

# Mero goals

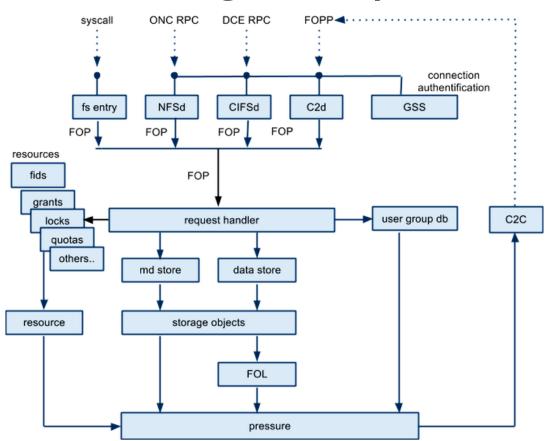
- extensible storage platform
  - observability
  - extensible scalable 3rd party plugins
- horizontally scalable (number of devices and nodes)
- vertically scalable (number of cores)
- flexible deployment

Warning: not everything in this presentation is fully implemented.

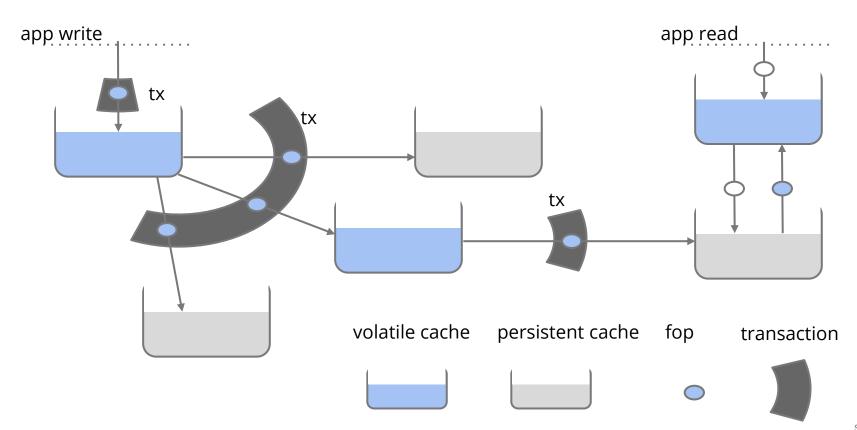
#### features

- 0-copy, 2-phase IO
- extensible meta-data: distributed key-value store
- layouts (data, meta-data, fault-tolerance)
  - network striping (parity de-clustering)
  - o composite: NBA, snapshot, multi-tier
- distributed transactions (consistency)
- resource manager (coherency)
- containers, function shipping
- threadless server design
- fdmi, addb
- user space, portable

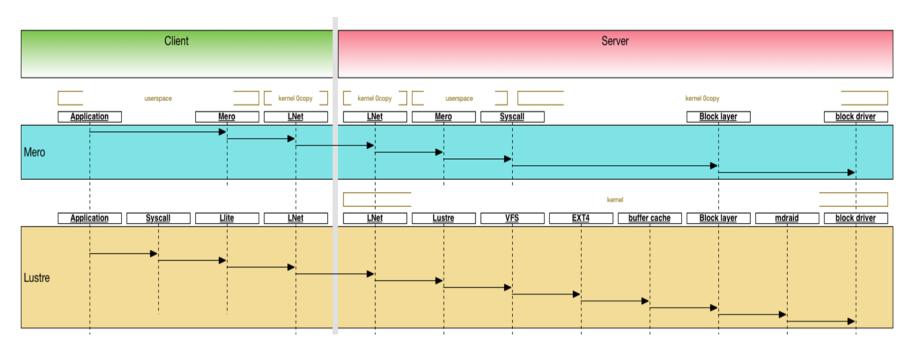
# structure: instance (single node, process or kernel)



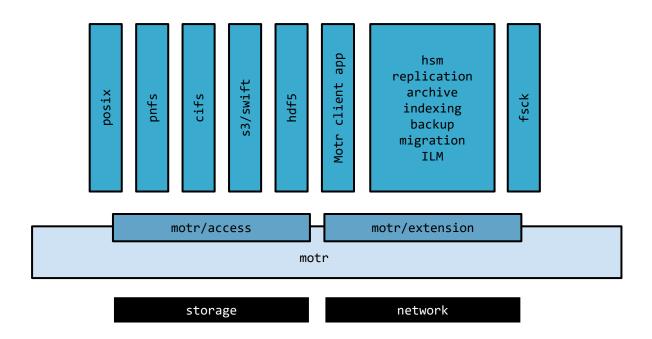
# structure: system view



# structure: (non-)layering



## structure: components



## components

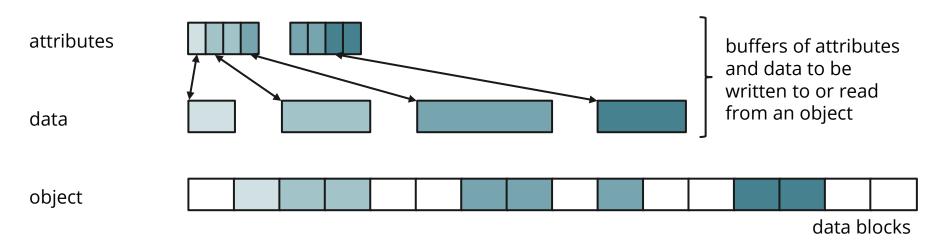
- Motr client
- transactions (dtm)
- resource manager
- fdmi
- addb
- network raid, layouts
- containers
- function shipping
- lingua franca
- integrity checking

- meta-data back-end
- network, rpc, fop, HA
- fom, reqh
- device io (stob)
- network raid repair
- security

### Motr client: overview

- object: network striped by default
- index: distributed key-value store
- operation: asynchronous, state machine, tracks progress
- transaction: atomic group of operations
- resource: ownership, coherent distributed caching
- layout: placement of data on storage
- 128-bit persistent identifiers, assigned by user
- all entry-points are asynchronous

## Motr client: object



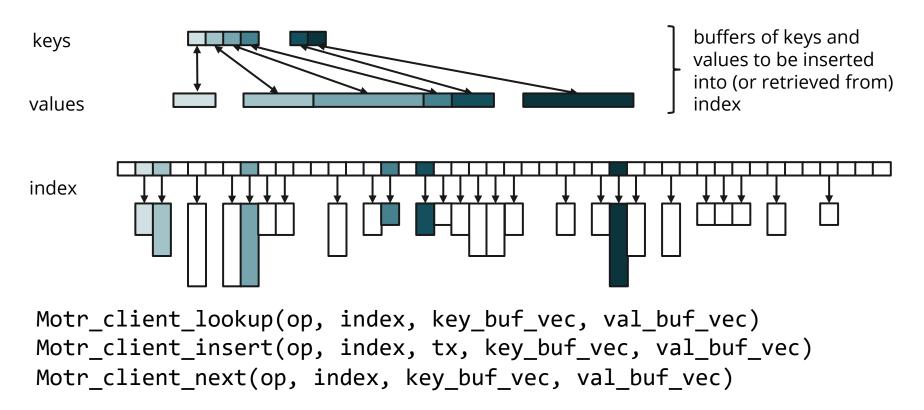
```
Motr_client_read(op, obj, data_buf_vec, attr_buf_vec, extent_vec)

Motr_client_write(op, obj, tx, data_buf_vec, attr_buf_vec, extent_vec)

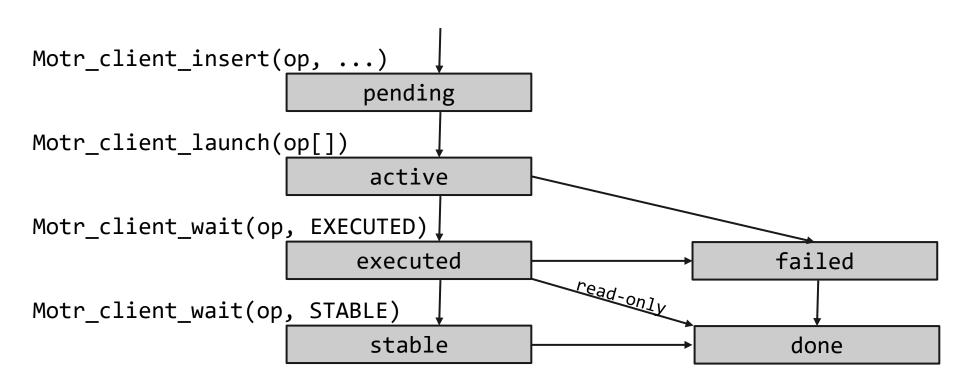
Motr_client_alloc(op, obj, tx, extent_vec)

Motr_client_free(op, obj, tx, extent_vec)
```

#### Motr client: index



# **Motr client: operation**



## components

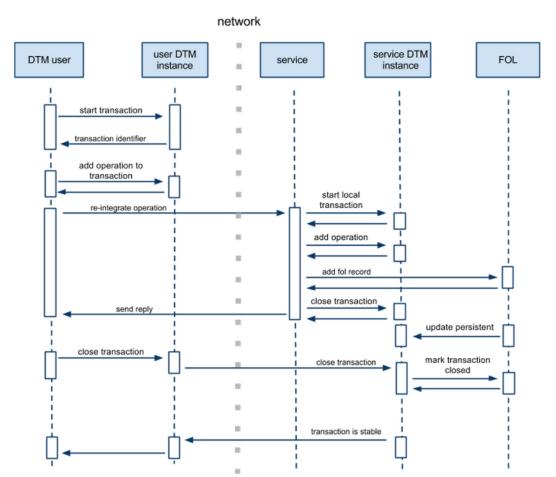
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### dtm: transactions

- a transaction is a group of operations
- **a**tomicity w.r.t. certain failures (network partitions, node restarts)
- consistency is defined by user
- guarantees neither isolation nor serialisability
- doesn't guarantee synchronous durability: too expensive for small transactions. Asynchronous stabilisation: the user is notified when the transaction becomes stable
- stabilisation ordering, liveness

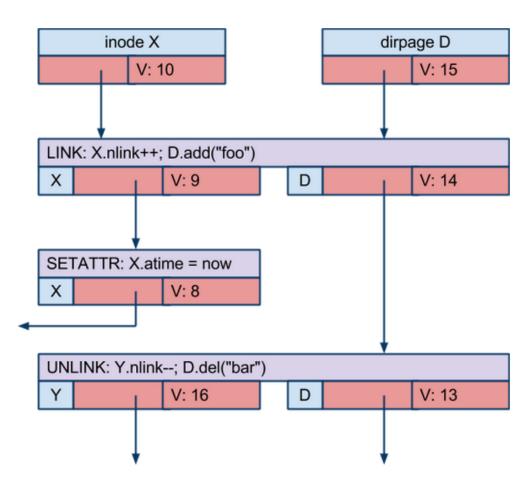
# dtm: sequence



### dtm: features

- masks certain transient failures
  - network partition, re-ordering, duplication
  - node restart
- write-ahead logging on each server
- undo for data
- redo for meta-data
- stabilisation: global logical clock (epochs)

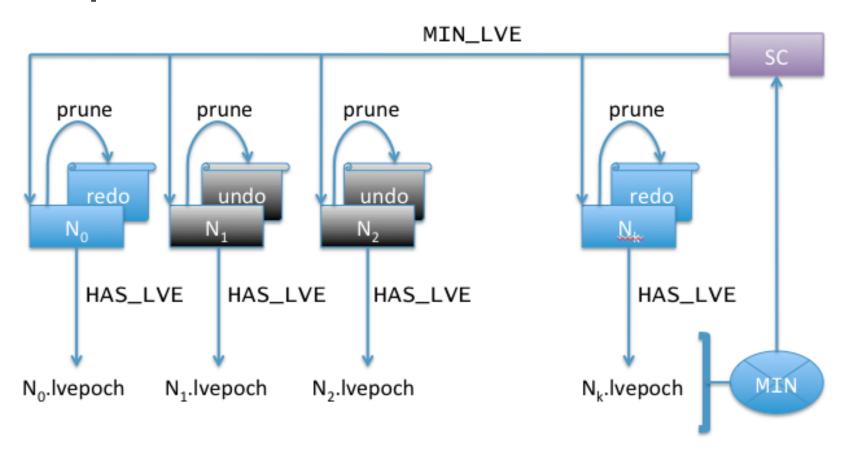
### dtm: versions



# dtm: epochs

- distributed clock to detect operation dependency and ordering (epoch, Fidge-Mattern, Lamport)
- messages are tagged with the logical timestamps:
   Event1 depends on Event0, then Event0.epoch <= Event1.epoch</li>
- any node can advance its clock independently
- operations are kept in the persistent log in epoch order, until epoch is stable. Then the log is pruned
- global coordination to determine when an epoch is stable

# dtm: epochs



## components

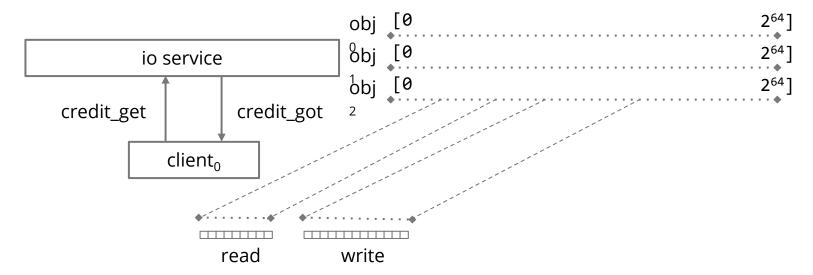
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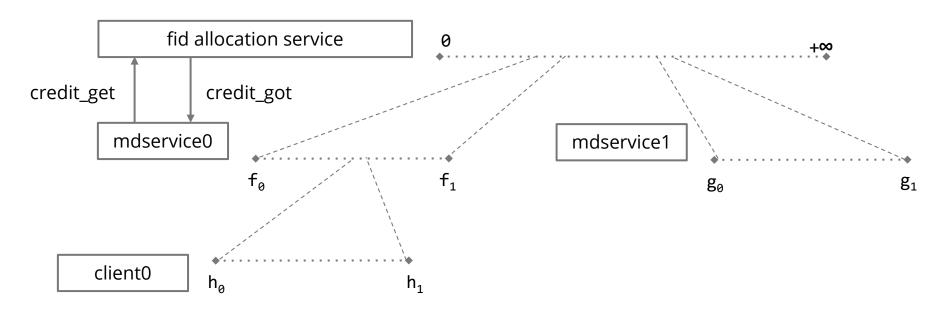
#### rm: definition

- resource: anything with ownership. An extent in an object, an entire object, a key in an index, etc.
- credit: a right to use a resource in a particular way
- credits control:
  - distributed caching
  - concurrency
- credits can be borrowed and sublet
- resource manager is separate from resource
- resource manager resolves conflicts
- user can define new resource types

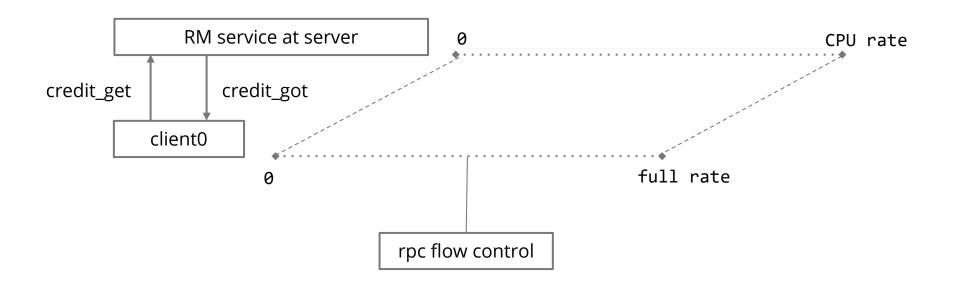
# Example: Block extent in a Motr client object



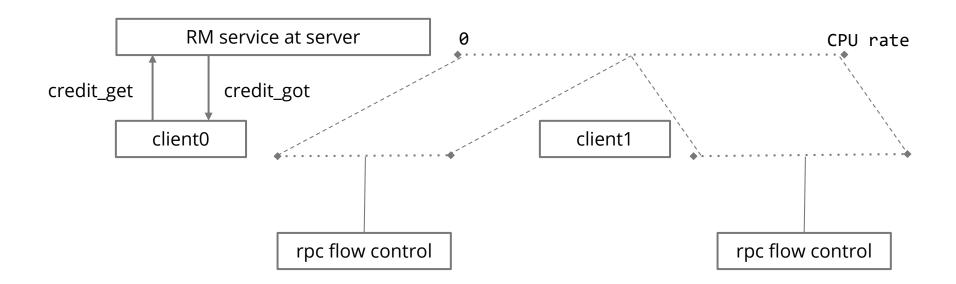
Example: fid extent allocation. Fid: 128 bit.



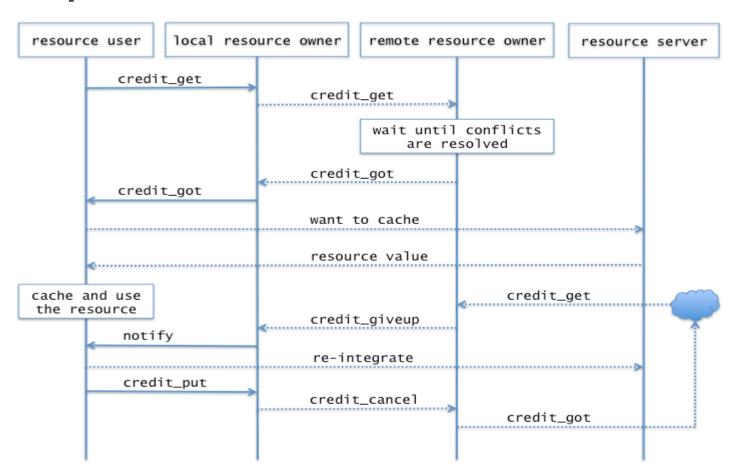
# Example: server CPU cycles



# Example: server CPU cycles



## rm: sequence



#### rm: resources

## Resource types

- open-files
- file extents
- disk space (grants)
- quotas
- server memory
- server cpu cycles
- file identifiers (fids)
- inode numbers

- network bandwidth
- storage bandwidth
- layout
- cluster configuration
- power

## rm: resource manager context

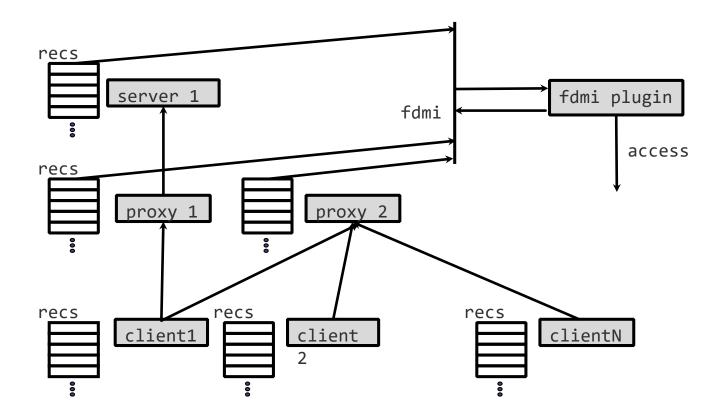
- generic infrastructure:
  - o RM service: BORROW, REVOKE, CANCEL
  - client (Motr client) interface
- specific resource types:
  - resource and credit names
  - conflict, credit ordering
- RM users:
  - resource acquisition and release logic
  - cache invalidation
  - assignment of RM services

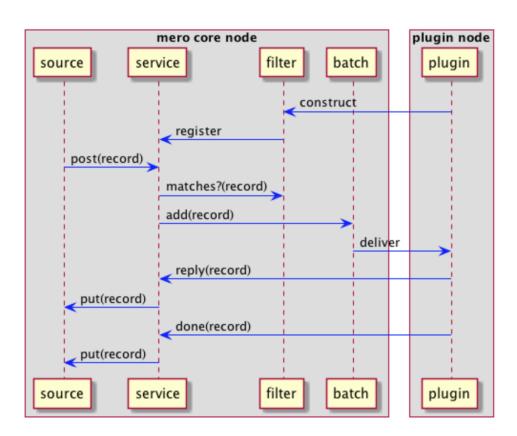
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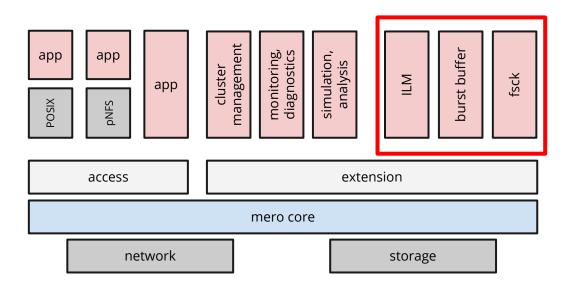
- meta-data back-end
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- operation log (fol):
  - record each operation on an object
  - log consists of records, log maintained by each node
- file-system definition and manipulation interface (fdmi)
  - scalable publish-subscribe interface
  - subscribe to records matching certain filter
  - map-reduce-style mechanism to deliver matching records to the subscribers
  - transactional delivery (EOS)
  - o delivers: fol records, addb records, HA events, others
- horizontal scalability
  - offload plugin processing and data-structures
  - asynchronous processing, batching





- ILM
  - replication
  - migration
  - o backup, archival,
  - o hsm
- indexing
- fsck
- data re-structuring
  - proxy de-staging
  - o RAID re-striping
- guided interfaces
  - o profiling
  - prefetching, destaging



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### addb

- systems grow larger and more complex
- how well the system is utilised?
- is it failure or expected behaviour?
- is it system or application behaviour?
- sources of data:
  - system logs
  - operating system
  - application traces
- very large amount of collected data
- ... or insufficiently detailed, or both
- difficult to analyse and correlate

### addb

- instrumentation on client and server
- data about operation execution and system state
- passed through network
- cross-referenced
- always on (post-mortem analysis, first incident fix)
- simulation (change configuration, larger system, load mix)

### addb: anatomy of a record

- measurement and context
- timestamped
- labels: identify context
- payload: up to 16 64-bit values,
- interpreted by consumer

### addb: sensors and histograms

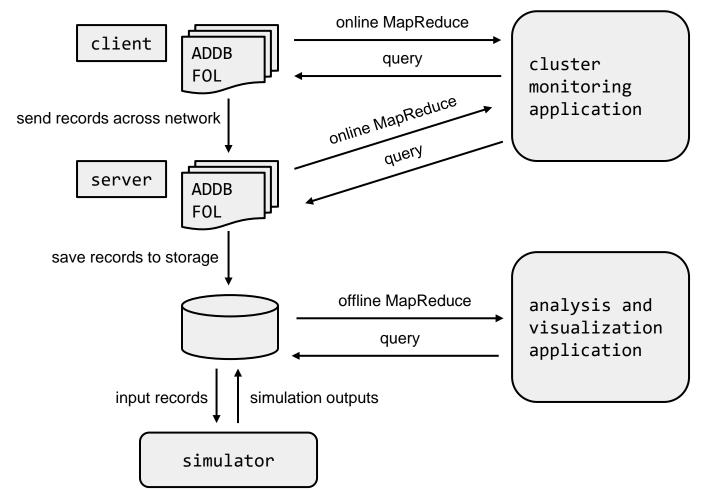
- some events are too frequent
- collapse them into counters
- count last events per locality
- automatically size buckets

### addb: state-machine transitions

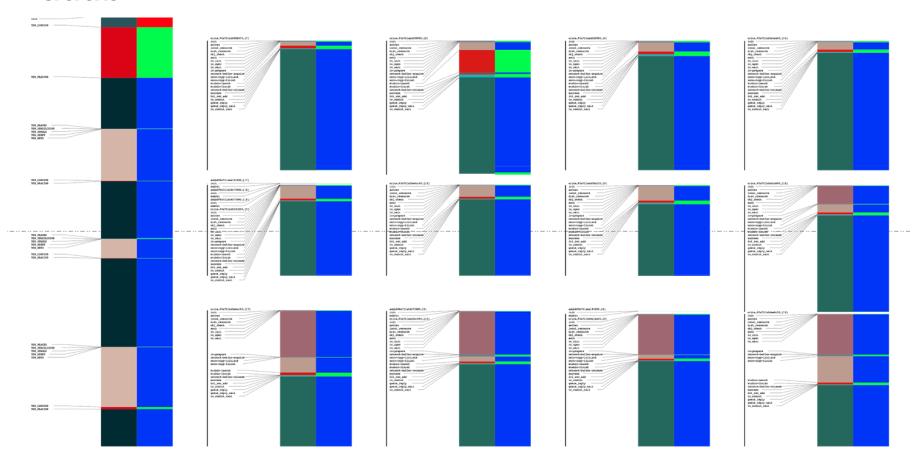
- request processing on clients and servers
- state machines
- state transition delays: in "binary microseconds", second >> 10

```
0: 25862
        1138
4132 :
                ***
         309
 6198 :
8264 :
10330 :
12396 :
14462 :
16528 :
          12
18594:
20660:
22726 :
24792 :
```

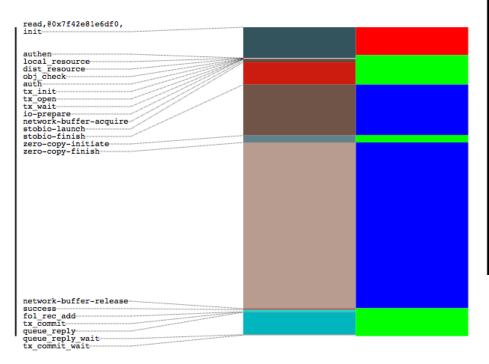
### addb

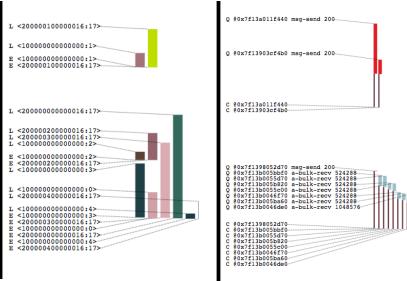


### addb



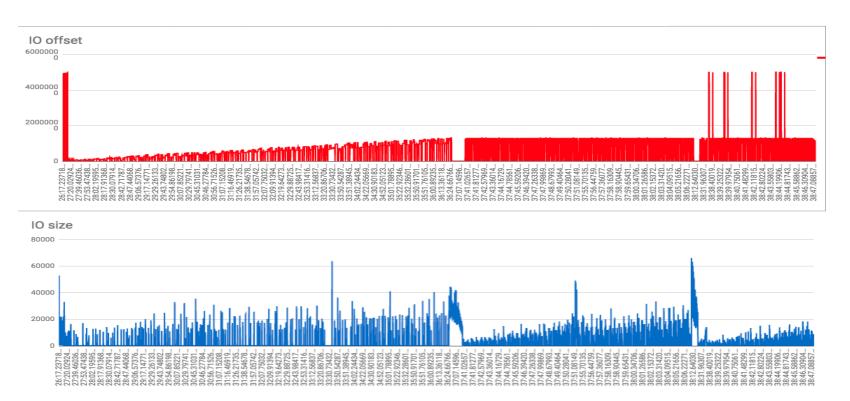
### addb





## addb: ad hoc profiling

\$ m0addb2dump ... | grep 'stob-io-launch' | awk '{print \$2, \$6, \$12}'



### addb: interface

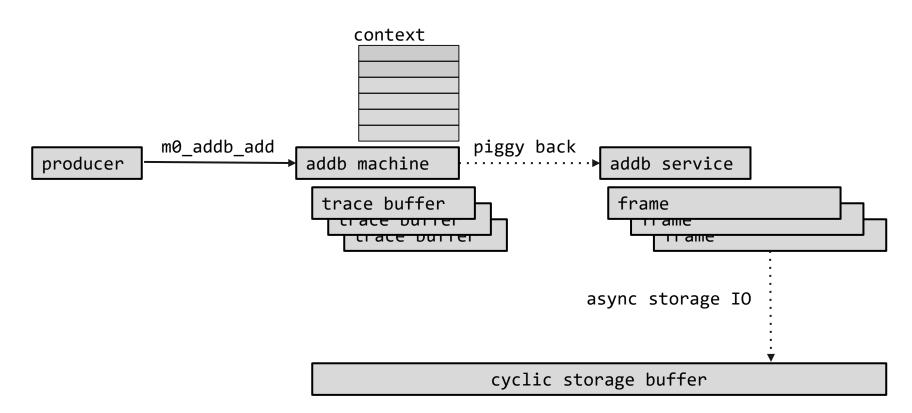
```
/**
 * Adds a label to the current context.
 * @param id

    label identifier

 * @param n - number of 64-bit values in label payload
 * @param value - payload
void m0 addb push(uint64 t id, int n, const uint64 t *value);
 * Removes the top-most label in the current context stack.
  @param id - label identifier
 * @pre "id" must the identifier of the top-most label.
void m0 addb pop(uint64 t id);
 * Adds one-time measurement in the current context.
 * @param id
                - measurement identifier
 * @param n
                - number of 64-bit values in measurement payload
 * @param value - payload
void m0 addb add(uint64 t id, int n, const uint64 t *value);
```

- very simple interface
- binary values only
- stack (LIFO) context
- context management is modular
- separate interface for sensors (not shown)
- usable from system and applications

### addb: data flow



### components

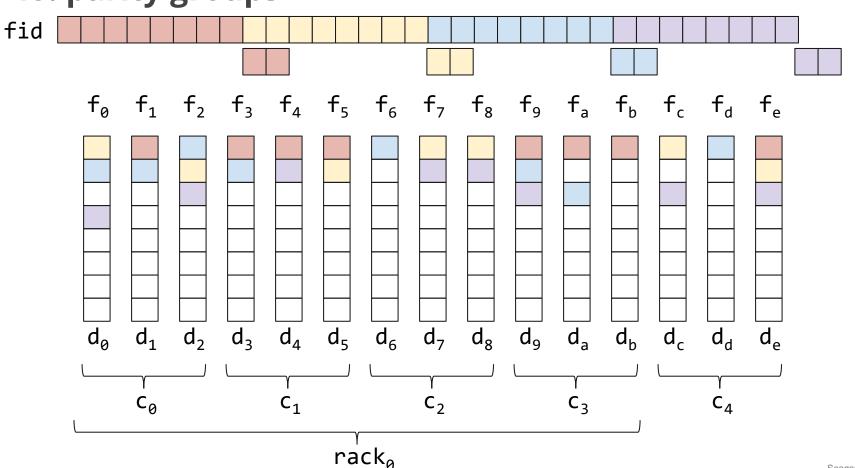
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### io: layout

- determines how an object is stored in underlying containers
- layouts for data and meta-data
- examples:
  - network striping with parity de-clustering (default)
  - compression
  - encryption
  - de-duplication
  - composite (NBA, small files, migration)

### io: parity groups



# io: permutations

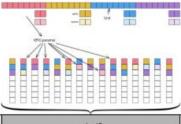
00	01	02	03	04	$\mathfrak{o}_p$	$0_q$	10	11	12	13	14	$_{1p}$	$1_q$	20
21	22	23	24	$_{2p}$	2q	30	31	32	33	34	$\mathfrak{z}_p$	$\mathfrak{z}_q$	40	41
42	43	44	$^{4}p$	4q	50	51	52	53	54	5 <sub>p</sub>	5q	60	61	62
63	64	6p	6q	70	71	72	73	74	7p	$7_q$	80	81	82	83
84	8 <sub>p</sub>	8q	90	91	92	93	94	9p	$9_q$	100	101	102	103	104
10p	10q	110	111	112	113	114	$11_p$	11q	120	121	122	123	124	12p
12q	130	131	132	133	134	13p	13q	140	141	142	143	144	14 <sub>p</sub>	14q
	PDRAID [15 (5+2+0)], 1 Tile													

03	11	01	$_{1q}$	12	0q	20	14	00	10	$\mathfrak{o}_p$	$1_p$	04	02	13
24	32	22	40	33	30	41	3 <sub>p</sub>	21	31	2q	3q	2p	23	34
4p	53	43	61	54	51	62	5q	42	52	50	60	4q	44	5p
6q	74	64	82	7 <sub>p</sub>	72	83	80	63	73	71	81	70	6p	$7_q$
90	$9_p$	8 <sub>p</sub>	103	9q	93	104	101	84	94	92	102	91	8q	100
111	11q	10q	124	120	114	12p	122	10p	11p	113	123	112	110	121
132	140	130	14p	141	13 <sub>p</sub>	14q	143	12q	13q	134	144	133	131	142
PDRAID [15 (5+2+0)], 1 Tile														

# **layout: composite**

#### Complex Layout

individual layout per extent

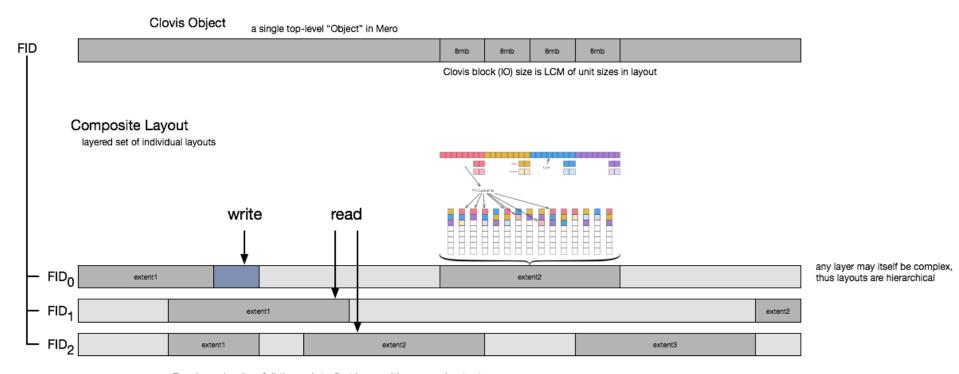


FID extent1 extent2

any layer may itself be complex, thus layouts are hierarchical

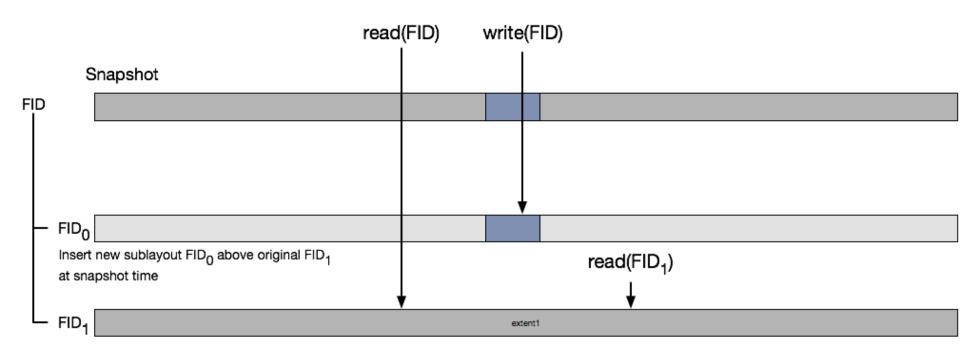
extent2

### layout: composite

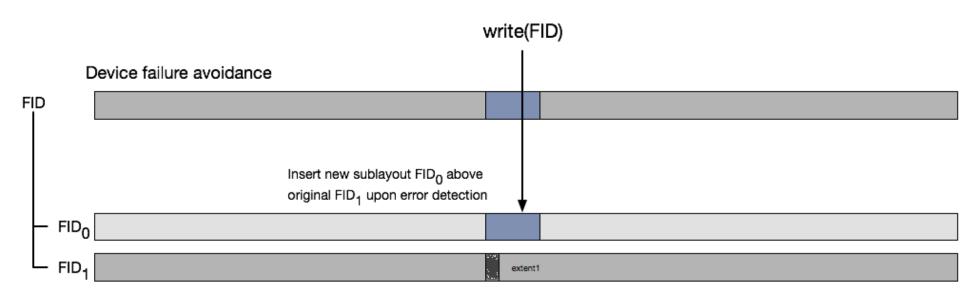


Reads and writes fall through to first layer with mapped extents. (Newly written blocks are added to read extents.) Layers can also be read/written directly using FID<sub>Sub</sub>

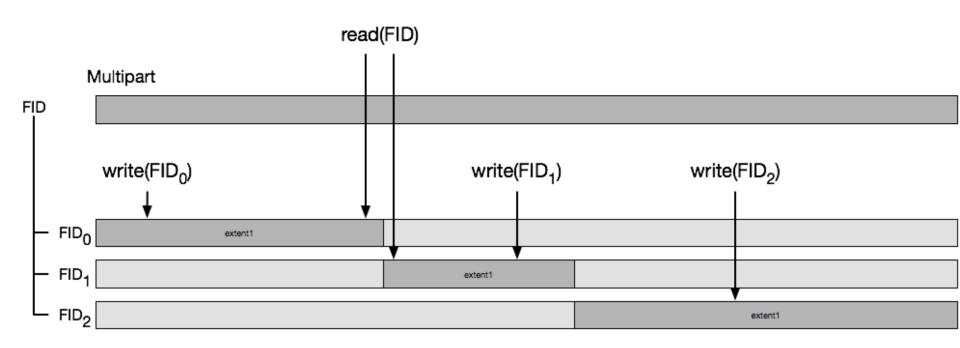
### layout: composite: snapshot



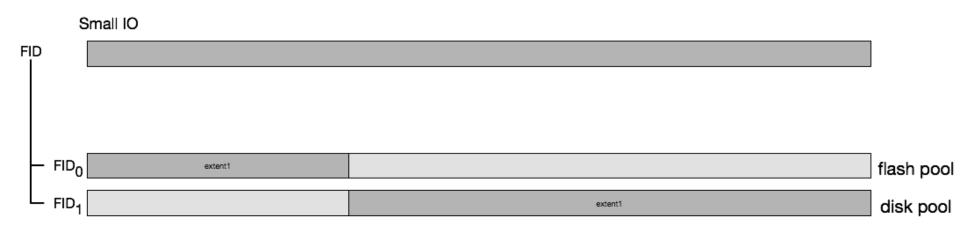
## layout: composite: nba



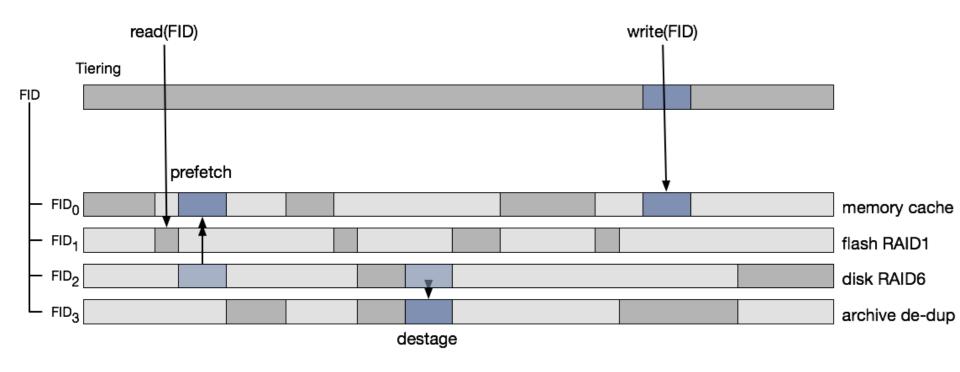
## layout: composite: s3 partial upload



# layout: composite: small files IO



## layout: composite: HSM, tiering



### components

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- function shipping
- lingua franca
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#### containers

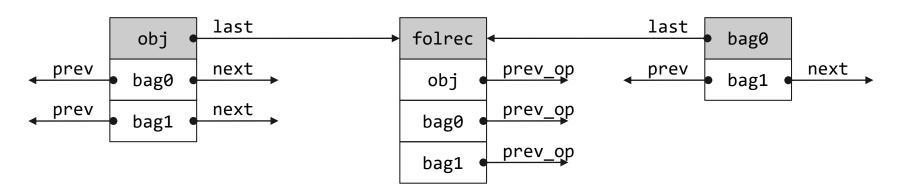
- container: an entity with a fid
- application fully controls containers
  - add an entity to a container;
  - remove an entity from a container;
  - list container elements;
  - execute bulk operation on container contents (compute-incontainer)
- arbitrary "topology" (only restriction: no duplicates)

#### Non-properties:

- separate fid-space for contents
- strong isolation guarantees

### containers implementation

- component containers, bag:
  - entities with fids,
  - lazily created when an entity is added to the container
- all local entities are linked together
- file operations log records (folrecs) for operations on container contents are linked



### containers rationale

- efficiently identify containers to which an object belongs
- flexible container nesting
- concurrent mass-operations on containers
- container history is traceable
- fdmi filters on containers
- function shipping to a container

#### Don't have:

- ordering of container contents
- ordered enumeration
- "implicit" containers (e.g., "all objects on this server")

#### containers use case

- application creates a container to track "entities of interest"
- entities are added to the container
- fdmi plugin is registered to receive notifications of all updates to the container
- examples:
  - hsm: fileset of hot objects
  - replication: source fileset

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### function shipping

- move computation closer to data (compute-in-storage)
- reduce network transmission overhead
- application structuring mechanism
- bulk computation on all Motr client entities:
  - object: function on data blocks
  - index: function on key-value records
  - container: function on member fids
- built-in fault tolerance

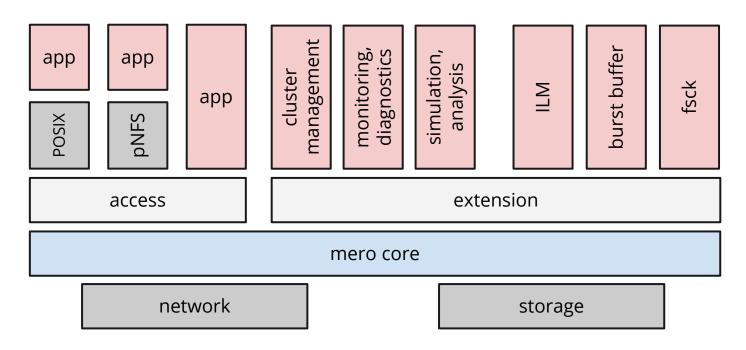
### function shipping implementation

- computations are first class Motr client entities
  - unique fid, globally addressable
  - registered dynamically by an application
- low level trusted mechanism:
  - dynamically load shared library into Mero service process
  - invoke computations remotely, argument-result passing
- untrusted mechanism:
  - run untrusted code (e.g., Python) in a separate address space
- client uses layouts to start execution and recover from failures

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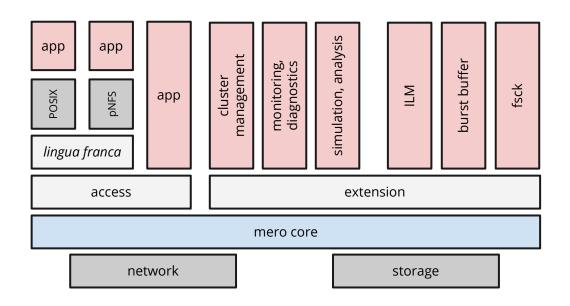
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Multiple front-ends: interoperability, common meta-data mechanism

- POSIX file system: via pNFS
  - straightforward semantics
  - concurrency:
    - concurrent reads
    - one writer
    - concurrent writers
- POSIX file system: via S3
  - interpret pathname as an URL (object key)
  - mapping of security and ownership attributes
- S3: *via* POSIX: parse URL as a pathname
- POSIX file system: via HDF5

- S3 object: "blue\_bucket/finance-9a84c723ee89f4b723b46cc5f1642b3"
- mount S3 store as POSIX
  - \$ cd blue\_bucket
  - \$ ls -1 finance-9a84c723ee8\*
  - -rw----- 1 satoshi satoshi 285 January 03 2009 finance-9a84c723ee89f4b723b46cc5f1642b3
- enumerate and locate objects
- access attributes
  - system attributes: layout, containers;
  - o common attributes: size;
  - front-end specific attributes: POSIX nlink
- efficiency: remote meta-data lookups are expensive



POSIX, pNFS, CIFS, S3/SWIFT, MPI Object IO, HDF5, MySQL, block device

### lingua franca implementation

- what are the *entities* managed by an FE?
- how entities are named?
- how entities are organised: tree, graph, array?
- how attributes are associated with entities?
  - different FEs have different sets of native attributes
  - o an FE wants to add attributes to foreign entities
  - an FE wants to interpret foreign attributes
  - some attributes are shared by multiple FEs
- core system has its own attributes

### lingua franca implementation

- a domain of files (entities), a file identified by a 128-bit fid
- two indices:
  - name-space (ns): records file attributes
  - object index (oi): maps from the file fid to all its names

#### ns index structure:

key	value
FRONTENDID+FNAME+ATTRCLASS+ATTRID	attribute value

#### oi index structure:

FID+FRONTENDID+FNAMEID	FRONTENDID+FNAME
------------------------	------------------

### lingua franca implementation: S3

• FRONTENDID: 3

ATTRCLASS: 3

FNAME: Object-URI: bucketname+NUL+object\_key+NUL

object: "app\_bucket/statement.xls", fid 0x60000:0x17ae76d0f

ns index:

key	value	comment
<pre>3app_bucket\0statement.xls\0!fid</pre>	0x60000:0x17ae76d0f	the fid of the object "statement.xls"
<pre>3app_bucket\0statement.xls\03content-length</pre>	3201526	object size in bytes
3app_bucket\0statement.xls\03content-md5	0x62ae2f12137738a9:0x173f5224f81446aa	md5 checksum
<pre>3app_bucket\0statement.xls\03last-modified</pre>	2017-04-26-15:29:30.85267	last modification timestamp
<pre>3app_bucket\0statement.xls\03</pre>		other S3 attributes

# lingua franca implementation: S3

object: "app\_bucket/statement.xls", fid 0x60000:0x17ae76d0f

oi index:

key	value	comment
3:00000000000060000:0000000017ae76d0:30	<pre>3app_bucket\0statement.xls</pre>	name of the object

# lingua franca implementation: POSIX

• FRONTENDID: P

ATTRCLASS: P

FNAME: parent\_directory\_fid+name\_in\_the\_directory

object: "/etc/passwd", fid 0x70000:0x322e1673fd

parent directory: "/etc", fid: 0x70000:0x18203a6485

ns index:

key	value	comment
P:70000:18203a6485:passwd!fid	0x70000:0x322e1673fd	the fid of the object "/etc/passwd"
P:70000:18203a6485:passwdPsize	16523	file size
P:70000:18203a6485:passwdPatime	2017-04-26-15:29:30.85267	access time
P:70000:18203a6485:passwdP		other POSIX attributes

# lingua franca implementation: POSIX

object: "/etc/passwd", fid 0x70000:0x322e1673fd

oi index:

key	value	comment
P:000000000000000000000000000000000000	P:70000:18203a6485:passwd	name of the object
P:000000000000000000000000000000000000	P:70000:18203a6485:passwd.1	another name: hard-link

# In /etc/passwd /etc/passwd.1

# lingua franca implementation: features

- each attribute is a separate key-value pair: flexibility
- new attributes can be added to existing files
  - without breaking compatibility
- attributes can be enumerated (NEXT operation)
  - o an FE can selectively handle attributes it understands
- contents of a directory can be enumerated
- new names can be added to a file

## components

- Motr client
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- fdmi
- addb
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- function shipping
- lingua franca
- integrity checking

- meta-data back-end
- network, rpc, fop, HA
- fom, reqh
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# integrity checking

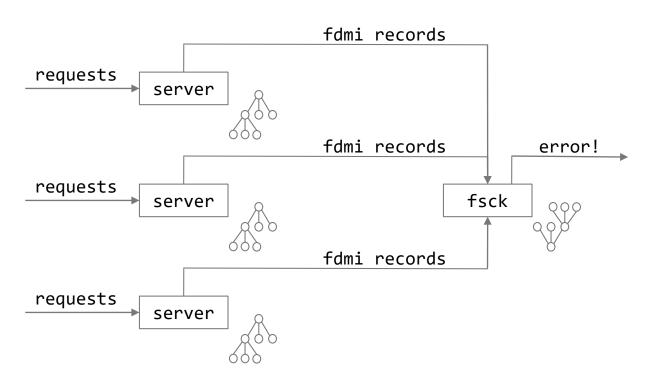
- redundancy, fancy metadata: not an answer (has been tried)
  - bugs (more important over time)
  - recovery from catastrophic failures
- traditional fsck
  - not distributed
  - specific to particular meta-data format
  - does not scale
    - time
    - space

# integrity checking

- need scalable integrity checking
- run it all the time
- on dedicated separate nodes (horizontal scalability)
  - maintain redundant "inverse" meta-data
  - update meta-data to match system evolution (fdmi)
  - detect inconsistencies
  - report, recover from redundancy
  - recover catastrophic failures
- usual redundancy: parity, checksums, background scrub
- fdmi: transactional coherence with the main state

parallel programming

# integrity checking



#### inverse meta-data

- block allocation
- pdclust structure
- key distribution
- b-tree structure
- application specific invariants
  - POSIX tree
  - o hdf5

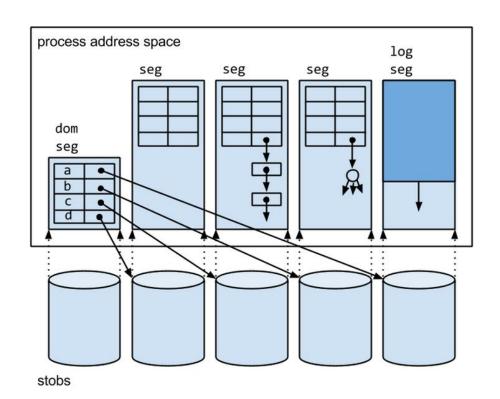
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#### backend

- segment
- transaction
- domain
- container
- WAL
- redo-only
- allocator
- btree



## components

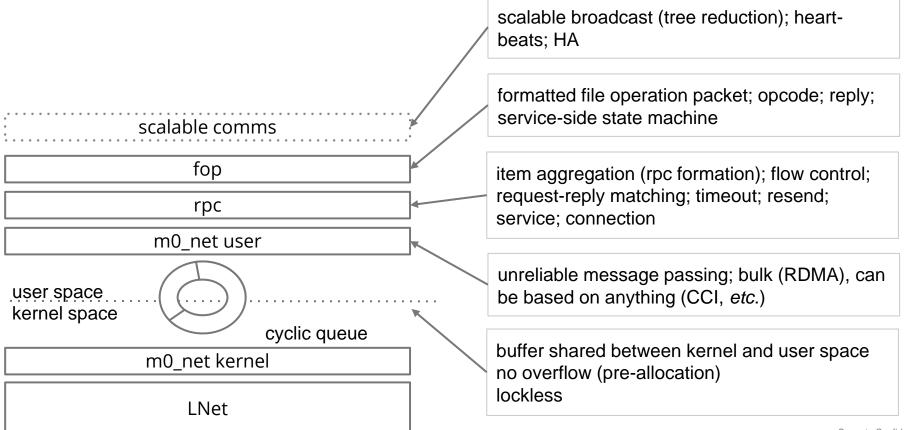
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#### comm

- network: LNet, 0-copy, unreliable message passing
- rpc
  - message packing (formation),
  - request-reply semantics
  - retransmit
- xcode: serialisation library
- fop: operation packet
- reduce-broadcast
  - communication and aggregation tree
  - provided by HA

#### comm



#### comm: xcode

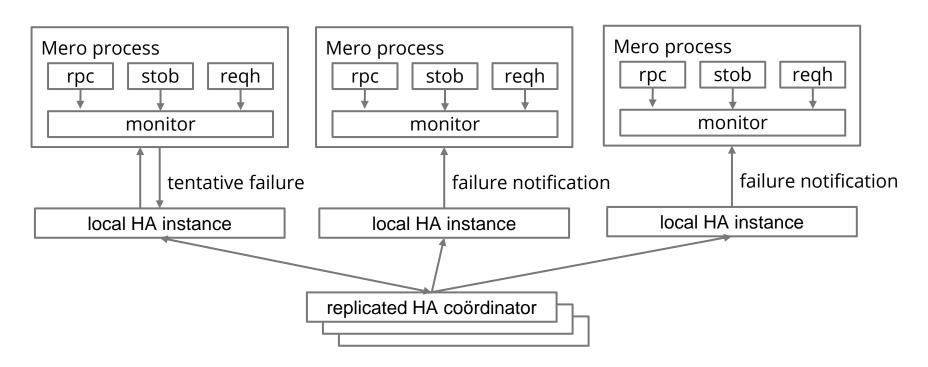
- a library adding a (modest) amount of reflection to C
- annotate header file foo.h
- gccxml: C parser
- generate foo\_xc.h, foo\_xc.c

```
struct m0_foo {
      uint32_t f_nr;
      struct m0_bar *f_bar;
} M0_XCA_RECORD;
```

```
static struct m0_xcode_type m0_foo_xc = {
    .xct_name = "m0_foo",
    .xct_sizeof = sizeof (struct m0_foo),
    .xct_children = {
        [0] = { "f_nr", &M0_XT_U32, offsetof(struct m0_foo, f_nr) },
        [1] = { "f_bar", &m0_bar_xc, offsetof(struct m0_foo, f_bar) }
};
```

m0\_xcode\_{en,de}code(), m0\_xcode\_{find,print,read}()

#### comm: HA



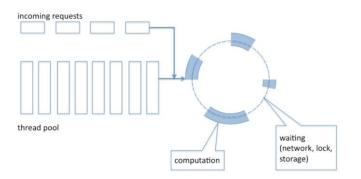
## components

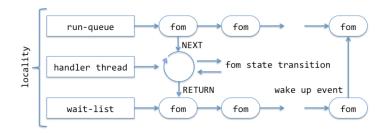
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#### fom

- thread-per-request:
  - multiple cores, NUMA,
  - locking,
  - cache ping-pong,
  - o c10K, many threads
- reqh:
  - thread per core
  - non-blocking scheduler
  - locality of reference
  - load balancing
  - long-term scheduling

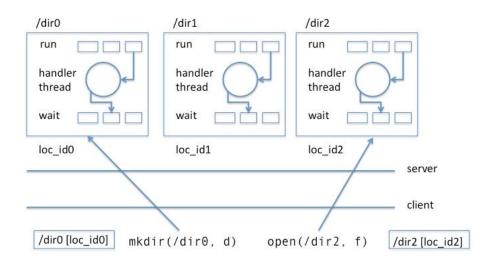




# fom: request handler

```
fom transition(locality *1) {
while (!locality->shutdown) {
                                             for each(fom, 1->run_queue) {
        while (have work) {
                network drain();
                                                     good = goodness(1, fom);
                stob drain();
                                                     if (good > max_goodness) {
                fom transition();
                                                                   best = fom;
                                                                   max_goodness =
        wait for work();
                                    fom good;
                                             best->state(); /* state transition */
goodness(locality *1, fom *f) {
        if (abs(fom->next block_nr, l->elevator_pos) < threshold)</pre>
                result += 1;
        if (l->pending rpc[fom->next nid] > 0)
                result += 1;
        result += take_deadline_into_account(1, f);
        result += take priority into account(1, f);
```

# fom: load balancing



- each locality is a small server
- clients talk to particular locality, using opaque identifier
- change identifier for load-balancing

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#### stob

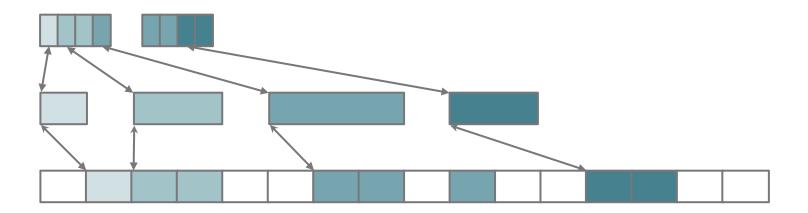
- array of data-blocks, [0, 2<sup>64</sup>), initially a hole
- create, delete, read, write, alloc, free operations
- IO at block granularity
- no usual meta-data (attributes, etc.)
- block attributes can be used for checksums, encryption keys, hash fingerprints
- scatter-gather-scatter operations: data and block attributes

# stob: IO

attributes

buffers

object



# stob: implementations

- linuxstob (aka devstob)
  - stob = file
  - aio
- adstob (allocation data stob)
  - multiple stobs stored in a backend stob
  - block allocator balloc
    - based on ext4 mballoc

## components

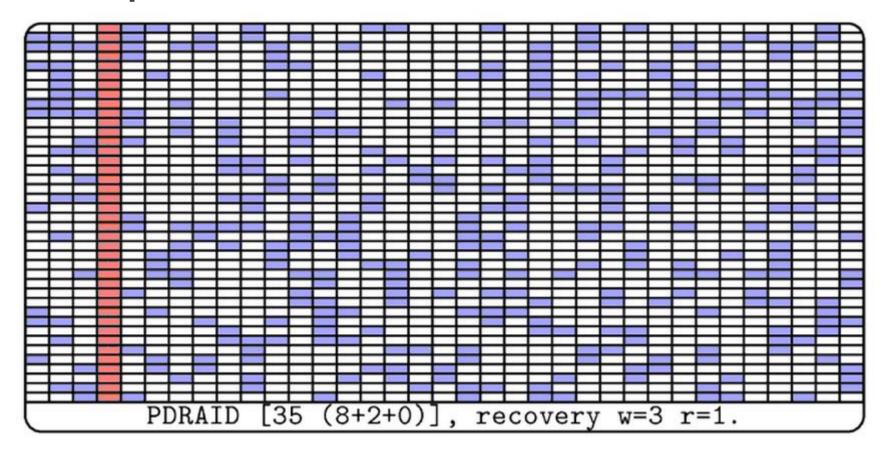
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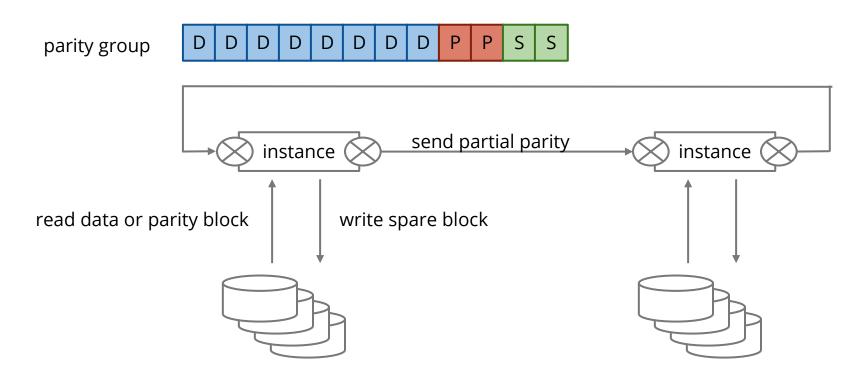
#### sns

- guaranteed IO performance during repair
- fast repair
- copy machine
- repair
- rebalance
- pool
- flattening

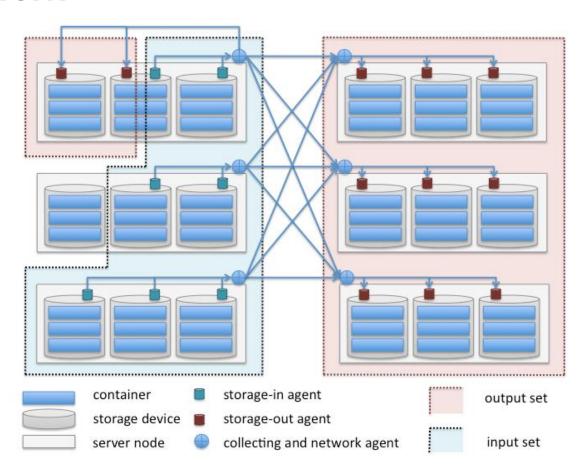
# sns: repair



# sns: copy machine

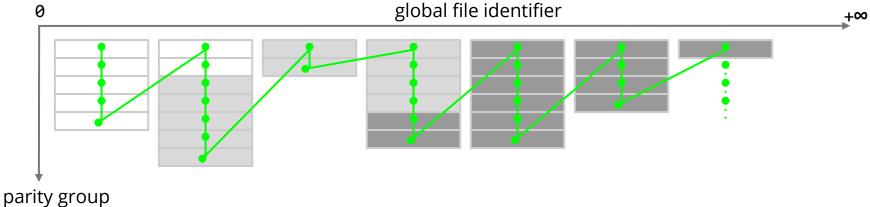


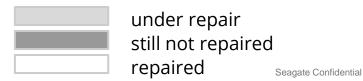
#### sns: network



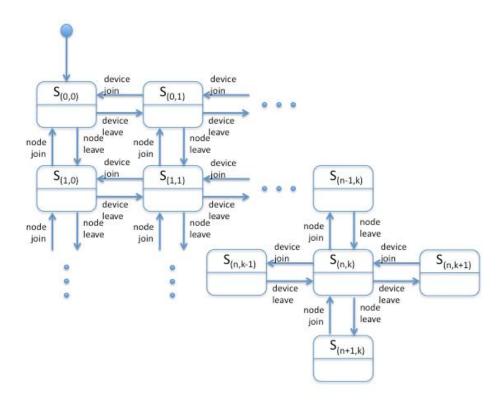
## sns: coördination

copy machine sliding window

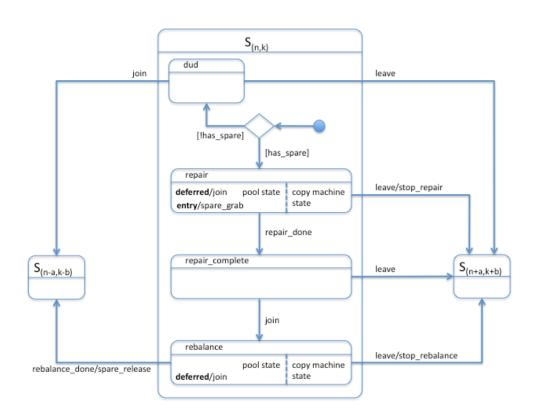




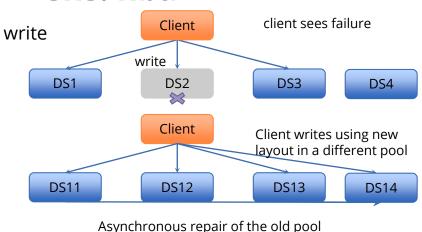
#### sns: failure state machine

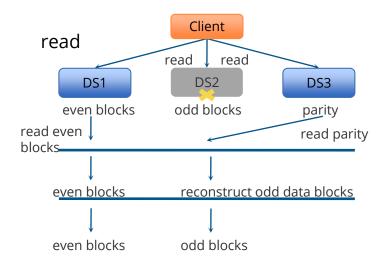


#### sns: failure in detail



#### sns: nba





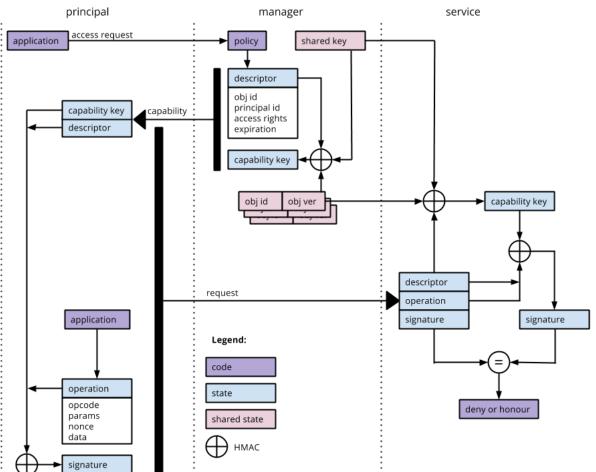
- client senses failure by timeout, notifies HA
- how new layout is selected? Layout formula
- composite layout with list of extents
- flattening: restore simple layout

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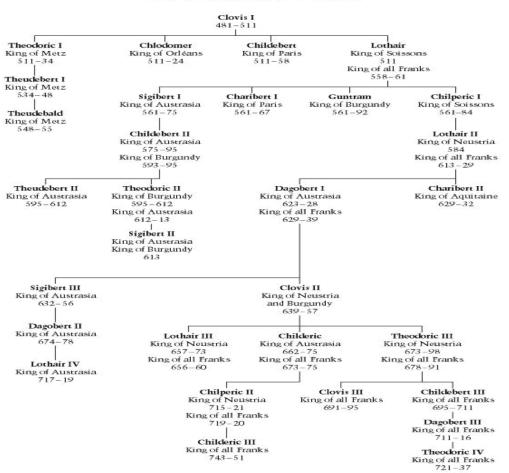
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security



# questions?

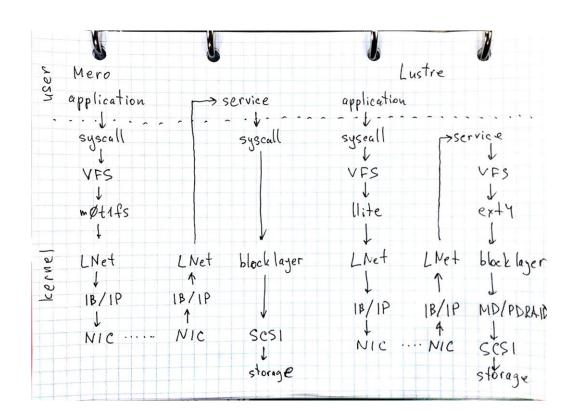
#### THE MEROVINGIAN KINGS



#### use cases

- IO data-flow end-to-end
- lookup path through all layers
- LOMO/WOMO: server and client failure

#### data-flow



#### lookup: m0t1fs

### lookup: rpc out

```
m0 rpc post(fop)
   o m0_rpc item send()
       m0_rpc_item_start_timer()
      m0_rpc_frm_enq_item()
        frm insert()
           o queue add [URGENT, BOUND]
        frm balance()
           o if (ready) frm fill packet()
              ■ m0 rpc packet add item()
           o frm packet ready()
              ■ m0 net buffer add()
```

#### lookup: rpc in

```
buf recv cb() (= [M0 NET QT MSG RECV])
   o net buf received()
       packet received()
        item received()
           o m0 rpc slot reply received()
              ■ item find(item xid(reply))
              ■ req->rio replied()
              ■ m0 rpc item change state(req,
                MØ RPC ITEM REPLIED)
```

m0\_rpc\_item\_wait\_for\_reply()

# lookup: network out

```
    m0_net_buffer_add()
    m0_net_tm_tlink_init_at_tail(buf, ql)
    nx_ops->xo_buf_add(buf) (= nlx_xo_buf_add)
    M0_NET_QT_MSG_SEND
    nlx_core_buf_msg_send()
    nlx_ucore_ioctl(M0_LNET_BUF_MSG_SEND
```

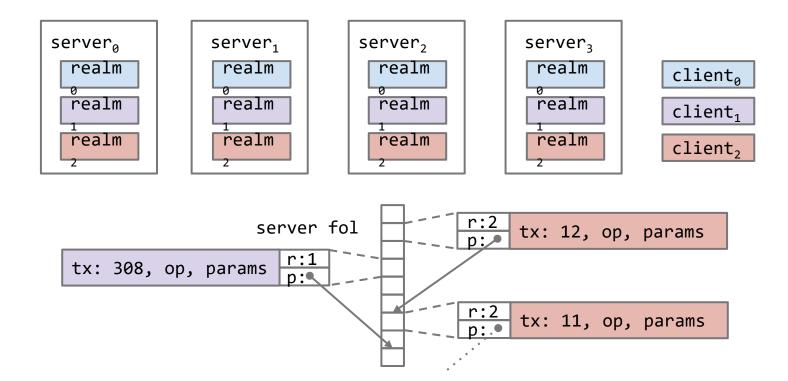
# lookup: network in

```
    nlx_tm_ev_worker()
    !queue.empty() or semaphore_down()
    queue get
    m0_net_tm_event_post()
    buf_recv_cb() (= [M0_NET_QT_MSG_RECV])
```

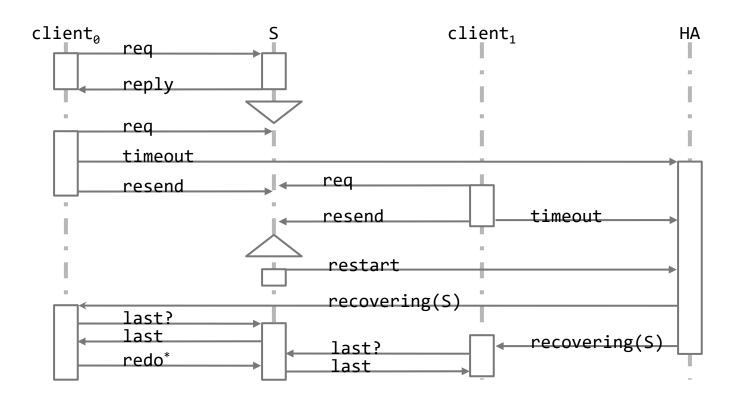
# lookup: service side

```
m0 rpc item dispatch()
   o m0 reqh fop handle()
       fop->fto_create(fop, &fom)
      ■ m0 fom queue(fom)
loc handler thread()
   o fom exec(fom)
       fom->fo ops->fo tick(fom)
         m0 md tick getattr()
            o m0 fom tick generic()
            o m0 cob locate()
               ■ m0 be btree lookup()
```

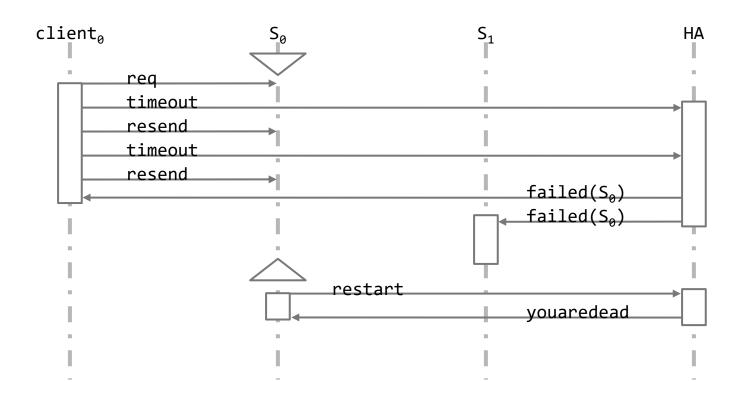
#### LOMO/WOMO realms



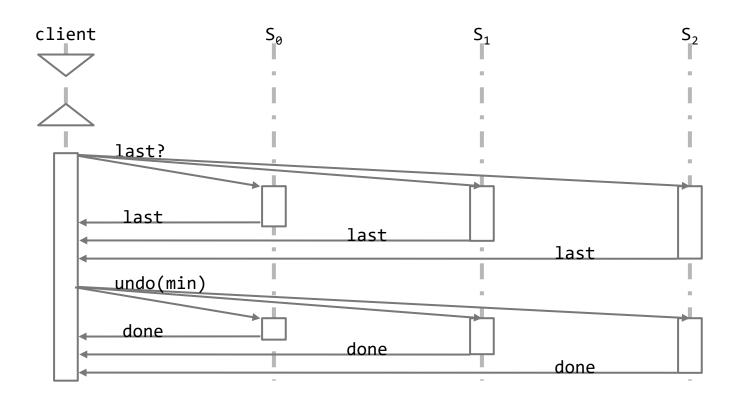
#### transient server failure



#### permanent server failure



#### client failure



# io:

