

Chapter 18: Sequential Quadratic Programming

- * One of the best for non-linearly constrained opt.
- * Solves quadratic substeps at
- * Works for both line-search and trust-region methods
- * Small or large problems
- * Active set methods for non-linear programming:
 - IQP approach: estimate p and active set together
 - EQP: estimate them separately.
- * Basic ideas:
 - 18.1 for local problems
 - Extend ^{to} "global" using line-search and trust-region
 - attack large problems.

18.1 Local SQP Method

Equality only in:
 $\min f(x)$

$$\begin{bmatrix} c_1(x) \\ c_2(x) \\ \vdots \\ c_m(x) \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \quad \text{18-2} \quad \frac{1}{9}$$

subject to $c(x) = 0$ (several constr.)

where $\left. \begin{array}{l} f: \mathbb{R}^n \rightarrow \mathbb{R} \\ c: \mathbb{R}^n \rightarrow \mathbb{R}^m \end{array} \right\} \text{smooth functions}$

Basic idea:

- Apply Newton to KKT.

KKT: Start with $\mathcal{L}(x, \lambda) = f(x) - \lambda^T c(x)$

Define Jacobian:

$$A(x)^T = [\nabla c_1(x) \quad \nabla c_2(x) \quad \dots \quad \nabla c_m(x)]$$

Note: $\nabla_x (\lambda^T c(x)) = \nabla \left[\sum_{i=1}^m \lambda_i c_i(x) \right]$

$$= \sum_{i=1}^m \lambda_i \nabla c_i(x) = A(x)^T \lambda$$

corrected

Note
the
transposition.

Apply the KKT conditions:

18-3
9

$$\mathcal{L}_x(x, \lambda) = \nabla f(x) - A(x)^T \lambda$$

For the second term, we have: $A^T(x)\lambda$.

$$\mathcal{L}_x(x, \lambda) = \nabla f(x) - A(x)^T \lambda = 0$$

In addition, from the constraints:

$$c(x) = 0.$$

Together, we define $F(x, \lambda)$ by:

$$F(x, \lambda) = \begin{bmatrix} \nabla f(x) - A(x)^T \lambda \\ c(x) \end{bmatrix} = 0. \quad (\Delta\Delta)$$

with $(n+m)$ equations

in $(n+m)$ unknowns x, λ .

For $\textcircled{A\Delta}$, we note that this is 18-4
9
a standard nonlinear problem for which we can use the multivariate Newton method for solving $F(x, \lambda) = 0$.

In 1-D, recall the Newton algorithm:

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

For M-D, it is simply (see chapter 11):

$$x_{k+1} = x_k - [\nabla f(x_k)]^{-1} f(x_k)$$

Or, rewrite this as:

$$x_{k+1} = x_k + p_k, \text{ where}$$

$$\textcircled{*} \left\{ \begin{aligned} [\nabla f(x)] p_k &= -f(x_k). \end{aligned} \right.$$

To apply $\textcircled{*}$ to $F(x, \lambda) = 0$, we have:

$$\nabla F(x, \lambda) = \begin{bmatrix} \nabla_x F_1(x, \lambda) & \nabla_\lambda F_1(x, \lambda) \\ \nabla_x F_2(x, \lambda) & \nabla_\lambda F_2(x, \lambda) \end{bmatrix}$$

We have:

18-5
9

$$\begin{aligned}\nabla_x F_1(x, \lambda) &= \nabla_x (L_x(x, \lambda)) \\ &= \nabla_{xx}^2 L(x, \lambda) \\ &= W(x, \lambda) \text{ by definition of } W(x, \lambda).\end{aligned}$$

$$\nabla_\lambda F_1(x, \lambda) = -A(x)^T$$

$$\nabla_x F_2(x, \lambda) = \nabla_x C(x) = A(x)$$

$$\nabla_\lambda F_2(x, \lambda) = \nabla_\lambda C(x) = 0.$$

We thus write:

$$\begin{bmatrix} x_{k+1} \\ \lambda_{k+1} \end{bmatrix} = \begin{bmatrix} x_k \\ \lambda_k \end{bmatrix} + \begin{bmatrix} p_k \\ p_\lambda \end{bmatrix},$$

where p_k and p_λ solve:

$$\underbrace{\begin{bmatrix} W_k & -A_k^T \\ A_k & 0 \end{bmatrix}}_{[\nabla F(x, \lambda)]} \underbrace{\begin{bmatrix} p_k \\ p_\lambda \end{bmatrix}}_{-F(x, \lambda)} = \underbrace{\begin{bmatrix} -\nabla f_k + A_k^T \lambda_k \\ -C_k \end{bmatrix}}_{-F(x, \lambda)}$$

In chapter 16, "non-singularity" of the solutions is proven, based on:

18- 6
9

$$W_k = \nabla_{xx}^2 L(x, \lambda)$$

Assumption 18.1

- (a) The constraint Jacobian A_k has full row rank.
- (b) The matrix W_k is positive definite for: $d^T W_k d > 0$ for $A_k d = 0, d \neq 0$.

(a) is LICQ, (b) is satisfied at the solution.

SQP Framework

We can define an SQP approach, based on what we just did.

At each iterate, compute (x_k, λ_k) .

Solve: $\min_P \frac{1}{2} P^T W_k P + \nabla f_k^T P$

18-7
9

subject to: $A_k P + C_k = 0.$

To solve this, form $\mathcal{L}(x, \lambda)$ locally:

$$\mathcal{L}(P_k, \mu_k) = \frac{1}{2} P_k^T W_k P_k + \nabla f_k^T P_k - \mu_k^T (A_k P_k + C_k)$$

$$\Rightarrow \mathcal{L}_{P_k}(P_k, \mu_k) = W_k P_k + \nabla f_k - A_k^T \mu_k = 0$$

giving:

$$W_k P_k + \nabla f_k - A_k^T \mu_k = 0$$

$$A_k P_k + C_k = 0$$

which is:

$$\begin{bmatrix} W_k & -A_k^T \\ A_k & 0 \end{bmatrix} \begin{bmatrix} P_k \\ \lambda_{k+1} \end{bmatrix} = \begin{bmatrix} -\nabla f_k \\ -C_k \end{bmatrix} \quad \text{--- } \textcircled{\square}$$

for $\lambda_{k+1} = \mu_k.$

Algorithm 18.1 (Local SQP Algorithm) 18-8 9

Choose an initial pair (x_0, λ_0)

for $k=0, 1, 2, \dots$ $W_k = \nabla_{xx}^2 L_k$

Evaluate $f_k, \nabla f_k, W_k = W(x_k, \lambda_k),$
 C_k , and A_k ;

Solve $\textcircled{\square}$ for P_k, M_k .

$$x_{k+1} = x_k + P_k;$$

If convergence test satisfied

STOP with solution (x_{k+1}, λ_{k+1}) ;

end Update : $x_{k+1} = x_k + P_k, \lambda_{k+1}$

Evaluate $F(x, \lambda)$, and check
if $\|F(x, \lambda)\| \approx 0$.

... Can be extended for
inequalities, line-search,
and trust-region methods.

For inequalities:

18-9
9

$$\min f(x)$$

subject to: $c_i(x) = 0, i \in E$

$$c_i(x) \geq 0, i \in I$$

Linearize the problem to:

$$\min_P f_k + \nabla f_k^T P + \frac{1}{2} P^T \nabla_{xx}^2 f_k P$$

Subject to:

$$\nabla C_i(x_k)^T P + c_i(x_k) = 0, i \in E$$

$$\nabla C_i(x_k)^T P + c_i(x_k) \geq 0, i \in I$$

Solve by methods in chapter 16.