

UF 2: Multi-thread programming

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Thread issues

Thread issues

- ❶ Non-atomic access
- ❷ Interference
- ❸ Memory consistency errors

Non-atomic access

Expressions can define complex actions that can decompose into other actions.

For example, `c++`, does not describe an atomic action:

1. Retrieve the current value of `c`.
2. Increment the retrieved value by 1.
3. Store the incremented value back in `c`.

Thread interference

It happens when two operations, running in different threads, but acting on the **same data**, interleave.

Thread interference Example

```
class Counter {  
    private int c = 0;  
    public void increment()  
    {  
        c++;  
    }  
    public void decrement()  
    {  
        c--;  
    }  
    public int value()  
    {  
        return c;  
    }  
}
```

Thread A invokes increment
Thread B invokes decrement

Thread A: Retrieve c=0.

Thread B: Retrieve c=0.

Thread A: Increment retrieved
value; result is 1.

Thread B: Decrement retrieved
value; result is -1.

Thread A: Store result in c; c is
now 1.

Thread B: Store result in c; c is
now -1.

Thread interference

Thread interference bugs are **difficult to detect** and fix, since the execution order of Threads is not controlled.

Memory consistency error

It occurs when different threads have inconsistent views of what should be the **same data**.

Memory consistency error Example

```
class Counter {  
    private int c = 0;  
    public void increment()  
    {  
        c++;  
    }  
    public void decrement()  
    {  
        c--;  
    }  
    public int value()  
    {  
        return c;  
    }  
}
```

Thread A invokes increment
Thread B invokes value

The obtained value might well be "0", because there's no guarantee that thread A's change to c will be visible to thread B.

Memory consistency error

Solution: To establish a happens-before relationship between related statements.

Happens-before relationship

The results of a write by one thread are **guaranteed** to be visible to a read by another thread only if the **write** operation **happens-before** the **read** operation.

Happens-before relationships

❑ start()

A call to start on a thread happens-before any action in the started thread

❑ join()

All actions in a thread happen-before any other thread successfully returns from a join on that thread.

Happens-before relationships

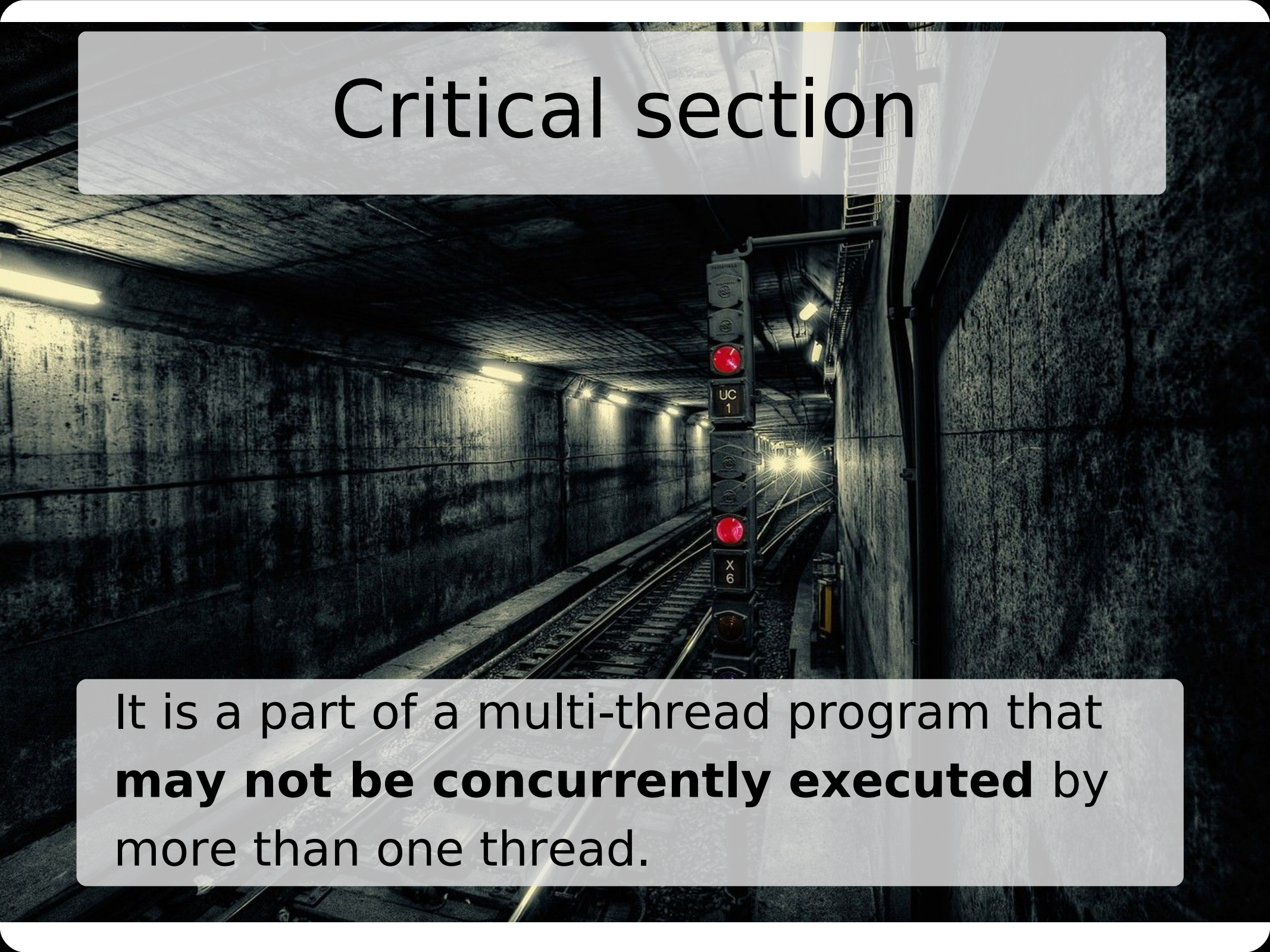
- ☐ synchronized
- ☐ volatile

A write to a volatile field happens-before every subsequent read of that same field. It does not entail mutual exclusion locking.



Synchronization

Critical section

A photograph of a subway tunnel. The walls are made of dark, textured concrete. The floor has metal tracks. A signal light is visible on the right side of the tracks, showing a red light. The tunnel is dimly lit with overhead lights.

It is a part of a multi-thread program that **may not be concurrently executed** by more than one thread.

Critical section

A photograph of a subway tunnel. The tunnel has concrete walls and a ceiling with recessed lighting. Tracks run down the center of the tunnel. A signal light is visible on the right side of the tracks, showing a red light. The signal light has labels 'UC 1' and 'X 6'.

Access to shared
resources as

- ☐ Files
- ☐ Shared variables or objects
- ☐ Data base resources

Mutual exclusion

It is the requirement of ensuring that no two concurrent Threads are in their critical section at the same time.

Synchronized code

Synchronized method

```
synchronized public void syncMethod(){  
    //...  
}
```

Synchronized statement

```
synchronized (this)  
{  
    //....  
}
```


Synchronized code

- ☐ It allows establishing happens before relation-ships.
- ☐ It ensures mutual exclusion.

Synchronized methods

They enable a simple strategy for **preventing** thread interference and memory consistency **errors**.

Synchronized methods

Reads or writes to an object's variables, visible to more than one thread, are done through synchronized methods.

Synchronized methods

Final fields can be safely read through non-synchronized methods, once the object is constructed.

Synchronized methods

When one thread is executing a synchronized method for an object, all other threads that invoke synchronized methods for the same object are **blocked** until the first is done.

Synchronized methods

This automatically establishes a **happens-before relationship** with any subsequent invocation of a synchronized method for the same object.

Synchronized methods

A Thread blocked by this means, is set to the **ready state** when the execution of the synchronized method

- ☐ ends
- ☐ A return occurs
- ☐ An Exception occurs

Synchronized methods

Example

```
class Counter {  
    private int c = 0;  
    public synchronized void increment() {  
        c++;  
    }  
    public synchronized void decrement() {  
        c--;  
    }  
    public synchronized int value() {  
        return c;  
    }  
}
```

Synchronized methods

Example

Thread A invokes increment
Thread B invokes decrement

Thread A: Retrieve $c=0$.

Thread A: Increment retrieved value; result is 1.

Thread A: Store result in c ; c is now 1.

Thread B: Retrieve $c=0$.

Thread B: Decrement retrieved value; result is -1.

Thread B: Store result in c ; c is now -1.

Intrinsic locks

- ❑ Synchronization is built around an internal entity known as the **intrinsic lock** or **monitor lock**.
- ❑ Every object has an intrinsic lock associated with it.

Static synchronized methods

When a **static synchronized method** is invoked, the thread acquires the intrinsic lock for the Class object associated with the class.

Synchronized statements

It must be specified the object that provides the **intrinsic lock**

synchronized (**this**)

```
{  
    //....  
}
```

Synchronized statements

```
public void addName(String name) {  
    synchronized(this) {  
        lastName = name;  
        nameCount++;  
    }  
    nameList.add(name);  
}
```

The method needs to synchronize changes to `lastName` and `nameCount`, but to avoid synchronizing invocations of other objects' methods.

Synchronized statements

```
public class MsLunch {  
    private long c1 = 0;  
    private long c2 = 0;  
    private Object lock1 = new Object();  
    private Object lock2 = new Object();  
    public void inc1() {  
        synchronized(lock1) {c1++;}  
    }  
    public void inc2() {  
        synchronized(lock2) {c2++;}  
    }  
}
```

They are also useful for improving concurrency with fine-grained synchronization.

Synchronized statements

```
public class MsLunch {  
    private long c1 = 0;  
    private long c2 = 0;  
    private Object lock1 = new Object();  
    private Object lock2 = new Object();  
    public void inc1() {  
        synchronized(lock1) {c1++;}  
    }  
    public void inc2() {  
        synchronized(lock2) {c2++;}  
    }  
}
```

Instead of using synchronized methods or otherwise using the lock associated with this, we create two objects solely to provide locks.

Reentrant synchronization

This describes a situation where synchronized code, directly or indirectly, invokes a method that also contains synchronized code, and both sets of code use the same lock.

Guarded block

```
public void guardedJoy() {  
    //Simple loop guard. Wastes processor time.  
    //Don't do this!  
    while(!joy) {}  
        System.out.println("Joy has been achieved!");  
}
```

joy is a shared variable
set by another thread

Such a block begins by polling a condition that must be true before the block can proceed.

Guarded block

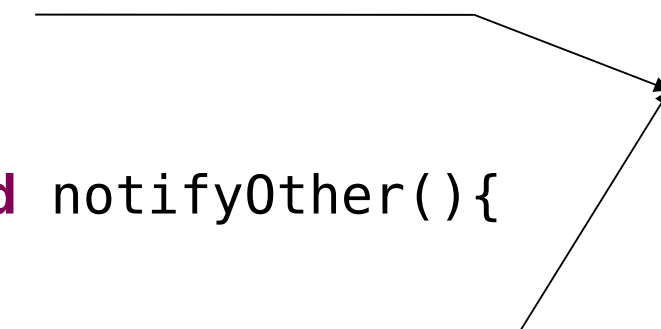
```
public synchronized void guardedJoy() {  
    //This guard only loops once for each special  
    //event, which may not be the event we're waiting for.  
    while(!joy) {  
        try {  
            wait();  
        }  
        catch (InterruptedException e) {}  
    }  
    System.out.println("Joy and efficiency have been  
    achieved!");  
}
```

A more efficient guard invokes `Object.wait` to suspend the current thread.

Guarded block

```
public synchronized notifyJoy(){  
    joy = true;  
    notifyAll();  
}
```

```
public synchronized notifyOther(){  
    other = true;  
    notifyAll();  
}
```



The thread which
invoked guardedJoy
becomes ready

The invocation of wait does not return until another thread has issued a notification that some special event may have occurred — though not necessarily the event this thread is waiting for

Synchronization issues

Liveness

It is the ability of a concurrent application to execute in a **timely** manner.

Liveness problems

- ❶ Deadlock
- ❷ Starvation
- ❸ Live-lock

Deadlock

It's a situation where two or more threads are **blocked forever**, waiting for each other.

Deadlock - Example

```
public class Deadlock {  
    static class Friend {  
  
        private final String name;  
        public Friend(String name)  
        {this.name = name;}  
        public String getName()  
        {return this.name;}  

```

```
        public synchronized void bow(Friend bower) {  
            System.out.format("%s: %s"+" has bowed to me!\n",  
                              this.name, bower.getName());  
            bower.bowBack(this);  
        }  
        public synchronized void bowBack(Friend bower) {  
            System.out.format("%s: %s"+" has bowed back to  
            me!\n", this.name, bower.getName());  
        }  
    }  
}
```



Deadlock - Example

```
public class Deadlock {  
    static class Friend {  
  
        private final String name;  
        public Friend(String name)  
        {this.name = name;}  
        public String getName()  
        {return this.name;}  

```

```
        public synchronized void bow(Friend bower) {  
            System.out.format("%s: %s"+" has bowed to me!\n",  
                              this.name, bower.getName());  
            bower.bowBack(this);}  
        public synchronized void bowBack(Friend bower) {  
            System.out.format("%s: %s"+" has bowed back to  
            me!\n", this.name, bower.getName());  
        }  
    }  
}
```

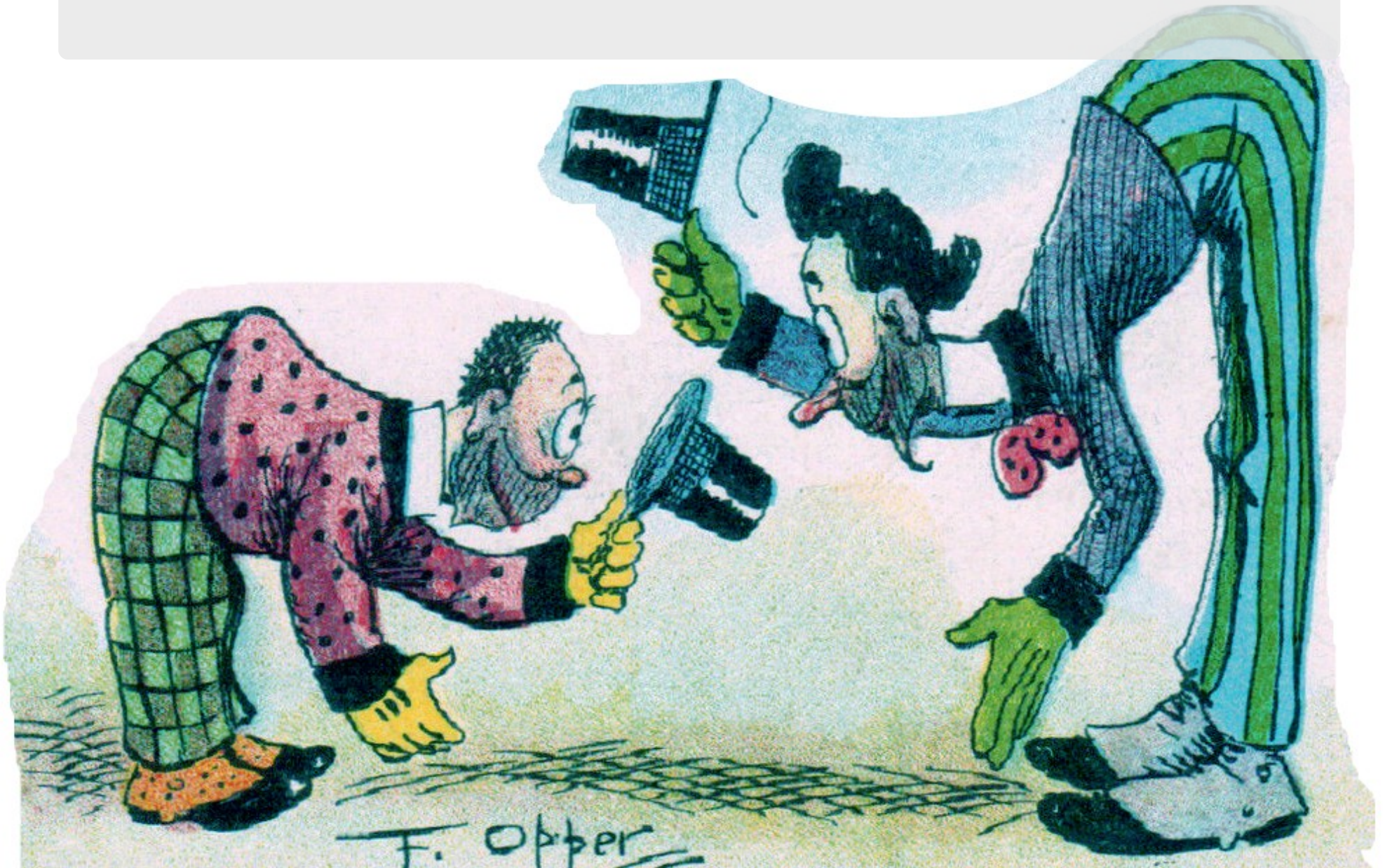


Deadlock - Example

```
public static void main(String[] args) {  
    final Friend alphonse =  
        new Friend("Alphonse");  
    final Friend gaston =  
        new Friend("Gaston");  
    new Thread(new Runnable() {  
        public void run() {  
            alphonse.bow(gaston); }  
    }).start();  
    new Thread(new Runnable() {  
        public void run() {  
            gaston.bow(alphonse); }  
    }).start();  
}
```



Deadlock



Starvation



It is a situation where a thread is unable to gain **regular access** to shared resources and is unable to make progress.

Starvation



This happens when shared resources are made unavailable for long periods by **greedy** threads.

Starvation



- ☐ An object provides a synchronized method that often takes a long time to return.
- ☐ Priority problems

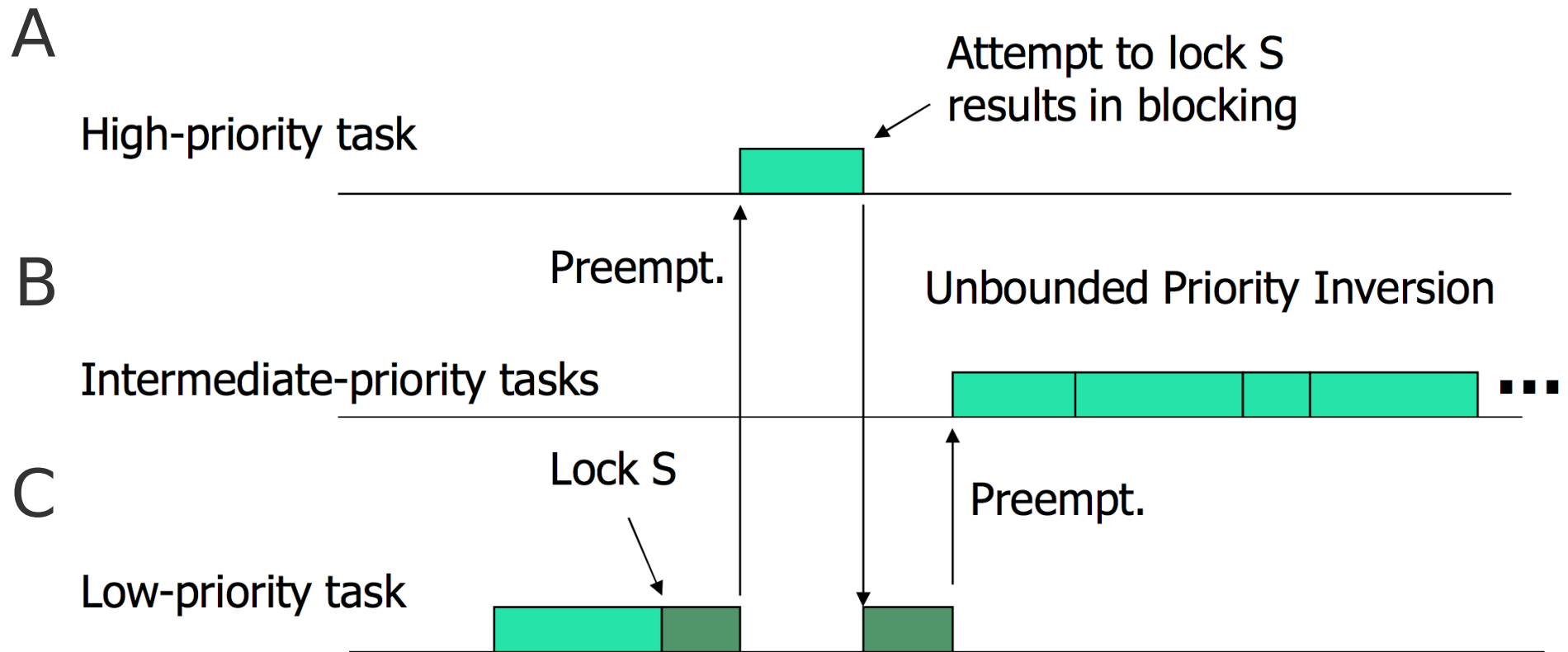
Starvation - Example

Priority inversion problem

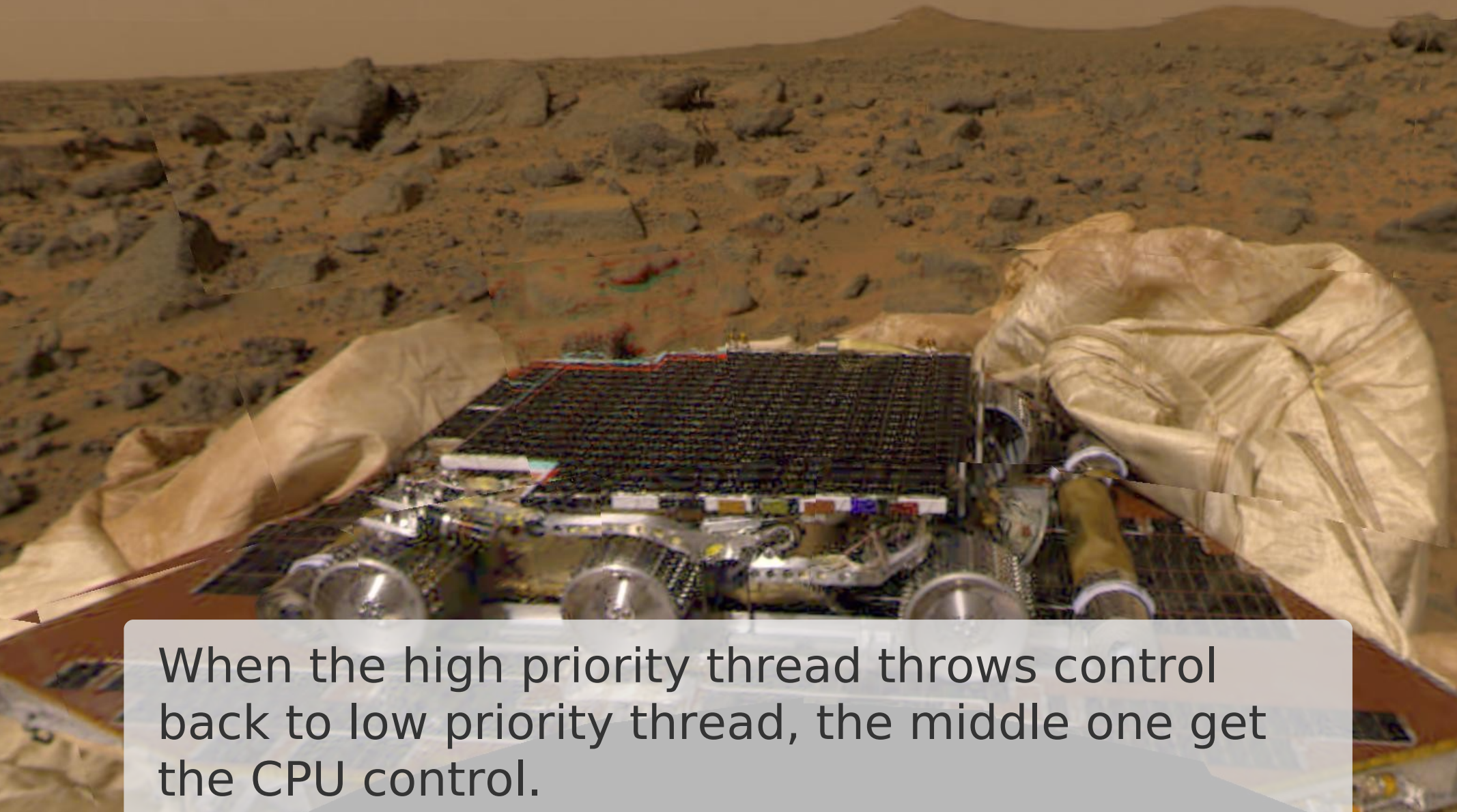
- ❑ Thread A is at high priority, waiting for result or resource from Thread C at low priority.
- ❑ Thread B at intermediate priority is CPU-bound.
- ❑ Thread C never runs, hence thread A never runs.

Starvation - Example

Priority inversion problem



Starvation - Example



When the high priority thread throws control back to low priority thread, the middle one get the CPU control.

Live-lock



Several threads spend all their time **trying to synchronize** instead of getting any work done.

Livelock



The threads are **not blocked**, they are too busy responding to each other to resume work.

Classical problems

- ➊ Producer-consumer
- ➋ Reader-writer
- ➌ Dining Philosophers
- ➍ Sleeping Barber's

Producer-Consumer problem

A photograph of a water bottle filling machine. Several clear plastic bottles with blue caps are positioned on a conveyor belt. The machine's metal components, including nozzles and structural frames, are visible in the background. The scene is industrial and brightly lit.

Two or more threads **communicate** with some threads “producing” data that others “consume”.

Producer-Consumer problem

A photograph of a water bottle filling machine. Several clear plastic bottles with blue caps are positioned in a row on a conveyor belt, waiting to be filled. The machine's metal components, including nozzles and structural frames, are visible in the background. The scene illustrates a buffer zone in a production line where items wait to be consumed.

- ☐ **Producer** generates things into a buffer
- ☐ **Consumer** takes those things and uses them

Producer-Consumer problem

A photograph of a water bottling machine. Several clear plastic bottles with blue caps are positioned on a conveyor belt. The machine's components, including metal pipes and mechanical parts, are visible in the background. The scene is set in a factory environment.

- ☐ **Producer** must wait if buffer is full
- ☐ **Consumer** mustn't extract data from empty buffer

Producer-Consumer

Guarded block - Example

```
public class Drop {  
    // Message sent from producer(P) to consumer(C).  
    private String message;  
    // true if C should wait for P to send message,  
    // false if P should wait for C to retrieve message.  
    private boolean empty = true;  
    public synchronized String take() {  
        // Wait until message is available.  
        while (empty) {  
            try {wait();}  
            catch (InterruptedException e) {}  
        }  
        // Toggle status.  
        empty = true;  
        // Notify producer that status has changed.  
        notifyAll();  
        return message;  
    }  
}
```

Producer-Consumer

Guarded block - Example

```
public synchronized void put(String message) {  
    // Wait until message has been retrieved.  
    while (!empty) {  
        try {wait();}  
        catch (InterruptedException e) {}  
    }  
    // Toggle status.  
    empty = false;  
    // Store message.  
    this.message = message;  
    // Notify consumer that status has changed.  
    notifyAll();  
}  
}
```


Producer-Consumer

Guarded block - Example

```
import java.util.Random;
public class Producer implements Runnable {
    private Drop drop;
    public Producer(Drop drop) {this.drop = drop;}
    public void run() {
        String importantInfo[] = {
            "Mares eat oats", "Does eat oats",
            "Little lambs eat ivy", "A kid will eat ivy too"};
        Random random = new Random();
        for (int i = 0; i < importantInfo.length; i++) {
            drop.put(importantInfo[i]);
            try {
                Thread.sleep(random.nextInt(5000));
            } catch (InterruptedException e) {}
        }
        drop.put("DONE");
    }
}
```

Producer-Consumer

Guarded block - Example

```
import java.util.Random;
public class Consumer implements Runnable {
    private Drop drop;
    public Consumer(Drop drop) {this.drop = drop;}
    public void run() {
        Random random = new Random();
        for (String message = drop.take();
            !message.equals("DONE");
            message = drop.take()) {
            System.out.format("MESSAGE RECEIVED: %s\n",
                message);
            try {
                Thread.sleep(random.nextInt(5000));
            } catch (InterruptedException e) {}
        }
    }
}
```

Producer-Consumer

Guarded block - Example

```
public class ProducerConsumerExample {  
    public static void main(String[] args) {  
        Drop drop = new Drop();  
        (new Thread(new Producer(drop))).start();  
        (new Thread(new Consumer(drop))).start();  
    }  
}
```