Textbook:

**Operating Systems – Internals and Design Principles,** by *William Stallings*.

Publisher: Pearson.

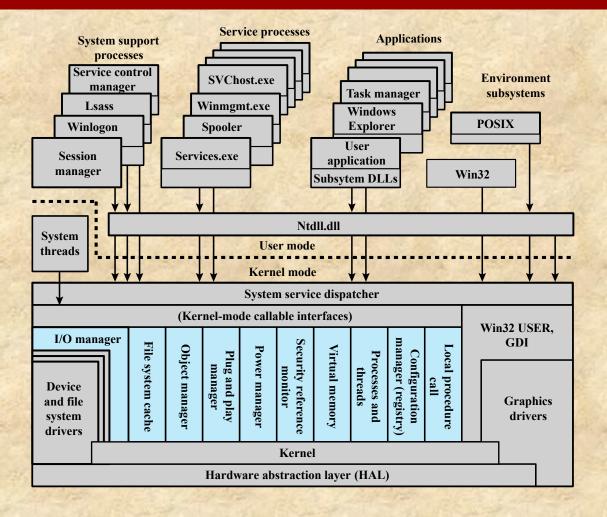
# Chapter 2 Operating Systems Overview



Part C

## "A short Walk through the Zoo"

- The various fundamental principles and theoretical techniques, which we have already looked at in parts A and B of this triple-lecture on Chapter 2, are combined and implemented differently in the various Operating Systems which are (or have been) available in academia and industry.
  - In this final part of our discussion of Chapter 2, which covers the remainder of this chapter from Section 2.7 onwards, we will take a glimpse at several different Operating Systems in order to see how the general theoretical concepts, which are generally valid, have been put to practice in different concrete forms and OS implementations.



Lsass = local security authentication server POSIX = portable operating system interface GDI = graphics device interface DLL = dynamic link libraries Colored area indicates Executive



## Kernel-Mode Components of Windows

#### Executive



■ Contains the core OS services, such as memory management, process and thread management, security, I/O, and interprocess communication

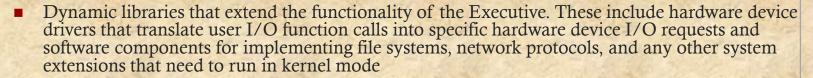
#### Kernel

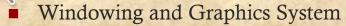
- Controls utilization of the CPUs. The Kernel manages thread scheduling, process switching, exception and interrupt handling, and multiprocessor synchronization
- Hardware Abstraction Layer (HAL)



■ Maps between generic hardware commands and responses and those unique to a specific platform and isolates the OS from platform-specific hardware differences

#### Device Drivers





■ Implements the GUI functions, such as dealing with windows, user interface controls, and drawing

### **SMP**

 "Windows" is also known for its support of symmetric multiprocessing (SMP)



- OS routines can run on any available CPUs, whereby different routines can run simultaneously on different CPUs
- Windows provides mechanisms for sharing data and resources between processes and flexible interprocess communication capabilities







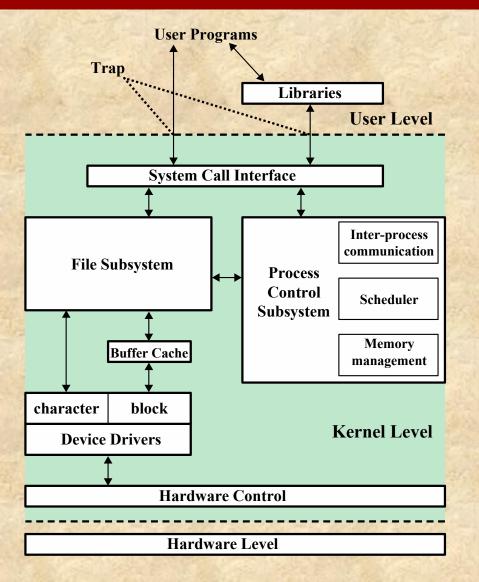




Figure 2.15 Traditional UNIX Kernel

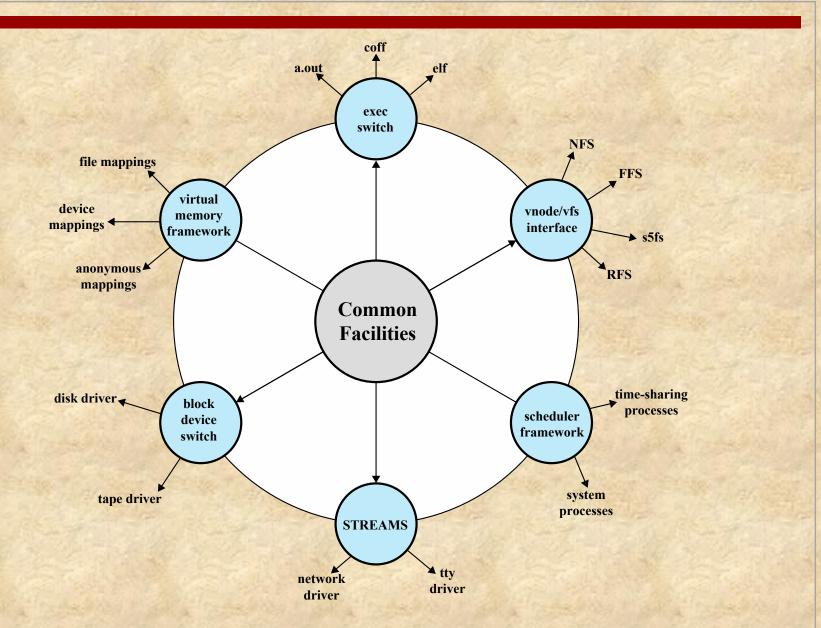


Figure 2.16 Modern UNIX Kernel



## LINUX Overview

- Started as a UNIX variant for the IBM PC
- Linus Torvalds, a Finnish student of computer science, programmed the initial version
- Linux was first posted on the Internet in 1991
- Today it is a full-featured UNIX system that runs on virtually all platforms
- Is free and the source code is available
- Key to the success of Linux has been the availability of free software packages under the auspices of the Free Software Foundation (FSF)
- Highly modular and easily configured

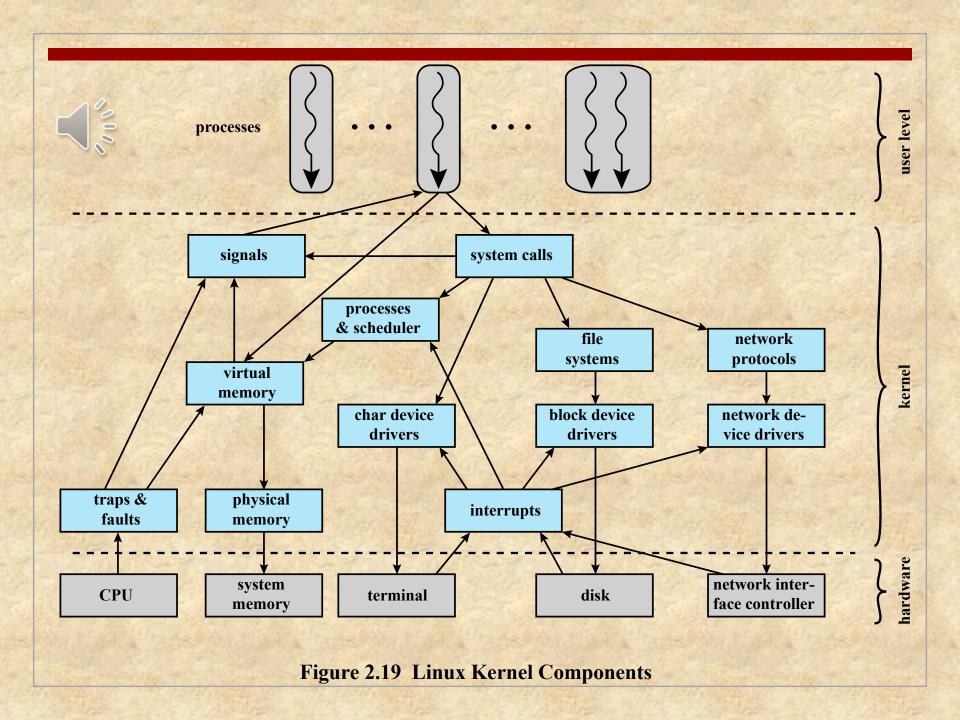


## Modular Structure

- Linux development is global and done by a loosely associated group of independent developers
- Although Linux does not use a microkernel approach, it achieves many of the potential advantages of the approach by means of its particular modular architecture
- Linux is structured as a collection of modules, a number of which can be automatically loaded and unloaded on demand

#### Loadable Modules

- Relatively independent blocks
- A module is an object file whose code can be linked to and unlinked from the kernel at runtime
- A module is executed in kernel mode on behalf of the current process
- Have two important characteristics:
  - Dynamic linking
  - Stackable modules



## Android Operating System

- A Linux-based system originally designed for mobile phones
- The most popular mobile OS
- Development was done by Android Inc., which was bought by Google in 2005
- 1<sup>st</sup> commercial version (Android 1.0) was released in 2008

- Most recent version is Android7.0 (Nougat)
- Android has an active community of developers and enthusiasts who use the Android Open Source Project (AOSP) source code to develop and distribute their own modified versions of the operating system
- The open-source nature of Android has been the key to its success



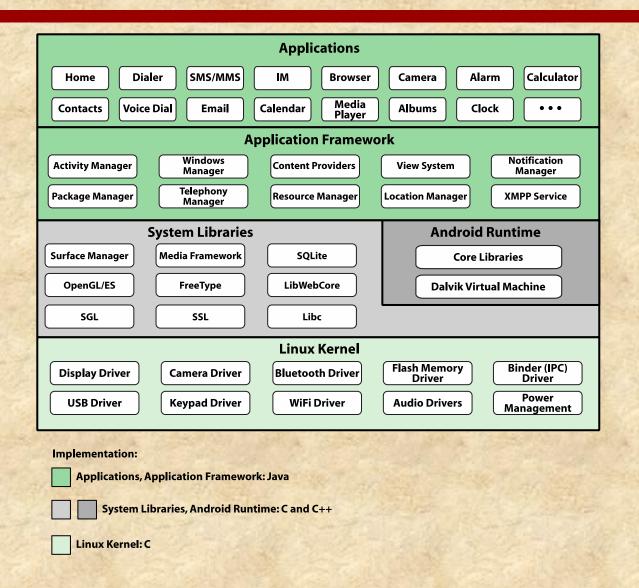




Figure 2.20 Android Software Architecture









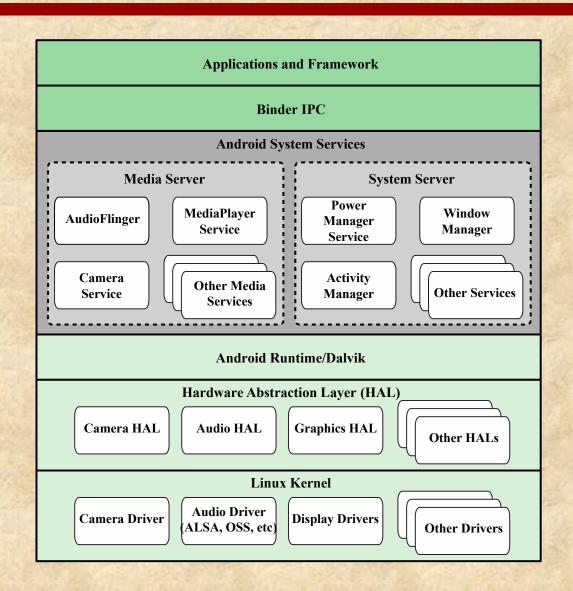


Figure 2.22 Android System Architecture

## Power Management: new features



#### **Alarms**

- Implemented in the Linux kernel and is visible to the app developer through the AlarmManager in the RunTime core libraries
- Is implemented in the kernel so that an alarm can trigger even if the system is in sleep mode
  - This allows the system to go into sleep mode, saving power, even though there is a process that requires a wake up

#### Wakelocks

- Prevents an Android system from entering into sleep mode
- These locks are requested through the API whenever an application requires one of the managed peripherals to remain powered on
- An application can hold one of the following wakelocks:
  - Full\_Wake\_Lock
  - Partial\_Wake\_Lock
  - Screen\_Dim\_Wake\_Lock
  - Screen\_Bright\_Wake\_Lock





## Conclusion of Chapter 2

■ Chapter 2 mainly served the purpose of an "Extended Table of Contents" in our textbook.



- It provided you with an overview of what an operating system is, of what most important components it is made, and what are the most crucial issues which the designer of an operating system must take into consideration.
- Many further details of all this are explained and discussed in the subsequent chapters of the book.