

Contents

1	Review of Propositional Logic	2
1.1	Connectives	2
1.1.1	Truth Table of the Connectives	2
1.2	Important Tautologies	2
1.3	Indirect Arguments/Proofs by Contradiction/Reductio as absurdum	3

1 Review of Propositional Logic

Task: Recall enough propositional logic to see how it matches up with set theory.

Definition: A proposition is any declarative sentence that is either true or false.

1.1 Connectives

	<u>Connectives</u>	<u>Notation in Maths</u>
and	\wedge	
or	\vee	"Inclusive or"
not	\neg	Sometimes denoted \sim
implies	\rightarrow	if/then; called implication \Rightarrow
if and only if	\leftrightarrow	Called equivalence \Leftrightarrow

1.1.1 Truth Table of the Connectives

Let P, Q be propositions:

P	Q	$P \wedge Q$	P	Q	$P \vee Q$	P	$\neg P$	P	Q	$P \rightarrow Q$	P	Q	$P \leftrightarrow Q$
F	F	F	F	F	F	F	T	F	F	T	F	F	T
F	T	F	F	T	T	F	T	F	T	T	F	T	F
T	F	F	T	F	T	T	F	T	F	F	T	F	F
T	T	T	T	T	T	T	T	T	T	T	T	T	T

Priority of the Connectives

Highest to Lowest: $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$

1.2 Important Tautologies

$$\begin{array}{ll}
 (P \rightarrow Q) & \leftrightarrow (\neg P \vee Q) \\
 (P \leftrightarrow Q) & \leftrightarrow [(P \rightarrow Q) \wedge (Q \rightarrow P)] \\
 \neg(P \wedge Q) & \leftrightarrow (\neg P \vee \neg Q) \\
 \neg(P \vee Q) & \leftrightarrow (\neg P \wedge \neg Q)
 \end{array}
 \left. \vphantom{\begin{array}{l} (P \rightarrow Q) \\ (P \leftrightarrow Q) \\ \neg(P \wedge Q) \\ \neg(P \vee Q) \end{array}} \right\} \text{De Morgan Laws}$$

As a result, \neg and \vee together can be used to represent all of $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$.

Less obvious: One connective called the sheffer stroke $P|Q$ (which stands for "not both P and Q" or "P nand Q") can be used to represent all of $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$ since $\neg P \leftrightarrow P|P$ and $P \vee Q \leftrightarrow (P|P) | (Q|Q)$.

Recall if $P \rightarrow Q$ is a given implication, $Q \rightarrow P$ is called the converse or $P \rightarrow Q$.
 $\neg Q \rightarrow \neg P$.

1.3 Indirect Arguments/Proofs by Contradiction/Reductio as absurdum

Based on the tautology $(P \rightarrow Q) \leftrightarrow (\neg Q \rightarrow \neg P)$

Example: Famous argument that $\sqrt{2}$ is irrational.

Proof:

Suppose $\sqrt{2}$ is rational, then it can be expressed as fraction form $\frac{a}{b}$. Let us **assume** that our fraction is in the lowest term, **i.e.** their only common divisor is 1.

Then,

$$\sqrt{2} = \frac{a}{b}$$

Squaring both sides, we have

$$2 = \frac{a^2}{b^2}$$

Multiplying both sides by b^2 yields

$$2b^2 = a^2$$

Since $a^2 = 2b^2$, we can conclude that a^2 is even because whatever the value of b^2 has to be multiplied by 2. If a^2 is even, then a is also even. Since a is even, no matter what the value of a is, we can always find an integer that if we divide a by 2, it is equal to that integer. If we let that integer be k , then $\frac{a}{2} = k$ which means that $a = 2k$.

Substituting the value of $2k$ to a , we have $2b^2 = (2k)^2$ which means that $2b^2 = 4k^2$. dividing both sides by 2 we have $b^2 = 2k^2$. That means that the value b^2 is even, since whatever the value of k you have to multiply it by 2. Again, if b^2 is even, then b is even.

This implies that both a and b are even, which means that both the numerator and the denominator of our fraction are divisible by 2. This contradicts our **assumption** that $\frac{a}{b}$ has no common divisor except 1. Since we found a contradiction, our assumption is, therefore, false. Hence the theorem is true.

qed