

Finite Differencing

$$\frac{\partial f}{\partial x_i} \approx \frac{f(x + \epsilon e_i) - f(x - \epsilon e_i)}{2\epsilon} \quad \text{--- (1)}$$

\*  $\epsilon > 0$  is "small"

\*  $e_i =$  1 in  $i$ -th position only.

$$e_i = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

\* (1) is called the central-difference formula.

Another formula is:

$$\frac{\partial f}{\partial x_i}(x) \approx \frac{f(x + \epsilon e_i) - f(x)}{\epsilon} \quad \text{--- (2)}$$

(2) is called the forward-difference formula.

Note that (1)+(2) refer to one-dimensional functions of  $x_i$ , where all other coordinates are fixed.

## Unit Roundoff (p. 60+)

ch7-2  
24

Represent floating point numbers using  $.d_1 d_2 \dots d_t$  binary digits.

The number is:

$$\sum_{i=1}^t d_i 2^{-i} \times 2^e,$$

where  $e$  is an exponent.

Define  $u = 2^{-t}$  to be the unit roundoff.

Now add 1 to a floating point number:

$$1 + \sum_{i=1}^t d_i 2^{-i} \times 2^e$$

and we don't have enough digits (in general) to represent the number.

Eg:  $t=2$ . Let  $A = (0.11)_2$

$$\Rightarrow A = 1 \times 2^{-1} + 1 \times 2^{-2} = \sum_{i=1}^2 d_i 2^{-i} \times 2^0$$

$$1 + A = (1.11)_2 = (0.111)_2 \times 2^1$$

If we truncate the least-significant bit:

$1+A \cong (0.11)_2 \times 2^{+1}$  with an error  
of  $(2^{-3}) \times 2^{+1} = \underline{\underline{2^{-2}}}$ .

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Now if we apply roundoff:

$1+A = (0.111)_2 \times 2^{+1} \cong 1 \times 2^{+1}$   
with an error of  $(2^{-3} \times 2^{+1}) = \underline{\underline{2^{-2}}}$  again.  
In this example:  $2^{-t}$  is the worst error  
for roundoff arithmetic.

For truncation arithmetic:

$2^{+1} \times (0.00 \underbrace{111 \dots}_{\substack{\text{infinite} \\ \text{number} \\ \text{of 1s}}})_2 = (0.01)_2 \times 2^{+1} = \underline{\underline{2^{-t+1}}}$   
is the worst error.

For determining  $t$ , Dennis & Schnabel  
use the concept of a machine-epsilon (machep)  
which is computed as the smallest  $\tau$   
such that  $1+\tau > 1$ . Here is the code:

```
 $\tau = 1;$   
while  $(1 + \tau * 0.5 > 1)$  do  
     $\tau = 0.5 * \tau;$ 
```

which exits with  $\tau$ .

Note that on some machines the first bit  
omitted, assumed to be 1.

# Floating-Point Arithmetic (Conte & de Boor,

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For an n-digit number:

$$fl(x) = \pm \underbrace{(\cdot d_1 d_2 \dots d_n)}_{\text{mantissa}} \beta^e$$

mantissa,  $\beta=2$  (except calculators!)

and  $e$  is the exponent.

Assume:  $m < e < M$ .

Overflow:  $|x| > \beta^M$ .

Underflow:  $|x| \leq \beta^{m-n}$

Write:  $fl(x) = x(1 + \delta)$

where:

\*  $|\delta| < \frac{1}{2} \beta^{1-n}$  in rounding arithmetic,

\*  $-\beta^{1-n} < \delta \leq 0$  in truncating (chopping).

$$\max_x |\delta| = \underline{\text{unit roundoff}} = u$$

Relationship between actual and floating-point approximations:

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$$x \square^* y = (x \square y) (1 + \delta)$$

for  $\square: +, -, *, /$  operations

and some  $\delta$  with  $|\delta| \leq u$ ,

$\square^*$  = floating point operation,

$\square$  = exact operation.

Eg.  $\underbrace{x + y}_{\text{floating point}} = \underbrace{(x + y)}_{\text{actual}} (1 + \delta), \quad |\delta| \leq u$

Equivalently:

$$\left\{ \begin{array}{l} x \square^* y = (x \square y) / (1 + \delta) \\ \text{for some } \delta \text{ with } |\delta| \leq u \end{array} \right.$$

# Application:

Consider the code:

$$s_0 := c$$

$$s_i := s_{i-1} - a_i b_i, \quad i=1, 2, \dots, r$$

$$s := s_r / a_{r+1}$$

Apply the approximation model:

$$\hat{s}_0 = c$$

$$\hat{s}_i = [\hat{s}_{i-1} - a_i b_i (1+\delta)] (1+\delta)$$

$i=1, 2, \dots, r$

for multiplying

for subtracting

$$\hat{s} = \hat{s}_r / [a_{r+1} (1+\delta)]$$

Note that we have used the same

$\delta$  w/out subscripts. This follows

by first using  $\delta_i$ , and then showing

that this can be simplified to  
this form.

# 1.4 Loss of Significance and Error Propagation

## Condition and instability

ch 7-8  
24

Write:

$$\text{Exact} = \text{approximation} + \text{Error}$$

$$x = x^* + e$$

Relative error is:  $(x - x^*) / x$ .

Loss of significance digits:

$x^*$  approximates  $x$  to  $(r)$  significant  $\beta$ -digits provided that:

$$|x - x^*| \leq \frac{1}{2} \beta^{s-r+1}$$

where: \*  $\beta = 2$  for non-calculators

\*  $s$  is given by:  $\max_s \beta^s \leq |x|$ ,

and  $s$  is integer.

\*  $(r$  is the number of significant digits to be determined.)

E.g.

For  $x^* = 3$ ,  $x = \pi$ ,  $\beta = 10$ .

$\frac{ch 7 \pm 9}{2d}$

$$\Rightarrow \max_s 10^s \leq \pi \Rightarrow s = 0.$$

$$\Rightarrow |3 - \pi| \leq \frac{1}{2} 10^{1-r}$$

$$\text{is } 0.14... \leq \frac{1}{2} 10^{1-r}$$

for  $r = 1$ , but fails for  $r = 2$ .

$$(0.14... \leq \frac{1}{2} \times \frac{1}{10} = 0.05 \text{ fails})$$

$\Rightarrow 3$  agrees to  $\pi$  to one-significant digit.

For  $x^* = 22/7$ , we have agreement to three significant digits (verify).



## Condition

ch 7<sup>+</sup>-10  
24

From:  $f(x) - f(x^*) \approx f'(x)(x - x^*)$ ,

consider the quotient of relative errors:

$$\frac{\left( \frac{f(x) - f(x^*)}{f(x)} \right)}{\left( \frac{x - x^*}{x} \right)} = \frac{f(x) - f(x^*)}{(x - x^*)} \cdot \frac{x}{f(x)}$$

$$\approx \frac{f'(x)}{f(x)} \cdot x \quad \text{for } |x - x^*| \text{ small.}$$

Formally, define:

$$\text{cond}(f)(x) = \max \left\{ \left| \frac{f(x) - f(x^*)}{f(x)} \right| / \left| \frac{x - x^*}{x} \right| \right. \\ \left. \text{such that: } |x - x^*| \text{ is small} \right\}$$

and write:

$$\text{cond}(f)(x) \approx \left| \frac{f'(x) \cdot x}{f(x)} \right|$$

Note that  $\text{cond}(f)$  measures how <sup>ch 7+1</sup> $\frac{24}{24}$  a small initial error gets amplified through a computation:

$$\text{max-relative-error-of}(f)(x) = \text{cond}(f) \cdot \left| \frac{x - x^*}{x} \right|$$

Ex:  $f(x) = \frac{10}{1-x^2}$

$$\Rightarrow \left| \frac{f'(x)x}{f(x)} \right| = \frac{2x^2}{1-x^2} \text{ which is } \underline{\text{unstable}} \text{ for } \underline{x = \pm 1}.$$

For stable computation, require that each step in the computation has a small condition number for every value of  $x$ . Else, we have instability.

Example:

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24

$f(x) = \sqrt{x+1} - \sqrt{x} \sim \otimes$   
is computed using:

$$x_0 := x \Rightarrow f_0(x) = x$$

$$x_1 := x_0 + 1 \Rightarrow f_1(x) = x + 1$$

$$x_2 := \sqrt{x_1} \Rightarrow f_2(x) = \sqrt{x}$$

$$x_3 := \sqrt{x_0} \Rightarrow f_3(x) = \sqrt{x}$$

$$x_4 := x_2 - x_3 \Rightarrow f_4(x, y) = x - y$$

Fix  $x_2$ , to get  $g(t) = x_2 - t$   
This gives:

$$\left| \frac{g'(t) \cdot t}{g(t)} \right| = \left| \frac{t}{x_2 - t} \right|$$

$$\text{For } x = 12345, \quad \left| \frac{t}{x_2 - t} \right| \sim \underline{\underline{22,222}}$$

$$\text{while } \text{cond}(f)(12345) = \left| \frac{f'(x)}{f(x)} \cdot x \right| \approx \underline{\underline{1/2}}$$

$\Rightarrow$  Original  $f$  is ok, but the  
algorithm is 44,000 worse  
in relative error.

This kind of error can usually  $\frac{ch7+-1}{24}$   
be fixed.

Observe:

$$f(x) = (\sqrt{x+1} - \sqrt{x}) \cdot \frac{(\sqrt{x+1} + \sqrt{x})}{(\sqrt{x+1} + \sqrt{x})} = \frac{x+1-x}{\sqrt{x+1} + \sqrt{x}}$$

which gives the following algorithm:

$$x_0 := x$$

$$x_1 := x_0 + 1$$

$$x_2 := \sqrt{x_1}$$

$$x_3 := \sqrt{x_0}$$

$$\begin{cases} x_4 := x_2 + x_3 \end{cases}$$

$$\begin{cases} x_5 := 1/x_4 \end{cases}$$

Now the last two steps compute:

$$g(t) = \frac{1}{(x_2+t)} \Rightarrow \left| \frac{g'(t)t}{g(t)} \right| = \left| \frac{t}{x_2+t} \right|$$

where  $\left| \frac{t}{x_2+t} \right| \approx \frac{1}{2}$  for  $t \approx x_2$ ,

which is more stable, as stable  
as  $f(x)!$  ( $\approx 1/2$ ).

For matrices, we use:

$$\text{cond}(A) = \|A\| \cdot \|A^{-1}\|$$

where:

$$\|A\| = \max_x \frac{\|Ax\|}{\|x\|}$$

if they exist!

For the Euclidean norm (p.173):

$$\|A\|_2 = \max_x \frac{\|Ax\|_2}{\|x\|_2} \quad \left( = \frac{\lambda_1 \leftarrow \text{largest}}{\lambda_n \leftarrow \text{smallest eigenvalue}} \right),$$

$$\|x\|_2 = \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}, \text{ while:}$$

the maximum norm:

$$\|A\|_\infty = \max_x \frac{\|Ax\|_\infty}{\|x\|_\infty}$$

$$\text{where } \|x\|_\infty = \max_{1 \leq i \leq n} |x_i|$$

Eg:  $\left\| \begin{bmatrix} 1 \\ -2 \end{bmatrix} \right\|_\infty = 2 //$

$$\|A\|_\infty = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|$$

p.173

Eg: For  $A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$

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$$\|A\|_{\infty} = \max(1+2, 3+4) = 7$$

Then, similar to before:

for  $Ax=b$ ,  $e = x - \hat{x}$  is the error for  $\hat{x}$  approximating  $x$  and

for the residual error: 
$$\begin{aligned} r &= Ax - A\hat{x} \\ &= Ae \end{aligned}$$

We can show:

$$\frac{1}{\text{cond}(A)} \cdot \frac{\|r\|}{\|b\|} \leq \frac{\|e\|}{\|x\|} \leq \text{cond}(A) \cdot \frac{\|r\|}{\|b\|}$$

↑  
relative  
error in solving  
 $Ax=b$ .

# Determining the step-size for ch 7 I-1b 24 finite-differencing

To provide optimal values for  $\epsilon$  for finite differencing, we will follow:

Conte and deBoor, "Elementary Numerical Analysis", 3rd ed.

Finite-differencing approximates  $f'$  using polynomials, and then bounds the error.

Eg:

Forward-differencing:

$$f'(x) \approx \frac{1}{\epsilon} [f(x+\epsilon) - f(x)]$$

with an error of:  $E(f) = -\frac{1}{2}\epsilon f''(\xi)$ ,  
for some  $\xi \in (x, x+\epsilon)$ .

This is an  $O(\epsilon)$  approximation.

So, if we have a bound for  $f''(\cdot)$ ,  $L$

then:  $|f'(x) - \frac{1}{\epsilon} [f(x+\epsilon) - f(x)]| < \frac{1}{2} \epsilon L$

For central differencing:

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$$f'(x) \approx \frac{1}{2\varepsilon} [f(x+\varepsilon) - f(x-\varepsilon)]$$

with an error of  $E(f) = -\frac{\varepsilon^2}{6} \times f'''(\xi)$

where  $\xi \in (x-\varepsilon, x+\varepsilon)$ . The method is of  $O(\varepsilon^2)$ .

For the second-derivative:

$$f''(x) \approx \frac{1}{\varepsilon^2} [f(x-\varepsilon) - 2f(x) + f(x+\varepsilon)]$$

with an error of:  $E(f) = -\frac{\varepsilon^2}{12} \times f^{(iv)}(\xi)$

where  $\xi \in (x-\varepsilon, x+\varepsilon)$ . It is  $O(\varepsilon^2)$ .

What should  $\varepsilon$  be?

\* We need to take quantization artifacts into account.

Replace  $f(x)$  by  $f(x) + E$ .

for the error in evaluating  $f$ .



Ex: Compute optimal  $\varepsilon$  for  
central-differencing:

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Set:

$$\underset{\substack{\uparrow \\ \text{computer} \\ \text{value}}}{f_c(x-\varepsilon)} = \underset{\substack{\uparrow \\ \text{true} \\ \text{value}}}{f(x-\varepsilon)} + \underset{\substack{\uparrow \\ \text{error in} \\ \text{computation.}}}{E_-}$$

and:  $f_c(x+\varepsilon) = f(x+\varepsilon) + E_+$

Then:

$$f'(x) \approx \frac{1}{2\varepsilon} \left[ (f(x+\varepsilon) + E_+) - (f(x-\varepsilon) + E_-) \right]$$

$$\approx \frac{1}{2\varepsilon} \left[ \underbrace{(f(x+\varepsilon) - f(x-\varepsilon))}_{\text{original error}} + (E_+ - E_-) \right]$$

For the error, consider the original error and the extra error term  $(E_+ - E_-)$ .

The maximum error is thus bounded by

$$\left| -\frac{\varepsilon^2}{6} f'''(\xi) \right| + \frac{1}{2\varepsilon} |E_+ - E_-|$$

Define:

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$$A = \max |E_+ - E_-|,$$

$$B = \max_x |f'''(x)|$$

For the optimal value  $\varepsilon^*$ , we have:

$$\frac{d}{d\varepsilon} \left[ \frac{A}{2\varepsilon} + \frac{B\varepsilon^2}{6} \right] = 0$$

$$\Rightarrow \frac{-A}{2\varepsilon^2} + \frac{2B\varepsilon}{6} = 0$$

$\Rightarrow$  For  $\varepsilon \neq 0$ :

$$-\frac{A}{2} + \frac{B\varepsilon^3}{3} = 0 \Rightarrow \boxed{\varepsilon = \left( \frac{3A}{2B} \right)^{1/3}} \quad (*)$$

\* For infinite precision:

$$A \rightarrow 0 \quad \text{and} \quad \varepsilon \rightarrow 0$$

\* For low precision, then we need to evaluate (\*).

\* Same for small  $f'''(\cdot)$ , occurring for noisy data.

① Approximating the Jacobian

For the residuals:

$$\frac{\partial r}{\partial x_i}(x) \approx \frac{r(x + \epsilon e_i) - r(x)}{\epsilon}$$

② Approximating the Hessian

$$\nabla^2 f(x) e_i = \frac{\partial(\nabla f)}{\partial x_i}(x) \approx \frac{\nabla f(x + \epsilon e_i) - \nabla f(x)}{\epsilon}$$

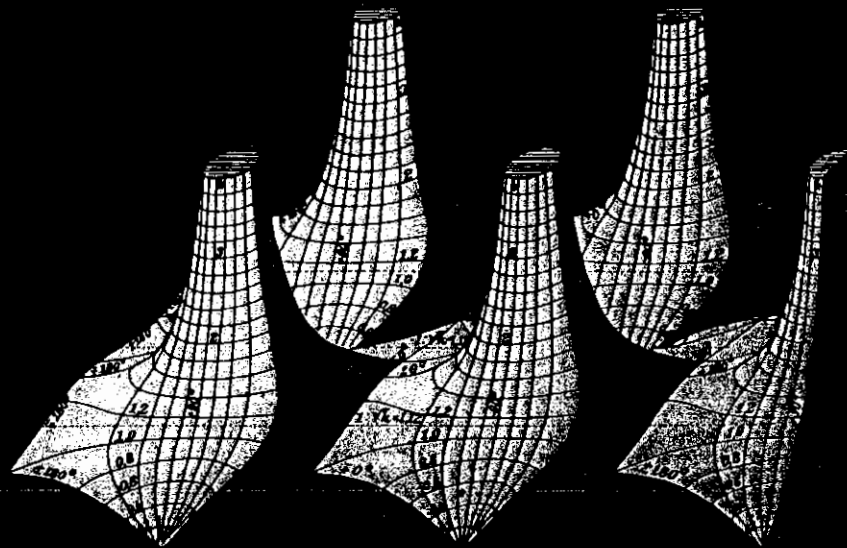
with an error that is  $O(\epsilon)$ Similarly:

$$\nabla^2 f(x) \underline{p} \approx \frac{\nabla f(x + \epsilon \underline{p}) - \nabla f(x)}{\epsilon}$$

that is  $O(\epsilon)$ .Also:

$$\frac{\partial^2 f}{\partial x_i \partial x_j}(x) = \frac{f(x + \epsilon e_i + \epsilon e_j) - f(x + \epsilon e_i) - f(x + \epsilon e_j) - f(x)}{\epsilon^2}$$

with an error of  $O(\epsilon)$



# HANDBOOK OF MATHEMATICAL FUNCTIONS

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Edited by Milton Abramowitz and Irene A. Segun

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25.3.3

$$R'_n(x) = \frac{f^{(n+1)}(\xi)}{(n+1)!} (\xi) \pi'_n(x) + \frac{\pi_n(x)}{(n+1)!} \frac{d}{dx} f^{(n+1)}(\xi)$$

$$\xi = \xi(x) \quad (x_0 < \xi < x_n)$$

Equally Spaced Abscissas

Three Points

25.3.4

$$f'_p = f'(x_0 + ph)$$

$$= \frac{1}{h} \left\{ (p - \frac{1}{2}) f_{-1} - 2p f_0 + (p + \frac{1}{2}) f_1 \right\} + R'_2$$

Four Points

25.3.5

$$f'_p = f'(x_0 + ph) = \frac{1}{h} \left\{ -\frac{3p^2 - 6p + 2}{6} f_{-1} \right.$$

$$+ \frac{3p^2 - 4p - 1}{2} f_0 - \frac{3p^2 - 2p - 2}{2} f_1$$

$$\left. + \frac{3p^2 - 1}{6} f_2 \right\} + R'_3$$

Five Points

25.3.6

$$f'_p = f'(x_0 + ph) = \frac{1}{h} \left\{ \frac{2p^3 - 3p^2 - p + 1}{12} f_{-2} \right.$$

$$- \frac{4p^3 - 3p^2 - 8p + 4}{6} f_{-1} + \frac{2p^3 - 5p}{2} f_0$$

$$- \frac{4p^3 + 3p^2 - 8p - 4}{6} f_1$$

$$\left. + \frac{2p^3 + 3p^2 - p - 1}{12} f_2 \right\} + R'_4$$

For numerical values of differentiation coefficients see Table 25.2.

Markoff's Formulas

(Newton's Forward Difference Formula Differentiated)

25.3.7

$$f'(a_0 + ph) = \frac{1}{h} \left[ \Delta_0 + \frac{2p-1}{2} \Delta_0^2 \right.$$

$$+ \frac{3p^2 - 6p + 2}{6} \Delta_0^3 + \dots + \frac{d}{dp} \binom{p}{n} \Delta_0^n \left. \right] + R'_n$$

25.3.8

$$R'_n = h^n f^{(n+1)}(\xi) \frac{d}{dp} \binom{p}{n+1} + h^{n+1} \binom{p}{n+1} \frac{d}{dx} f^{(n+1)}(\xi)$$

$$(a_0 < \xi < a_n)$$

25.3.9  $hf'_0 = \Delta_0 - \frac{1}{2} \Delta_0^2 + \frac{1}{3} \Delta_0^3 - \frac{1}{4} \Delta_0^4 + \dots$

25.3.10  $h^2 f''_0 = \Delta_0^2 - \Delta_0^3 + \frac{11}{12} \Delta_0^4 - \frac{5}{6} \Delta_0^5 + \dots$

25.3.11

$$h^3 f'''_0 = \Delta_0^3 - \frac{3}{2} \Delta_0^4 + \frac{7}{4} \Delta_0^5 - \frac{15}{8} \Delta_0^6 + \dots$$

25.3.12

$$h^4 f^{(4)}_0 = \Delta_0^4 - 2\Delta_0^5 + \frac{17}{6} \Delta_0^6 - \frac{7}{2} \Delta_0^7 + \dots$$

25.3.13

$$h^5 f^{(5)}_0 = \Delta_0^5 - \frac{5}{2} \Delta_0^6 + \frac{25}{6} \Delta_0^7 - \frac{35}{6} \Delta_0^8 + \dots$$

Everett's Formula

25.3.14

$$hf'(x_0 + ph) \approx -f_0 + f_1 - \frac{3p^2 - 6p + 2}{6} \delta_0^2 + \frac{3p^2 - 1}{6} \delta_1^2$$

$$- \frac{5p^4 - 20p^3 + 15p^2 + 10p - 6}{120} \delta_0^4 + \frac{5p^4 - 15p^2 + 4}{120} \delta_1^4$$

$$+ \dots - \left[ \binom{p+n-1}{2n+1} \right]' \delta_0^{2n} + \left[ \binom{p+n}{2n+1} \right]' \delta_1^{2n}$$

25.3.15

$$hf'_0 \approx -f_0 + f_1 - \frac{1}{3} \delta_0^2 - \frac{1}{6} \delta_1^2 + \frac{1}{20} \delta_0^4 + \frac{1}{30} \delta_1^4$$

Differences in Terms of Derivatives

25.3.16

$$\Delta_0 \approx hf'_0 + \frac{h^2}{2!} f''_0 + \frac{h^3}{3!} f'''_0 + \frac{h^4}{4!} f^{(4)}_0 + \frac{h^5}{5!} f^{(5)}_0$$

25.3.17

$$\Delta_0^2 \approx h^2 f''_0 + h^3 f'''_0 + \frac{7}{12} h^4 f^{(4)}_0 + \frac{1}{4} h^5 f^{(5)}_0$$

25.3.18

$$\Delta_0^3 \approx h^3 f'''_0 + \frac{3}{2} h^4 f^{(4)}_0 + \frac{5}{4} h^5 f^{(5)}_0$$

25.3.19

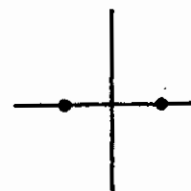
$$\Delta_0^4 \approx h^4 f^{(4)}_0 + 2h^5 f^{(5)}_0$$

25.3.20

$$\Delta_0^5 \approx h^5 f^{(5)}_0$$

Partial Derivatives

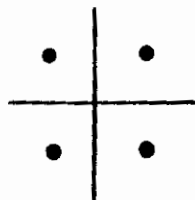
25.3.21



$$\frac{\partial f_{0,0}}{\partial x} = \frac{1}{2h} (f_{1,0} - f_{-1,0}) + O(h^2)$$

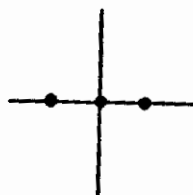
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25.3.22



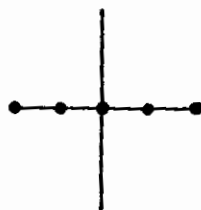
$$\frac{\partial f_{0,0}}{\partial x} = \frac{1}{4h} (f_{1,1} - f_{-1,1} + f_{1,-1} - f_{-1,-1}) + O(h^2)$$

25.3.23



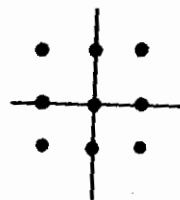
$$\frac{\partial^2 f_{0,0}}{\partial x^2} = \frac{1}{h^2} (f_{1,0} - 2f_{0,0} + f_{-1,0}) + O(h^2)$$

25.3.24



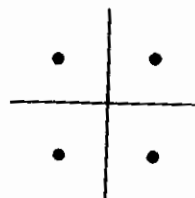
$$\frac{\partial^2 f_{0,0}}{\partial x^2} = \frac{1}{12h^2} (-f_{2,0} + 16f_{1,0} - 30f_{0,0} + 16f_{-1,0} - f_{-2,0}) + O(h^4)$$

25.3.25



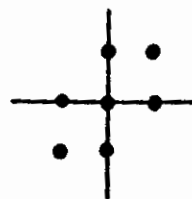
$$\frac{\partial^2 f_{0,0}}{\partial x^2} = \frac{1}{3h^2} (f_{1,1} - 2f_{0,1} + f_{-1,1} + f_{1,0} - 2f_{0,0} + f_{-1,0} + f_{1,-1} - 2f_{0,-1} + f_{-1,-1}) + O(h^2)$$

25.3.26



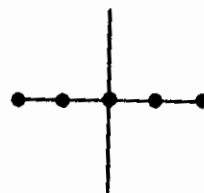
$$\frac{\partial^2 f_{0,0}}{\partial x \partial y} = \frac{1}{4h^2} (f_{1,1} - f_{1,-1} - f_{-1,1} + f_{-1,-1}) + O(h^2)$$

25.3.27



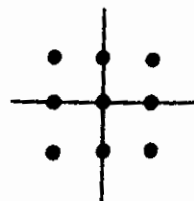
$$\frac{\partial^2 f_{0,0}}{\partial x \partial y} = \frac{-1}{2h^2} (f_{1,0} + f_{-1,0} + f_{0,1} + f_{0,-1} - 2f_{0,0} - f_{1,1} - f_{-1,-1}) + O(h^2)$$

25.3.28



$$\frac{\partial^4 f_{0,0}}{\partial x^4} = \frac{1}{h^4} (f_{2,0} - 4f_{1,0} + 6f_{0,0} - 4f_{-1,0} + f_{-2,0}) + O(h^2)$$

25.3.29



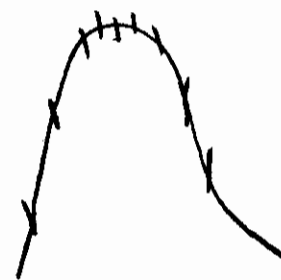
$$\frac{\partial^4 f_{0,0}}{\partial x^2 \partial y^2} = \frac{1}{h^4} (f_{1,1} + f_{-1,1} + f_{1,-1} + f_{-1,-1} - 2f_{1,0} - 2f_{-1,0} - 2f_{0,1} - 2f_{0,-1} + 4f_{0,0}) + O(h^2)$$

# Suggested Approacher

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- \* Experiment with different precisions: float, double, etc.
- \* Experiment with different values for  $\epsilon$ , guided by (\*)
- \* Consider adaptive step methods, using smaller  $\epsilon$  for high curvature points:

(see (\*) ...)



- \* Use Maple/Mathematica for exact derivative evaluation.