Information Theory and Computation Exercise 9

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Theory

In this exercise we are gonna diagonalize the Quantum Ising model with the Hamiltonian:

$$H = \lambda \sum_{i=1}^{N} \sigma_i^z + \sum_{i=1}^{N-1} \sigma_i^x \sigma_{i+1}^x$$
 (1)

As a basis for the one particle Hilbert space we chose the one diagonalizing σ^z :

$$\sigma^z |0\rangle = |0\rangle \tag{2}$$

$$\sigma^z |1\rangle = -|1\rangle \tag{3}$$

An N-particles wave function is defined by a family of $\{n_i\}n_i = 0, 1$ i.e.

$$|n_1\rangle |n_2\rangle \dots |n_M\rangle$$
 (4)

We can observe that there is a one to one correspondence between a wavefunction and an integer number form 0 to $2^N - 1$; from the coefficients above we obtain the number by computing:

$$n_1 + n_2 \times 2 + n_3 \times 2^2 + \dots + n_N \times 2^{N-1}$$
 (5)

In other words, we treat the sequence $n_N n_{N-1} \dots n_1$ as a binary number. Since we chose the diagonal basis the action of $\sum_{i=1}^N \sigma_i^z$ on a N particles state is trivial; instead $\sigma_i^x \sigma_{i+1}^x$ as as:

$$\sigma_i^x \sigma_{i+1}^x |n_1\rangle |n_2\rangle \dots |n_M\rangle = |n_1\rangle |n_2\rangle \dots |1 - n_i\rangle |1 - n_{i+1}\rangle \dots |n_M\rangle$$
 (6)

i.e. flips the *i*-th and the i+1-th spin. Denoting $|m\rangle = |m_1 m_2 \dots m_N\rangle$ and $|n\rangle = |n_1 n_2 \dots n_N\rangle$ with m and n constructed as above said, the matrix element $\langle m | \sigma_i^x \sigma_{i+1}^x | n \rangle$ is non zero if and only if the numbers m and n, in binary, have the same digits apart from the *i*-th and i+1-th, where they must be opposite; one can express this property by the following predicate:

$$\langle m | \sigma_i^x \sigma_{i+1}^x | n \rangle \neq 0 \leftrightarrow XOR(n, 3 \times 2^{i-1}) == m$$
 (7)

The XOR operation is done logically, bit-by-bit. The above property can be used to compute the Hamiltonian matrix elements:

$$H_{m,n} = \langle m | H | n \rangle = \lambda \delta_{m,n} + \sum_{i=1}^{N-1} \delta_{m,XOR(n,3 \times 2^{i-1})}$$
 (8)

The number $3 \times 2^{i-1}$ is a *mask* number, useful to compare the bits of each index m and n: it is equal to the binary 11_2 shifted towards the left in order to force the i-th nd i+1-th bits to be opposite (an the others to be equal). Practically, given some m, look for all the n's that have non zero $\delta_{m,XOR(n,3\times 2^{i-1})}$, for $i=1\ldots N-1$.

Code Development

Essentially we need two function to build the Hamiltonian; the first one is the function computing the magnetization for a given state (i.e. the first part of the Hamiltonian); basically, for a state indexed by the number n, the magnetization is given by the number of zeros minus the number of ones in the binary representation:

The second function we need is the one that actually computes the Hamiltonian matrix elements:

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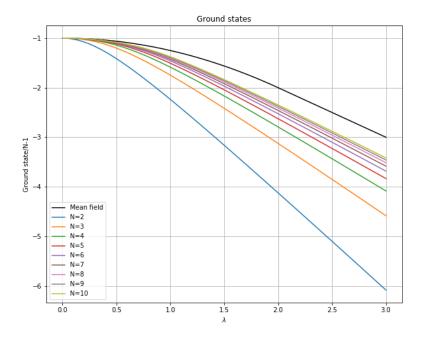
end do end function

The remaining part of the program had the job of diagonalizing the obtained matrix (via *zheev*) and saving the computed energies.

Results

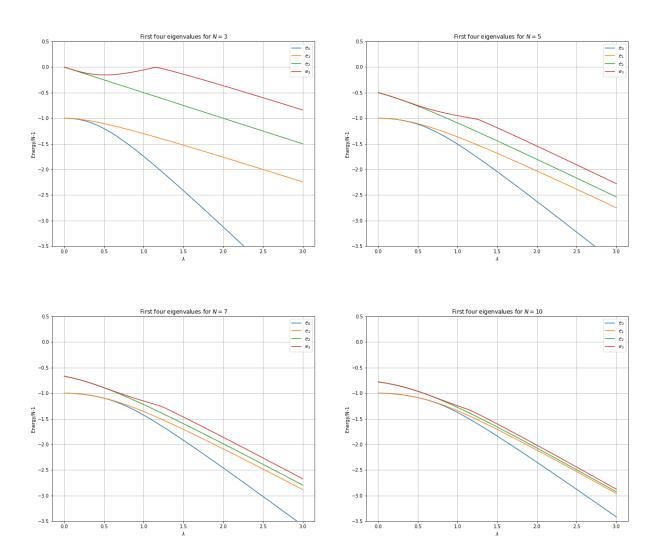
We run the code for a number of spins going from 2 to 10. Below we plot the ground state (divided by N-1) of each system and the mean field energy density which has the expression:

$$e_0(\lambda) = \begin{cases} -1 - \frac{\lambda^2}{4} & 0 < \lambda < 2\\ -\lambda & \lambda > 2 \end{cases} \tag{9}$$



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The computed ground states seem to look, as N highers, more similar to the mean field solution. Below, instead, we show, for some N's, the first four eigenvalues:



We can notice that as N highers the ground state is more and more distinct form the other energy levels which tend to come closer, especially at higher λ 's. Another related interesting phenomenon that we observe is the degeneracy removal, due to leaving $\lambda=0$.

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