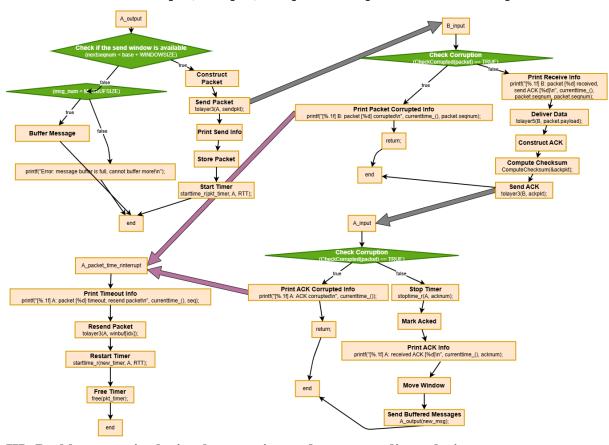
## IERG3310 Lab1 Report

## SU Caiyi 1155191405

I. The environment for compiling my codes

```
PS C:\Users\苏采奕> gcc --version
gcc.exe (MinGW.org GCC Build-2) 9.2.0
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This is free software; see the source for copying conditions. There is NO warranty; not even for MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE.
```

II. Flow chart of A\_output, A\_input, B\_input and A\_packet\_time\_rinterrupt



## III. Problems met in the implementation and corresponding solutions

**Problem**: After my code can successfully test the two examples given, I tried to test other Enter channel pattern strings. I found that if I set the Enter channel pattern string input to put x and - at the beginning of the information transmission, such as "xxxx----oooooooooo", my output will show that my packet cannot be sent out completely (as shown in the figure below)

```
IERG3310 labl: Reliable Data Transfer Protocol-SR, implemented by 1155191405

Enter the number of messages to simulate:

Seter time_ between messages from sender's layer5 [ > 0.0]:

Enter channel pattern string

XXXX———osoposoposo

Enter channel pattern string

XXXX——osoposoposo

Enter sender's window size

Packet Transmission Log

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1.0. A: send packet [0] base [0]

1.0. A: send packet [1] base [0]

1.0. A: send packet [2] base [0]

1.0. A: buffer packet [3] base [0]

1.1. A: packet [1] time out, resend packet

1.1. A: buffer packet [3] base [0]

1.2. A: packet [1] time out, resend packet

1.2. B: packet [99999] corrupted
```

```
(32.0) A: packet [0] time_out, resend packet
[34.0) A: packet [1] time_out, resend packet
[35.0] A: packet [2] time_out, resend packet
[35.0] A: packet [999999] corrupted
[40.5] B: packet [999999] corrupted
[40.5] B: packet [2] received, send ACK [2]
[47.0] A: packet [0] time_out, resend packet
[40.0] A: packet [1] received, send ACK [1]
[40.0] A: received ACK [2]
[53.5] B: packet [1] received, send ACK [1]
[55.5] B: packet [1] received, send ACK [1]
[60.0] A: received ACK [0]
[60.0] A: received ACK [0]
[60.0] A: received ACK [1]
[62.0] A: received ACK [1]
[62.0] A: received ACK [1]
[63.0] A: received ACK [1]
[64.0] A: received ACK [1]
[65.0] A: received ACK [1]
[66.0] A: received ACK [1]
[67.0] A: received ACK [1]
[68.0] A: received ACK [1
```

**Solution**: After I carefully checked the code and the principle of SR, I found that it was because the A side only base++ when acknum == base (the window only moves one packet at a time), but did not continue to check whether the subsequent sequence number was also ACKed, or did not update the status of the corresponding window cache in the "out-of-order" scenario.

The correct approach should be:

- ◆ When A\_output() sends a new packet, put it into the send window array winbuf[i] and write the status:
  - Sent + Waiting for ACK.
- ◆ When A\_input() receives a certain ACK=acknum:
  - Find the item with seq=acknum in the send window and mark it as "ACKed".
  - If acknum == base, move base forward, but not just base++, but continuously check: Is the next sequence number marked as ACKed? If so, continue ++ until you encounter a seq that has not been ACKed.
  - When the window moves forward, if there are still cached messages (in buffer[]) and there is space in the current window, continue to send new packets.

After my modification, the input of "xxxx----oooooooooooo" can be transmitted completely, but when I re-enter the two examples given, I find that the base of the packet at the 32nd second in the output of the second example is different from my output. I checked SR's interactive animation and it was the same as my output. I immediately asked TAs and the professor and got a positive answer.