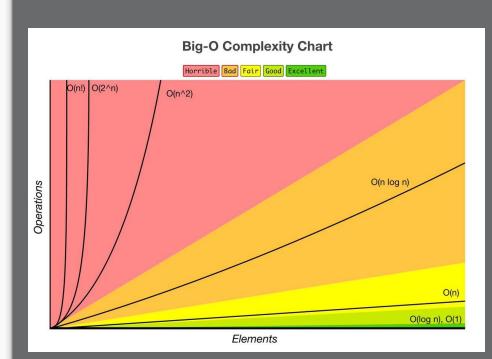
Algorithmic Analysis (Big-O) Sorting

Frank Burkholder
Taryn Heilman
Joe Gartner





Learning Objectives



- Understand Big-O notation and how to use it to describe algorithms
- Be ready to write implementations of
 - Insertion Sort
 - Bubble Sort
 - Heap Sort
- Optional: Python Review (or maybe New!) Define a decorator

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Algorithmic Time Complexity

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"Big-O" Notation



- <u>Big O Notation</u> Used to describe how the runtime (time complexity) and size (space complexity) of an algorithm increases as the size of the input array of **length** *n* increases. We'll be focusing on time complexity.
- This is an <u>order of magnitude</u> approximation, meaning we only worry about the leading term. In big-O, a process that requires *n* computations is the same as one that requires *3*n* computations, both are *O(n)*.
- Expressed as a function of n (e.g, C^n , n^3 , n^2 , $n \log n$, n, $\log n$)

Order of magnitude approximation examples



$\operatorname{Number} N$	Expression in $N=a imes 10^b$	Order of magnitude \boldsymbol{b}
0.2	2×10^{-1}	-1
1	1×10^{0}	0
5	0.5×10^{1}	1
32	0.32×10^2	2
999	0.999×10^3	3

Person (L~10⁰ m)

Source: Wikipedia



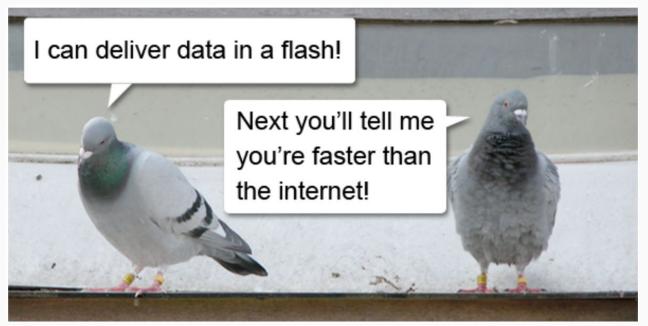
"Big-O" Notation - why?

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This topic often shows up in software interviews, as it is a good test of several things:

- Estimating the growth in time for an algorithm as the input size grows
- Evaluate your knowledge of data structures
- They test your ability to think about what you are doing





Source another example

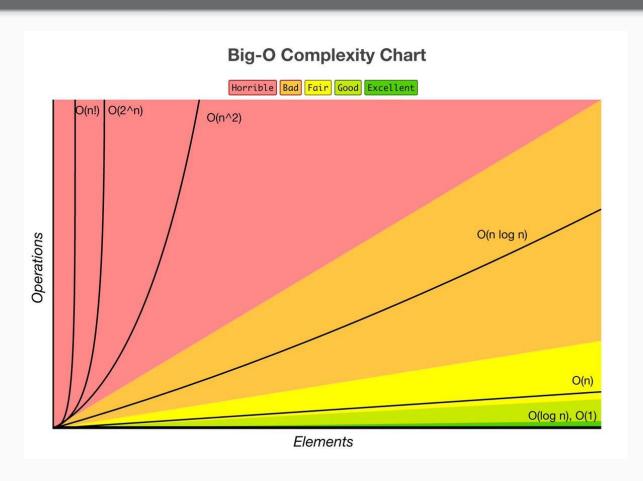
O(N) of common routines (more complete on Wikipedia)



Time Complexity (best to worst)	Name	Description - Given an input of size n	Example operation or algorithm
O(1)	Constant	Only a single step required to complete the task.	Lookup in a set. Appending to a list.
O(log n)	Logarithmic	The number of steps it takes to accomplish the task are decreased by some factor with each step.	Binary search
O(n)	Linear	The number of of steps required is directly related to <i>n</i> .	Lookup in a list. Kmeans algorithm.
O(n log n)	Log-linear	The number of of steps required is directly related to <i>n</i> multiplied by some factor that is a factor of <i>n</i> (but much less than <i>n</i>).	Merge sort
O(n²)	Quadratic	The number of steps it takes to accomplish a task is square of <i>n</i> (bad).	Double for loop. Create covariance matrix. Hierarchical clustering.
O(C^n)	Exponential	The number of steps it takes to accomplish a task is a constant to the <i>n</i> power (very bad).	Graph colouring.

O(n) rough comparison of number of operations





```
def check_list_elements(x, lst):
    for i in lst:
        if element==x:
            return x
```

Assume there are *n* items in lst.

How many comparisons does this function make?

What is the run time complexity?

What is a way to reduce the complexity?

```
def check_list_elements(x, lst):
    for i in lst:
        if element==x:
            return x
```

Assume there are *n* items in lst.

How many comparisons does this function make? Each item 1st with x.

What is the run time complexity? 1st is n elements, so n, or O(n)

What is a way to reduce the complexity? Maybe use a set instead of a list?

```
def print_pairs(lst):
    for x in lst:
        for y in lst:
            print(x,y)
```

Assume there are *n* items in lst. How many pairs does this function make? What is the run time complexity?

```
def print_pairs(lst):
    for x in lst:
        for y in lst:
            print(x,y)
```

Assume there are *n* items in lst.

How many pairs does this function make? Each item in 1st with every other item in 1st. What is the run time complexity? 1st is n elements, so n * n, or $O(n^2)$

```
def print_pairs_two_lists(lst1, lst2):
    for a in lst1:
        for b in lst2:
            print(a, b)
```

Assume there are *n* items in lst1 and m items in lst2 How many pairs does this function make? What is the run time complexity?

```
def print_pairs_two_lists(lst1, lst2):
    for a in lst1:
        for b in lst2:
            print(a, b)
```

Assume there are n items in lst1 and m items in lst2 How many pairs does this function make? Each item in lst1 with each item in lst2. What is the run time complexity? so n * m, or O(nm)

What is the O(n) complexity of this algorithm?



```
def max min(lst):
    current max = 0
    current min= 0
    for i in 1st:
        current max = max(current max, i)
    for i in 1st:
        current min = min(current min, i)
    print(current max, current min)
def max min2(lst):
    current max = 0
    current min= 0
    for i in 1st:
        current max = max(current max, i)
        current min = min(current min, i)
    print(current max, current min)
```

Assume there are *n* items in lst.

What is the run time complexity for each function?

```
def max min(lst):
    current max = 0
    current min= 0
    for i in 1st:
        current max = max(current max, i)
    for i in 1st:
        current min = min(current min, i)
    print(current max, current min)
def max min2(lst):
    current max = 0
    current min= 0
    for i in 1st:
        current max = max(current max, i)
        current min = min(current min, i)
    print(current max, current min)
```

Assume there are *n* items in lst.

What is the run time complexity for each function? O(n) for both



```
def find anagrams (lst):
    result = []
    for word1 in 1st:
        for word2 in 1st:
            if word1 != word2 and sorted (word1) == sorted (word2):
                 if word1 not in result:
                     result .append (word1)
                 if word2 not in result:
                     result .append (word2)
     return result
```

Assume there are *n* words in 1st.

How many comparisons does this function make? Each word in 1st with each word in 1st. What is the run time complexity? 1st is n elements, so n * n, or $O(n^2)$

```
def find_anagrams (lst):
    result = []
    d = defaultdict(list)
    for word in lst:
        d[tuple(sorted(word))].append(word)
    for key, value in d.iteritems():
        if len(value) > 1:
            result.extend(value)
    return result
```

How many comparisons does this function make?

What is the run time complexity? *O(n)*



```
def find_anagrams (lst):
    result = []
    d = defaultdict(list)
    for word in lst:
        d[tuple(sorted(word))].append(word)
    for key, value in d.iteritems():
        if len(value) > 1:
            result.extend(value)
    return result
```

How many comparisons does this function make? **Doesn't make any. Checks for** more than one word for a sorted character key in d, which is much smaller than length *n*. The lst is iterated through only once.

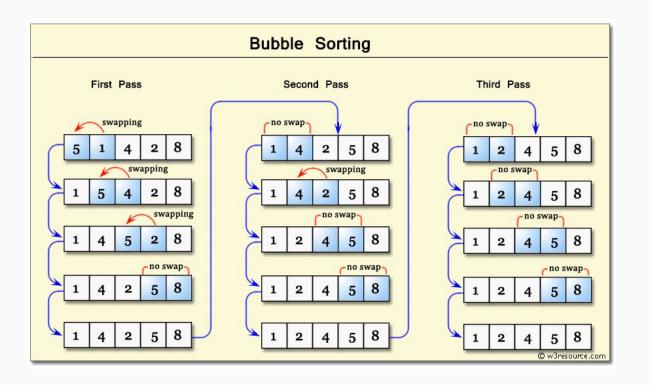
What is the run time complexity? O(n)

find_anagrams.py

Sorting Algorithms

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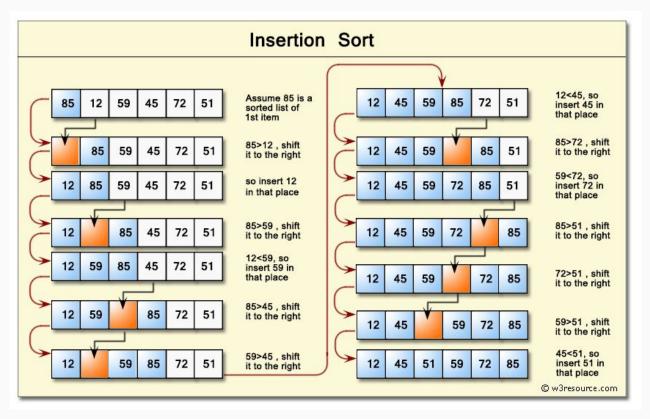
Wikipedia visualization



What is the run time complexity?

Insertion Sort





Wikipedia visualization.

```
# psuedocode
def insertionSort(lst):
   for i in range(1, len(lst)):
    val = lst[i]
     while i > 0 and lst[i - 1] > val:
         move lst[i - 1] to the right
         decrement i
     assign val to lst[i]
```

What is the run time complexity?

Pair up and sort a shuffled suit from a deck of cards.

Part I:

Try <u>Bubble Sort</u> and <u>Insertion Sort</u> something at home (5-6 playing cards or maybe some coins). Can you explain and demonstrate each sorting algorithm to your partner?

Part II:

There are two text files in the lecture repo:
unsorted num 1.txt and unsorted num 2.txt

Write a Python script that you run from the command line, using argument parsing, to sort the numbers in the text file. Use Bubble Sort.

You should be able to run your script like this:

\$ python bubblesort.py --in unsorted num 1.txt --out sorted num 1.txt

There are better sorting algorithms!



- Heap Sort
- Merge Sort
- Quick Sort
- All of the above have complexity O(n log n)

Read more at https://wikipedia.org/wiki/Sorting_algorithm

Learning Objectives



High level takeaways:

- Why data structures matter (do you need to use a list? maybe tuple, dict, set is better?)
- When referring to time it takes to run code... our favorite answer, it depends!!!! (size of data, computational power, memory, etc.) so we use Big O notation instead
- One common place to talk about Big O is in sorting algorithms
- Keep in mind 'premature optimization is the root of all evil'

Review: Decorators



```
16 class SortTester(ABC):
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       Abstract base class for our sorting classes
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30
        Allows for the calling of a sort, a stack count sort, and a timed sort
       def init (self):
            pass
        @classmethod +
        def this sort(cls, input list):
                                                      Decorators
            pass
        @timeit ←
        def timed sort(self, input list):
31
            pass
32
33
```

Learning Objectives



- Understand Big-O notation and how to use it to describe algorithms
- Be ready to write implementations of
 - Insertion Sort
 - Bubble Sort
 - Heap Sort
- Python Review (New!) Define a decorator