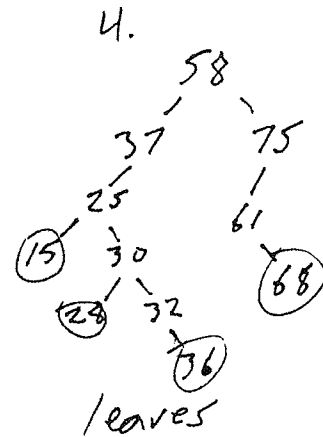
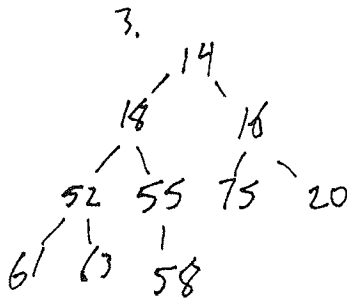
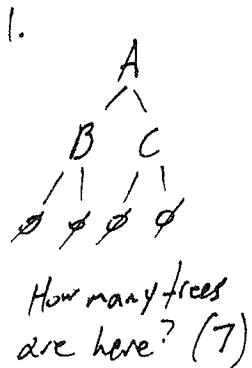
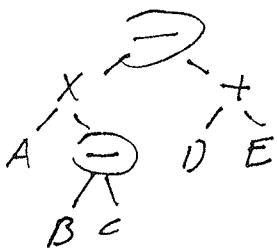


Example trees



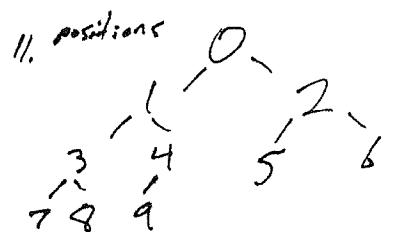
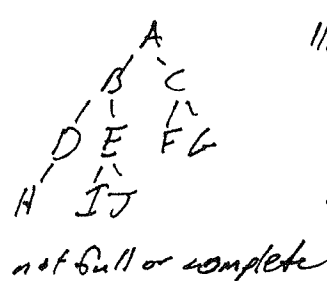
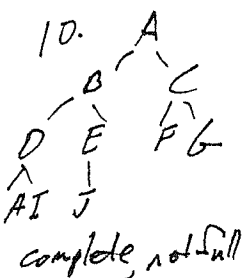
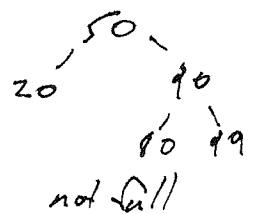
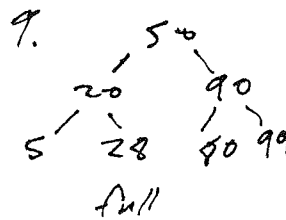
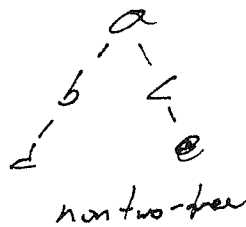
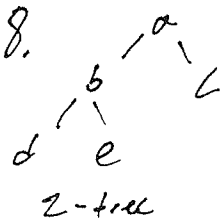
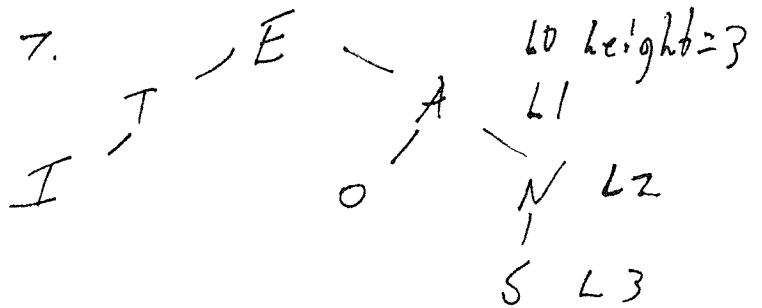
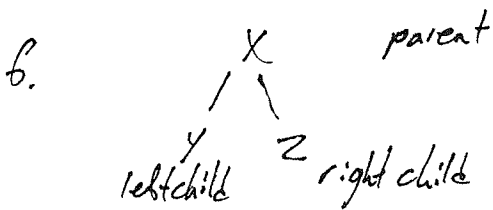
path 77-3
77, 25, 30, 32

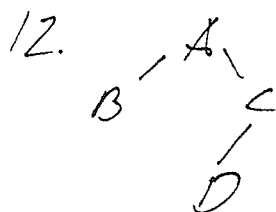
5.



element value = "-" at root (t)

element value = "-" at $\text{root}(\text{rightTree}(\text{leftTree}(t)))$





$$n(t) = 4$$

$$t \text{ is a } \rightarrow \text{floor} \left(\frac{4+1}{2} \right) = 2 = \text{leaves}(t)$$

∴

t is a tree? No!

instead

$$\text{out vision} \rightarrow \frac{4+1}{2} = 2.5 \geq \text{leaves}(t)$$

True

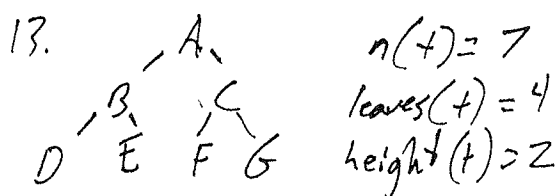
$$\frac{n(t)+1}{2.0} \leq 2^{\text{height}(t)}$$

$$\text{height}(t) = 2$$

$$\frac{4+1}{2.0} \leq 2^2$$

$$2.5 \leq 4$$

True



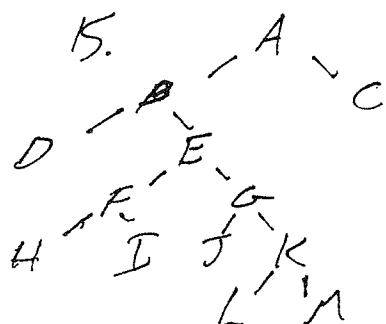
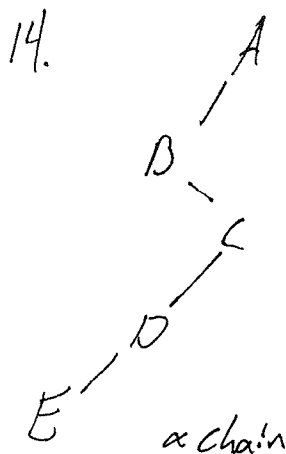
$$n(t) = 7$$

$$\text{leaves}(t) = 4$$

$$\text{height}(t) = 2$$

$$1, 3) \quad \frac{7+1}{2} = \text{True}$$

$$2, 4) \quad \frac{7+1}{2} = 4 = 2^2 = \text{True}$$



$$2 + 4 + 4 + 4 + 5 + 5 + 1 = 25$$

inOrder(+)

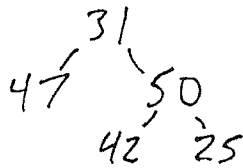
16. Tree + has n elements

Each element has 2 subtrees

\therefore There will be $2n$ recursive calls for inOrder(+)

$\therefore O(n)$

Tree +:



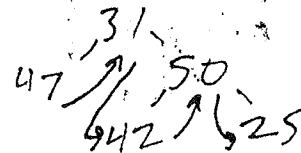
leftTree(+): (2) 47 \rightarrow
 (1) ~~31~~ ~~50~~

Back to where we were before

process root:

31 \rightarrow

processing order



rightTree(+): 50
 42 25

leftTree(rightTree(+))

42 \rightarrow

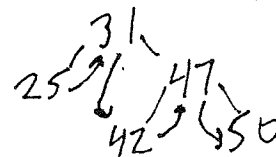
root(rightTree(+)) \rightarrow

50

rightTree(rightTree(+))

25 \rightarrow

17. BST



Output

25

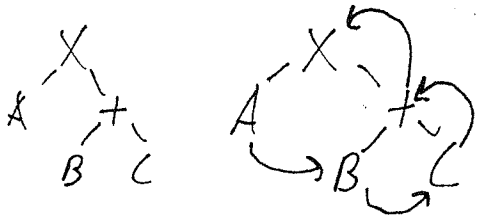
31

42

47

50

18. postOrder(+)
O(n)



Output

A

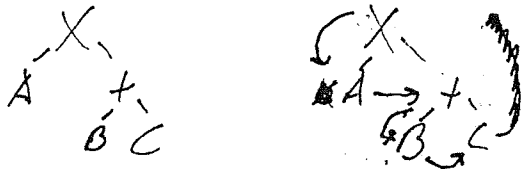
B

C

+

X

19. preOrder(+)
O(n)



Output

X

A

+

B

C

20. A

Breadth-First

Output

A

B

C

D

E

~~F~~

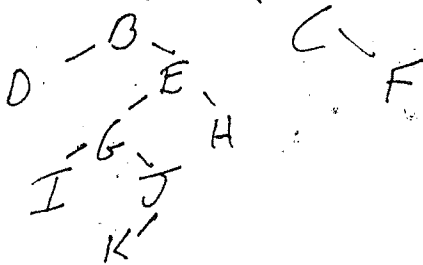
G

H

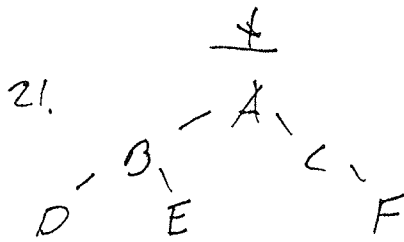
I

J

K



steps are ordered in numbers
X = sequence



Output

A	4
B	9
C	14
D	18
E	21
F	24

