# File Organization

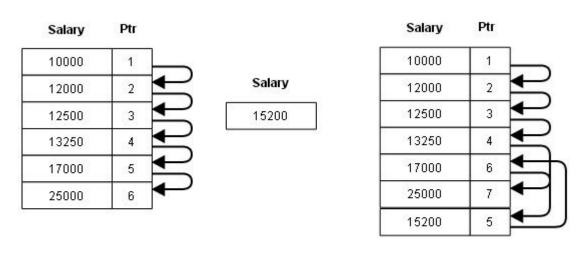
And Hashing

How do you physically store the data?

### Objectives of file organization

- 1. Records should be accessed as fast as possible
- 2. Any insert, update or delete transaction on records should be easy, quick and should not harm other records.
- 3. No duplicate records should be induced as a result of insert, update or delete
- 4. Records should be stored efficiently so that cost of storage is minimal.

### 1. Sequential File Organization



(a) Initial status of table

(b) Record to be inserted

(c) After insertion of new record

Sequential (sorted) file organization - Record insertion

#### Insertion:-

- 1. Insert the record at the end of the file
- 2. Set appropriate pointer for the inserting record
- 3. Update other record pointers accordingly
- Reorganize all the records to arrange the records in sorting order of the ordering attribute

#### Deletion:-

- 1. Locate the record to be deleted
- 2. Update all other records' pointer

#### Advantages:

1. Efficient query retrieval if the query uses sorting attribute as the search key

#### Disadvantages:

- 1. Insertion and deletion are expensive
- 2. Updating the sorting attribute values is expensive
- 3. Retrieval of records on non-ordering attributes is not efficient

### 2. Heap file organization (unordered)

Simplest type of file organization

#### Insertion:

- 1. New records are inserted at the end of the file
- 2. No order is required in storing the database

#### Deletion:

 Frequent deletion of records would create much free space, thus needs periodic reorganization of the file

#### Advantages:

- 1. Simple file organization
- 2. Insertion is easy
- 3. Best method for bulk loading of data

#### Disadvantages:

- 1. Because of no particular ordering, finding a record would require to do a linear search only.
- 2. Frequent deletion would required periodic file reorganization

### 3. Hash file organization

It is a file organization technique where a hash function is used to compute the address of a record. It uses value of an attribute or set of attributes and gives the location where the record could be stored

#### Points:

- If bucket(s) is/are full, then overflow buckets can be used to store more records.
- 2. The attribute(s) that is frequently used for data manipulation can be chosen as the input for the hash function.
- Same hash function that was used to store the records has to be used for deletion, modification or selection of records.

### Advantages:

- 1. Quick access to records in terms of selection. [If queried on the attribute that was used for hashing]
- 2. Easy to insert, delete, or update a record.

#### Disadvantages:

- Records are randomly stored in scattered locations. May waste a lot of space in case of small files.
- 2. Not efficient for range based queries
- 3. If querying attribute is not the hashed attribute, you may need to scan the entire table for retrieval.

## Two types of hashing

- 1. Static Hashing
- 2. Dynamic Hashing

### Static hashing

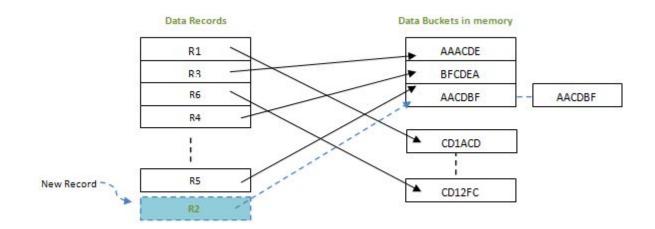
In this method, the resultant data address bucket will always be the same.

Problem: Bucket overflow

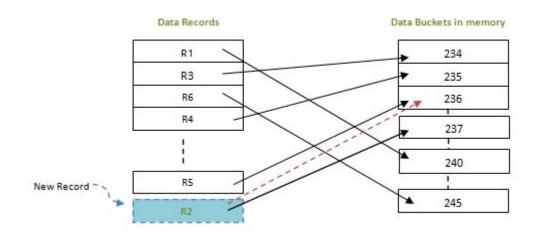
What if the data bucket address generated by the hash function is already full?

Some ways to resolve it.

### Closed hashing (overflow chaining)



### Open hashing (linear probing)



### Dynamic Hashing (extendible hashing)

In this method of hashing, data buckets grows or shrinks as the records increases or decreases.

Continuing with an example.