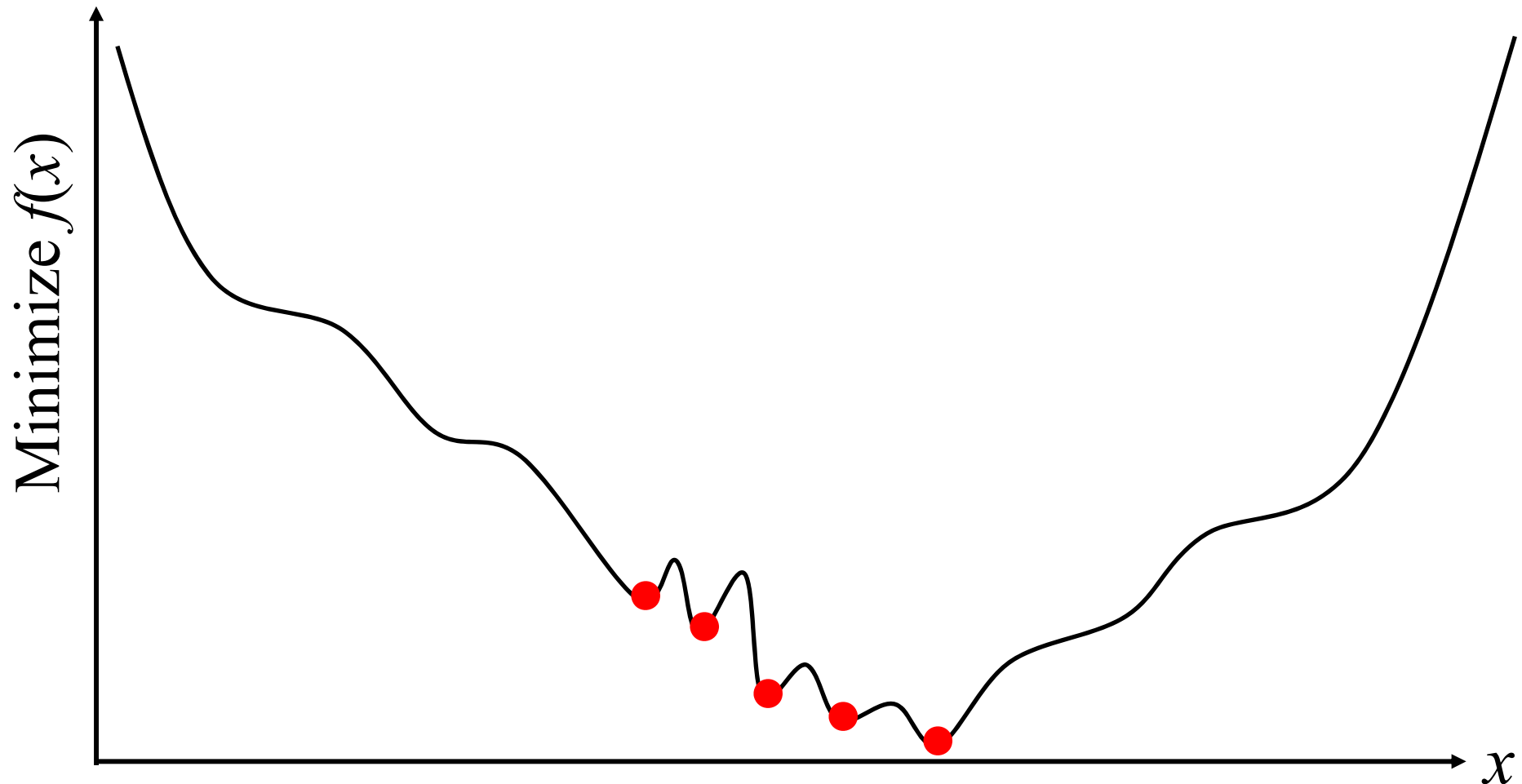
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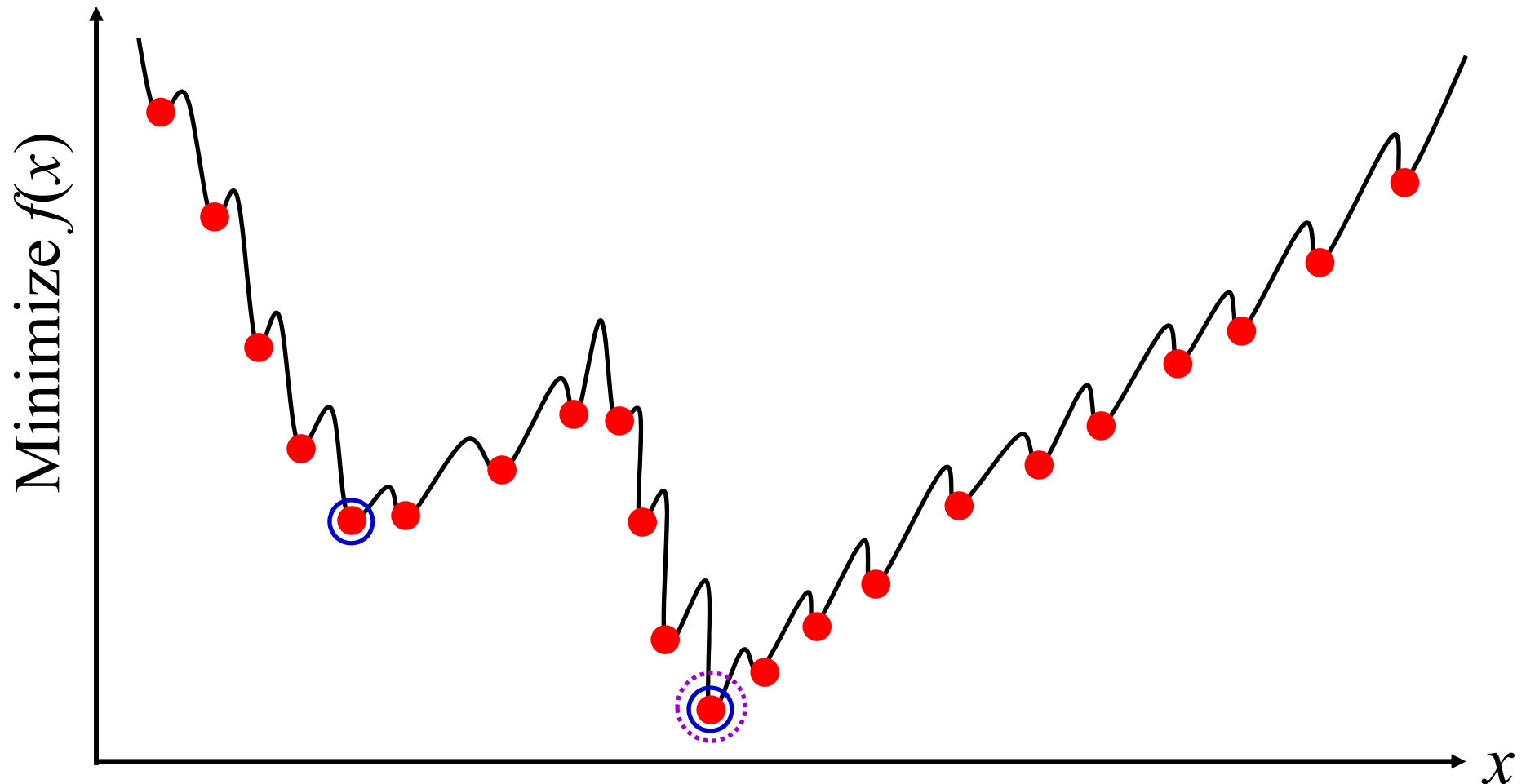
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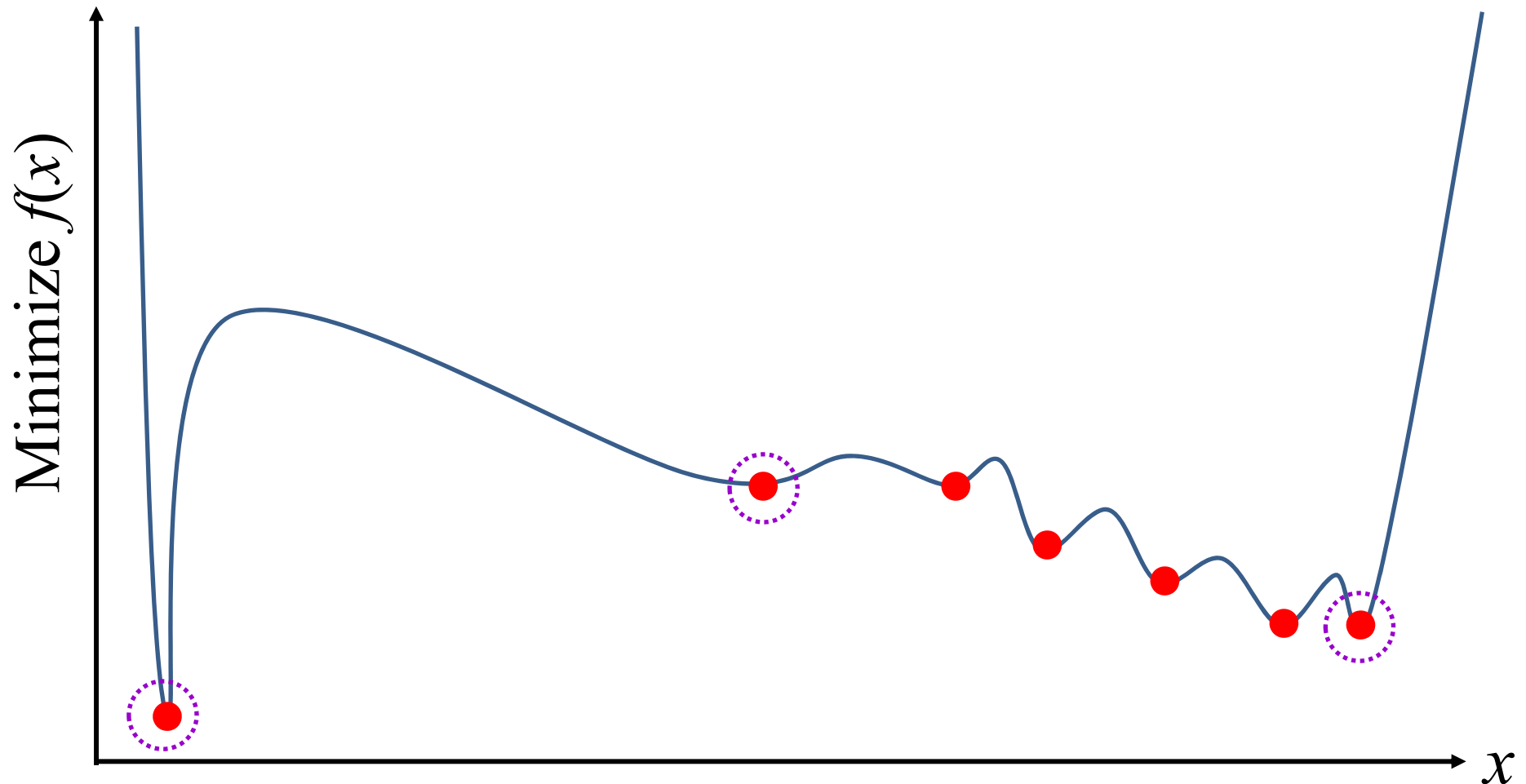
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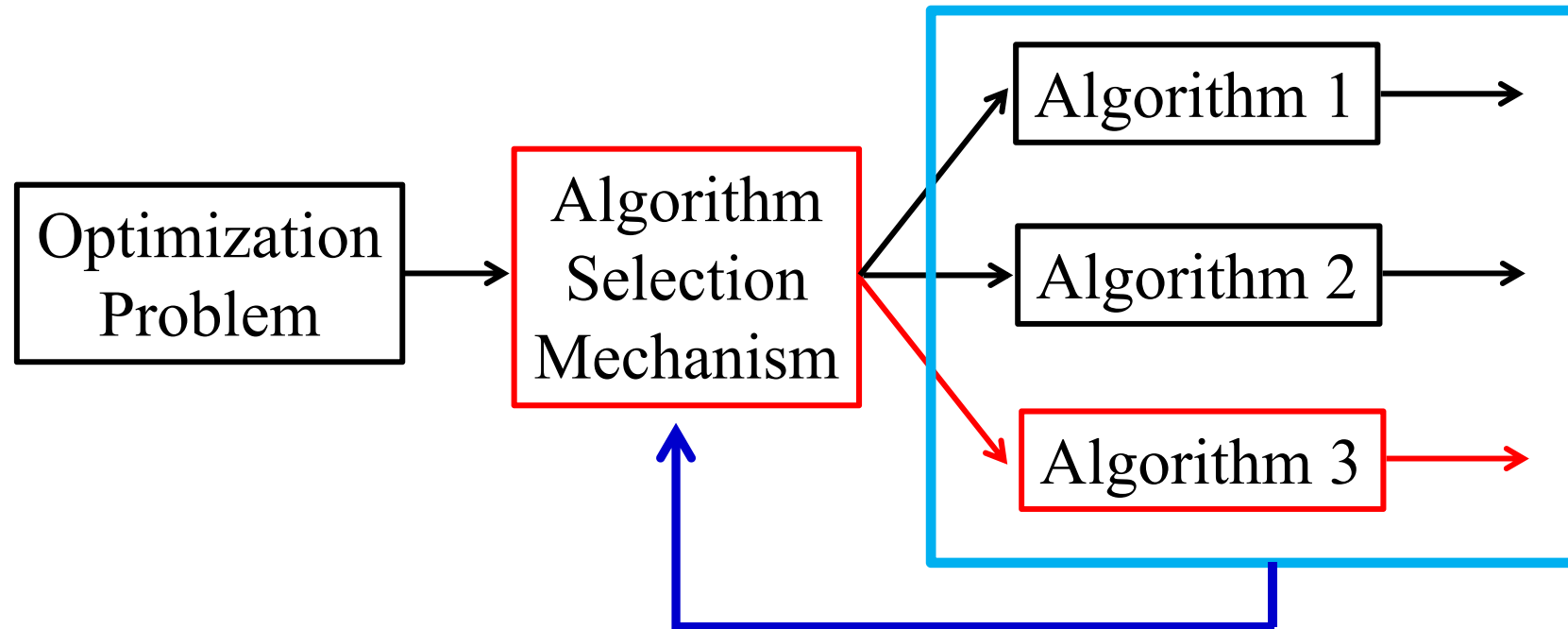
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Variable Neighborhood Search

Algorithm Ensemble

Based on the characteristics of the problem, an appropriate algorithm will be selected.

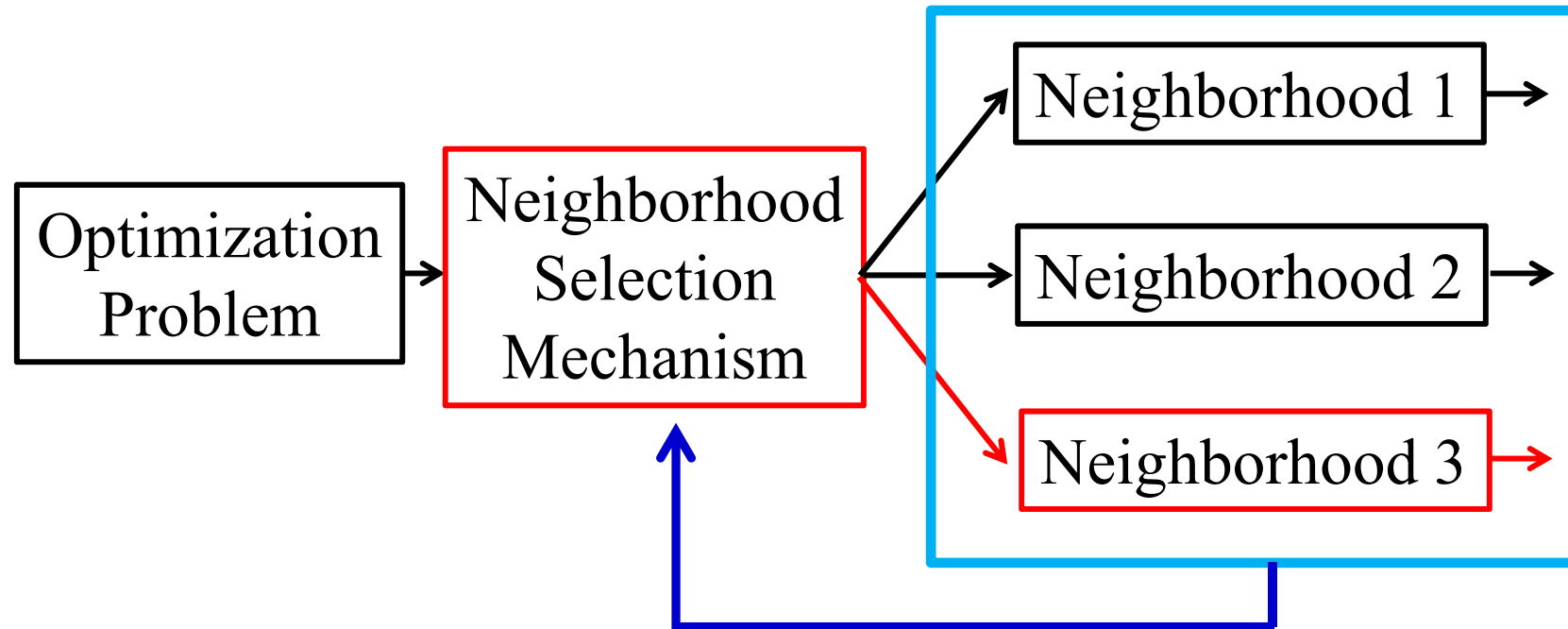


Depending on the progress of the search, the selection of an appropriate algorithm will be changed.

Variable Neighborhood Search

Algorithm Ensemble

Based on the characteristics of the problem, an appropriate algorithm will be selected.



Depending on the progress of the search, the selection of an appropriate neighborhood will be changed.

Lab Session Task 3

Specify neighborhood structures $\mathcal{N}_1, \mathcal{N}_2, \mathcal{N}_3, \dots$ for designing a **Variable Neighborhood Descent (VND)** algorithm for a TSP problem.