

An operadic model structure on compact closed categories

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Compact Closed Categories

A compact closed category is a symm. mon. cat $(\mathcal{X}, \otimes, I)$ such that

- each object $x \in \mathcal{X}$, has a dual $x^* \in \mathcal{X}$:

- there are distinguished morphisms

- $\cap_x : I \rightarrow x \otimes x^*$

- $\cup_x : x^* \otimes x \rightarrow I$ satisfying

$$\begin{array}{c} x \\ \text{id}_x \otimes \cap_{x^*} \\ \cup_{x^*} \quad \otimes \quad \cap_{x^*} \\ x \end{array} = \begin{array}{c} x \\ | \quad \text{id}_x \\ x \end{array} = \begin{array}{c} x \\ \cap_x \otimes \text{id}_x \\ \text{id}_x \otimes \cup_x \\ x \end{array}$$

Let x_0 denote object monad of x

Trace for $f: x \otimes z \rightarrow y \otimes z$ in \mathcal{X} write

$$\text{tr}_z(f) = (\text{id}_y \otimes \cup_{z^*})(f \otimes \text{id}_z)(\text{id}_x \times \cap_z) : x \rightarrow y$$

write $\text{tr}(x) = \text{tr}_x(\text{id}_x)$.

Note:

In this talk, we assume, for convenience,
that categories $X \in \text{Comp}$ satisfy

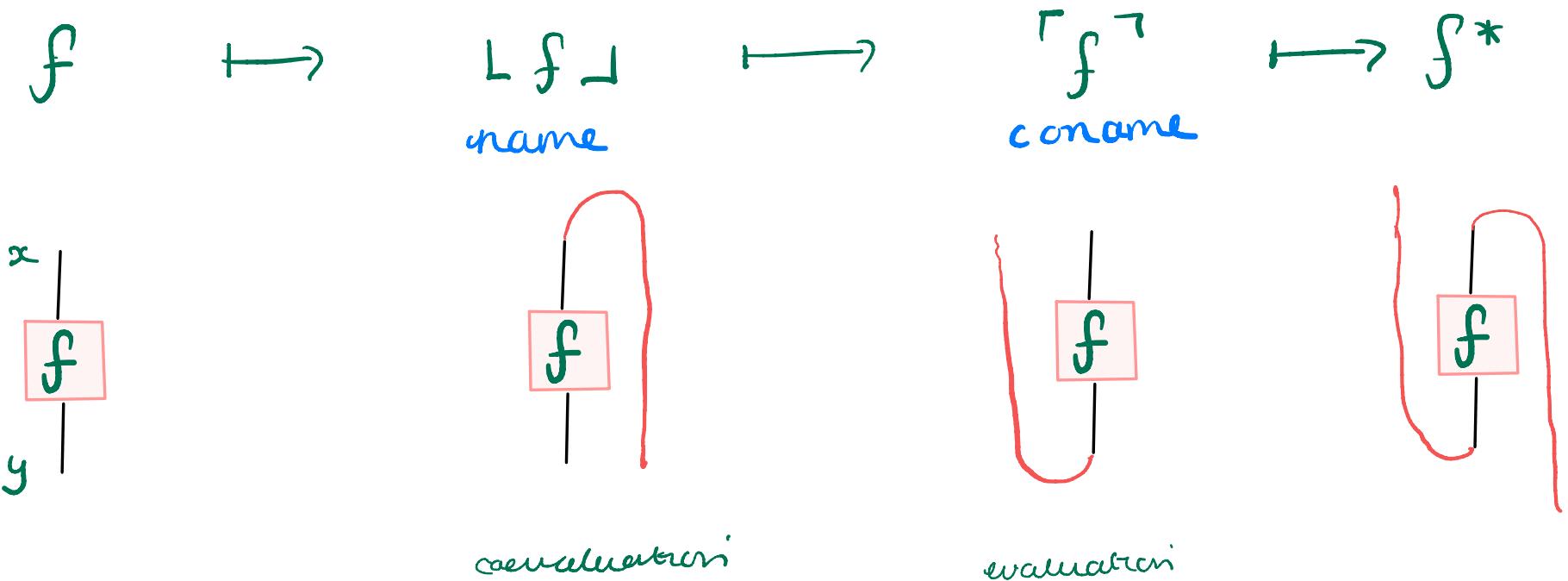
- X strict monoidal
- $\forall x \in X_0, x^{**} = x$.

Internal homs

$$[x,y] = yx^*$$

X compact closed with objects x, y

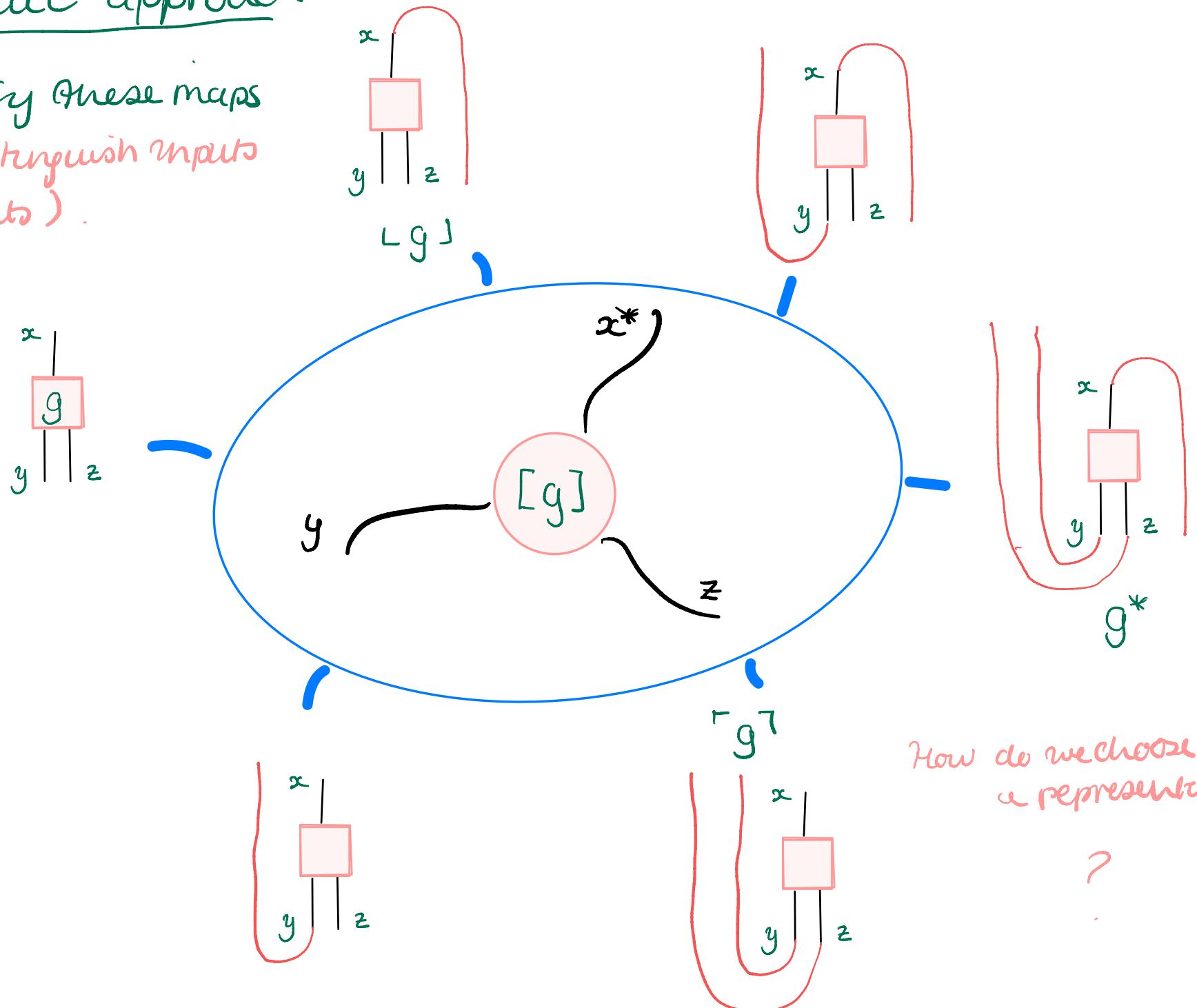
$$X(x, y) \cong X(I, yx^*) \cong X(y^*x, I) \cong X(y^*, x^*)$$



Operadic approach

Identify these maps

(don't distinguish inputs
& outputs)

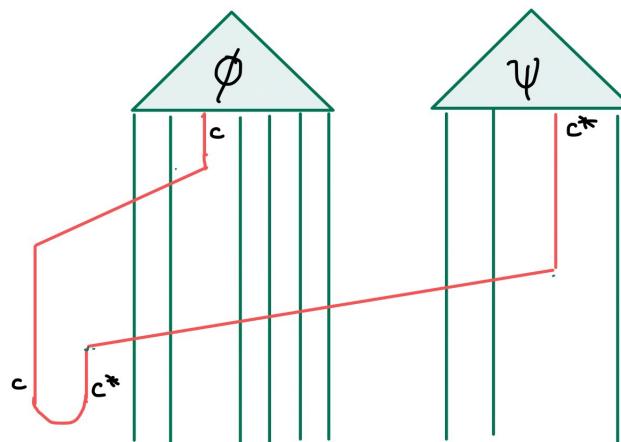
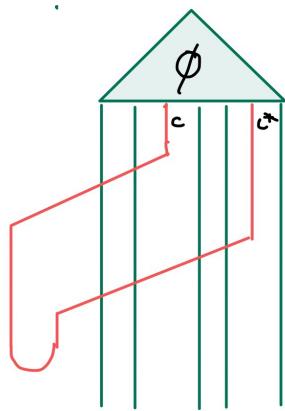


Choice of Representative

We defined the monadic functor $\text{Comp} \rightarrow \text{C}^0$

in terms of the slice I/X of $X \in \text{Comp}$.

Operations defined in terms of composition with cups.



$$\phi, \psi \in I/X.$$

- Which morphisms $F: X \rightarrow Y$ can we rebuild by looking at

$$I_F: I/X \rightarrow I/Y$$

- What does it mean if $I/X \cong I/Y$?

Structure

1. Introduce 'étale' and 'torsion-free' morphisms of compact closed categories
2. Localization: Cat. of fractions from \cap, \cup , trace morphisms + 250s
3. Model Structure - factorisations + lifts.

Étale morphisms

Defn Morphism $F: X \rightarrow Y$ in Comp is étale if

- for all morphisms $g: y_1 \rightarrow y_2$ in Y ,
if $l_{g^{-1}}: I \rightarrow y_2 y_1^*$ is in $\text{Im}^{es}(F)$, then
 g is in $\text{Im}^{es}(F)$.
-

Lemma: F is étale if and only if, for all $y \otimes z \in \text{Im}^{\text{es}}(F)$, there are $w, x \in X_0$ s.t $wx = a$

$\underbrace{y \otimes z}_{\text{if } F(a)}$ and $F(w) \cong y, F(x) \cong z$.

Proof Say $y \otimes z \in \text{Im}^{\text{es}}(F)$. Then so are

$$\eta_{y \otimes z} \cong L[\text{id}_y \otimes \eta_z] \cong L[\eta_y \otimes \text{id}_z].$$

Hence, if F is étale $\text{id}_y \otimes \eta_z$ and $\eta_y \otimes \text{id}_z$, and hence y and z are in $\text{Im}^{\text{es}}(F)$.

Conversely, assume $y \otimes z \in \text{Im}^{\text{es}}(F)$. $\Rightarrow y, z \in \text{Im}^{\text{es}}(F)$ & $y, z \in Y$.

Let $g \in Y(y, z)$ be such that $Lg \in \text{Im}^{\text{es}}(F)$. So
 By assumption $\exists w, x \in X$ and $f \in X(I, xw^*)$ s.t
 $F(f) \cong g$, and $F(w) \cong y, F(x) \cong z$.

But, since $X(I, xw^*) \cong X(w, x)$, there is $f \in X(w, x)$ s.t $F(f) \cong g$.

□

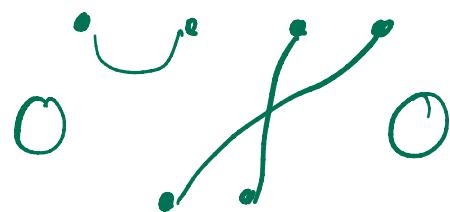
Ex

Let B have objects $n \in \mathbb{N}$

morphisms $m \rightarrow n$

isotopy classes of 1-manifolds

M together w/ $\partial M \xrightarrow{\sim} m+n$



Let $B_2 \hookrightarrow B$ be subcategory on even objects

$\cap \in \text{Im}(\text{incl.})$ but \sqcup is not.

The inclusion $B_2 \hookrightarrow B$ is NOT étale.

Examples

• MOTIVATING EXAMPLE

Category CO of circuit operads embeds in Kleisli category CO_M

of monad $\text{CO} \xrightarrow{\cong} \text{Comp} \simeq \text{CO}^M$.

Morphisms in the image are precisely étale morphisms between free compact closed categories on circuit operads.

These are the arity-preserving morphisms (called étale in \langle Joyal-Kock,R \rangle modular operad terminology.)

- compact closed categories of ring modules

q: $R \rightarrow S$ homomorphism of finitely presented comm. rings

$$\varphi_!: R\text{Mod} \rightarrow S\text{Mod}$$
$$M \mapsto M \otimes_R S \quad \bar{w} \quad (m \cdot r, s) \sim (m, \varphi(r)s)$$

If $\varphi_!$ is étale on Comp then

- φ preserves prime ideals (unramified)

- To do: Study other properties of these morphisms

flatness

infinitesimal lifting etc

Prop

If $F: X \rightarrow Y$ is étale, then

i) F is isofibration $\Leftrightarrow I/F: I/X \rightarrow I/Y$ is isofibration

other

ii) " " equivalence \Leftrightarrow " " " equivalence of categories

Proof

i) Immediate from definitions: $g \in \text{Im}^{\text{es}}(F) = \text{Im}(F) \Leftrightarrow g \circ \in \text{Im}(F)$.

ii) If $F: X \rightarrow Y$ étale and $I/F: I/X \xrightarrow{\cong} I/Y$, then

f.f:

for all $x_1, x_2 \in X$

$$\begin{array}{ccc} I & & \\ \nearrow & \searrow & \\ I & \xrightarrow{\cong} & I \end{array}$$

$$X(I, x_2 x_1^*) \cong Y(I, F(x_2 x_1^*))$$

$\cong Y(F(x_1), F(x_2))$

so F is fully faithful.

I/F is equiv.

ess. surj.: for all $y \in Y$, $\gamma_y = L^{\text{id}_y} \downarrow$ is in $\text{Im}^{\text{es}}(I/F)$.

ob but then id_y , and hence $y \in \text{Im}^{\text{es}}(F)$. $\leftarrow F$ étale.

□.

Trace Localisation

- I) A calculus of left fractions
- II) Trace localisation

I A calculus of left fractions

X compact closed

$W = W_X \subset X$ is wide subcategory with

- all isomorphisms
- all cups, caps, trace: $\forall x \in X_0 \cap_x, \cup_x, \text{tr}(x) \in W$.

Up to iso, a morphism in W has the form

$$w = \text{id}_t \otimes U_x \otimes \cap_y \otimes \text{tr}(z): t \otimes x^* \otimes x \rightarrow t \otimes y \otimes y^*$$

$t, x, y, z \in X_0$

$\left. \begin{matrix} U_x \\ \cap_y \\ \text{tr}(z) \end{matrix} \right\} = \cap_{z^*}$

NOTATION: I'll denote morphisms $w \in W_X(x, y)$ by $x \rightsquigarrow y$.

PROP (X, W) defines a 'calculus of left fractions'

In other words:

F1) For all $f \in X(a, c)$ and $w \in W(a, b)$

there exists • an object $d \in X$.

• morphisms $f' \in X(b, d)$, $w' \in W(b, d)$

such that

$a \xrightarrow{w} b$ in X .

$$\begin{array}{ccc} f & \downarrow & f' \\ a & \xrightarrow{w} & b \\ c & \xrightarrow{w'} & d \end{array}$$

F2) For all $w: a \rightsquigarrow b$, $f_1, f_2: b \rightrightarrows c$ s.t $f_1 w = f_2 w$

there is a $v: c \rightsquigarrow d$ such that $v f_1 = v$

$$a \xrightarrow{w} b \xrightarrow{\frac{f_1}{f_2}} c \rightsquigarrow d$$

PROP (X, W) defines a 'calculus of left fractions'

PROOF

F1) For all $f \in X(a, c)$ and $w \in W(a, b)$, want commuting square f

$$\begin{array}{ccc} a & \xrightarrow{w} & b \\ f \downarrow & \lrcorner & \downarrow f' \text{ in } X. \\ c & \xrightarrow{w'} & d \end{array}$$

given such $f \in X(a, c)$ and $w \in W(a, b)$,

there are objects $t, x, y, z \in X_0$ s.t.

- $a = t x^* x \bullet b = t y y^*$
- $w = \text{id}_t \otimes U_x \otimes \eta_y \otimes \text{tr}(z)$.

And -

$$\begin{array}{ccc} a = t x^* x & \xrightarrow{\text{id}_t \otimes U_x \otimes \eta_y \otimes \text{tr}(z)} & t y y^* = b \\ \downarrow & \lrcorner & \downarrow f \otimes \eta_{x^*} \otimes y^* \\ c & \xrightarrow{\text{tr}(x^* y \otimes z)} & c \end{array}$$

F2) For all $w: a \rightsquigarrow b$, $f_1, f_2: b \rightrightarrows c$ s.t. $f_1 w = f_2 w$ }
want $v: c \rightsquigarrow d$ such that $v f_1 = v f_2$.

$$\begin{array}{ccc} a & \xrightarrow{w} & b \\ & \xrightarrow{f_1} & \xrightarrow{v} c \\ & \xrightarrow{f_2} & \end{array} \rightsquigarrow c \xrightarrow{v} d$$

If $w: a \rightsquigarrow b$, $f_1, f_2: b \rightrightarrows c$ s.t., $f_1 w = f_2 w$ then $f_1 = f_2$ \square .

Trace Localization

Hence, there is a category $X[W^{-1}]$:

- Objects: $X[W_x^{-1}]_0 = X_0$
- Morphisms $[f, w]: a \rightarrow b$: represented by cospans
 $a \longrightarrow b \quad \text{up to equivalence:}$

$(f_1, w_1) \sim (f_2, w_2): a \rightarrow b$ if there is a commuting diagram

$$\begin{array}{ccc} & b_1 & \\ f_1 \nearrow & \downarrow w_1 & \\ a & h_1 \downarrow & b \\ & \cong & \\ & h_2 \uparrow & \\ f_2 \searrow & \swarrow w_2 & \end{array}$$

where

$$h_1 f_1 = h_2 f_2 : a \rightarrow c$$

$$h_1 w_1 = h_2 w_2 : b \rightsquigarrow c$$

$\Rightarrow h_1, h_2 \in W$ since W satisfies 2 out of 3.

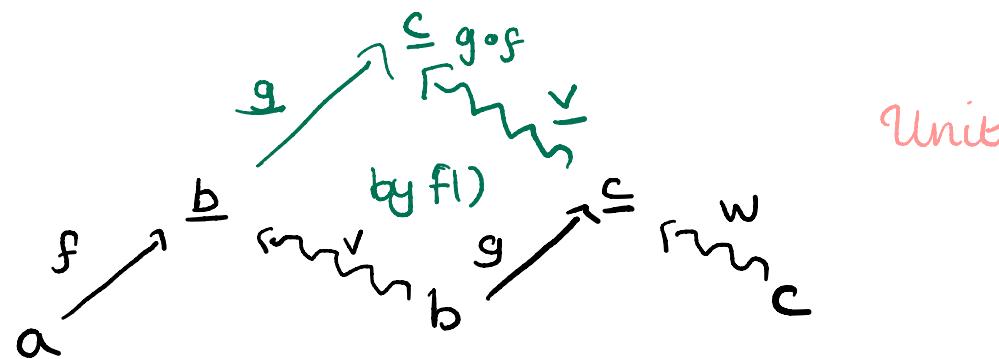
So $(f_1, w_1) \sim (f_2, w_2): a \rightarrow b$ if b_1, b_2 differ up to cups, caps & traces.

Write $[f, w]$ for morphism $a \rightarrow b$ represented by (f, w) .

$X[W^{-1}]$

composition

$$[g, w] \circ [f, v] = [g \circ f, v \circ w]$$



Units

$$a \xrightarrow{id} a = a$$

Trace localisation: $\text{Tr}_X: X \rightarrow X[W_X^{-1}], f \mapsto (f, \text{id})$

Lem
PF

$$\text{For all } x \in X_0, \quad [\cap_x, \text{id}_{xx^*}]^{-1} = [\cup_{x^*}, \text{id}_I]: xx^* \rightarrow I$$

$$I \xrightleftharpoons{\cap_x} I \xrightleftharpoons{xx^*} I \xrightleftharpoons{\cup_{x^*}} I \xrightleftharpoons{I} I$$

since $\cap_x, \cup_{x^*} \in W$. \square .

$X[w^{-}]$ is compact closed

for $[f, w] : a \rightarrow b$, construct

$$\begin{array}{ccc} & b & \\ f \nearrow & \swarrow w & \\ a & & b \end{array}$$

$[f, w]^* = [f^*, w^*] : b^* \rightarrow a^*$ using F1:

$$\begin{array}{ccccc} & b^* & & w^* & \\ & \downarrow & & \swarrow & \\ f^* & & (F1) & & b^* \\ \downarrow & & & & \downarrow \\ a^* & & f_* & & w_* \\ & & \searrow & & \\ & & a^* & & \end{array}$$

$L[f, w] : I \rightarrow ba^*$

$$\begin{array}{ccccc} & ba^* & & w \otimes id_{a^*} & \\ Lf \nearrow & \swarrow & & & \\ I & & ba^* & & \end{array}$$

Lem Let $F: X \rightarrow Y$ be a monomorphism. Then the induced morphism $F_{W^{-1}}: X[W_X^{-1}] \rightarrow Y[W_Y^{-1}]$ is étale $\iff F$ is.

PF • By previous Lemma: If F is étale then $y \otimes z \in \text{Im}^{\text{es}}(F)$, y and z are in the essential image.

So $F: X \rightarrow Y$ étale $\Rightarrow F_{W^{-1}}: X[W_X^{-1}] \rightarrow Y[W_Y^{-1}]$ is.

Conversely. Assume, $y \otimes z \in \text{Im}^{\text{es}}(F_{W^{-1}}) \Rightarrow y, z \in \text{Im}^{\text{es}}(F_{W^{-1}}) \cap \text{gr}(Y)$.

And let $y \otimes z \in \text{Im}^{\text{es}}(F)$. Then $y \otimes z \in \text{Im}^{\text{es}}(F_{W^{-1}})$
 $\Leftrightarrow y, z \in \text{Im}^{\text{es}}(F_{W^{-1}})$. But then $id_y, id_z \in \text{Im}^{\text{es}}(F_{W^{-1}})$

$\Rightarrow id_y, id_z$ and hence $y, z \in \text{Im}^{\text{es}}(F)$
 by definition of localization D

"Torsion-free" morphisms

Let X be compact closed

MX compact closed w

objects $(MX)_0$ is the free monoid $M(X_0)$ on X_0

morphisms $(MX)((x_1, \dots, x_m), (y_1, \dots, y_n)) =$
 $x(x_1 \otimes \dots \otimes x_m, y_1 \otimes \dots \otimes y_n)$

Defn A morphism $F: X \rightarrow Y$ of compact closed categories
is torsion free

if there is a free monad Z_0 and a monoid morphism

$$F_0: MX_0 \xrightarrow{\text{free}} Z_0$$

w.t.

$$\begin{array}{ccc} MX_0 & \xrightarrow{F_0} & Z_0 \\ \downarrow & & \downarrow \\ X_0 & \xrightarrow{F_0} & Y_0 \end{array}$$

as a pushout of
monoid morphisms.

Operadic model structure on Comp

Cofibrations

$\Phi: X \rightarrow Y$ injective on objects
 torsion free

Fibrations

$\Psi: X \rightarrow Y$ isofibration
 étale
 reflect trace morphisms:
 $\text{tr}_z(g) \in \text{Im}^{\text{es}}(\Psi) \Rightarrow g \in \text{Im}^{\text{es}}(\bar{\Psi}).$

Weak equivalences

$F: X \rightarrow Y$ induced functor
 $I/F_{W^{-1}}: I/X[W_X^{-1}] \xrightarrow{\cong} I/Y[W_Y^{-1}]$
 equivalence of categories.

Operadic model structure on Comp

Cofibrations

$I: X \rightarrow Y$ injective on objects
 torsion free

Cofibrant objects:

Compact closed categories X
 ω free object monoid
 (circuit operads!)

Fibrations

$\Psi: X \rightarrow Y$ isofibration
 étale
 reflects trace morphisms

All objects are fibrant

Weak equivalences

$F: X \rightarrow Y$ induced functor
 $I/F_{W^{-1}}: I/X[W_X^{-1}] \xrightarrow{\cong} I/Y[W_Y^{-1}]$
 equivalence of categories.

Prop Trivial fibrations are surjective equivalences of categories.

Proof $F: X \xrightarrow{\cong} Y$ since F étale and reflects trace.

$F: X \rightarrow Y$ surjective since F isofibration.

Factorisation

Let $F: X \rightarrow Y$ be any morphism.

I) There exists a factorisation

$$\begin{array}{ccc} & c & \\ i \nearrow & \swarrow p & \\ X & \xrightarrow{F} & Y \end{array}$$

PROOF (up to symmetry asos).

MY_0 the free object monoid on Y_0 - object set of Y

c has objects $(x, (y_1, \dots, y_n)) \in X_0 \times MY_0$

$$\begin{array}{ccc} MX_0 & \xrightarrow{(-, I)} & MX_0 \times MY_0 \\ \downarrow & & \downarrow \\ X_0 & \xrightarrow{\quad} & X_0 \times MY_0 \end{array}$$

given by pushout

morphisms

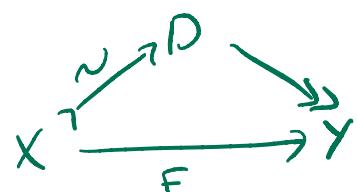
$$c((x, y_1, \dots, y_m), (x', y'_1, \dots, y'_n)) = Y(F(x) \otimes y_1 \otimes \dots \otimes y_m, F(x') \otimes y'_1 \otimes \dots \otimes y'_n)$$

Obvious *, \otimes , etc.

$i: X \rightarrow c: x \mapsto (x, I_Y)$ injective on objects to a free

$p: c \xrightarrow{\sim} Y$: projection surjective equivalence by construction

$$\begin{aligned} p \circ i &= F & x \mapsto (x, I_Y) &\mapsto F(x) \\ (f: x \rightarrow x') &\mapsto (F(f): (x, I_Y) \rightarrow (x', I_Y)) \\ &&&\downarrow \\ &&F(f): F(x) \rightarrow F(x') & \end{aligned}$$



The construction of the trw.cof - fib factorisation is more involved.

I) For $F: X \rightarrow Y$, define $\text{Fib}(F) \hookrightarrow Y$ to be the smallest compact closed subcategory of Y s.t for all morphisms $g \in Y(y_1, y_2)$ of Y , if $Lg \in Y(I, y_2 y_1^*)$ is in $\text{Im}^{es}(F)$, then g is in $\text{Fib}(F)$, and if $\text{tr}_z(g) \in \text{Im}^{es}(F)$, then $g \in \text{Fib}(F)$.

Lemma

The inclusion $\text{Fib}(F) \hookrightarrow Y$ is a fibration

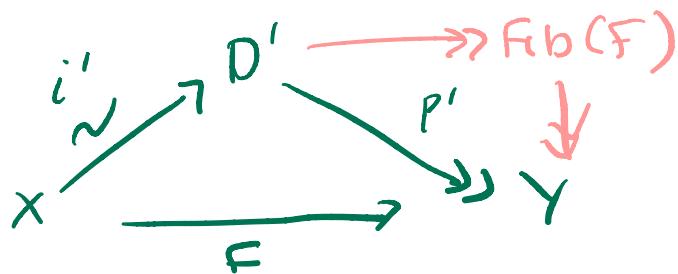
Proof

Definition of $\text{Fib}(F)$

□

We'll build D from $\text{Fib}(F)$ in two stages.

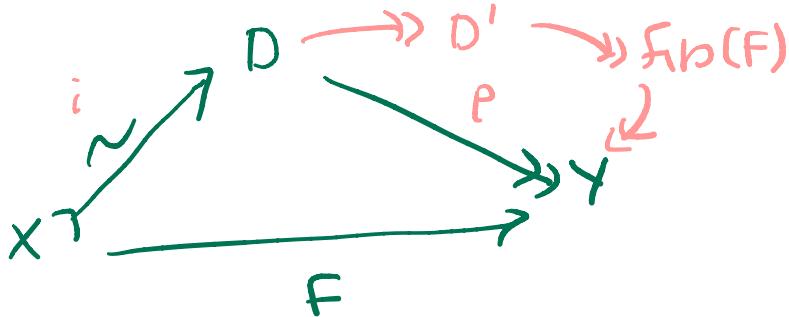
i)



i' weak equivalence
injective on objects

But not necessarily torsion-free

ii)



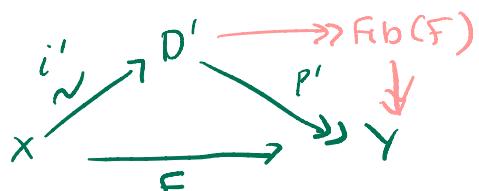
i trivial fibration

RMK

By definition of $\text{Fib}(F)$, for all $y \in \text{Fib}(F)_0$

there exist $x \in X_0$, $\tilde{y} \in \text{Fib}(F)_0$ s.t. $F(x) \cong y \tilde{y}$.

1) Construct factorisation:



D' has

objects: pairs (x, y)

where $x \in X_0$, $y \in \text{Fib}(F)_0$

$\exists \tilde{y} \in \text{Fib}(F)_0$, $y \tilde{y} \cong F(x)$.

morphisms $(x, y) \rightarrow (x', y')$ are pairs $((f_i)_{i=1}^k, g)$ s.t.

- $g \in \text{Fib}(F)(y, y')$
- $f_i \in \text{Mor}(X)$ $1 \leq i \leq k$ and $\text{dom}(f_i) = x$, $\text{cod}(f_k) = x'$
- there are $w, v \in \text{Mor}(W_{\text{Fib}(F)}) \subset \text{Mor}(W_B)$
 • $g(f) \stackrel{\text{def}}{=} w \circ F(f) \circ v \in \text{Fib}(F)(y, y')$

and • $g \cong g(f)$ in $I/\text{Fib}(F)[W_{\text{Fib}(F)}^{-1}]$

- k, q is minimal for these properties:

None of the f_i are (vertically) composable in X s.t. properties preserved.

$\text{dom}(v) = y$,
 $\text{cod}(w) = y'$ s.t.

where

$f \stackrel{\text{def}}{=} \bigotimes_{i=1}^k f_i \in \text{Mor}(X)$

Sorry, what?

morphisms $(x, y) \rightarrow (x', y')$ are pairs $((f_i)_{i=1}^k, g)$ s.t

- $g \in \text{Fib}(F)(y, y')$
- $f_i \in \text{Mor}(X) \quad 1 \leq i \leq k \text{ and } \text{dom}(f_i) = x, \text{ cod}(f_k) = x'$

- there are $w, v \in \text{Mor}(W_{\text{Fib}(F)}) \subset \text{Mor}(W_B)$ s.t
 $\text{dom}(v) = y, \text{ cod}(w) = y'$

$$\bullet g(f) \stackrel{\text{def}}{=} w \circ F(\underline{f}) \circ v \in \text{Fib}(F)(y, y')$$

$$f \stackrel{\text{def}}{=} \bigotimes_{i=1}^k f_i \in \text{Mor}(X)$$

and • $g \cong g(f)$ in $I/F_{\text{fib}}(F)[W_{\text{Fib}(F)}^{-1}]$

- $k \geq 1$ is minimal for these properties:

None of the f_i are (vertically) composable in X s.t properties preserved.

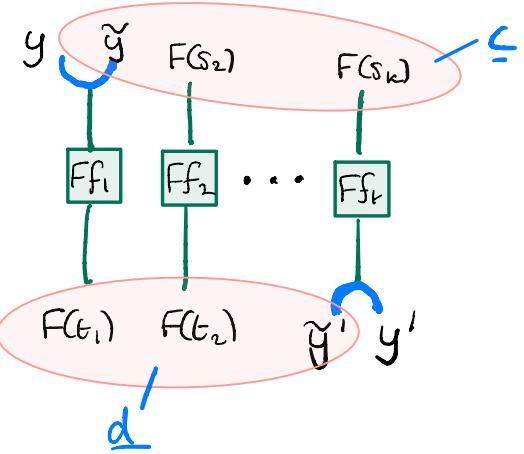
- there are $w, v \in \text{Mor}(W_{\text{Fib}(F)}) \subset \text{Mor}(W_B)$ $\text{dom}(v) = y$, s.t
 $\text{cod}(w) = y'$
 $\bullet g(f) \stackrel{\text{def}}{=} w \circ F(f) \circ v \in \text{Fib}(F)(y, y')$
 - and $\bullet g \cong g(f)$ in $I/\text{Fib}(F)[W_{\text{Fib}(F)}^{-1}]$
- $f \stackrel{\text{def}}{=} \bigotimes_{i=1}^k f_i \in \text{Mor}(X)$

Exercise to check at $Lf \in I/X$, $Lw \circ F(f) \circ v \in F/Y$.

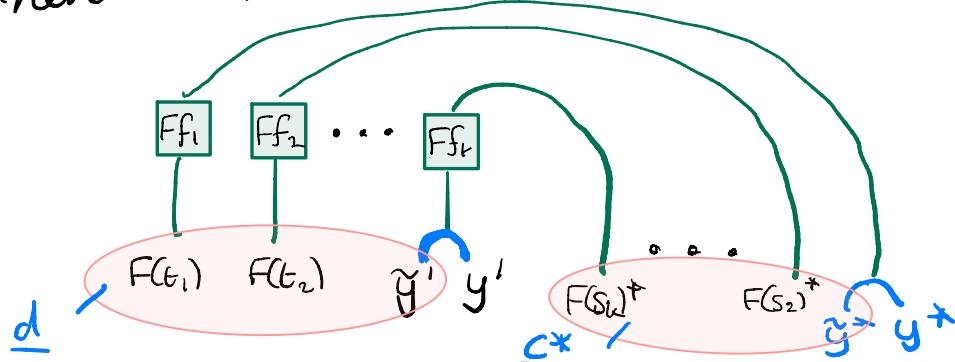
for $1 \leq i \leq k$ let $s_i = \text{dom}(f_i)$, $t_i = \text{cod}(f_i)$ $s_1 = x$, $t_k = x'$
let $y \tilde{y} \cong F(x)$, $F(x') \cong \tilde{y}' y'$ in $\text{Fib}(F) \circ$.
We'll assume $y \tilde{y} = F(x)$, $F(x') = \tilde{y}' y'$ for convenience.

Let $\underline{c} = \tilde{y} \otimes (\bigotimes_{i=2}^k F(s_i))$, $\underline{d} = \left(\bigotimes_{j=1}^{k-1} F(t_j) \right) \otimes \tilde{y}'$ in $\text{Fib}(F) \circ$.

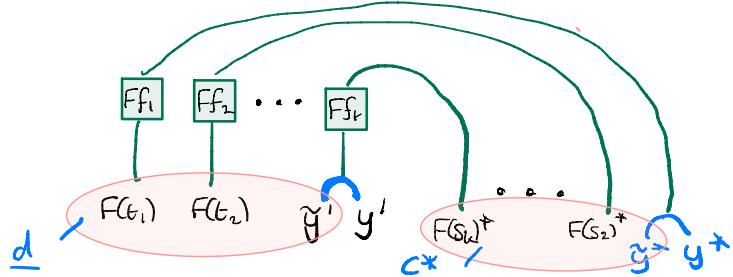
Then $F(f) : y \otimes \underline{c} \longrightarrow \underline{d} \otimes y$ in $\text{Fib}(F)$



and hence $L F(f) : I_Y \longrightarrow \underline{d} \otimes y \otimes \underline{c}^* \otimes y^*$



$$L^F(\underline{f}) \downarrow : I_y \longrightarrow \underline{d} \otimes y \otimes \underline{c}^* \otimes y^*$$



Since $g(\underline{f}) \stackrel{\text{def}}{=} w \circ F(\underline{f}) \circ v \in \text{Fib}(F)(y, y')$, it must be the case that
 $Lg(F) \downarrow = (id_{y, y^*} \otimes V_b \otimes \text{tr}(z)) \circ L^F(\underline{f}) \downarrow \in \text{Fib}(F)(I, y' y^*)$

hence

$$\underline{d} \underline{c}^* = \underline{b}^* \underline{b}$$

Can rewrite

- there are $w, v \in \text{Mor}(W_{\text{Fib}(F)}) \subset \text{Mor}(W_B)$ $\text{dom}(v) = y$, s.t $\text{cod}(w) = y'$
- $g(\underline{f}) \stackrel{\text{def}}{=} w \circ F(\underline{f}) \circ v \in \text{Fib}(F)(y, y')$
- and • $g \cong g(\underline{f})$ in $I/\text{Fib}(F)[W_{\text{Fib}(F)}^{-1}]$

$$\underline{f} \stackrel{\text{def}}{=} \bigotimes_{i=1}^k f_i \in \text{Mor}(X)$$

as

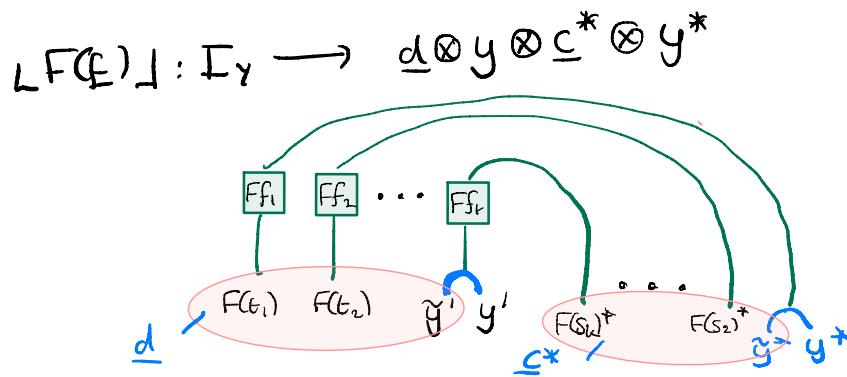
- there is $b \in \text{Fib}(F)_0$, and $g(\underline{f}) \cong g$ in $I/\text{Fib}(F)[W_{\text{Fib}(F)}^{-1}]$
such that $Lg(\underline{f}) \downarrow \cong (id_{y, y^*} \otimes V_b) \circ L^F(\underline{f}) \downarrow$.

$((f_i)_{i=1}^k, g) : (x, y) \rightarrow (x', y')$ in D'

- $g \in \text{Fib}(F)(y, y')$
- $f_i \in \text{Mor}(X)$ $1 \leq i \leq k$ and $\text{dom}(f_i) = x$, $\text{cod}(f_k) = x'$
- there is $b \in \text{Fib}(F)_0$, and $g(f) \cong g$ in $I/\text{Fib}(F)[W_{\text{Fib}(F)}^{-1}]$
such that $\lfloor g(f) \rfloor \cong (id_y \otimes \cup_b) \circ F \lrcorner$.

• $k \geq 1$ is minimal for these properties:

None of the f_i are (vertically) composable in X s.t properties preserved



If any \star

$$\begin{cases} F(s_i) = F(s_j)^* & 1 < i < j \leq k \\ F(s_i) = F(t_j) & 1 < i < j < k \\ F(t_i) = F(s_j) & 1 \leq i \leq j \leq k \\ F(t_i) = F(t_j)^* & 1 \leq i < j < k \end{cases}$$

then can compose f_i, f_j in X .

So condition says that there are no i, j s.t \star .

actually, it says a bit more ...

Lem Let $y, y' \in \text{Im}^{\text{es}}(F)$. Then

$((f_i)_{i=1}^k, g) : (x, y) \rightarrow (x', y')$ in D' has form
 $\cdot (f, g)$ with $F(f) \cong g$ in $\text{Fib}(F)[\text{W}_{\text{Fib}(F)}]$.

PF

Let $y\tilde{y} \cong F(x)$, then $\tilde{y} \in \text{Im}^{\text{es}}(F)$:

Namely $\text{id}_{F(x)} \cong \text{id}_{y\tilde{y}}$

$$\begin{array}{ccc} \text{tr}_y & \text{id}_{F(x)} & \cong \\ \nearrow & \curvearrowright & \downarrow \\ \text{Im}^{\text{es}}(F) & & \text{Im}^{\text{es}}(F) \end{array}$$

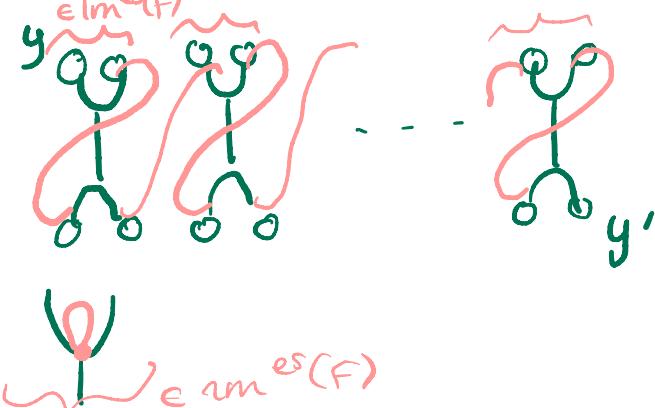
$\cong \text{tr}(y) \otimes \text{id}_{\tilde{y}}$.

Likewise \tilde{y}' .

Hence if $\lfloor g(f) \rfloor \cong (\text{id}_{y'} y^* \otimes \cup_b) \circ F \vdash \rfloor$

then $b \in \text{Im}^{\text{es}}(F)$ and so f_i 's composable in X .

Some examples of morphisms in D' :



Morphism $p': D' \rightarrow Y$ $p': D' \rightarrow \text{Fib}(F) \hookrightarrow Y$
 $(x, y) \mapsto y, (f_i)_i, g \mapsto g$

étale, iso-fibration, reflects trace by construction
 fibration

Morphism $i': X \rightarrow D'$ $x \mapsto (x, F(x))$ $f \mapsto (f, F(f))$

clearly injective on objects.

Must show

$\Gamma = I/i'_{w^{-1}} : I/X[w_x^{-1}] \rightarrow F/D[w_0^{-1}]$ is equivalence

PROP

$$\Gamma = I/i'_{W^{-1}} : I/X[W_X^{-1}] \rightarrow I/D[W_D^{-1}] \text{ is equivalence}$$

PF

faithful: Let $[I \xrightarrow{\alpha_i} x_i \xleftarrow{v_i} x_i] \in X[W_X^{-1}] (I, x_i)$ for $i=1,2$

A morphism $I/X[W_X^{-1}] ([\alpha_1, v_1], [\alpha_2, v_2])$ has a representative of the form:

$$\begin{array}{ccc} \alpha_1 & \downarrow I & \alpha_2 \\ x_1 & \nearrow f & \searrow x_2 \\ v_1 & & v_2 \\ x_1 & & x_2 \end{array} \quad \text{by}$$

$$\begin{array}{ccccc} \alpha_2 & \nearrow \alpha'_1 & x_2 & \nearrow x_2 \\ \alpha_1 & \nearrow x_1 & f & \nearrow x_1 \\ I & & f' & \nearrow x_2 \\ & & & x_2 \end{array}$$

In particular it follows that distinct morphisms in $I/X[W_X^{-1}] ([\alpha_1, v_1], [\alpha_2, v_2])$ have distinct images in $I/D[W_D^{-1}] ([(\alpha_1, F(\alpha_1)), (v_1, F(v_1))], [(\alpha_2, F(\alpha_2)), (v_2, F(v_2))])$.

fullness

Let

$$[f, g] \in I/D[W_D^{-1}] ([(\alpha_1, F(\alpha_1)), (v_1, F(v_1))], [(\alpha_2, F(\alpha_2)), (v_2, F(v_2))])$$

$$\begin{array}{ccc} (\alpha_1, F(\alpha_1)) & \downarrow I & (\alpha_2, F(\alpha_2)) \\ (x_1, F(x_1)) & \nearrow f & (x_2, F(x_2)) \\ (v_1, F(v_1)) & & (v_2, F(v_2)) \\ (x_1, F(x_1)) & \nearrow g & (x_2, F(x_2)) \end{array}$$

be any monomorphism. Then $g = f$ in $I/D[W_D^{-1}]$
by definition

PROP

$$\Gamma = I/i_{w^{-1}} : I/X[w_x^{-1}] \rightarrow F/D[w_0^{-1}] \text{ is equivalence}$$

Pf

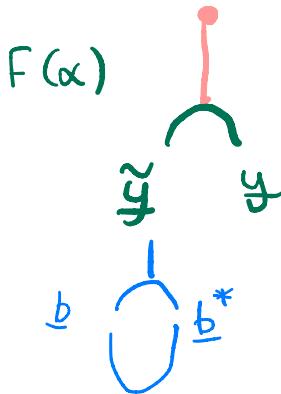
ess. says on obj.

Let $I \xrightarrow{\quad} (\underline{x}, y) \xleftarrow{\text{ess.}} (\underline{x}, y)$ represent an object of $I/D[w_0^{-1}]$

Since $\text{dom } (\alpha_i) = I_x$ for all $1 \leq i \leq k$, $k=1$ and $(\alpha_i)_{i=1}^k = \alpha$.

Then there is a $\tilde{y} \in \text{Fib}(F)$ with $\tilde{y} \cong F(\underline{x})$

Moreover, there is $\underline{b} \in X$. $(U_{F(\underline{b})} \otimes \text{id}_{\tilde{y}}) \circ F(\alpha) \cong \beta : I \rightarrow \tilde{y}$ in $D'[w_0^{-1}]$
 therefore $\tilde{y} \cong F(\underline{b}\underline{b}^*)$.



But then $(\alpha, \beta) \cong (\alpha, F(\alpha))$ in $I/D'[w_0^{-1}]$ by

□

So we have factorisation

$$\begin{array}{ccc} & \nearrow D' & \\ X & \xrightarrow[F]{} & Y \end{array}$$

Then construct $D \cong D'$
 objects $(x, \underbrace{y_1, \dots, y_n}_{\text{"M(Fib(F))}})$ s.t. $F(x) \cong y_1 \otimes \dots \otimes y_n \otimes \tilde{y}$
 for some $\tilde{y} \in \text{Fib}(F)$.

morphisms $D((x, y_1, \dots, y_n), (x', y'_1, \dots, y'_n))$
 $= D'((x', y_1 \otimes \dots \otimes y_n), (x', y'_1 \otimes \dots \otimes y'_n))$.

Then D_0 is given by pushout

$$\begin{array}{ccc} Mx_0 & \longrightarrow & Mx_0 \times M\text{Fib}(F)_0 \\ \downarrow & & \downarrow \\ x_0 & \longrightarrow & D_0 \end{array}$$

so have factorisation

$$\begin{array}{ccc} & \nearrow D & \\ X & \xrightarrow[F]{} & Y \end{array}$$

Lifts

Let

$$\begin{array}{ccc} X & \xrightarrow{F} & A \\ \Phi \downarrow y_2 & \nearrow \text{red dashed} & \downarrow \Psi \\ Y & \xrightarrow{g} & B \end{array} \quad \text{commute in } \underline{\text{Comp}}$$

Let $g \in Y(y_1, y_2)$ not in $\text{Im}(\Phi)$. Since Φ is w.e., there exist $x, x' \in X$ and $f' \in X(I, x)$, $v \in X(x, x')$ s.t. there is a commuting diagram in Y :

$$\begin{array}{ccccc} I & \xrightarrow{\Phi(f')} & \Phi(x) & \xleftarrow{\Phi(v)} & \Phi(x') \\ \parallel & & & & \parallel z \\ I & \xrightarrow{g} & y_2 y_1^* & = & y_2 y_1^{*\dagger} \end{array}$$

In particular, by defn of W_y , there is a z (possibly $z = I$) in Y such that $\Phi(f') \cong L \text{tr}_z(g) \perp$.

It follows that

$\Psi F(f') \cong GLtr_z(g) \hookrightarrow$ and therefore, since $\bar{\Psi}$ is a fibration, $g \in \text{Im}(\bar{\Psi})$, and there are $a_1, a_2 \in A$ and $f \in A(a_1, a_2)$ s.t

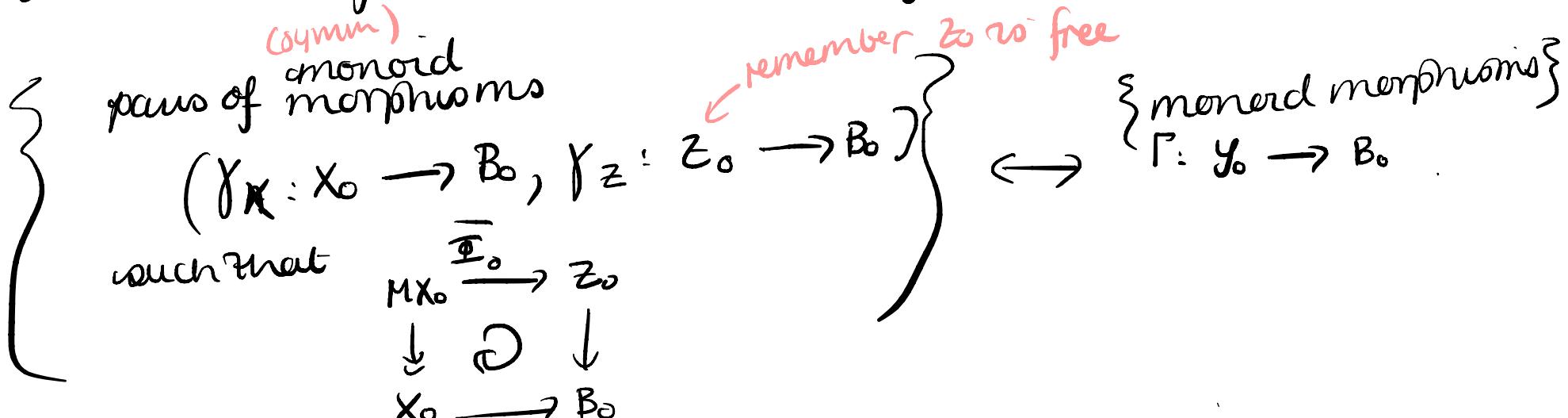
$$B(\bar{\Psi}(a_1), \bar{\Psi}(a_2)) \ni \bar{\Psi}(f) = g \in B(g(y_1), g(y_2))$$

In particular, for each $y \in Y$, may take $g = \text{id}_y$.

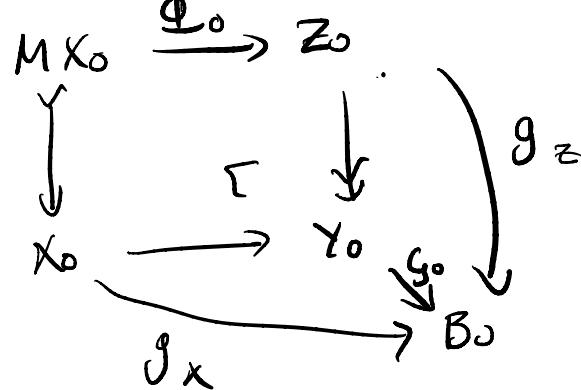
So, we can construct a lift $Y \rightarrow A$ in the category $\underline{\text{CAT}}$ of small categories

However Still need to prove that there is a lift in Comp.

Since Φ is a cofibration, there is a bijection



Let $g_x: x_0 \rightarrow B_0$, $g_z: Z_0 \rightarrow B_0$ be the pair of morphisms such that



commutes.

We must construct a monoidal functor

$$Z_0 \xrightarrow{L_2} A_0 \text{ such that } \bar{\Psi}_0 \circ L_2 = g_{M\bar{A}}: Z_0 \rightarrow B_0$$

Since Z_0 is free and $\bar{\Psi}$ reflects monoidal products, we may define L_2 on singletons such that duals are preserved. Hence have you got morphisms?

\therefore desired not

Let $\begin{array}{ccc} X & \xrightarrow{F} & A \\ \Phi \downarrow & s \downarrow \bar{\Psi} & \text{commute in } \underline{\text{Comp}} \\ Y & \xrightarrow{g} & B \end{array}$

Since Φ is a trivial fibration, it is a surjective equivalence of categories.

Hence for each $g \in \mathcal{Y}(y_1, y_2)$ $g \notin \text{Im}(\Phi)$, and all $a_1, a_2 \in A$
 $\text{s.t. } \bar{\Psi}(a_i) = g(y_i) \quad (i \in \{1, 2\})$, there is a unique $f \in A(a_1, a_2)$

s.t $g(g) = \bar{\Psi}(f)$.

Then, since Φ is a cofibration, we can construct a lift
 $Z_0 \rightarrow A_0$ as before, and hence the desired lift

$$L: Y \rightarrow A$$

Weak equivalences of cofibrant objects

Since all objects of Comp are fibrant,
fibrant-cofibrant objects for this model structure
are precisely those that have free object monad.

By a previous talk, these are in the image
of the free functor $\text{Co} \xrightarrow{\sim} \text{Comp}$.

Indeed, straightforward to show that a w.e. of cofibrant
objects is induced by an equivalence of circuit
operads.

Example

If $X \in \text{Comp}$ has a 0-object 0 , then
 $X \xrightarrow{*}$ is a weak equivalence in the
trivial compact closed category
operadic model structure

Next time

- Homotopy category Ho Comp
- Enriching the model structure ..

Thank you!