Homework 1

Submitted by:

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ROIL NO: EE 18 MTECH 11030

Subject: Introduction to Modern AI Ams 9: Desire the Bellman's equation for: (3) State-rame function Va(3). Ne know, $V_{\pi}(s) = E_{\pi} \left[R_{t} / S_{t} = 6 \right], R_{t} = \sum_{k=0}^{2^{k}} l_{t+k+1}$ $= E_{\pi} \left[\begin{array}{ccc} z & z^{k} & \text{RE+k+1} & \text{S}_{4} = 6 \end{array} \right] - 0$ We use the law of total enfectation to desire the Bellonan's equation. The law of total enfectation is given as, E[X] = SE[X|A;] P(A;), LA;; is a comtable fastition of the sample of the Now, IP (at=a | St=s) is the foling of (s,a) $\mathcal{L}_{A}(s) = \sum_{a} \pi(s,a) \; \mathcal{E}_{\pi} \left[\sum_{k=0}^{\infty} 2^{k} \, \mathcal{E}_{t+k+1} \, \middle| \, S_{t} = S, \; A_{t} = a \right]$ Now, conditions 3 on the future state of we get,

$$(3) \Rightarrow \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \sum_{n} \sum_{t=0}^{s'} \chi^{k} \chi_{t+k+1} | s_{t}=s, a_{t}=a, s_{t+n} s' | \mathcal{N}_{t+n} s' | s_{t}=s, a_{t}=a)$$

$$\Rightarrow \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \sum_{t=0}^{s'} \chi^{k} \chi_{t+k+1} | s_{t}=s, a_{t}=a, s_{t+n} s' | \mathcal{N}_{(s,n)} s' | \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \mathcal{N}_{(s,n)} | \mathcal{N}_{\pi}(s) | \mathcal{N}_{\pi}(s) | \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \mathcal{N}_{(s,n)} | \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \mathcal{N}_{\pi}(s) | \mathcal{N}_{\pi}(s) | \mathcal{N}_{\pi}(s) = \sum_{n} \pi(s,n) \sum_{s'} \mathcal{N}_{\pi}(s) | \mathcal$$

(29) Action-value function: PA(3,a) = ET Rt | St=8, 9t= a], Rt = = 24 EC+R+1 Now, conditioning (4) on the future states, we obtain, $\Rightarrow 2\pi(s,a) = \sum_{s,r} E_{\pi} \left[\sum_{k=0}^{\infty} 2^{k} \mathcal{L}_{k+k+1} \left[s,a,s^{r} \right] \mathbb{P}(s' \mid \beta,a) \right]$ Denotes Isia) = IP(s'(s,a) $\Rightarrow \Re(S_{i}a) = \sum_{c'} E_{\pi} \left[\sum_{k=0}^{\infty} 2^{k} \mathcal{L}_{c+k+1} \left[S_{i}a, s' \right] \right] \mathcal{P}_{(S_{i}a)}^{S'}$ = E PGIA) Ex [441 | SI A, S'] + 2 = PGA ET = 2KALLEL [SIAIS]

Denote, Retti = Ex [Etti | 2,9,5']
and observe mad, by Markov foopers,

ET [Z 2 KALLEL [S19,5]] = ET [Z 2 KALLELZ [S']

 $= \sum_{s'} \Re(s, a) = \sum_{s'} \Re(s, a) \Re(s, a) + 2 \sum_{s'} \Re(s, a) \exp(s') + 2 \sum_{s'} \Re(s') + 2 \sum_{s'} \Re(s') \exp(s') + 2 \sum_{s'} \Re(s') + 2$

Now, conditioning \mathfrak{D} on a' (forsible actions when) we have, $\mathfrak{B} = \mathfrak{I}_{R}(s,a) = \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}^{A+1}_{(s,a)} + \mathfrak{I}_{S'}^{S'} \mathbb{R}^{A+1}_{(s,a)} \mathbb{R}$

$$2\pi(s,a) = \sum_{s'} \mathbb{P}'_{(s,a)} \left[\mathbb{R}^{k+1}_{(s,a,s')} + 2 \sum_{a'} 9_{\pi}(s',a') \pi(s',a') \right]$$

Ams 2 (?) State transition table for enample 1 description: A mobile rootest has the job of firking cans and folding them on ben. The bet owns on battery. The decision on how to search for a can is made by reinforcement learning agent leased on the charge level of the botts leathery. The agent should devide whether the bot Should (i) actively search for ears, (ii) wait for comeone to bring the came to it (399) nechange the battery. The states are: 3= [lon, high (change levels of bot's battery) The actions connexponding to states are: X (low) = { seauch, want, nechange} A (high) = { seauch, won't } The comes fording rewards are assumed to be: Kreanch: average vervard during reach. kwart: average remand during wast. -3: when sechange after suming out 0: donning rechange

5	S'	a	f(s' s,a)	14 (0,3,5)
high	ligh	Search	\mathcal{L}	Pseanch
high	low	search	1-2	Reauch
high	high	wast	1	Carit
high	low	wast	0	Knearet
lon	low	senuch	P	Reamh
low	high	search	1-B	-3
low	low	wast	<u> </u>	ywart
low	ligh	sort	0	Knaft
lon	low	recharge	D	0
low	high	sechange	4	0

$$g = \{ q, B \}, G = good extre$$

$$B = End extre$$

The actions connexponding to states are:
$$X(G) = \{Stary, move\}$$

Legal = 3 (for staying in state G)

Equal = -1 (for soing to state B)

Equal = -1 (for soing to state B)

Equal = -1 (for staying in state B)

Equal = -1 (for staying in state B)

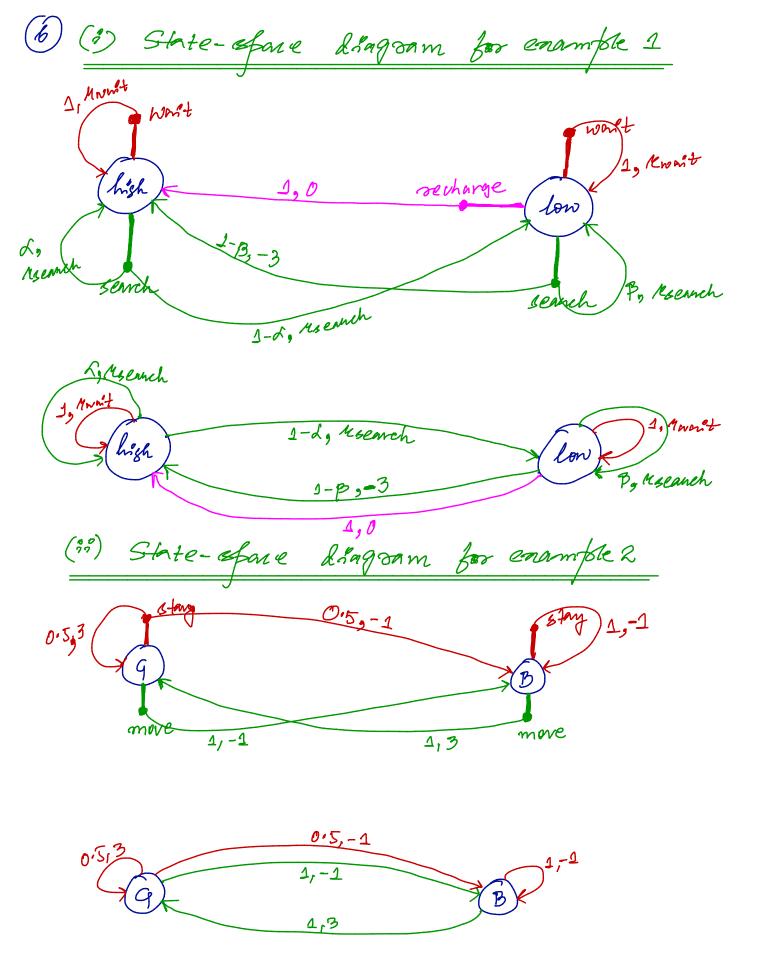
Equal = 3 (for moving to state G from)

Equal = -1 (for moving to state G from)

Equal = -1 (for moving to state B)

State transition table:

8	6	a	B(g/15,a)	4 (A, S, S)
9	9	stang	0.2	3
9	В	Stay	0.2	-1
9	9	move	0	0
9	75	move	1	<u>-1</u>
お	B	Story	4	-1
B	Q	Hay	0	0
B	B	move	O	O
13	9	more	4	3



After Emplementing the value iteation algorithm for the can collecting notot enample, the following results are obtained: Considering $\Delta = 100$, $E = 1 \times 10^{9}$, oftemal values V*(S=High) = 3-176470587 v*(s= Low) = 1.9999999 Iteration needed = 31 Policy: State = High, Action = search Policy: State = Low, Action = won't Python file: EE 18 MTECH 11 030_ Question_ 2-C- final. itynto After Emplementing the value iteation algorithm for the second MDP enample, the following nesults are obtained: Considering $\Delta = 100$, $E = 5 \times 10^{-9}$ oftemal values V*(S=Hogh) = 2079999999 v*(s= Low) = 4.3999999 Iteration needed = 30 Policy: State = Good, Action = Stay Policy: State = Bad, Action = More Kython file: EE 18 MTECH 11 030_ Question_ 2-d-bonal. itynto

Ams: 1)

Deriving the Bellman's equation for state-

Non,

$$\frac{1}{2} \int_{k=0}^{\infty} \frac{1}{2^{k}} \left[\frac{1}{2^{k}} \right] dk = \frac{1}{2^{k}} \int_{k=0}^{\infty} \frac{1}{2^{k}} \frac{1}{2^{$$

Now,

$$E_{\pi}[E_{\pi}[q_{t+1}|s_{t+1}|c_{t}] = E_{\pi}[q_{t+1}|s_{t}]$$
 $food:$

Let $G_{t+1} = g'$, $S_{t+1} = 8'$, $G_{t} = g$, $S_{t} = 8$

Then,

 $E_{\pi}[E_{\pi}[q_{t+1}|s_{t+1}]|c_{t}] = E_{\pi}[E_{\pi}[q'|s',s]]$
 $= E_{\pi}[S_{g'}[g',s']]$

=
$$\sum_{s'} \sum_{g'} g' p(g' | g', s) p(s' | s)$$

= $\sum_{s'} \sum_{g} g' \times p(g', g', s) p(g', s) p(g', s)$
= $\sum_{s'} \sum_{g} g' \times p(g', g', s) p(g', s) p(g', s)$
= $\sum_{s'} \sum_{g} g' p(g' | g', s) p(g', s) p(g', s)$

=
$$\sum_{s'} \sum_{g'} \frac{1}{f(g', g', s)} \frac{1}{f(s)}$$

= $\sum_{g'} \frac{1}{f(g', s)} \frac{1}{f(s)}$
= $\sum_{g'} \frac{1}{f(g', s)} \frac{1}{f(s)}$
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