

# A\* Search Algorithm

28/07/19

- informed, guided search
- uses heuristics
- selects path with least cost at each iteration

$$f(n) = g(n) + h(n)$$

total cost of path if node "n" is chosen      cost from starting node to node "n"      heuristic function

$$\text{Objective: } \underset{n}{\operatorname{argmin}} f(n)$$

- A\* returns least cost path iff heuristic is **admissible** (never overestimates actual cost to get to goal) ↑
- implemented using priority queue
- Dijkstra can be seen as a special case of A\* where  $h(n) = 0 \forall \text{ nodes}$
- time complexity is polynomial iff:
  - there is a single goal state
  - search space is a tree
  - $|h(x) - h^*(x)| = O(\log h^*(x))$   
     $\hat{=}$   $h^*$  is the optimal heuristic

## Algorithm

procedure  $A^*$ (source, destination,  $h$ ):

    pqqueue :=  $\{\}$  // priority queue

    predecessor := nil // mapping from node to node

$g := \{\infty \text{ for all nodes}\}$

$g[\text{source}] = 0$

$f := \{\infty \text{ for all nodes}\}$

$f[\text{start}] = h(\text{start})$

    while pqqueue is not empty:

        current = pqqueue.extract\_min()

        if current == goal:

            return reconstruct\_path(predecessor, current)

        pqqueue.remove(current)

        for each neighbour of current:

$g' := g[\text{current}] + d(\text{current}, \text{neighbour})$

            if  $g' < g[\text{current}]$ :

                predecessor[neighbour] = current

$g[\text{neighbour}] = g'$

$f[\text{neighbour}] = g' + h(\text{neighbour})$

if neighbour not in pqueue:  
pqueue.add(neighbour)

return error

procedure **reconstruct\_path** (predecessor, current):  
path := { current }

while current in predecessor:  
current = predecessor[current]  
path = [current] + path

return path