### COS221 L14 - (E)ER to Relational Mapping (Chapter 9 in Editions 6 and 7)

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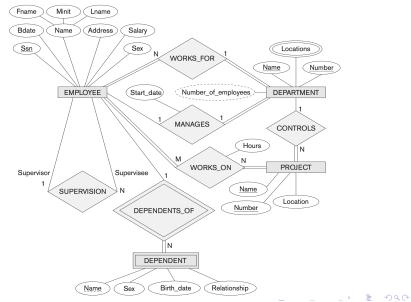
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### (E)ER-to-Relational Mapping Algorithm

- **Step 1:** Mapping of regular (strong) entity types
- **Step 2:** Mapping of weak entity types
- **Step 3:** Mapping of binary 1:1 relationships
- **Step 4:** Mapping of binary 1:N relationships
- **Step 5:** Mapping of binary M:N relationships
- **Step 6:** Mapping of multivalued attributes
- **Step 7:** Mapping of N-ary relationships
- **Step 8:** Mapping specialisation and generalisation
- **Step 9:** Mapping unions

#### The example

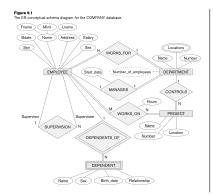
Figure 9.1
The ER conceptual schema diagram for the COMPANY database.

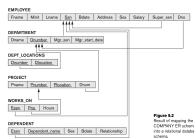


#### **Step 1:** Mapping of regular (strong) entity types

For each regular entity type, E, in the ER schema

- Create a relation R
- Add simple attributes of E to R
- Include the simple component attributes of composite attributes
- Choose one of the key attributes of E for R. If the attribute is a composite, then use all the simple attributes of the composite for the primary key of R.

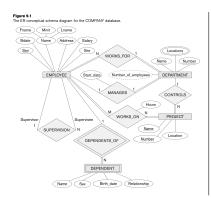


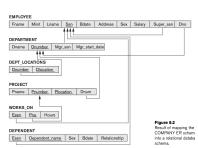


#### **Step 2:** Mapping of weak entity types

For each weak entity type, W, in the ER schema with owner entity E

- Create a relation R
- Add the simple attributes of W to R
- Include as foreign key attributes of R, the primary key of the relations that correspond to the owner entity type
- The primary key of R is a combination of the primary key(s) of the owner(s) and the partial key of W, if any
- If there is a weak entity type E<sub>1</sub>, whose owner is also a weak entity type, E<sub>2</sub>, then E<sub>1</sub> should be mapped before E<sub>2</sub> to determine the primary keys first.





### **Step 3:** Mapping of binary 1:1 relationships

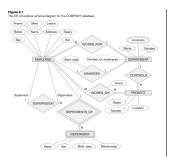
For each binary 1:1 relationship type R, identify the relations S and T that correspond to the entity types participating in R

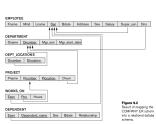
- 1. Foreign key approach most useful and followed unless special conditions exist
- Merged relation approach Works when both relations are total and can be merged. That is, when both relations have the same number of tuples at all times.
- 3. Cross-reference or relationship relation approach Sets up a third relation, *U*, for the purpose of cross-referencing the primary keys of *S* and *T*. The relation *U* becomes a look-up table because it relates one tuple from *S* with one tuple from *T*.

### **Step 3:** Mapping of binary 1:1 relationships (Approach 1)

For each binary 1:1 relationship type R, identify the relations S and T that correspond to the entity types participating in R, Choose one of the relations, say S - preferably the one with the total participation in R

- include as foreign key in S, the primary key of T
- include all the simple attributes (or simple components of the composite attributes) of the 1:1 relationship type R in S





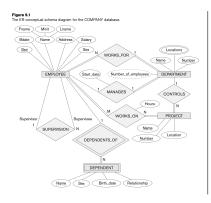
Choose DEPARTMENT as S because total in the MANAGES relation (every department has a manager). If EMPLOYEE was chosen as S, Dept\_managed would have been included in EMPLOYEE. Where an employee is not

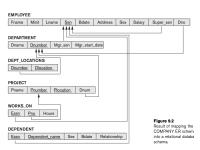
a manager, a NULL value would need to have been included.

#### **Step 4:** Mapping of binary 1:N relationships

For each binary 1:N relationship type R, identify the relation S on the N side

- include in S the primary key of relation T (the other entity on the relationship).
- include any simple attributes (or simple components of composite attributes) of the 1:N relationship type as attributes of S





# **Step 4:** Mapping of binary 1:N relationships (Alternative approach)

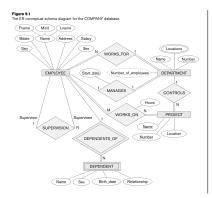
An alternative approach is to use the relationship relation (cross-reference) approach, discussed as the third option of the 1:1 relationships. In this approach

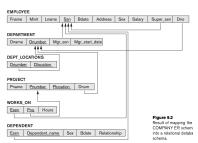
- Create a separate relation with attributes being the primary keys of the two relations involved in the relationship which will act as foreign keys for the respective relations.
- The primary key in this relation is the same as that of the first relation.
- This approach can lead to NULL values in the foreign keys.

#### **Step 5:** Mapping of binary M:N relationships

For each binary M:N relationship type R,

- create a new relation S to represent R.
- include as foreign key attributes in S, the primary keys of the relations that represent the participating entity types
- make the combination of these primary keys the primary key
- include the simple attributes (or the composite as simple components of the relationship

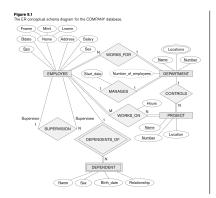


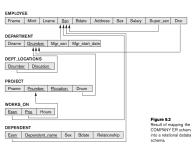


#### **Step 6:** Mapping of multivalued attributes

For each multivalued attribute A, create a new relation R

- R will include an attribute corresponding to A and a primary key K as foreign key in R of the relation that represents the entity type that has A as a multivalued attribute
- the primary key R is a combination of A and K (if the multivalued attribute is composite, include its simple components).

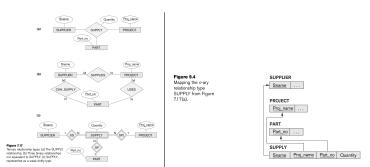




#### **Step 7:** Mapping of N-ary relationships

For each N-ary relationship type R, where n  $\geq$  2, create a new relation S to represent R

- include as foreign attributes in S as primary key attributes that participate in the relationship
- include any simple attributes (or simple components of composites) as attributes of S
- the primary of S is a combination of all foreign keys that reference the relations participating in the relationship
- of the cardinality constraints on any of the entity types E participating in relation R is 1, then the primary key of S should not include the foreign key attribute that reference E' corresponding to E



#### **Step 8:** Mapping specialisation and generalisation

Convert each specialisation with m subclasses  $\{S_1, S_2, ..., S_m\}$  and generalised superclass  $C(k, a_1, ..., a_n)$ , where k is the primary key into relation schemas using one of the following options:

#### Option 8A: Multiple relations - superclass and subclasses

- create a relation L for C with attributes of C and the primary key k
- create a relation  $L_i$  for each subclass  $S_i$  with the attributes  $\{k\} \cup \{\text{attributes of } S_i\}$  and k as primary key
- works for total or partial, disjoint or overlapping specialisations

#### Option 8B: Multiple relations - subclass relations only

- create a relation  $L_i$  for each subclass  $S_i$  with attributes {attributes of  $S_i$ }  $\cup$  {k,  $a_1$ , ...,  $a_n$ } with primary key of k
- only works for specialisations where the subclasses are total, also recommended for specialisations with a disjointedness constraint

#### Option 8C: Single relation with one type attribute

- create a relation L with attributes  $\{k, a_1, ..., a_n\} \cup \{\text{attributes of } S_i\} \cup .... \{\text{attributes of } S_m\} \cup \{t\}$  and primary key of k
- t is a type or discriminating attribute whose value indicates the subclass to which the tuple belongs
   used for subclasses that are disjoint
- can generate NULL values

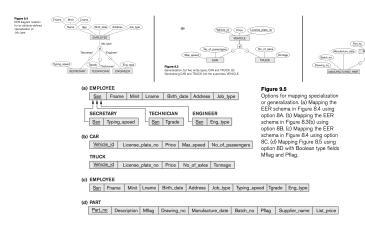
#### Option 8D: Single relation with multiple type attributes

- create a relation L with attributes  $\{k, a_1, ..., a_n\} \cup \{\text{attributes of } S_i\} \cup .... \{\text{attributes of } S_m\} \cup \{t_1, t_2, ..., t_m\}$  and primary key of k, where  $t_i$  is a Boolean type attribute indicating whether the tuple belongs to subclass  $S_i$
- used for specialisations whose subclasses are overlapping, also works for disjoint

#### Step 8: Mapping specialisation and generalisation

Convert each specialisation with m subclasses  $\{S_1, S_2, ..., S_m\}$  and generalised superclass  $C(k, a_1, ..., a_n)$ , where k is the primary key into relation schemas using one of the following options:

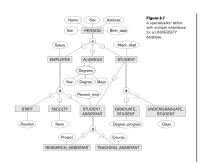
- Option 8A: Multiple relations superclass and subclasses
- Option 8B: Multiple relations subclass relations only
- Option 8C: Single relation with one type attribute
- Option 8D: Single relation with multiple type attributes



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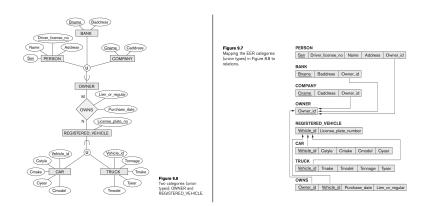
## **Step 8:** Mapping specialisation and generalisation - Lattices





#### Step 9: Mapping unions

When mapping classes whose superclasses have different keys, a surrogate key is defined when creating a relation which corresponds to the category. This key is the foreign key in the other relations participating in the relationship.



Owner\_id is the surrogate key.