

**Computer Science and Engineering**  
**IIIT Kalyani, West Bengal**

**Compilers Design Laboratory (Spring: 2017 - 2018)**

*3rd Year CSE: 6th Semester*

Assignment - 5

Marks: 20

Assignment Out: 9<sup>th</sup> February, 2018

Report on or before: 9<sup>th</sup> March, 2018

1. Consider the following *context-free grammar*  $G$ : the *non-terminals* are  
 $N = \{ AE, BE, D, DL, E, F, ES, IOS, IS, NE, P, PE, RE, S, SL, T, TY, VL, WS \}$ ,  
the *terminals* are  
 $\Sigma = \{ + - * / = < > ( ) \{ \} := ; \text{ and else end ic id if int do fc float not or print prog scan str then while } \}$ ,  
the *start symbol* is  $P$  and the production rules are,

```
P → prog DL SL end
DL → D DL | ε
D → TY VL ;
TY → int | float
VL → id VL | id
SL → S SL | ε
S → ES | IS | WS | IOS
ES → id := E ;
IS → if BE then SL end |
    if BE then SL else SL end
WS → while BE do SL end
IOS → print PE | scan id
PE → E | str
BE → BE or AE | AE
AE → AE and NE | NE
NE → not NE | { BE } | RE
RE → E = E | E < E | E > E
E → E + T | E - T | T
T → T * F | T / F | F
F → ( E ) | id | ic | fc
```

Most of the terminals are self-explanatory. **id** is an (identifier), **ic** is an integer constant, **fc** is a floating-point constant, **str** is a string of characters, **:=** is an assignment etc.

2. Transform the grammar  $G$  to an equivalent  $LL(1)$  grammar  $G_1$  by removing *left recursion* and *left factoring*. Compute  $First(A \rightarrow \alpha)$  for every production rule and  $Follow(A)$  of every non-terminal producing an  $\epsilon$ .

In the report clearly present this computation to establish that the grammar is  $LL(1)$ .

3. You already have the scanner written for these non-terminals in the assignment 3. The function `yylex()` of the scanner returns tokens. Attributes are available in global variables.
4. Write a C program to implement a *recursive descent predictive parser* for the modified grammar  $G_1$ . The parser calls `int yylex()` for token and uses the global variables if required.

Given an input string  $x$  over the alphabet  $\Sigma$ , the parser either *accepts* or *rejects* the string depending on whether  $x \in L(G) = L(G_1)$  or not. If there is an error, it may also print the *line number* and *character position* where the error has occurred.

5. Take the scanner files `myLex.c++` and `myLex.h` from the assignment 3. The name of the parser file should be `myRDPparser.c++`.

6. Prepare a `Makefile` to compile both the files. Finally prepare a `.tar` file  
(`$ tar cvf <rollNo>.tar Makefile myLex.c myRDPparser.c myLex.h`)  
and return the `.tar` file.