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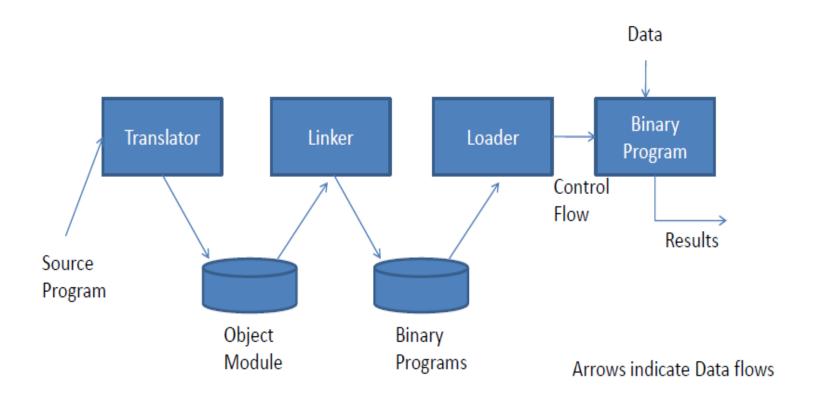
Linkers



Linker

- Steps for Execution:
- 1. Translation of program
- 2. **Linking** of program with other programs needed for its execution.
- 3. **Relocation** of the program to execute from the specific memory area allocated to it.
- 4. Loading of the program in the memory for the purpose of execution.







Linkers

- Links with other programs needed for its execution
- Processes a set of object modules to produce a ready-to-execute program form called binary program
- Loader loads this program into memory for execution



Linker

 A Linker is a system program that combines the code of a target program with codes of other programs and library routines



Object Module and Binary Program

Object Module:

 Contains target code (Machine language) of the program and information about other programs and library routines that it needs to invoke during its execution

Binary Program:

 Target code and other program routines combined together to form Binary program



Addresses

Address of a program entity may vary at different times:

- Translation time address: Address assigned by translator (ORIGIN or START)
- 2. Linked address: Address assigned by linker
- Load time address: Address assigned by loader



Addresses

Origin of a program may have to be changed by linker or loader because

- Same set of translated addresses may have been used in different object modules, resulting in memory conflicts
- 2. OS MM module may require that a program should be allocated specific area of memory



Example:

	Stateme	nt	Address	Code					
	START	500							
	ENTRY	TOTAL	The trans	lated					
	EXTRN	MAX, ALPHA	origin						
	READ	А	500)	+ 09 0 540					
LOOP			501)						
				ranslated Time					
			Add	ress of LOOP					
	MOVER	AREG, ALPHA	518)	+ 04 1 000					
	BC	ANY, MAX	519)	+ 06 6 000					
	BC	LT, LOOP	538)	+06 1 501					
	STOP		539)	+00 0 000					
Α	DS	1	540)						
TOTAL	DS	1	541)						
	END								



Relocation concept

- Program relocation is the process of modifying the addresses used in the address sensitive instructions of a program such that the program can execute correctly from the designated area of memory
- If linked origin ≠ translated origin, relocation must be performed by the linker
- If load origin ≠ linked origin, relocation must be performed by the loader

Sample Assembly Language Program and generated code (Program P)

	START	500						
	ENTRY	TOTAL						
	EXTERN	MAX,ALPHA						
	READ	Α	500 +400))		09	00	540+400=940
LOOP			501+400	90:	1			
	MOVER	AREG, ALPHA	518+400	918	}	04	01	000
	BC	ANY, MAX	519+400	919)	07	06	000
	BC	LT, LOOP	538)	07	01	501		
	STOP		539)	00	00	000		
Α	DS	1	540)					
TOTAL	DS	1	541)					
	END							



Performing Relocation

```
t_origin<sub>p</sub> - translated origin of program P

l_origin<sub>p</sub> - linked origin of program P

t<sub>symb</sub> - translation time address of a symbol symb
```

I_{symb} - link time address of a symbol symb

Relocation factor of P is defined as $relocation_factor_p = l_origin_p - t_origin$ ______1



Performing Relocation

```
relocation\_factor_p = l\_origin_p - t\_origin
t_{symb} = t\_origin_p + d_{symb}
t_{symb} = l\_origin_p + d_{symb}
t_{symb} = l\_origin_p + d_{symb}
```

From 1 and 3

$$I_{symb} = t_origin_p + relocation_factor_p + d_{symb} -----4$$

From 2 and 4

$$I_{symb} = t_{symb} + relocation_factor_p$$
5



For Program P:

- Translated origin is 500
- Suppose I_origin = 900
- Relocation_factor = 900- 500 = 400
- Relocation will be performed for instructions with translated time address 500 and 538
 - For instruction with translated time address 500:address 540 in operand field will change to (540+400)= 940
 - For instruction with translated time address 538:address 501 in operand field will change to (501+400)= 901



Linking

- Program interacts with another program unit using its instructions & data in its own instructions
- Public definitions & external references required
 - ENTRY: Public definitions
 - A symbol defined in a program unit that may be referenced in other program unit.
 - EXTRN: External references
 - A reference to a symbol that is not defined in the program unit containing the reference (defined in other program)
- Linking is the process of binding an external reference to the correct link time address



Program Q

START 200

ENTRY ALPHA

__

ALPHA DS 25 231) 00 00 025

END



Linkers

Object module contains all info necessary to relocate & link program with other programs

- 1. Header: has translated origin, size & execution start address
- 2. Program: has machine code
- Relocation table (RELOCTAB): each entry contains translated address of an address sensitive instruction
- 4. Linking table (LINKTAB): contains PD/EXT symbols

Linker generates the linked addresses of all the symbols & instructions



RELOCTAB and LINKTAB

- Program P:
- RELOCTAB:
 - **-500)** 09 00 540
 - -538) 07 01 501
- LINKTAB:
 - TOTAL PD
 - MAX EXT
 - ALPHA EXT



RELOCTAB and LINKTAB

- Program Q:
- LINKTAB:
 - ALPHA PD

NTAB:

Symbol	Linked addr
Р	900
TOTAL	941
Q	942
ALPHA	973



Linking of Program P and Q

Program P:

-0 -	START 500					
	ENTRY	TOTAL				
	EXTERN	MAX,ALPHA				
	READ	A	900)	09	00	940
	LOOP		901)			
	MOVER	AREG, ALPHA	918)	04	01	973
	BC	ANY, MAX	919)	07	06	000
	-					
	ВС	LT, LOOP	938)	07	01	901
	STOP		939)	00	00	000
	A DS	1	940)			
	TOTAL DS	51	941)			
	END					
Program Q	:					
C	START	200				
	ENTRY	ALPHA				
			942)			
	ALPHA DS	5 25	973)	00	00 (025
	END	-	/	2.3		



Self-Relocating Programs

Programs can be classified into

- 1. Non relocatable programs
- cannot be executed in any memory area other than its translated origin
- -due to lack of information pertaining to address sensitive instructions in program
- 2. Relocatable programs
- Has info available related to address sensitive instructions in program
- 3. Self-relocating programs



Self-Relocating Programs

- Performs relocation of its own address sensitive instructions
- 2 provisions for this
- 1. Table containing address sensitive instructions exists as part of program
- 2. Relocating logic: Code to perform relocation of address sensitive instructions also exists as part of program
- Can execute in any area of memory
- Useful in time sharing operating systems



Static Link Libraries

- Is the process of copying all library modules used in the program into the final executable image.
- This is performed by the linker and it is done as the last step of the compilation process.
- The linker combines library routines with the program code in order to resolve external references, and to generate an executable image suitable for loading into memory.
- When the program is loaded, the operating system places into memory a single file that contains the executable code and data.
- This statically linked file includes both the calling program and the called program.
- Statically linked files are significantly larger in size because external programs are built into the executable files.



Static Link Libraries

- If any of the external programs change then they have to be recompiled and re-linked again else the changes won't reflect in existing executable file.
- Statically linked program takes constant load time every time it is loaded into the memory for execution
- Programs that use statically-linked libraries are usually faster
- In statically-linked programs, all code is contained in a single executable module. Therefore, they never run into compatibility issues.



Dynamic Link Libraries

- In dynamic linking the names of the external libraries (shared libraries) are placed in the final executable file while the actual linking takes place at run time when both executable file and libraries are placed in the memory.
- Several programs use a single copy of an executable module.
- Is performed at run time by the operating system
- Only one copy of shared library is kept in memory. This significantly reduces the size of executable programs, thereby saving memory and disk space
- Individual shared modules can be updated and recompiled. This is one of the greatest advantages dynamic linking offers.



Dynamic Link Libraries

- Load time might be reduced if the shared library code is already present in memory.
- Programs that use shared libraries are usually slower than those that use statically-linked libraries.
- Dynamically linked programs are dependent on having a compatible library. If a library is changed (for example, a new release may change a library), applications might have to be reworked to be made compatible with the new version of the library. If a library is removed from the system, programs using that library will no longer work.



MIT-WPU Final Year (B.Tech)

System Software and Compiler Design

5/3/2021



Module II

- Macro processor: Macro Definition, Macro expansion and nested macros
- Loaders: Loader schemes: Types of loaders, direct linking loaders.
- Linkers: Relocation and linking concepts, self-relocating programs, Static and dynamic link libraries.

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Loader

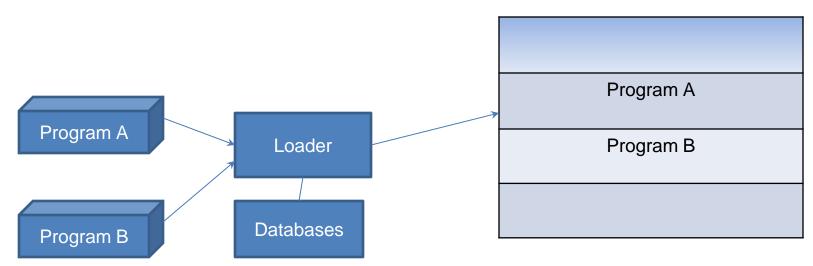


Loaders

Loader is a program that accepts the object program, prepares these programs for execution by the computer and initiates the execution of the program.

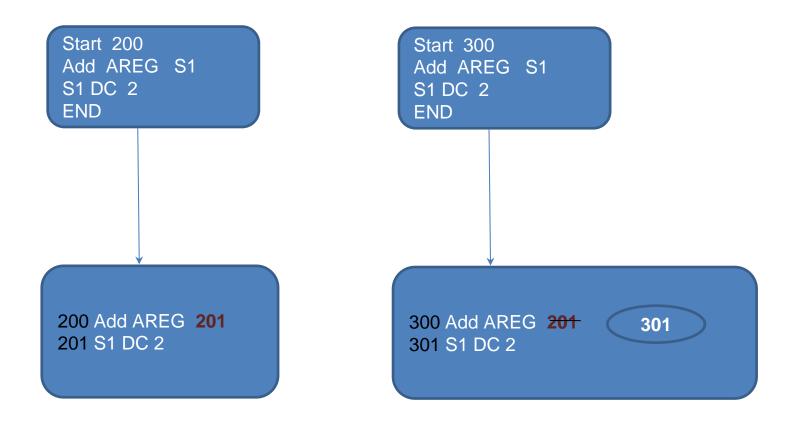
Functions of the loader

- 1. Allocate space in memory for the programs.(Allocation)
- 2. Resolve symbolic references between object decks.(Linking)
- 3. Adjust all addr dependent locations, such as addr constants, to correspond to the allocated space.(**Relocation**)
- 4. Physically place the m/c instr and data into memory.(**Loading**)



Programs loaded in memory ready for execution

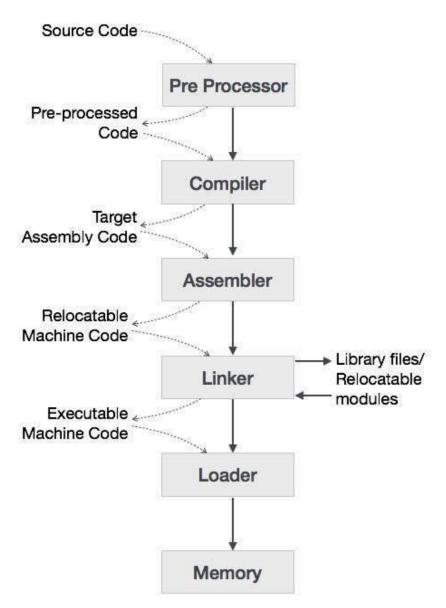
Concept of Relocation



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Language Processing System



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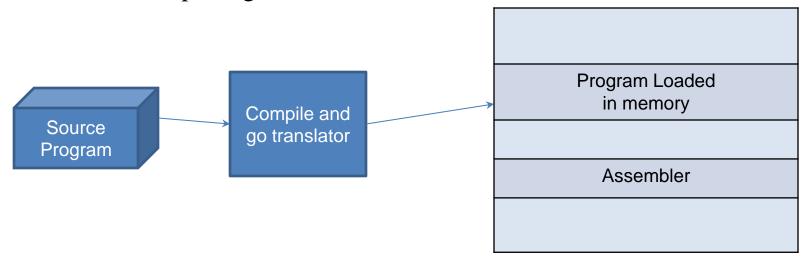
Different Types of Loader Schemes

- 1. Compile and Go Loaders
- 2. General Loader Scheme
- 3. Absolute Loaders
- 4. Relocating Loaders
- 5. Direct Linking Loaders



1. Compile and Go Loaders

- -Assembler places the code into core
- -Loader consists of one instr that transfers to the starting instr of the newly assembled program
- -easy to implement
- -portion of memory is wasted because of assembler
- -every time the program is run it has to be retranslated
- -difficult to handle multiple segments

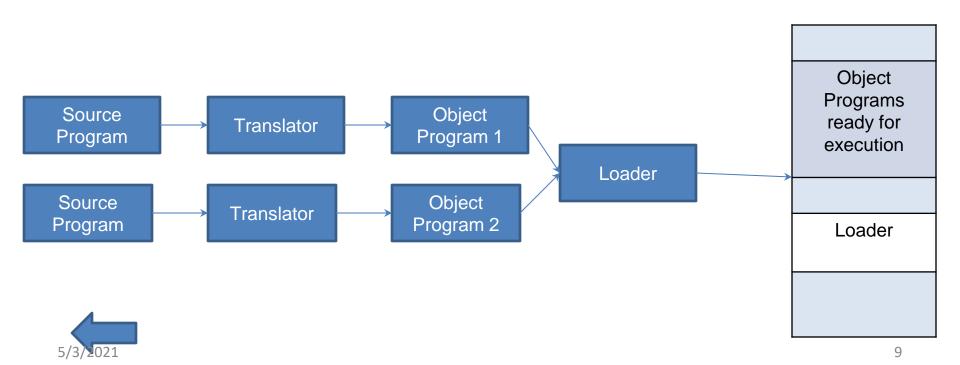






2. General Loader Scheme

- As loader is smaller than assembler more memory is available
- Reassembling of program is not required to run the program later.
- Loader is present in memory.





3. Absolute Loader

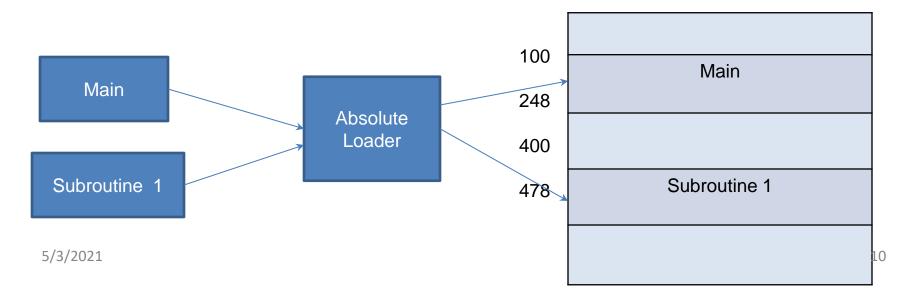
- -Same as "compile and go" loader except data is punched on cards instead of memory.
- -Loader accepts m/c language text and places it into memory at the location specified by the assembler

Advantages

-More memory is available, -simple to implement

Disadvantages

- -Programmer must specify address to the assembler where the program is loaded.
- -In case of multiple subroutines, programmer has to remember address of each subroutine.





Relocating Loader (Binary Symbolic Subroutine)

- -To avoid possible reassembling of all subroutines when a single subroutine is changed.
- -To perform **task of allocation and linking** for the programmer.
- -Allows many procedure segments but only one data segment.
- -Translated code segments and the information regarding relocation and intersegment references is passed to the loader.

Information provided by the assembler to the BSS loader.

- Transfer Vector
- Contains the address and names related to the subroutines referenced in the program.
- Total length of the program
- length of transfer vector
- Relocation Bits
- relocation bit is associated with every instruction
- Relocation bits can be 0 or 1.
- If 1 then address field needs relocation
- If 0 then address field does not need relocation

ST	14 SAVE			
ST	14	SAVE		
Relocation Bit=0		Relocation Bit=1		

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4. Relocating Loader (Binary Symbolic Subroutine)

- In BSS
- All four functions of loader (allocation, linking, relocation and loading) are performed automatically by the BSS loader.
- **Relocation bits** are used to solve the problem of relocation.
- The **transfer vector** is used to solve the problem of linking.
- The **program length** information is used to solve the problem of allocation.

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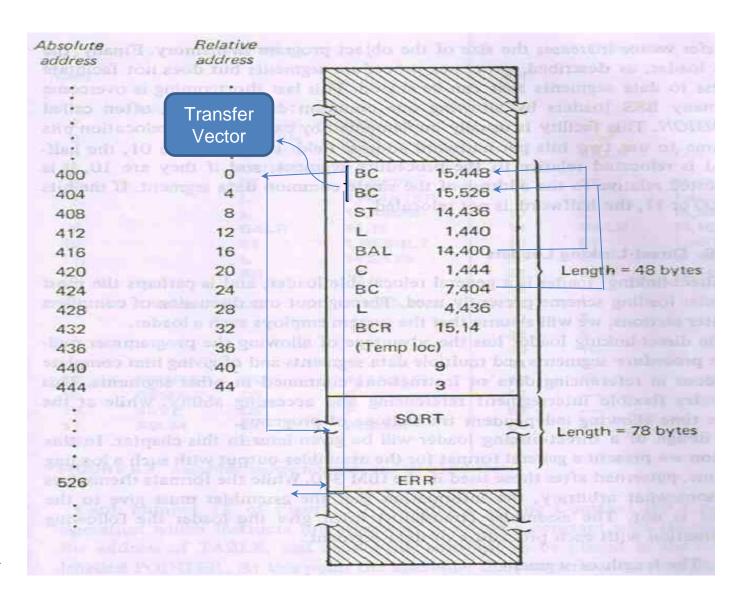
4. Relocating Loader (Contd..)

Source program Pro			Program	Length =	48 bytes	
			Transfer	Vector =	8 bytes	
			<u>Rel.</u>	<u>Rel</u>	Object Co	<u>ode</u>
MAIN	START		<u>Addr.</u>	<u>Bits</u>		
	EXTRN	SQRT	0	00	'SQRT'←	
	EXTRN	ERR	4	00	'ERRb' ←	
	ST	14,SAVE	8	01	ST	14,36
	L	1,=F'9'	12	01	L	1,40
	BAL	14,SQRT	16	01	BAL	14,0
	С	1,=F'3'	20	01	С	1,44
	BNE L BR	ERR 14,SAVE 14	24 28 32	01 01 0	BC L BCR	7,4 14,36 15,14
SAVE	DS	F	36	00	(Temp lo	cation)
	END		40	00	9	
			44	00	3	

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4. Relocating Loader (Contd..)





Disadvantages of Relocating Loader

- The transfer vector linkage is only useful for transfers and **not well** suited for loading or storing external data.
- The transfer vector increases the size of the object program in memory.
- BSS loader processes procedure segments but does not facilitate access to data segments that can be shared.

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5. Direct Linking Loader

- Flexible intersegment referencing and accessing ability.
- Allows independent translation of programs.

Information provided by the assembler with each procedure or data segment

- Length of the segment.
- List of symbols and relative locations.
- List of symbols not defined but referenced.
- Information where address constants are located.
- M/c translation of source program and relative address assigned.

Assembler produces 4 types of cards in the object deck.

$ESD \square$	External Symbol Dictionary.
$TXT \square$	Actual Object Code.
$RLD \square$	Relocation and Linkage Direct

RLD □ Relocation and Linkage Directory.

END _ End of object deck.



5. Direct Linking Loader(contd...)

ESD cards

- Contains info related to all the symbols defined and referenced in the program.
- Values for ESD cards are

SD (Segment Definition)	name on START card
LD (Local Definition) □	Specified on ENTRY card

ER (External Reference)

specified on EXTRN card

TXT cards

Contains actual object code translated version of program.

RLD cards

- The location constant that needs to be changed due to relocation
- By what is has to be changed
- The operation to be performed(+/-)

END cards

End of object deck and specifies the starting address for execution if the assembled routine is the main program.

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5. Direct Linking Loader(contd...)

Card No	ALP	Rel Loc	Translation
1.	JOHN START	Official	Findow need - Flaces need
2.	ENTRY RESULT		[index reg] + [base reg] + [12] = 54 + 0 + 2 = 56
3.	EXTRN SUM		1
4.	BALR 12, 0	0	BALR 12,0
5.	USING *, 12		[12]< current value of LC
6	ST 14, SAVE	2	ST 14, 54(0,12)
7.	L 1, POINTER	6	L 1, 46(0,12)
8.	L 15, ASUM	10	L 15, 58(0,12)
9.	BALR 14, 15	14	BALR 14, 15
10.	ST 1, RESULT	16	ST 1, 50(0,12)
11.	L 14, SAVE	20	L 14, 54(0,12)
12.	BR 14	24	BCR 15, 14
13.	TABLE DC F '1, 7, 9, 10, 3'	28 32 36 40 44	1 7 9 10 3
14.	POINTER DC A(TABLE)	48	28
15.	RESULT DS F	52	-
16.	SAVE DS F	56	-
17.	ASUM DC A(SUM)	60	?
18.	END	64	



ESD And RLD Cards

ESD Cards				
Ref No	Symbol	Type	Relative Loc	Length
1.	JOHN	SD	0	64
2.	RESULT	LD	52	-
3.	SUM	ER	-	-

RLD Cards				
Ref No	Symbol	Flag	Length	Rel Loc
14	JOHN	+	4	48
17	SUM	+	4	60

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TXT Cards

TXT Cards			
Ref No	Rel Loc	Object Co	de
4	0	BALR	12,0
6	2	ST	14, 54(0,12)
7	6	L	1, 46(0,12)
8	10	L	15, 58(0,12)
9	14	BALR	14, 15
10	16	ST	1, 50(0,12)
11	20	L	14, 54(0,12)
12	24	BCR	15, 14
13	28	1	
13	32	7	
13	36	9	
13	40	10	
13	44	3	
14	48	28	
17	60	0	

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1 2 3 4	PG1 PG1ENT1	START ENTRY PG1ENT1, PG1ENT2 EXTRN PG2ENT2, PG2
5 6 7 8 9 10	PG1ENT2	DC A (PG1ENT1) DC A (PG1ENT2+15) DC A (PG1ENT2-PG1ENT1-3) DC A (PG2) DC A (PG2ENT1+PG2-PG1ENT1+4) END
12 13 14	PG2	START ENTRY PG2ENT1 EXTRN PG1ENT1, PG1ENT2
15 16 17 18 19	PG2ENT1	DC A (PG1ENT1) DC A (PG1ENT2+15) DC A (PG1ENT2-PG1ENT1-3) END

Source Card Ref no	Relat Addr			
1 2 3 4 5 6 7 8 9 10 11	0 20 30 40 44 48 52 56 60	PG1ENT1 PG1ENT2	DC A (P DC A (P DC A (P	PG1ENT1, PG1ENT2 PG2ENT2, PG2 G1ENT1) G1ENT2+15) G1ENT2-PG1ENT1-3) G2) G2ENT1+PG2-PG1ENT1+4)
12 13 14 15 16 17 18	0 16 24 28 32 36	PG2 PG2ENT1	DC A (P	PG2ENT1 PG1ENT1, PG1ENT2 G1ENT1) G1ENT2+15) G1ENT2-PG1ENT1-3)

OBJECT DECK FOR PG1

ESD Cards

Source Card Ref No	Name	Туре	ID	Relative Address	Length
1	PG1	SD	01	0	60
2	PG1ENT1	LD	01	20	
2	PG1ENT2	LD	01	30	
3	PG2	ER	02		
3	PG2ENT1	ER	03		

TXT Cards (Those having address constants)

Source	Relative	Contents	Comments
Card	Address		
Ref No			
6	40-43	20	
7	44-47	45	30 +15
8	48-51	7	30-20-3
9	52-55	0	UNKNOWN TO PG
10	56-59	-16	-20+4

RLD Cards

Source Card Ref No	ESD-ID	Length in bytes	FLAG + or -	Relative Address
6	01	4	+	40
7	01	4	+	44
9	02	4	+	52
10	03	4	+	56
10	02	4	+	56
10	01	4	-	56

PG2

ESD Cards

Source Card Ref No	Name	Туре	ID	Relative Address	Length
12	PG2	SD	01	0	36
13	PG2ENT1	LD		16	
14	PG1ENT1	ER	02		
14	PG1ENT2	ER	03		

Txt Cards (Those having address constants)

1	Relative Address	Contents	Comments
16	24-27	0	
17	28-31	15	
18	32-35	-3	

RLD Cards

Source Card Ref No	ESD-ID	Length in bytes	FLAG + or -	Relative Address
16	02	4	+	24
17	03	4	+	28
18	03	4	+	32
18	02	4	-	32



MIT-WPU Final Year (B.Tech)

System Software and Compiler Design



Module II

- Macro processor: Macro Definition, Macro expansion and nested macros
- Loaders: Loader schemes: Types of loaders, direct linking loaders.
- Linkers: Relocation and linking concepts, self-relocating programs, Static and dynamic link libraries.

Introduction to Macro

- In the mid-1950s, when assembly language programming was commonly used to write programs for digital computers, the use of **macro instructions** was initiated for two main purposes:
 - to reduce the amount of program coding that had to be written by generating several assembly language statements from one macro instruction and
 - to enforce program writing standards, e.g. specifying input/output commands in standard ways.
 - Macro instructions were effectively a middle step between assembly language programming and the high-level programming languages that followed, such as FORTRAN and COBOL
 - — Ед.
 - **#define** PI 3.14159 //"PI" to be replaced with "3.14159" wherever it occurs
 - #define pred(x)((x)-1) // What this macro expands to depends on what argument x is passed to it. Here are some possible expansions:
 - pred(2) \rightarrow ((2) -1)
 - pred(y+2) \rightarrow ((y+2) -1)
 - pred(f(5)) \rightarrow ((f(5))-1)



Introduction

- Macro instructions (macros) are single -line abbreviations for group of instructions.
- Macros are used to provide a program generation facility through macro expansion.
- Many languages provide built in facilities for writing macros.
 e.g. PL/I, C, Ada and C++, Assembly languages.
- Generating preprocessors or software tools like Awk of Unix has an equivalent effect.
- A macro is a unit of specification for program generation through expansion.

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Introduction (contd...)

Macros are defined at the start of the program.

A Macro Definition consists of

- MACRO Pseudo opcode
- Name of macro
- List of formal parameters
- Body of macro(instruction) or Sequence to be abbreviated
- MEND Pseudo opcode

Macro Definition Syntax:-

- 1) Macro header :- It contains keyword 'MACRO'.
- 2) Macro prototype statement syntax :-
 - < Macro Name > [& < Formal Parameters >]
- 3) Model Statements: It contains 1 or more simple assembly statements, which will replace MACRO CALL while macro expansion.
- 4) MACRO END MARKER: It contains keyword 'MEND'.

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Introduction

- Macro name with a set of actual parameters, is replaced by some code, generated from macro body. This is called macro expansion.
- Two types of expansions:
- **Lexical expansion:** It implies replacement of a character string by another character string during program generation.
- **Semantic expansion:** Generation of type specific instructions for manipulation of byte and word operands.

Example 4.1 (Benefits of semantic expansion) The following sequence of statements is used to increment the value stored in a memory word by a given constant:

- 1. Move the value from the memory word into a CPU register.
- 2. Increment the value in the CPU register.
- Move the new value into the memory word.



Macro-processor

- Many times some blocks of code is repeated in the course of a program
- They may consists of code
 - -- to save or exchange set of registers
 - -- to set up linkages
 - -- to perform a series of arithmetic operations
- In this situation macro instruction facility is useful.
- Programmer defines a single instruction to represent a block of code.
- For every occurrence of this one line instruction the assembler will substitute the entire block.

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Macros and Functions

Macros

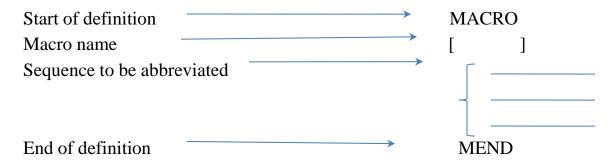
- Macro is Preprocessed
- No Type Checking is done in Macro
- Using Macro increases the code length
- Speed of Execution using Macro is Faster
- Before Compilation, macro name is replaced by macro value
- Macros are useful when small code is repeated many times
- Macro does not check any Compile-Time Errors
- Does not require CALL and RETURN

Functions

- Function is Compiled
- Type Checking is done in functions
- Using Function keeps the code length unaffected
- Speed of Execution using Function is Slower
- During function call, transfer of control takes place
- Functions are useful when large code is to be written
- Function checks Compile-Time Errors
- Requires CALL and RETURN



Macro Instruction



Macro Example

MACRO

INCR

ADD 1,DATA

ADD 2,DATA

ADD 3,DATA

MEND

:

INCR

:

INCR

Example of macro

* Source Program > mACRO	Taxget Trogram with Marro expansions
my macro	a chay (9)
15 ADD AREG, X	START 100
ADD BREG,X	ADD AREG,X
MEND	HDD BREG,X
START 100 imput Macro output	ADD AREG,X
my macro processor	ADD BREG,X
20 MYMACRO	ADD AREG,X
MyMACRO	ADD BREG,X
STOP	STOP



Example-Macro

 Macro instructions are single line abbreviations for group of instructions

e.g.

A 1,DATA

A 2,DATA

A 3,DATA

• • • • • • • •

A 1,DATA

A 2,DATA

A 3,DATA

DATA DC F'5'

	EXPANDED SOURCE
MACRO	
INCR	
A 1,DATA	
A 2,DATA	
A 3,DATA	
MEND	
	A 1,DATA
INCR	A 2,DATA
	A 3,DATA
	A 1,DATA
INCR	A 2,DATA
	A 3,DATA
DATA DC F'5'	



Contd...

Source program with Macro definitions & macro calls

expanded source without definition & call

MACRO

INCR &ARG

ADD AREG, & ARG

ADD BREG,&ARG

MEND

MOVER CREG, LABEL1

INCR DATA1

SUB CREG,LABEL1

MOVER CREG,LABEL1
ADD AREG,DATA1
ADD BREG,DATA1
SUB CREG ,LABEL1



Macro features

- 1. Macro instruction arguments
- 2. Conditional macro expansion
- 3. Macro calls within macros
- **4.**Macro instructions defining macros



Macro Instruction Arguments

1. Macro instruction arguments

- Macro facility lacks in flexibility: there is no way for a specific macro call to modify the coding that replaces it.
- So extension of this facility consists of providing for arguments or parameters in macro call
 - dummy arguments
 - positional arguments
 - keyword arguments
 - label arguments

these arguments are preceded by &.



Macro Instruction Arguments Contd...

e.g.	Source	Expanded source
•••••	MACRO	
•••••	INCR &ARG	
A 1,DATA1	A 1,&ARG	
A 2,DATA1	A 2,&ARG	
A 3,DATA1	A 3,&ARG	
	MEND	
•••••		
A 1,DATA2		A 1,DATA1
A 2,DATA2	INCR DATA1	A - 2,DATA1
A 3,DATA2		A 3,DATA1
••••		
••••	INCR DATA2	A [1,DATA2
DATA1 DC F'5'		A - 2,DATA2
DATA2 DC F'10'		A \(\lambda \),DATA2
	DATA1 DC F'5'	DATA1 DC F'5'
	DATA2 DC F'10'	DATA2 DC F'10'
	·	



Macro Instruction Arguments (Label Arguments)

MACRO		Loop1 A	1,DATA1
&LAB INCR	&ARG1,&ARG2,&ARG3	A	2,DATA2
&LAB A	1, &ARG1	A	3,DATA3
A	2, &ARG2	7.1	3,0711713
A	3, &ARG3	• • • • • •	
MEND		Loop2 A	1,DATA3
•••••		A	2,DATA2
LOOP1 INCR	DATA1,DATA2,DATA3	A	3,DATA1
	DATA3,DATA2,DATA1	DATA1 D	C F'5'
DATA1 DC F'	5'	DATA2 D	C F'10'
DATA2 DC F'10'		DATA3 DC F '15'	
DATA3 DC F	'15'		
••••			



Macro Instruction Arguments

Positional arguments

arguments are matched with dummy arguments according to the order in which they appear

e.g. INCR A B C

Keyword arguments:

it allows reference to dummy arguments by name as well as by position

e.g. INCR &ARG1=A,&ARG3=C,&ARG2=B



Macro Instruction Arguments (Positional Arguments)

MACRO

INCR &ARG1,&ARG2,&ARG3,&LAB

&LAB A 1, &ARG1

A 2, &ARG2

A 3, &ARG3

MEND

.

INCR DATA1,DATA2,DATA3,LOOP1

.

INCR DATA3,DATA2,DATA1,LOOP2

DATA1 DC F'5'

DATA2 DC F'10'

DATA3 DC F'15'

.



2 Pass Macroprocessor

- Recognize macro definitions
- Save the definitions
- Recognize calls
- Expand calls and substitute arguments



2 Pass Macroprocessor structure

Pass I

- 1. **Input** macro source deck
- 2. Output macro source deck copy for use by pass 2
- 3. Macro definition table (MDT) ,used to store body of macro def.
- 4. Macro name table (MNT) ,used to store names of defined macros
- 5.MDTC-MDT counter, used to indicate next available entry in MDT
- 6.MNTC-MNT counter, used to indicate next available entry in MNT
- 7. Argument List Array(**ALA**) used to substitute index markers for dummy arguments before storing macro definition

Example

78. /	- A		_	$\overline{}$
		\mathbb{C}	К	
			V 4	w

M1 &ARG1,&ARG2

ADD AREG & ARG1

ADD BREG & ARG2

MEND

MACRO

M2 &ARG3,&ARG4

SUB AREG & ARG3

SUB BREG & ARG4

MEND

START 300

MOVER AREG S1

MOVEM BREG S2

M1 D1 D2

MOVER AREG S1

M2 D3 D4
PRINT S1
PRINT S2
S1 DC 5
S2 DC 6
END



Contd...

• Pass 2

- 1. Copy of **input** macro source deck
- 2. **Output** expanded source deck to be used as input to assembler
- 3. **MDT**, created by Pass 1
- 4. **MNT**, created by pass 1
- 5. **MDTP-** MDT pointer used to indicate next line of text to be used during macro expansion
- 6. **ALA**, used to substitute macro call arguments for the index markers in the stored macro definition.



- Macro definition table(MDT)
- Used to store macro definition
- Created by pass-1
- Pass 1 identifies and stores all macro definitions in MDT
- Pass 2 can identify macro calls and expand these calls by using their macro definitions stored in MDTs
- Every line of macro definition except MACRO is stored in MDT as MACRO is not used to expand macro
- MEND indicates end of macro definition so it is stored in MDT
- MDT has 80 bytes per entry

Index	instruction



MNT

- Macro name table(MNT)
- Used to store name of macros and corresponding MDT index
- Created by Pass 1 and used by Pass 2
- So that pass 2 can decide whether opcode of this source instruction is macro call or not by searching through name field of MNT.

MNT index	Macro Name	MDT index

MDTC(Macro def table counter)

this variable stores the last count from the MDT table

MNTC(Macro name table counter)

this variable stores the no. of macros defined in the program



ALA

- Argument List Array(ALA)
- Used for association of dummy arguments and actual arguments
- Partially table is created by pass 1 where pass 1 associates an unique integer number with each dummy arguments in the order in which they appear
- Macros are stored in MDT by using these numbers
- It is partially constructed and used by pass 2
- It keeps track of dummy parameters when the macro is being defined and maintain the actual arguments when expanding the call

Integer index	dummy argument	Actual argument



Example

MACRO

INCR & ARG1 MOVER AREG,& ARG1 ADD AREG,& ARG1 MOVEM AREG,& ARG1

MEND

START
MOVEM BREG, A
ADD CREG,='1'
SUB CREG,A
INCR DATA1
MUL CREG,='1'
END

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START 200		
MACRO		
INCR &ARG1,&ARG2		
MOVER AREG, A		
ADD AREG,B		
MOVEM AREG,A		
MEND		
MOVER AREG,='1'		
MOVEM BREG,M		
INCR DATA1,DATA2		
DATA1 DC 5		
DATA2 DC 10		
END		

MDT

MDT index	Instruction
1	

MNT

MNT Index	Macro Name	MDT index
1		

ALA

ALA Index	Formal Argument
#1	

START 200 MACRO INCR &ARG1,&ARG2 MOVER AREG, A ADD AREG,B MOVEM AREG,A **MEND** MOVER AREG,='1' MOVEM BREG,M **INCR DATA1, DATA2** DATA1 DC 5 DATA2 DC 10 **END**

MDT

MDT index	Instruction
1	

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	

START 200 MACRO INCR &ARG1,&ARG2 MOVER AREG, A ADD AREG,B MOVEM AREG,A **MEND** MOVER AREG,='1' MOVEM BREG,M **INCR DATA1, DATA2** DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

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ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

START 200

MACRO

INCR &ARG1,&ARG2

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

MEND

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1,DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

START 200

MACRO

INCR &ARG1,&ARG2

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

MEND

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1, DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

START 200

MACRO

INCR &ARG1,&ARG2

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

MEND

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1, DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A
3	ADD AREG,B

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

START 200

MACRO

INCR &ARG1,&ARG2

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

MEND

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1, DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A
3	ADD AREG,B
4	MOVEM AREG,A

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

START 200

MACRO

INCR &ARG1,&ARG2

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

MEND

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1, DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A
3	ADD AREG,B
4	MOVEM AREG,A
5	MEND

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

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ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

Output of Pass 1

Output file

START 200

MOVER AREG,='1'

MOVEM BREG,M

INCR DATA1,DATA2

DATA1 DC 5

DATA2 DC 10

END

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A
3	ADD AREG,B
4	MOVEM AREG,A
5	MEND

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

Output of Pass 2

MDT

MDT index	Instruction
1	INCR &ARG1,&ARG2
2	MOVER AREG, A
3	ADD AREG,B
4	MOVEM AREG,A
5	MEND

MNT

MNT Index	Macro Name	MDT index
1	INCR	1

ALA

ALA Index	Formal Argument
#1	&ARG1
#2	&ARG2

Final Expansion

START 200

MOVER AREG,='1'

MOVEM BREG,M

MOVER AREG, A

ADD AREG,B

MOVEM AREG,A

DATA1 DC 5

DATA2 DC 10

END



Contd...

```
MACRO
```

&LAB INCR &ARG1,&ARG2,&ARG3

&LAB MOVER AREG, &ARG1

A DD AREG, &ARG2

MOVEM BREG, &ARG3

MEND

• • • • • •

LOOP1 INCR DATA1,DATA2,DATA3

• • • • • • • • • •

LOOP2 INCR DATA3,DATA2,DATA1

• • • • • • • •

DATA1 DC F'5'

DATA2 DC F'10'

DATA3 DC F'15'



Example Contd...

MDT

Index	instruction
1	&LAB INCR &ARG1,&ARG2,&ARG3
2	#0 MOVER AREG,#1
3	ADD AREG,#2
4	MOVEM BREG,#3
5	MEND

ALA

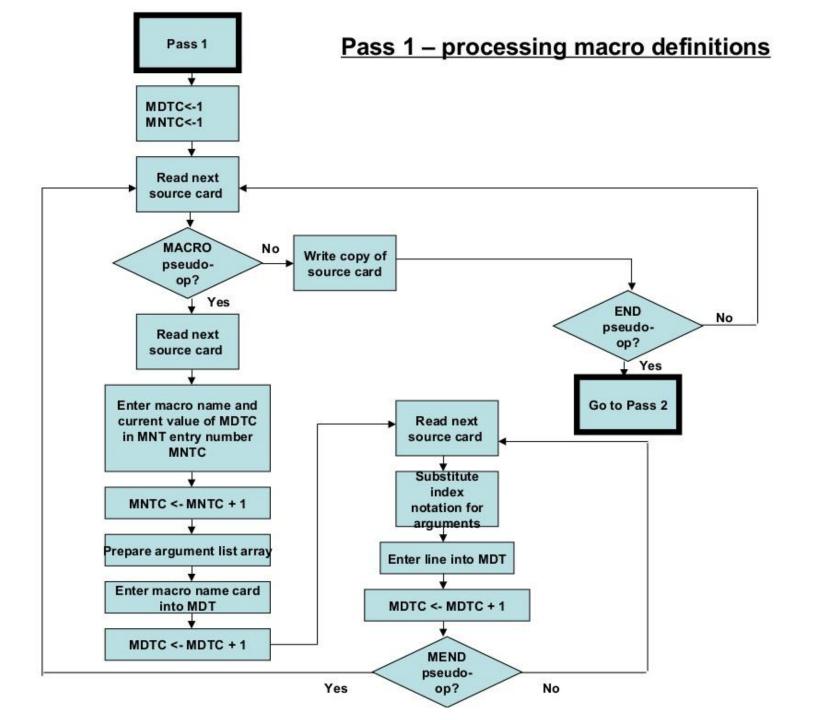
Index	arguments
#0	"LOOP1BBB"
#1	"DATA1BBB"
#2	"DATA2BBB"
#3	"DATA3BBB"

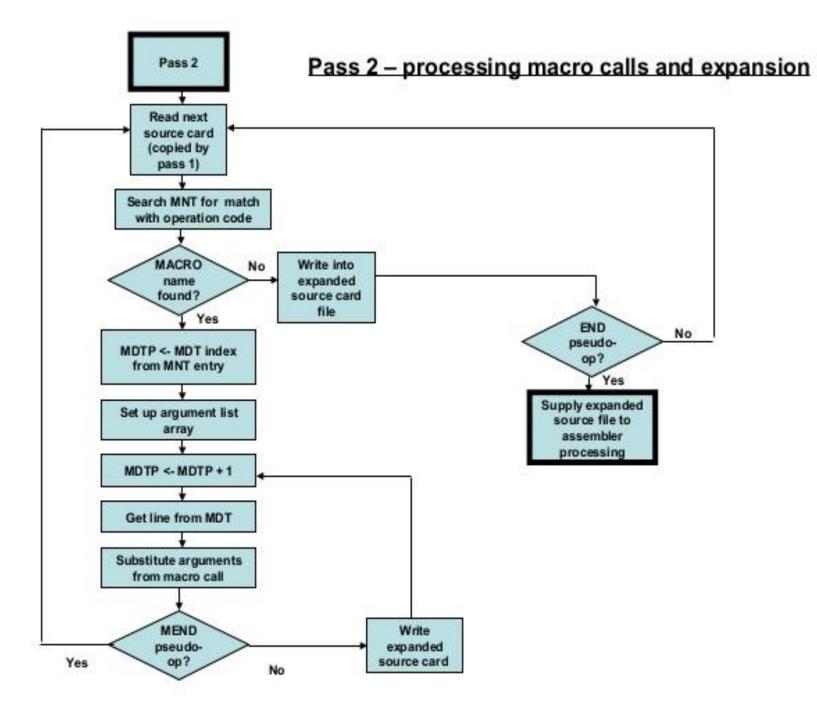
MNT

Index	Name(8 bytes)	MDT index (4 bytes)
1	"INCRbbbb"	1

ALA

Index	arguments
#0	"BBBBBBB"
#1	"DATA3BBB"
#2	"DATA2BBB"
#3	"DATA1BBB"







Contd...

```
START
MACRO
INCR &ARG1,&ARG2,&ARG3
      1, &ARG1
A
A 2, &ARG2
      3, &ARG3
Α
MEND
INCR DATA1,DATA2,DATA3
 INCR DATA3,DATA2,DATA1
 DATA1 DC F'5'
 DATA2 DC F'10'
 DATA3 DC F'15'
 END
```



Contd...

```
START
 MOVER AREG,A
 MOVEM AREG,B
 . . . . . .
 MACRO
 VARY &ARG1,&ARG2
 L 1,F'5'
 A 1,&ARG1
 A 1,&ARG2
 MEND
 MACRO
 INCR &ARG1,&ARG2,&ARG3
         1, &ARG1
  Α
         2, &ARG2
         3, &ARG3
 MEND
  INCR DATA1,DATA2,DATA3
  . . . . . . . . . .
  VARY DATA3,DATA2
  . . . . . . . . . .
  DATA1 DC F'5'
  DATA2 DC F'10'
  DATA3 DC F'15'
   . . . . . . .
   END
```



Conditional macro expansion

```
MACRO
        VARY &COUNT,&ARG1,&ARG2,&ARG3
&ARG0
&ARG0
        A 1, &ARG1
        A IF (&COUNT EQ 1).FINI
                                               TEST IF &COUNT=1
        A 2, &ARG2
        A IF (&COUNT EQ 2).FINI
                                                TEST IF &COUNT=2
          3, &ARG3
.FINI
          MEND
LOOP1 VARY 3, DATA1, DATA2, DATA3
 . . . . . . . . . .
LOOP2 VARY 2, DATA3,DATA2
LOOP3 VARY 1, DATA1
DATA1 DC F'5'
DATA2 DC F'10'
DATA3 DC F'15'
. . . . . . .
```



Expanded code

```
Loop1 A
         1,DATA1
         2,DATA2
      A 3,DATA3
Loop2 A 1,DATA3
      A 2,DATA2
      . . . . . . . . .
Loop1 A 1,DATA1
 . . . . . . .
DATA1 DC F'5'
DATA2 DC F'10'
DATA3 DC F '15'
```

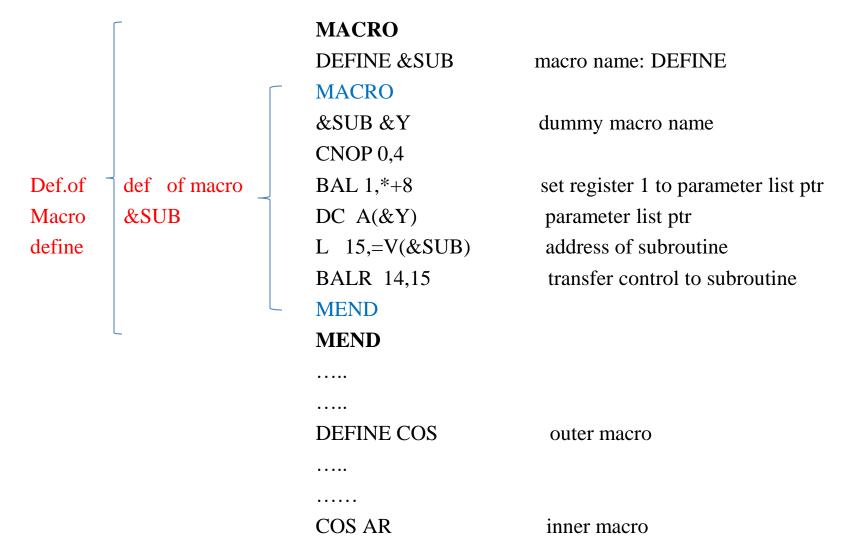


Contd...

- Pseudo opcodes AIF & AGO
- .FINI (labels starting with a period(.) are macro labels)and do not appear in the output of the macro processor
- AIF (conditional branch pseudo opcodes)
- AGO (unconditional branch pseudo opcodes) or goto statement.
- AIF & AGO control the sequence in which the macro processor expands the statements in macro instruction



Macro instructions defining macros





A single pass algorithm for macro processor

 MDI: macro definition input indicator works like switch keeps track of macro call

MDI is ON

- -during expansion of a macro call
- -lines are read from MDT

MDI is OFF

all other times

- --the reading of MEND line indicates end of macro and terminates expansion of a call ,MDI is set off next line is obtained from regular i/p stream.
- MDLC : macro definition level counter
 - -- keeps track of macro definition.
 - -- MDLC is incremented by 1 when a MACRO is encountered and decremented by 1 when MEND occurs
 - -- MDLC is used to insure that the entire macro def. including MACROs & MENDs get stored in MDT



	Index	instructions
START	1	
MACRO		
DEFINE &SUB		
MACRO		
&SUB &Y		
CNOP 0,4		
BAL 1,*+8		
DC A(&Y)		
L $15,=V(\&SUB)$		
BALR 14,15		
MEND		
MEND		
 DEFINE COS		
•••••		
COS AR		
END		

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1		

ALA

ALA Index	Formal Argument	Actual Argument
#1		



		.,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	
	Index	instructio	ns
START	1	DEFINE	&SUB
MACRO	2		
DEFINE &SUB			
MACRO			
&SUB &Y			
CNOP 0,4			
BAL 1,*+8			
DC A(&Y)			
L 15,=V(&SUB)			
BALR 14,15			
MEND			
MEND			
••••			
 DEFINE COS			
•••••			
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



		14161	
	Index	instructio	ons
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3		
MACRO			
&SUB &Y			
CNOP 0,4			
BAL 1,*+8			
DC A(&Y)			
L 15,=V(&SUB)			
BALR 14,15			
MEND			
MEND			
••••			
DEFINE COS			
••••			
•••••			
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



START
••••
MACRO
DEFINE &SUB
MACRO
&SUB &Y
CNOP 0,4
BAL 1,*+8
DC A(&Y)
L 15,=V(&SUB)
BALR 14,15
MEND
MEND
••••
DEFINE COS
COS AR
END

MDI			
Index	instructio	ons	
1	DEFINE	&SUB	
2	MACRO		
3	#1	&Y	
4			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5		
BAL 1,*+8			
DC A(&Y)			
L 15,=V(&SUB)			
BALR 14,15 MEND			
MEND			
 DEFINE COS			
 COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6		
DC A(&Y)			
L 15,=V(&SUB) BALR 14,15			
MEND			
MEND			
 DEFINE COS			
 COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7		
L 15,=V(&SUB) BALR 14,15			
MEND			
MEND			
 DEFINE COS			
••••			
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB)	8		
BALR 14,15 MEND			
MEND			
 DEFINE COS			
 COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



START	Index	instructions	
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9		
MEND	10		
••••	11		
DEFINE COS	12		
	13		
	14		
COS AR	15		
END	16		
	17		

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10		
DEFINE COS			
••••			
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2		



START	Index	instructions	
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
••••	11		
DEFINE COS	12		
	13		
	14		
COS AR	15		
END	16		
	17		

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	COS	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2		



START	Index	instructio	ns
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
	11		
 DEFINE COS			
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2		

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2		



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
	11	COS	&Y
DEFINE COS	12		
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	COS	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



START	Index	instructions	
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
••••	11	cos	&Y
 DEFINE COS	12	CNOP	0,4
	13		
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	cos	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



START	Index	instructions	
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
	11	cos	&Y
 DEFINE COS	12	CNOP	0,4
	13	BAL	1,*+8
	14		
COS AR			
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	cos	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



START	Index	instructio	ns
	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
••••	11	COS	&Y
DEFINE COS	12	CNOP	0,4
	13	BAL	1,*+8
	14	DC	A(#1)
COS AR	15		
END			

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	cos	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
	11	cos	&Y
 DEFINE COS	12	CNOP	0,4
••••	13	BAL	1,*+8
	14	DC	A(#1)
COS AR	15	L	15,=V(COS)
END	16		

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	COS	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
	11	COS	&Y
DEFINE COS	12	CNOP	0,4
	13	BAL	1,*+8
	14	DC	A(#1)
COS AR	15	L	15,=V(COS)
END	16	BALR	14,15
	17		

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	cos	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	



	Index	instructio	ns
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC $A(&Y)$	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
••••	11	COS	&Y
DEFINE COS	12	CNOP	0,4
	13	BAL	1,*+8
	14	DC	A(#1)
COS AR	15	L	15,=V(COS)
END	16	BALR	14,15
	17	MEND	

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	COS	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	

Expansion of Define Cos call



	Index	instructions	
START	1	DEFINE	&SUB
MACRO	2	MACRO	
DEFINE &SUB	3	#1	&Y
MACRO	4	CNOP	0,4
&SUB &Y CNOP 0,4	5	BAL	1,*+8
BAL 1,*+8	6	DC	A(&Y)
DC A(&Y)	7	L	15,=V(#1)
L 15,=V(&SUB) BALR 14,15	8	BALR	14,15
MEND	9	MEND	
MEND	10	MEND	
••••	11	COS	&Y
DEFINE COS	12	CNOP	0,4
	13	BAL	1,*+8
	14	DC	A(#1)
COS AR	15	L	15,=V(COS)
END	16	BALR	14,15
	17	MEND	

MNT

MNT INDEX	MACRO NAME	MDT INDEX
1	DEFINE	1
2	cos	11

ALA

ALA Index	Formal Argument	Actual Argument
#1	&SUB	COS
#2	&Y	AR

Expansion of COS AR call

CNOP	0,4
BAL	1,*+8
DC	A(#1)
L	15,=V(COS)
BALR	14,15



Macro calls within macros

```
MACRO
ADD1
      &ARG
MOVER AREG, & ARG
ADD AREG, ='1'
MOVEM AREG,&ARG
MEND
MACRO
ADDS
       &ARG1,&ARG2,&ARG3
ADD1
       &ARG1
ADD1
       &ARG2
ADD1
       &ARG3
MEND
ADDS DATA1,DATA2,DATA3
. . . . . .
DATA1 DC F'5'
DATA2 DC F'10'
DATA3 DC F'15'
```

.



Contd...

गान्तिर्थुंचे थुवा ।।		expanded source	expanded source
		level 1	level 2
MACRO	0		
ADD1	&ARG		
L	1,&ARG		
A	1,F'1'		
ST	1,&ARG		
MEND			
MACRO	O		
ADDS	&ARG1,&ARG2,&ARG3	}	
ADD1	&ARG1		
ADD1	&ARG2		
ADD1	&ARG3	expansion of ADD	expansion of ADD1
MEND			
			L 1,DATA1
		ADD1 DATA1	Γ A 1,F'1'
			ST 1,DATA1
ADDS	DATA1,DATA2,DATA3	ADD1 DATA 2	L 1,DATA2
		_	A 1,F'1'
		ADD1 DATA3	ST 1,DATA2
			L 1,DATA3
DATA1	DC F'5'		Γ A 1,F'1'
DATA2	DC F'10'		ST 1,DATA3
DATA3	DC F'15'		
END			

3/27/2020



START

• • • • • •

.

MACRO

ADD1 &ARG

L 1,&ARG

A 1,F'1'

ST 1,&ARG

MEND

. . . .

MACRO

ADDS &ARG1,&ARG2,&ARG3

ADD1 &ARG1

ADD1 &ARG2

ADD1 &ARG3

MEND

• • • •

ADDS DATA1,DATA2,DATA3

.

DATA1 DC F'5'

DATA2 DC F'10'

DATA3 DC F'15'

.....

END



START

.

• • • • • •

MACRO

ADD1 &ARG

L 1,&ARG

A 1,F'1'

ST 1,&ARG

MEND

• • • •

MACRO

ADDS &ARG1,&ARG2,&ARG3

ADD1 &ARG1

ADD1 &ARG2

ADD1 &ARG3

MEND

....

ADDS DATA1,DATA2,DATA3

.

DATA1 DC F'5'

DATA2 DC F'10'

DATA3 DC F'15'

.....

END

Index	instructions
1	ADD1 &ARG
2	L 1,#1
3	A 1,F'1'
4	ST 1 ,#1
5	MEND
6	ADDS &ARG1,&ARG2,&ARG3
7	ADD1 #1
8	ADD1 #2
9	ADD1 #3
10	MEND

MNT INDEX	MACRO NAME	MDT INDEX
1	ADD1	1
2	ADDS	6



Contd...

- MDTP= 6 when ADDS is called
- And MDI is set ON.
- Then READ function increments MDTP, gets line from the MDT(line 7)
- Then MDTP = MDTP+1 =7
- So it is ADD1 DATA1

Then MDTP =1 so here previous value of MDTP i.e. 7 will be lost.

This the problem with macro calls within macros

- So it will work recursively.
- means to process one macro before it is finished with another then to continue with the previous or outer.
- Recursive procedures usually operate by means of stack.
- Each stack frame is associated with each recursive call
- Here status of unfinished computations is preserved.

References

System Programming & Operating System TE: Macro Definition and call, Nested Macro
 Calls, (wikinote.org)