

4) AVL

Insert $O(\log_2 n)$

AVL is slower since AVL is balanced more strictly than RB trees so more rotations are required

Delete $O(\log_2 n)$

Slower than RB for same reason as insert

Search $O(\log_2 n)$

AVL is faster because the balancing is more strict. The number of nodes checked to the furthest node is less than RB tree because of that.

Red Black

$O(\log_2 n)$

Faster than AVL since fewer rotations are done to insert. The black height only has to be the same on each side so the balancing is less strict and less rotations are made.

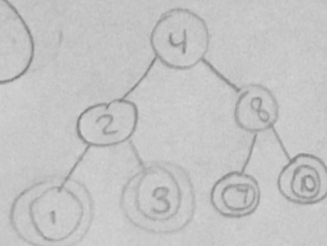
$O(\log_2 n)$

Faster than AVL for same reason as insert.

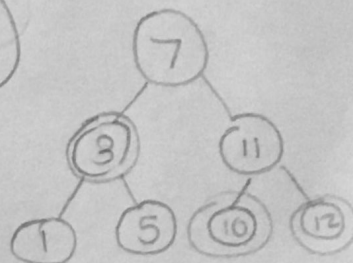
$O(\log_2 n)$

Since the black nodes are what ensures the tree is balanced, the tree can be more skewed if more red nodes are on one side than the other so searching can take a little longer than AVL.

5)



6)



7) `rbtree insert (rbtree* root, int data) {`

`if (root == root->root) {`

`make new node w/ data;`
`node color = black;`
`return new node;`

`}`

`else if (root == NULL) {`

`make new node w/ data;`
`node color = red;`
`return new node;`

`}`

`else {`

`if (root->data > data) {`

`root->left = insert(root->left, data);`

`}`

`else {`

`root->right = insert(root->right, data);`

`}`

`// recolor`

`if (previous color == red && previous->previous color == black) {`

`previous is set to black;`

`previous sibling is set to black;`

`grandparent is set to red;`

`}`

`if (this is root && a ko not black) {`
`set to black;`

`}`

// continued up here

`if (sibling is red and root is red) {`

`parent color = black;`

`grandparent color = red;`

`sibling color = black;`

`}`

`else if (parent is left of grandparent) {`

`if (root is right of parent) {`

`parent color = black;`

`grandparent color = red;`

`rotate right (grandparent);`

`}`

`else if (parent = right of grandparent) {`

`if (root is left of parent)`

`set color of parent to black;`

`set color of parent to red;`

`rotate left (grandparent);`

`}`

`return root;`

`}`