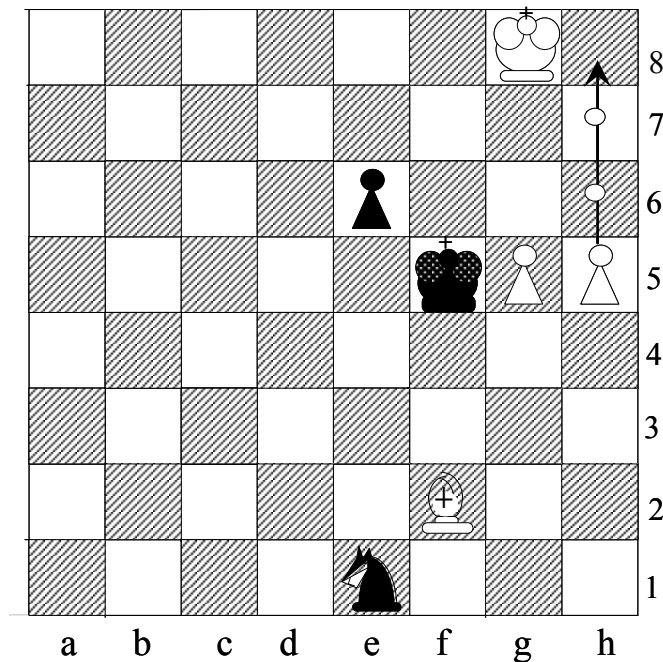


Midterm Exam

Oct 28, 8:00 am - 12:00 pm, LW-836

Midterm Review

1. Represent the system shown below as an ABG.
2. Generate Zone with the main trajectory $a(h5)a(h6)a(h7)a(h8)$. Use grammars of Trajectories and Zones, show the generations with important values of the functions and sets including $DOCK$, med_i , $lmed_i$, SUM , ST_k , $MOVE$, $next_i$, f , h_i^o , h_j , g , v , w , $TIME$, and $NEXTTIME$.
3. Show your understanding of Translations. Show a translation table (analogous to table in handout No. 14) for the main variation (your choice) assuming that Black moves first. Include only the most important trajectories.



Note. On the exam it is going to be a meaningful ABG (which may include non-chess-like pieces). It will be stated as a gaming problem so that generating a zone and moving in it would allow you to solve this ABG.

Class of Problems

Abstract Board Game

is the following eight-tuple

$$\langle X, P, R_p, \{ON\}, v, S_i, S_t, TR \rangle$$

$X = \{x_i\}$ is a finite set of *points*;

$P = \{p_i\}$ is a finite set of *elements*; $P = P_1 \cup P_2, P_1 \cap P_2 = \emptyset$;

$R_p(x, y)$ is a family of binary relations of *reachability* in X
 $(x \in X, y \in X, p \in P)$; y is *reachable* from x for p ;

$ON(p)=x$ is a partial function of *placement* of elements P into X ;

$v > 0$ is a real function, $v(p_i)$ are the *values* of elements;

S_i is a set of *initial* states of the system,
 a certain set of formulas $\{ON(p_i) = x_i\}$;

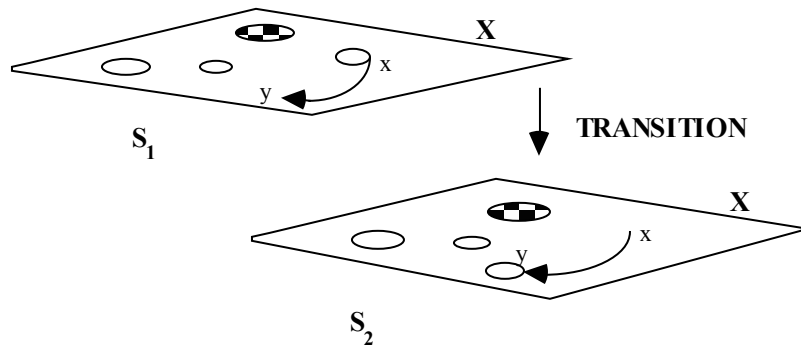
S_t is a set *target* states of the system (as S_i);

TR is a set of operators **TRANSITION**(p, x, y) for transition of the system from one state to another described as follows

precondition: $ON(p) = x \wedge R_p(x, y)$

delete: $ON(p) = x, ON(q) = y$

add: $ON(p) = y$



1	Q_1	$S(u, v, w) \rightarrow A(u, v, w)$	two	\emptyset
2 _i	Q_2	$A(u, v, w) \rightarrow t(h_i^o(u), l_o+1)$ $A((0, 0, 0), g(h_i^o(u), w), zero)$	$TIME(z) = DIST(z, h_i^o(u))$	3 \emptyset
3	Q_3	$A(u, v, w) \rightarrow A(f(u, v), v, w)$	$NEXTTIME(z) =$ $init(u, NEXTTIME(z))$	$four$ 5
4 _j	Q_4	$A(u, v, w) \rightarrow t(h_j(u), TIME(y))$ $A(u, v, g(h_j(u), w))$	$NEXTTIME(z) =$ $ALPHA(z, h_j(u), TIME(y) - l_j+1)$	3 3
5	Q_5	$A(u, v, w) \rightarrow A((0, 0, 0), w, zero)$	$TIME(z) = NEXTTIME(z)$	3 6
6	Q_6	$A(u, v, w) \rightarrow e$		\emptyset \emptyset

$Q_1(u) = (ON(p_0) = x) \wedge (MAP_{x, p_0}(y) \leq l \leq l_0) \wedge (\exists q ((ON(q) = y) \wedge (OPPOSE(p_0, q))))$

$Q_2(u) = T$; $Q_3(u) = (x \neq n) \vee (y \neq n)$ $Q_5(w) = (w \neq zero)$; $Q_6 = T$;

$Q_4(u) = \exists p ((ON(p) = x) \wedge (l > 0) \wedge (x \neq x_0) \wedge (x \neq y_0)) \wedge [(\neg OPPOSE(p_0, p) \wedge (MAP_{x, p}(y) = 1))$
 $\vee (OPPOSE(p_0, p) \wedge (MAP_{x, p}(y) \leq l))]$

$init(u, r) = \begin{cases} 2n, & \text{if } u = (0, 0, 0), \\ r, & \text{if } u \neq (0, 0, 0). \end{cases}$

$f(u, v) = \begin{cases} (x+1, y, l), & \text{if } ((x \neq n) \wedge (l > 0)) \vee ((y = n) \wedge (l \leq 0)) \\ (1, y+1, TIME(y+1) \times v_{y+1}), & \text{if } (x = n) \vee ((l \leq 0) \wedge (y \neq n)). \end{cases}$

Let $t_0 \in L_t^{l_0}(S)$, $t_0 = a(z_0)a(z_1)...a(z_m)$, $t_0 \in t_{p_0}(z_0, z_m, m)$;

If $((z_m = y_0) \wedge (p = p_0) \wedge (\exists k (1 \leq k \leq m) \wedge (x = z_k))) \vee$
 $((z_m \neq y_0) \vee (p \neq p_0)) \wedge (\exists k (1 \leq k \leq m-1) \wedge (x = z_k))$
then $DIST(x, p_0, t_0) = k+1$ else $DIST(x, p_0, t_0) = 2n$

$ALPHA(x, p_0, t_0, k) = \begin{cases} max(NEXTTIME(x), k), & \text{if } (DIST(x, p_0, t_0) \neq 2n) \\ & \wedge (NEXTTIME(x) \neq 2n); \\ k, & \text{if } DIST(x, p_0, t_0) \neq 2n \wedge \\ & (NEXTTIME(x) = 2n); \\ NEXTTIME(x), & \text{if } DIST(x, p_0, t_0) = 2n. \end{cases}$

$g_r(p_0, t_0, w) = \begin{cases} 1, & \text{if } DIST(r, p_0, t_0) < 2n, \\ w_r, & \text{if } DIST(r, p_0, t_0) = 2n. \end{cases}$ $TRACKS_{p_0} = \{p_0\} \times (\cup L[G_t^{(2)}(x, y, k, p_0)])$

If $TRACKS_{p_0} = e$ $1 \leq k \leq l$
then $h_i^o(u) = e$

else $TRACKS_{p_0} = \{(p_0, t_1), (p_0, t_2), \dots, (p_0, t_b)\}, (b \leq M)$ and $h_i^o(u) = \begin{cases} (p_0, t_i), & \text{if } i \leq b, \\ (p_0, t_b), & \text{if } i > b. \end{cases}$

$TRACKS = \cup TRACKS_p$, where $TRACKS_p$ is the same as for h_i^o

If $TRACKS = e$ $ON(p)=x$
then $h_i(u) = e$

else $TRACKS = \{(p_1, t_1), (p_1, t_2), \dots, (p_m, t_m)\}, (m \leq M)$ and $h_i(u) = \begin{cases} (p_i, t_i), & \text{if } i \leq m, \\ (p_m, t_m), & \text{if } i > m. \end{cases}$

At the beginning: $u = (x_0, y_0, l_0)$, $w = zero$, $v = zero$, $x_0 \in X$, $y_0 \in X$, $l_0 \in \mathbf{Z}_+$,

$p_0 \in P$, and $TIME(z) = 2n$, $NEXTTIME(z) = 2n$ for all z from X .

Grammar $G_t^{(2)}$ of shortest and admissible trajectories.

L	Q	Kernel, π_k	π_n	F_T	F_F
1	Q_1	$S(x, y, l) \rightarrow A(x, y, l)$		two	\emptyset
2_i	Q_2	$A(x, y, l) \rightarrow A(x, \text{med}_i(x, y, l), \text{lmed}_i(x, y, l))$ $A(\text{med}_i(x, y, l), y, l - \text{lmed}_i(x, y, l))$		three	three
3_j	Q_3	$A(x, y, l) \rightarrow a(x)A(\text{next}_j(x, l), y, f(l))$		three	4
4	Q_4	$A(x, y, l) \rightarrow a(y)$		three	5
5	Q_5	$A(x, y, l) \rightarrow e$		three	\emptyset

$$V_T = \{a\},$$

$$V_N = \{S, A\},$$

$$V_{PR}$$

$$Pred = \{Q_1, Q_2, Q_3, Q_4, Q_5\}$$

$$Q_1(x, y, l) = (\text{MAP}_{x,p}(y) \leq l < 2\text{MAP}_{x,p}(y)) \wedge (l < 2n)$$

$$Q_2(x, y, l) = (\text{MAP}_{x,p}(y) \neq l)$$

$$Q_3(x, y, l) = (\text{MAP}_{x,p}(y) = l) \wedge (l \geq 1)$$

$$Q_4(y) = (y = y_0)$$

$$Q_5(y) = (y \neq y_0)$$

$$Var = \{x, y, l\};$$

$$Con = \{x_0, y_0, l_0, p\};$$

$$Func = Fcon \cup Fvar;$$

$$Fcon = \{f, \text{next}_1, \dots, \text{next}_n, \text{med}_1, \dots, \text{med}_n,$$

$$\text{lmed}_1, \dots, \text{lmed}_n\} \quad (n = |X|),$$

$$f(l) = l-1, \quad D(f) = \mathbf{Z}_+ \setminus \{0\}$$

functions $\text{next}_i, \text{med}_i$ and lmed_i are defined below.

$$Fvar = \{x_0, y_0, l_0, p\}$$

$E = \mathbf{Z}_+ \cup X \cup P$ is the subject domain;

Parm: $S \rightarrow Var, \quad A \rightarrow Var, \quad a \rightarrow \{x\};$

$$L = \{1, 4\} \cup \text{two} \cup \text{three}, \quad \text{two} = \{2_1, 2_2, \dots, 2_n\}, \quad \text{three} = \{3_1, 3_2, \dots, 3_n\}$$

At the beginning of derivation: $x = x_0, y = y_0, l = l_0, \quad x_0 \in X, y_0 \in X, \quad l_0 \in \mathbf{Z}_+, p \in P.$

Definition of functions *med*, *lmed*, *next*

med_i is defined as follows:

$$D(med_i) = X \times X \times \mathbf{Z}_+ \times P$$

$$DOCK = \{v \mid v \text{ from } X, MAP_{x_0,p}(v) + MAP_{y_0,p}(v) = l\},$$

If

$$DOCK_l(x) = \{v_1, v_2, \dots, v_m\} \neq \emptyset$$

then

$$med_i(x, y, l) = v_i \text{ for } 1 \leq i \leq m \text{ and}$$

$$med_i(x, y, l) = v_m \text{ for } m < i \leq n,$$

otherwise

$$med_i(x, y, l) = x.$$

lmed_i is defined as follows:

$$D(lmed_i) = X \times X \times \mathbf{Z}_+ \times P$$

$$lmed_i(x, y, l) = MAP_{x,p}(med_i(x, y, l))$$

next_i is defined as follows:

$$D(next_i) = X \times \mathbf{Z}_+ \times X^2 \times \mathbf{Z}_+ \times P$$

$$SUM = \{v \mid v \text{ from } X, MAP_{x_0,p}(v) + MAP_{y_0,p}(v) = l_0\},$$

$$ST_k(x) = \{v \mid v \text{ from } X, MAP_{x,p}(v) = k\},$$

MOVE_l(x) is an intersection of the following sets:

$$ST_1(x), ST_{l_0-l+1}(x_0) \text{ and } SUM.$$

If

$$MOVE_l(x) = \{m_1, m_2, \dots, m_r\} \neq \emptyset$$

then

$$next_i(x, l) = m_i \text{ for } i \leq r \text{ and}$$

$$next_i(x, l) = m_r \text{ for } r < i \leq n,$$

otherwise

$$next_i(x, l) = x.$$