

Assignment-3Submitted by

Sayam Kumar

S20180010158

Problem 3.58

Parameters  $x$ ,  $y$  and  $z$  are passed in registers  $\%rdi$ ,  $\%rsi$  and  $\%rdx$ . As the return value is stored in  $\%rax$ . we will use a variable result to store result value.

decode 2:

```

subq    %rdx, %rsi
imulq   %rsi, %rdi
movq    %rsi, %rax
salq    $63, %rax
sarq    $63, %rax
xorq    %rdi, %rax
ret

```

Explanation $y = y - z$ ; Dest = Dest - Src $x = x * y$ 

result = y

result = result  $\ll$  63result  $\ll$  = 63result = result  $\wedge$  x

The corresponding C code is -

```

long decode2(long x, long y, long z)
{
    long result = y - z;
    return ( result << 63 >> 63 )  $\wedge$  ( result * x );
}

```

subq
sal
sar
xor

Now explaining the C code -

First,  $y - z$  is stored in result.

Then corresponding left and right shift are done and result is xor with result  $* x$ . That's why the C code is written.

Final ans -

long decode2 (long x, long y, long z)

```
{  
    long result = y - z;  
    return (result << 63 >> 63) ^ (result * x);  
}
```

Problem 3.60 x in %rdi, n in %esi

loop:

movl %esi, %ecx

Explanation

Storing n in %ecx

movl \$1, %edx

A variable (call mask) = 1

movl \$0, %eax

result = 0

jmp .L2

.L3

movq %rdi, %rbx

A variable (call temp) = x

andq %rdx, %rbx

temp = temp & mask

orq %rbx, %rax

result |= temp

salq %cl, %rdx

mask <<= n

.L2

testq %rdx, %rdx

mask & mask

jne .L3

mask != 0

rep; ret

Filling in the code -

long loop (long x, long n)

```
{  
    long result = 0;  
    long mask;  
    for (mask = 1; testq mask != 0; mask = mask << n) {  
        result |= mask & x; salq  
        orq andq  
    }  
    return result;  
}
```



A Which registers hold program values  $x$ ,  $n$ , result and mask?

Sol- It is given in question that  $x$  is stored in `%rdi`,  $n$  in `%esi`. Through the assembly code, it is clear that result is stored in `%rax` as `%rax` stores return value and mask is stored in `%rdx`.

B What are the initial values of result and mask?

Sol- Result = 0, mask = 1

because `movl $1, %edx`  $\Rightarrow$  mask = 1

`movl $0, %eax`  $\Rightarrow$  result = 0

C What is the test condition for mask?

Sol- mask  $\neq 0$

because `testq %rdx, %rdx`  $\Rightarrow$  mask  $\neq 0$   
`jne .L3`

D How does mask get updated?

Sol- mask = mask  $\ll n$

because `salq %cl, %rdx`  $\Rightarrow$  mask = mask  $\ll n$

E How does the result get updated?

Sol- result  $\mid$  = mask &  $x$

because  
`andq %rdx, %r8`  $\Rightarrow$  result  $\mid$  = mask &  $x$   
`orq %r8, %rax`

F Already written on previous page.

### Problem 3.62 Understanding the assembly Code line by line

p1 in %rdi, p2 in %rsi, action in %edx

.L2

```
movl    $27, %eax
ret
```

Explanation

result = 27  
break

.L3

```
movq    (%rsi), %rax
movq    (%rdi), %rdx
movq    %rdx, (%rsi)
ret
```

result = \*p2  
%rdx = \*p1  
\*p2 = %rdx  
break

.L5

```
movq    (%rdi), %rax
addq    (%rsi), %rax
movq    %rax, (%rdi)
ret
```

result = \*p1  
result += \*p2  
\*p1 = result  
break

.L6

```
movq    $59, (%rdi)
movq    (%rsi), %rax
ret
```

\*p1 = 59  
result = \*p2  
break

.L7

```
movq    (%rsi), %rax
movq    %rax, (%rdi)
movl    $27, %eax
ret
```

result = \*p2  
\*p1 = result  
result = 27  
break

.L9

```
movl    $12, %eax
ret
```

result = 12  
break

Now we have understood what each line is doing,  
we can put them all in C code on  
next page

long switch 3 (long \*p1, long \*p2, mode-t action)

```
{ long result = 0;
  switch (action) {
    case MODE-A:
      result = *p2;
      *p2 = *p1;
      break;
    case MODE-B:
      result = *p1 + *p2;
      *p1 = result;
      break;
    case MODE-C:
      *p1 = 59;
      result = *p2;
      break;
    case MODE-D:
      *p1 = *p2;
    case MODE-E:
      result = 27;
      break;
    case
    default:
      result = 12;
      break;
  }
  return result;
}
```

The reason why we have written break in case MODE-D because even MODE-E is calculating the same thing of result = 27. So no break statement in MODE-D.



Problem ~~6.63~~ 6.63 3.63 x in %rdi, n in %rsi  
for the jump table

0x4006f8: 0x4005a1 (60) 0x4005c3 (61)  
0x400708: 0x4005a1 (62) 0x4005aa (63)  
0x400718: 0x4005b2 (64) 0x4005bf (65)

Now understanding the assembly line by line  
Explanation

sub \$0x3c, %rsi n = n - 60  
cmp \$0x5, %rsi n-5 calculating  
ja 4005c3 <switch-prob  
+0x33> jump above  
⇒ switch cases are  
60, 61, 62, 63, 64, 65.

# So the address in jump table now map to 60, 61,  
62, 63, 64, 65 sequentially.

lea 0x0(, %rdi, 8), %rax	result = %rdi * 8
retq	break
mov %rdi, %rax	result = x
sar \$0x3, %rax	result >>= 3
retq	break
mov %rdi, %rax	result = x
shl \$0x4, %rax	result <<= 4
sub %rdi, %rax	result -= x
mov %rax, %rdi	x = result
imul %rdi, %rdi	x * = x
lea 0x4b(%rdi), %rax	result = x + 75
retq	break

By looking at the labels of addresses given in front of  
these assembly instructions. the following C code can  
be written.

```

long switch-prob (long x, long n) {
    long result = x;
    switch (n) {
        case 60:
        case 62:
            result *= 8;
            break;
        case 63:
            result >>= 3;
            break;
        case 64:
            x = (x < 4) - x;
        case 65:
            x *= x;
        default:
            result = x + 75;
            break;
    }
    return result;
}

```

Explanation - As there are no set statements in assembly for labels 64, 65, we do not write break statement in these cases. Moreover for case 61, it maps to default case as it is evident from the address of jump tables for label 61 ( $0x4005c3$ ).

---

# Problem 3.64

long store\_ele (long i, long j, long k, long \*dest)  
 i in %rdi, j in %rsi, k in %rdx, dest in %rcx

store\_ele:

```
leaq    (%rsi, %rsi, 2), %rax
leaq    (%rsi, %rax, 4), %rax
movq    %rdi, %rsi
salq    $6, %rsi
addq    $rsi, %rdi
addq    %rax, %rdi
addq    %rdi, %rdx
movq    A(, %rdx, 8), %rax
movq    %rax, (%rcx)
movl    $3640, %eax
ret
```

Explanation

$rax = 3rsi$   
 $rax = 13rsi$   
 $rsi = i$   
 $rsi = 64i$   
 $rdi = 65i$   
 $rdi = 65i^2 + 13rsi$   
 $k = k + 13j + 65i$   
 $result = A[k]$   
 $*dest = A[8(k + 13j + 65i)]$   
 set size of A as 3640

A) Extending the solution of 2D for 3D matrices -

$$A[i][j][k] = x - A + \text{size of (long)} (S * T * i + T * j + k)$$

B)  $A[8(13(Si + j) + k)]$  from \*dest

$\uparrow$   $\nwarrow$   
 $T=13$   $S=5$

Size of A = 3640

$$\Rightarrow R = \frac{A}{8 \times T \times S} = \frac{3640}{8 \times 13 \times 5} = 7 \quad (8 = \text{size of (long)})$$

So  $R=7$ ,  $S=5$ ,  $T=13$



Problem 6-45 3.65

The assembly code for transpose of a matrix is given below →

.L6:

movq (%rdx), %rcx

movq (%rax), %rsi

movq %rsi, (%rdx)

movq %rcx, (%rax)

addq \$8, %rdx

addq \$120, %rax

cmpq %rdi, %rax

jne .L6

Explanation

$$rcx = A[i][j]$$

$$rsi = A[j][i]$$

$$A[i][j] = rsi$$

$$A[j][i] = rcx$$

} swap

taking next element } size of  
going to next row } long = 8

rax - rdx compare

for loop

~~A) %rdx holds %.~~

A) %rdx holds  $A[i][j]$

B) %rax holds  $A[j][i]$

C) At size of long = 8 and we are 120 bytes ahead of  
A to get next row  $\Rightarrow M = \frac{120}{8} = 15$

### Problem 3.61

for this trial implementation of a condition move instruction, this can lead to segmentation fault -

```
long cread(long *xp) {  
    return (xp? *xp : 0);  
}
```

cread:

movq (%rdi), %rax

testq %rdi, %rdi

movl \$0, %edx

cmovle %rdx, %rax

ret

Explanation

$v = *p \leftarrow$  can lead to segmentation fault

Test x

Set  $ve = 0$

If  $x = 0$ ,  $v = ve$

Return v

To avoid this,

we can write

```
long cread-alt(long *xp)  
{  
    long tmp = 0L;  
    return * (xp? xp : &tmp);  
}
```

Hence, null value reference can be resolved with respect to adding a temporary variable and returning its address.

### Problem 3-68

for the function setval

max 8 (%rsi), %rax will store q→t. As int is 4 bytes long

$$4 < B \leq 8 \quad \text{--- (1)}$$

Similarly

32 (%rsi) will fetch q→u. now as long it is 8 bytes long.

$$\Rightarrow 2A \leq 20 \\ A \leq 10 \quad \text{--- (2)}$$

Now (%rdi + 16) represents y

No of elements in AB  $\Rightarrow$

$$4AB \leq 192 \Rightarrow AB \leq 48 \quad \text{--- (3)}$$

With equations (1), (2) and (3)

B=5, A=9 is the solution

Problem 3-71 The C program is as follows

```
#include <stdio.h>
```

```
#include <assert.h>
```

```
#define BUFFER 50
```

```
void good_echo()
```

```
{ char string[BUFFER];
```

```
while (1){
```

```
char *p = fgets(string, BUFFER, stdin);
```

```
if (p == NULL){
```

```
printf("No string entered.");
```

```
break;
```

```
}
```

```
printf("%s\n", p);
```

```
}
```

```
return;
```

```
}
```

```
int main() {  
    good_echo();  
    return 0;  
}
```

The error handled is in when no string is entered or there is a buffer overflow. fgets throws error on buffer overflow.



# Problem 3-66

$n$  in  $\%rdi$ ,  $A$  in  $\%rsi$ ,  $j$  in  $\%rdx$

sum\_col:

```
leaq 1(, %rdi, 4), %r8
leaq (%rdi, %rdi, 2), %rax
movq %rax, %rdi
testq %rax, %rax
jle .L4
scaq $3, %r8
leaq (%rsi, %rdx, 8), %rcx
movl $0, %eax
movl $0, %edx
```

Explanation

$t_1 = 4n + 1$   
 $t_2 = 3n$   
 $t_3 = 3n$   
 $t_{ext} = 3n$   
 $3n \leq 0$   
 $t_1 = 8t_1 = 8(4n + 1)$   
 $t_4 = 8j + A$   
 $t_2 = 0$   
 $t_5 = 0$

.L3

```
addq (%rcx), %rax
addq $1, %rdx
addq %r8, %rcx
cmpq %rdi, %rdx
jne .L3
up ret
```

$t_2 = *t_4 = *(A + 8j)$   
 $t_5 = t_5 + 1$   
 $t_4 + t_1 = t_4$   
 Cmp  $t_5$  &  $t_3$   
 if  $t_5 \neq 3n$  loop

.L4:

```
movl $0, %eax
ret
```

$t_5$  return 0  $\%rax = 0$

```
cmpq %rdi, %rdx
jne .L3
```

$t_5, t_3$  compare  
 $t_5 \neq 3n$

(NR)  $n = 3n$

```
leaq 1(, %rdi, 4), %r8
scaq $3, %r8
addq %r8, %rcx
```

$t_1 = 4n + 1$   
 $t_1 = 8t_1$   
 $t_4 = t_1 + t_4$

3n every loop, pointer move  $\mathcal{O}(4n + 1)$

$N(n) = 4n + 1$