

* Operating System * [Module 1].

- Computer System
 - OS manages hardware resources of a computer system.
 - Manage all the task in execution with respect to computer system.
 - Manages hardware resources and has eye on execution unit.
 - Manages Jobs, Tasks, processes in computer system.
 - Operating System is a Software. It is an interface between Hardware and user.
 - Operating system is a resource manager.
 - It can take input from user and process that input and give output to the user.
 - OS acts as an interface between Hardware and User.
 - manages, processes and converts input to desire output.
- Operating System
- Manually all the things were done before OS.
 - User need to do all the things.
 - Efficiency is improved after using OS.
 - Speed gets increased.
 - System is now multi-tasked.
- Operating System
- OS is a Software.
 - Software is Hardware dependent.
 - Software gets installed in Hardware.
 - Need of physical memory.
 - When a program runs into ram to produce intermediate results which is stored in registers.

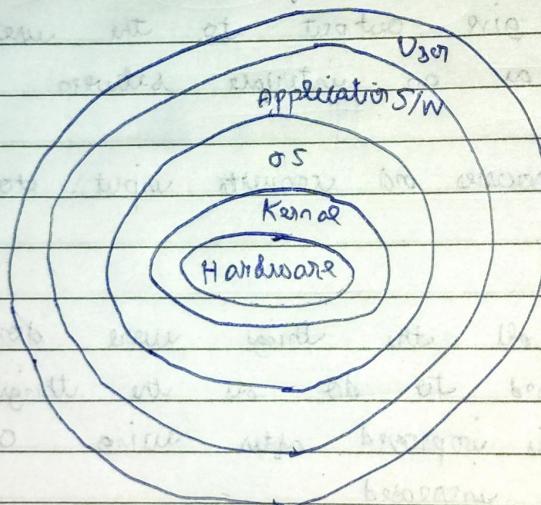
- Cache is a high speed memory than a RAM.

* What is OS?

- Hardware Manager: It manages all the hardware resources or components of computer.
- Process Manager: It supervises all the task which is being executed by processor.

* How OS is different from other application software?

- OS is installed forever in hard-drive.
- Applications are also installed over OS. Application are runned inside OS.

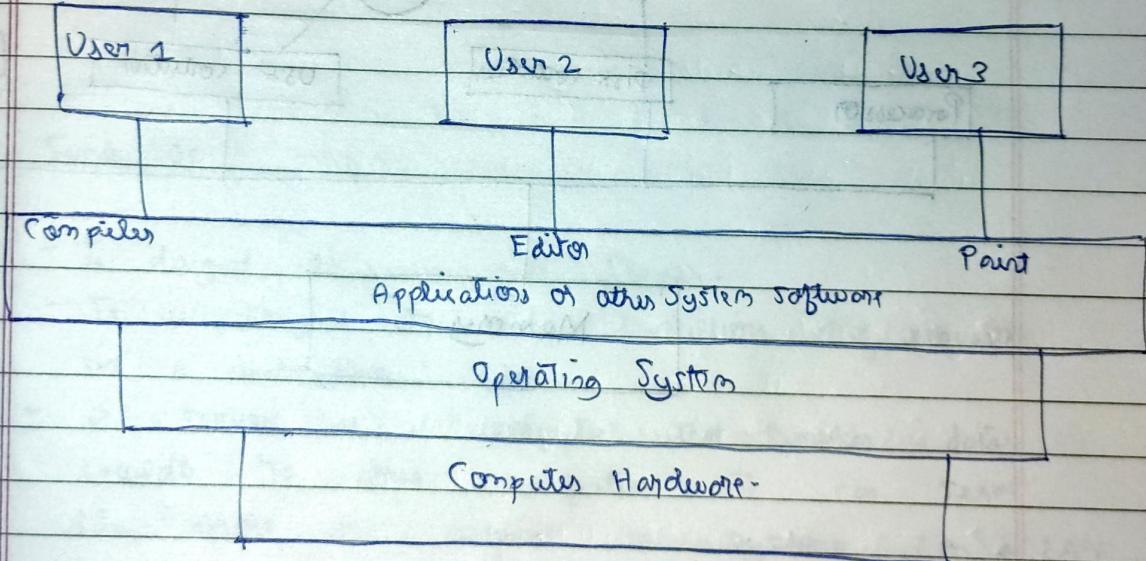


* Operating System Layered Architecture

OS runs over computer, Application runs over OS.

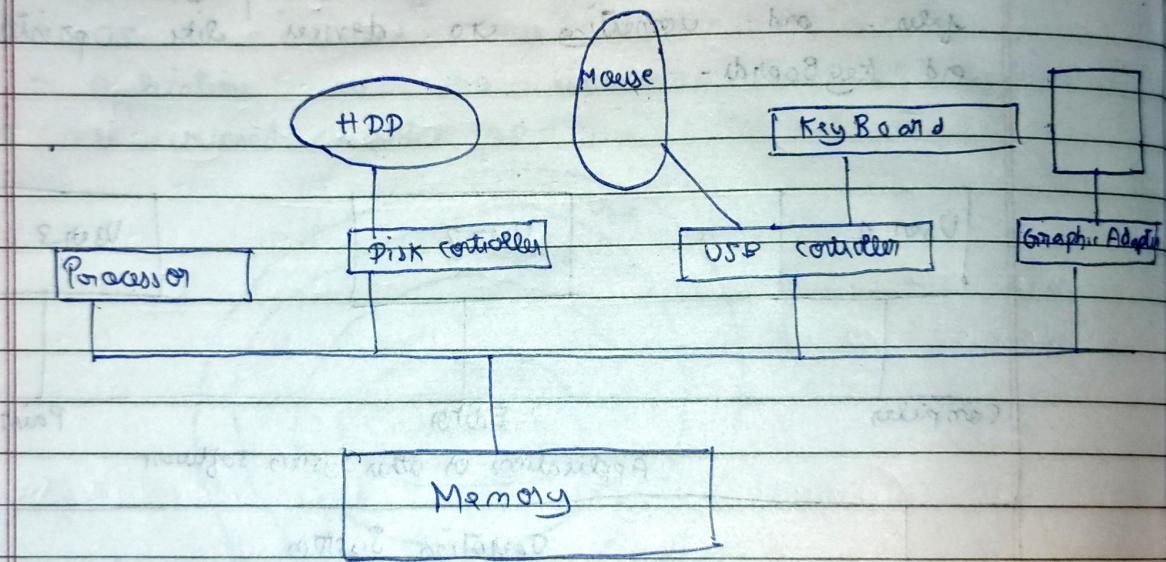
* Kernel in OS

- A Kernel is the core part of an Operating system.
- It acts as a Bridge between software applications and the hardware of a computer.
- Kernel manages system resources, such as the CPU, memory and devices, ensuring everything works together smoothly and efficiently.
- It Handles tasks like running programs, accessing files and connecting to devices like printer and keyboards.



* Basic computer organization required for OS

- Examples of well-known OS
- Mobile OS: Android, iOS, Windows
- Embedded Systems
- Real-time OS: RTOS & SRTOS
- Desktop OS: Personal Computer
- Server machine OS



* Booting is of two types:

- 1) Cold Booting: When the computer is started after being switched off.
- 2) Warm Booting: When the operating system does is restarted after system crash or freeze.

* Real-Time Operating System (RTOS)

- A real-time operating system (RTOS) is a special kind of operating system designed to handle tasks that need to be completed quickly and on time.
- Two types of RTOS
 - 1) HRT
 - 2) SRT

* Server OS

- It is designed to run over Server.
- It is designed to handle millions of requests in a single time.
- A server is a computer that makes data available to other computers. It can serve data across the internet to systems on a LAN or WAN.
- Examples of Server OS are:
 - 1) Apache HTTP Server
 - 2) Microsoft IIS

* Functions of OS

1) Process Management (Process scheduling Algorithms)

- It controls how processes are carried out and controls how your computer runs by handling the active process.
- This includes stopping processes, setting which processes should get more attention and many more.
- Responsible for managing the Start, Stop and scheduling of processes which are programs running on System.

2) Memory Management

3) Device Management (Manages multiple Devices HDD, Monitor, mouse, Printer, Speaker, Webcam).

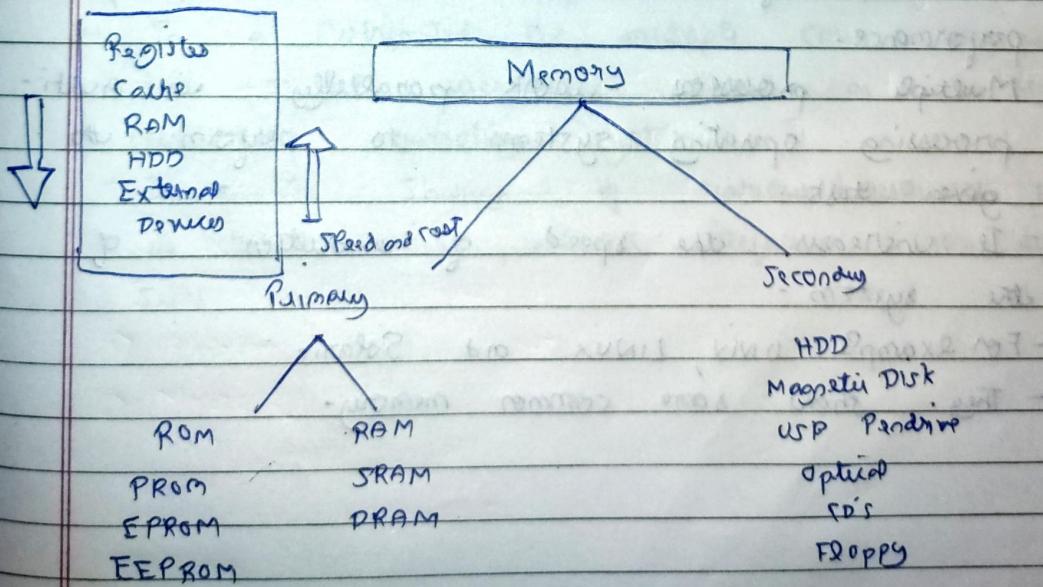
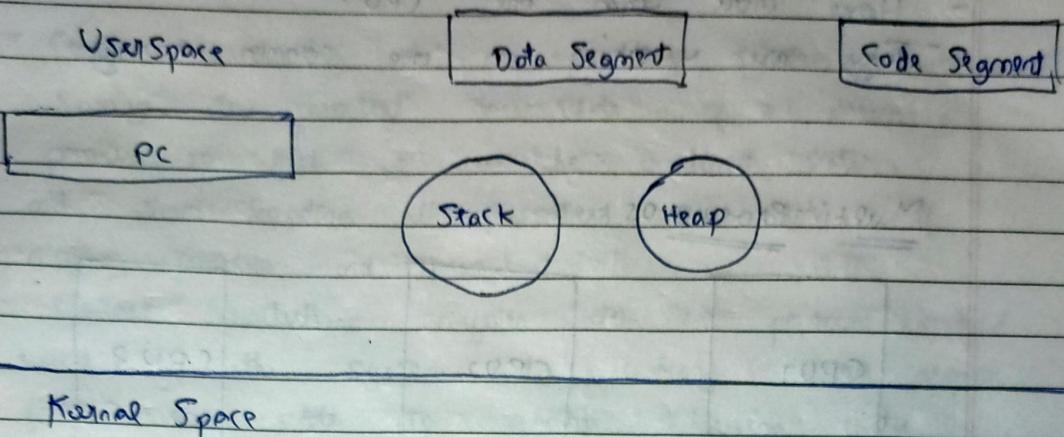
4) Disk Management (Disk scheduling Algorithms)

5) Network Management (Network Card/ controller)

6) File Management

7) Security Management (Firewall, Anti-virus, Anti-Spyware).

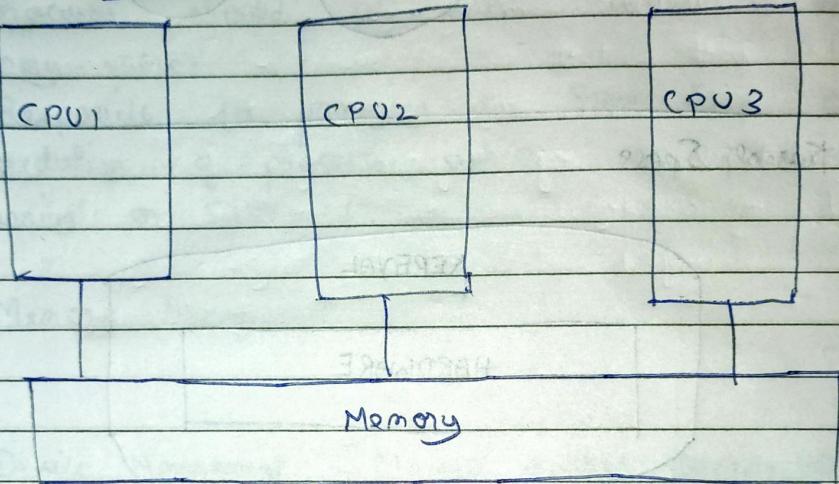
* User Space and Kernel Space



* Multi-processing Operating System.

- Have more than one CPU.
- They are working on one common memory.

Multi-Processors



- A multi-processing operating system is defined as a type of operating system that makes use of more than one CPU to improve performance.
- Multiple processors work parallelly in multi-processing operating systems to perform the given task.
- It increases the speed of execution of the system.
- For example UNIX, LINUX and Solaris.
- They share some common memory.

* Context switching

- Context switching in an operating system involves saving the context or state of a running process so that it can be restored later, and then loading the context or state of another process and run it.
- Context Switching refers to the process method used by the system to change the process from one state to another using the CPUs present in the system to perform its job.

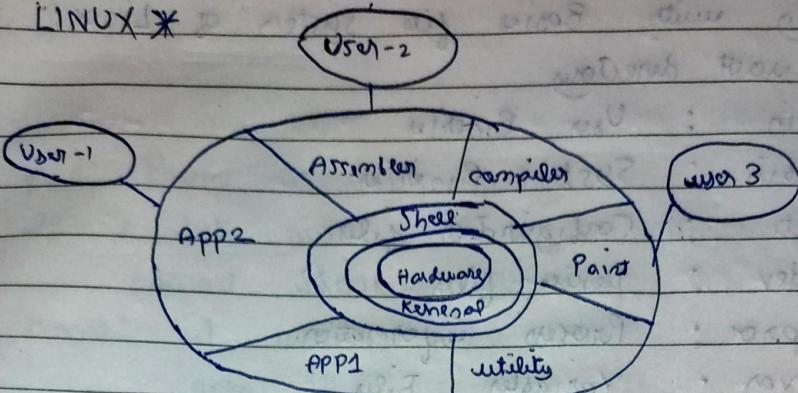
* Distributed Operating System -

- A Distributed OS refers to a model in which applications run on multiple interconnected computers, offering enhanced communication and integration capabilities composed to a network of
- In a Distributed OS, multiple CPUs are utilized but for end user, it appears as a typical centralized operating system.
- It enables sharing of various resources such as CPUs, disks, network interfaces across different sites.

* Batch Operating System

- A Batch operating system (BOS) is a type of computer system that allows multiple users to work on a system at the same time without communicating directly with each other.
- This is done by keeping users in separate batches or groups so that they cannot interact directly.
- This makes it useful for applications where users need to work on different parts of project without interfering with each other.

* LINUX *



Linux operating system

- Linux is an Open-source operating System.
- It is available free to use and user can modify it according to their need.
- The founder of Linux is Linus Torvalds. It is available since 1991.
- An Open source community is working behind the updation and upgradation of Linux code.
- Features of Linux:
 - 1) No cost/ Low cost
 - 2) Multi-Tasking
 - 3) Security
 - 4) Multi-user
 - 5) Stable and Scalable
 - 6) Networking
 - 7) CLI as well as GUI.

- Working with Basic file system of Linux
- / is root directory
- 1) /bin : User Binaries
- 2) /sbin : System Binaries
- 3) /etc : Configuration files
- 4) /dev : Device files
- 5) /proc : Process information
- 6) /var : Variables Files
- 7) /tmp : temporary files
- 8) /usr : User Programs
- 9) /home : Parent directory of user friendly directory
- 10) /boot : Boot loader Files
- 11) /opt : Apps
- 12) /lib : System Libraries

* Ubuntu for Windows

Present

- 1) \$ pwd → Prints working directory
/home/malkeet
- 2) \$ nano abc.txt
\$ ls
- 3) ls : It lists out all the files and directory of current working directory

\$ touch file1.txt
\$ ls

\$ touch file2.txt
\$ ls

\$ mkdir

- 1) pwd: Present working directory
- 2) ls: It lists out all the files and directories of current working directory
- 3) nano: it actually runs the nano editor and open the specified file
- 4) touch: It is used to create a new file
- 5) mkdir: It is used to create new directory
- 6) chmod: to give and revoke file and directory permission
- 7) rm: Used to remove a file

\$ ls -l



Give details about files in directory -

\$ whoami

Tells the currently working user.

owner	Group	Other	Read - R Write - W Execute - E
-	-	-	

u - owner

g - group

o - other

a - all users

chmod u+x file2.txt

\$ chmod

↓

It is used to give read, write and execute permission

\$ chmod a+wx file2.txt

ls -l

\$ chmod o-wx file2.txt

\$ chmod a-x file2.txt

8) cd :- This command is used to change the directory.

9) cd.. :- Goes one step in directory Back.

10) rmdir : Used to remove the directory

11) rm : to remove file and recursive directory

12)

rm -r

==

txt, jpg, xhtml, html

remove -r

quar -e

auto -a

\$ sudo adduser user2

- This command is used to create a user.
- sudo is known as root privileges.

\$ sudo chown user2 file1.txt

- The chown command in Linux is used to change the owner of a file/directory or link.
- Chown is short for "change owner".

cd /etc/
↓
ls

sudo nano sudoers (with root priv.)

* Display content of the file to the screen.

\$ nano file.txt

↓

\$ cat file.txt

↓

\$ head -5 file.txt (Prints first 5 lines of file)

↓

\$ tail -5 file.txt (Prints last 5 lines of file)

\$ head file.txt (Prints first 10 lines)

\$ tail file.txt (Prints last 10 lines)

*> chown: To change the owner of file

2) su: To switch the user upon current to any specified user.

3) cat: To display content of file on console.

4) head: To display top n lines of the file on console. By default it will print first 10 lines.

5) tail: To display bottom n lines of the file on console.

By default it will print last 10 lines.

6) sudo: To give some specific privilege to the user other than root.

\$ cat file.txt

\$ cat > file3.txt

hello

hi

Bye

You

\$ cat file3.txt

\$ head -5 file3.txt

* Redirection (>)

- Output of one command is used as input

\$ tail -5 file4.txt > file5.txt

* System call

- A system call is a way in which a computer program requests a service from the Kernel of the operating system it is executed on.
- A System call is a way for programs to interact with the Operating System.
- A computer program makes a system call when it requests the Operating System's Kernel.
- It is a mechanism used by programs to request services from the Operating System (OS).

* System call Types

File related calls have been omitted -

1. Read(), Write(), Delete(), Open(), Close(), Create()
2. Process related calls: New(), Fork(), Exit(), Wait(), Running(), etc.
3. Device related calls: Read(), ioctl(), etc.
4. Information related: getpid(), gettimeofday(), sysinfo(), etc.
5. Communication related: wait(), Signal(), status(), etc.

* Process

- In OS, a process is a sequence of instructions that are executed in a specific order.
 - A process is essentially a running program that represents the fundamental unit of work that must be implemented in a system.
 - It is information or code or data which help the processor to execute or complete the user task.
1. Code Segment: It consist of compiled code or instructions to be executed.
 2. Data Segment: It consist of data required for the execution.
 3. Information Segment: It have metaData about the system variables or system which helps in process execution.

* Process Control Block (PCB)

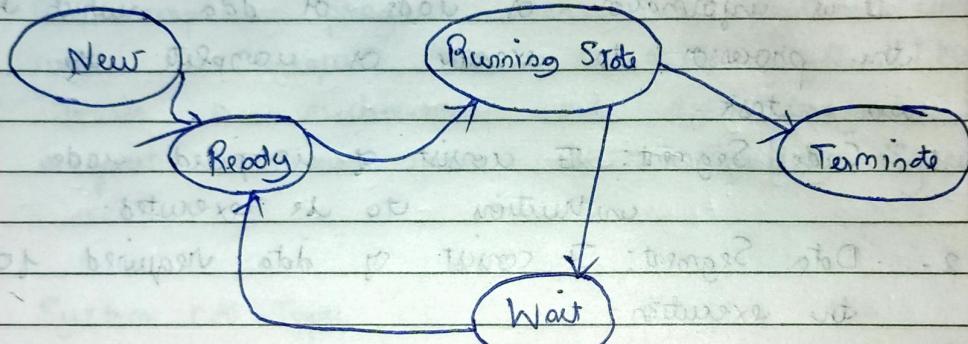
PID	1001
Pointer	0x1000
State	Ready
Allocated Registers	A, C, D
Program counter (PC)	0xdfff0
Allocated Hardware	KB, MS, PR
Accounting information	file.txt, ABC.txt
Priority	10
	:

PCB P1

- A process control block (PCB) is a data structure used by the operating system to manage information about a process. The process control keeps track of many important pieces of information needed to manage processes efficiently.

* Process Life Cycle

1. New State (Creation)
2. Ready State (Ready Queue)
3. Running State - Process is executing the process.
4. Terminated State
5. Wait State



Process Life Cycle

* Scheduler

- A scheduler is a software that helps schedule the processes in an operating system.
- It helps to keep all computer resources busy and allows multiple users to share system resources effectively.
- Types of Schedulers in OS:
 - 1) Short-term scheduler
 - 2) Long-term scheduler (LTS)
 - 3) Medium-term scheduler

* Scheduling Algorithm

- 1) Pre-Emptive Scheduling
- 2) Non-Pre-emptive scheduling

Pre-emptive Scheduling

- In this preemptive scheduling, a higher priority task takes precedence over a lower priority tasks, even if the lower priority task is already running.
- For example, if a high priority process enters the queue and is allocated to the CPU, the current process switches from the running queue to ready queue. The lower priority task waits for some time and continues its executing process until the higher priority job gets completed.

Non-preemptive Scheduling

- Non-preemptive scheduling is a CPU scheduling method where a process keeps control of a resource until it ends or becomes a waiting state.
- In this method, new processes must wait for the current process to finish its CPU cycle before they can begin.

D) FCFS (First come First Serve)

- Waiting time of process = CPU Allocation - Arrival time
- Avg Waiting Time = Sum WT of All process / no. of processes.

E) SJF (Shortest Job first) :- Shortest process first with smallest Burst time

F) Round Robin

Round Robin is a CPU scheduling algorithm where each process is cyclically assigned a fixed time slot.

- It is the preemptive version of the First

come First serve CPU Scheduling Algorithm

- Giving CPU to each process everytime named as

Gianton time slot.

Time slot during which CPU is given to process.

* Memory Management:

- Fixed size partition
- Variable size partition
- The memory of computer system can be divided into block of fixed size or variable size
- This process is known as Fixed size partitioning and variable size partitioning respectively.
- Fixed size partitioning: Here the memory is divided into fixed size of Block like 2 Bytes, 4 Bytes etc.
- The process which will get the memory must be more than the partitioning size -

* Internal Fragmentation

- Internal fragmentation is a phenomenon that occurs when a memory block is allocated to a process that is smaller than the amount of memory requested.
- This results in unused space within the memory block, which is then considered internal fragmentation.

* Variable size partitioning

- Here the memory is divided into variable size blocks in which blocks may 2 Bytes, 3 Bytes or any size.
- In variable size partitioning during contiguous memory allocation some time total memory available can not give to the new process which leads to extend fragmentation.

* Compaction

- Compaction is a technique to collect all the free memory present in the form of fragments into one large memory, which can be used to run other processes.
- It does not by moving all the processes towards one end of the memory and move the available free space towards the other end of the memory so that it becomes contiguous.

Before compaction

Main Memory

Occupied Space
Free Space
Occupied Space
Occupied Space
Free Space

After compaction

Main Memory

Occupied Space
Occupied Space
Occupied Space
Free Space
Free Space

* Process scheduling Revision

- Process creation using fork
- Waitpid and exec system calls (fork-exe-wait)
- Examples on process creation
- Parent and child processes
- Orphan and Zombie processes

* Process creation using fork

- fork(): It is method / system call which is used to create child processes of the current process executed
- What a child process does: It does exactly the same for which you have designed the parent process

\$ nano program.c

\$ gcc -o Program Program.c

\$./Program

\$

#include < stdio.h >

int main()

{

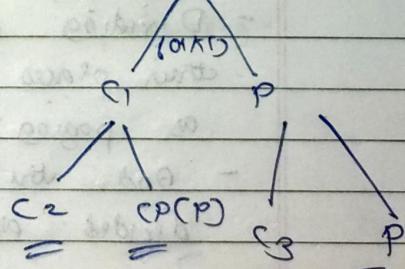
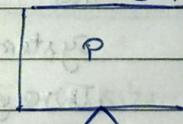
fork();

fork();

priority("Hello ~~with~~ Admin\n");

return 0;

No of parent process $\leftarrow 2 \times n = 1, n=2$
 No of child process $\leftarrow 2 \times n - 1 = 3, n=2$



Fourth time Hello All
 will be printed

* Orphan and Zombie process in OS

- A zombie process is a process that has finished execution but still has an entry in the process table.
- An orphan process is a process that is still running but its parent process has finished or terminated.

* Paging

- The process of retrieving process in the form of pages from the secondary storage into the main memory is known as Paging.
- The basic purpose of paging is to separate each procedure into pages.
- Additionally frames will be used to split the main memory.
- Instead of loading entire program in main memory, just load the needed one.
- This will increase the efficiency of your system and it will increase the throughput.
- Dividing process into number of pages and then placing which is on Hardrive is known as paging.
- And the sections to which they are divided are known as Pages.

Page 1	1B
Page 2	2B

PI of size 3B come

- The program is divided into pages and memory is divided into frames -

- * To support paging the size of page must be equal to the size of frames -
- * All the pages must be divided into equal size -

CPU --->

Byte 3

Logical Address

P1 is of 48

Page 0

0	1
2B	

Page 1

2	3
2B	

Frame No

0

0

1

1

2

3

2

—

—

3

—

—

Physical

Address

4

8

9

5

—

—

6

—

—

7

—

—

8

—

—

Page Table of P1

MMU \rightarrow Memory Management Unit

Page 0	Frame 1
Page 1	Frame 2

Fixed size
partitioning

* Page replacement Algorithms

1) LRU

2) MRU

- If the page is not in the frame, it's known as miss of the page and it's a page fault.

<u>P1 →</u>	2	3	2	4	3	5	3	7	*
Frame 1	8	8	8	5	4	5	5	4	7
Frame 2	3	3	3	3	3	3	3	3	3
Frame 3	2	2	2	2	2	2	2	2	2
Frame 4	1	1	1	6	6	6	5	5	7
init	HIT	HIT	HIT	HIT	MISS	HIT	MISS	HIT	MISS

* Paging

- Dividing the process into fixed size pages is known as paging.
- Why paging: Instead of loading the whole process which can't be loaded into main memory, it loads few pages of the process into main memory according to frames available and other pages loaded from storage device on demand by CPU.

* Paging Table

- It is a special table maintained by MMU (Memory management unit) to map the logical address of the demanded page with the physical address of the page in main memory where actually the page is placed.

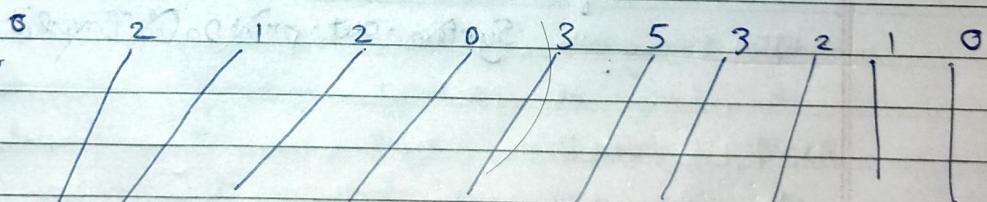
* Demand paging

- If the page is loaded in memory from OS per the demand of CPU i.e known as demand paging.

* Page Fault:

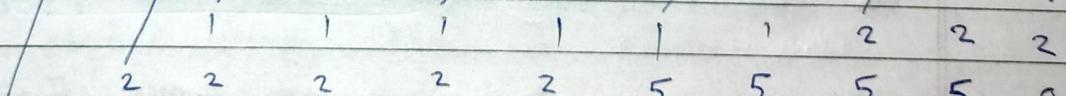
- If CPU demands a page and it's not there in main memory.
- That time an interrupt generated to load the page from virtual memory.
- Page

P1 →

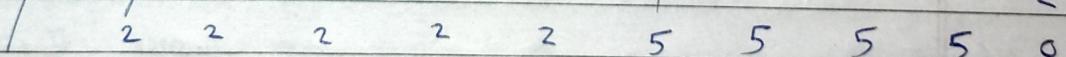


FrameSet

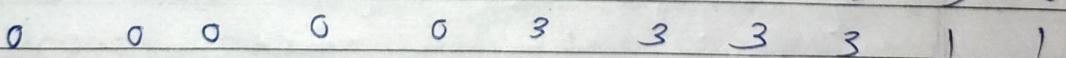
Frame 3



Frame 2



Frame 1



MISS MISS MISS HIT HIT MISS MISS HIT MISS MISS MISS

Hit Ratio = 3/11

Miss Ratio = 8/11

Paged WA

	0	2	1	2	0	3	5	3	2	1	0
Frame 4						3	3	3	3	3	3
Frame 3						1	1	1	1	1	1
Frame 2			2	2	2	2	2	2	2	2	0
Frame 1	0	0	0	0	0	0	5	5	5	5	5
	MISS	MISS	MISS	HIT	HIT	MISS	MISS	HIT	HIT	HIT	MISS

* Belady's Anomaly:

- Belady's anomaly is the phenomenon in which increasing the number of page frames results in an increase in the number of page faults for certain memory access patterns -

* Virtual Memory

- Virtual Memory is a memory management technique used by Operating Systems to give the appearance of a large, continuous block of memory to applications, even if the physical memory (RAM) is limited.
- It allows the system to compensate for physical memory shortages, enabling larger applications to run on Systems with less RAM.
- It is Space in HardDrive

- Hardware required for paging (virtual + cache)
- Translation look aside buffer
 - It is a page table represented in Cache memory to fasten the process of finding physical address for a given logical address
- In paging we load only few pages in main memory of a particular process.

* Virtual Memory

- It is a memory space in Hard-drive, which works like physical memory to certain large processes whose size is bigger than RAM.
- In virtual memory process is divided into fixed sized partitions known as pages.
- OS loads the process page from virtual memory to physical memory on demand of CPU i.e. Known as Demand paging.
- While working with virtual memory pages are saved as per their logical address.
- While reading the pages from virtual memory to physical memory the pages get physical address.
- If a page loaded from Virtual memory to physical memory i.e. Known as Swap-in process.
- If the page is shifted from physical memory to logical memory by the replacement process that is Known as Swap out memory.

* What is Shell?

- Shell is an interface between user and the Kernel.
- The user can interact with Kernel by using shell commands or shell script/ program.
- It takes input from user and passes it to the Kernel.

* What are different Shells in Linux?

cd / etc

name p1

#!/bin/bash

etc \$ ls

echo "Enter your name"

read str1

cat/etc/shells

echo Welcome, \$str1

x = 100

*

Echo " Enter num1"

read num1

if [\$num1 -eq 5]

then

Echo " Number is equal to 5"

else

Echo "Number is not equal to 5"

fi

O/P:= 100

* echo Enter Num 1

read Num 1

echo Enter Num 2

read Num 2

if [\$Num1 -gt \$Num2]

then

echo Num 1 is greatest

else

echo Num 2 is greatest

fi

* echo Enter Num 1

read Num 1

echo Enter Num 2

read Num 2

echo Enter Num 3

read Num 3

if [\$Num1 -gt \$Num2]

if [\$Num1 -gt \$Num3]

then

echo Num 1 is greatest

else

echo Num 3 is greatest

fi

else

if [\$Num2 -gt Num3]

then

echo Num 2 is greatest

else

echo Num 3 is greatest

fi

fi

* While loop in Shell programming

Syntax:

`while [condition]`

`do`

`statement`

`increment`

```
#!/bin/bash
```

`a=0`

`while [$a -lt 10]`

`do`

`echo $a`

`a=`expr $a + 1``

`done`

O/P:

0

1

2

3

4

5

6

7

8

9

10

* `cat file1 | grep lh.n`

`cat file1 | grep lh.*n`

- Piping allows us to cat and filter the content

`cat file2 | grep ^h`

Piping allows us to filtration among the result of the command.

Eg `cat file2 | grep ^h`

* Redirection (>)

It allows you to use output of one command as an input of other command.

* Tilde (~)

Used for going back to the user's home directory

* Dollar Symbol (\$)

It is also used to matching the end of the string for a line in a file.

* Caret (^)

It is used to match the starting of the string within given file.

```

echo Number of Param, $#
echo Script Name, $0
echo Hello, $1
echo Hi, $2
echo OK, $3
echo Bye, $*

```

Command line arguments

- Specify parameters in Shell while passing command line arguments

1. \$# : No. of parameters passed
2. \$0 : Script Name
3. \$i : where i can be any number, it will return that parameter
4. \$* : Gives all parameters passed - along with it

* Regular expression

- `grep ^S`: It will filter the pipeline and get the words which are starting with S.
- `grep ^n*f$`: It will filter the pipeline and get the lines which are starting with n and ending with f.
- `ls f*`: It will return all the files starting with f.

* Arithmetic expression -

#!/bin/bash

echo Enter a Number

read Num1

echo Enter a Number

read Num2

RES='expr \$Num1 * \$Num2'

echo Result, \$RES

Basic Arithmetic Operators

1) + Addition

2) - Subtraction

3) * Multiplication

4) / Division

5) % Modulo

6) += increment by constant

7) -= Decrement by constant

* Network Commands

- Telnet : It is used to take the control of remote server.
- SSH : It is also used to take control of remote server / machine -
- FTP : File Transfer protocol
It allows you to send or receive file from remote server / machine -
- SFTP : Secure File Transfer protocol -