13 November, 2014



Team Number: 4 Consultation Day/Hour: Tuesday, 1pm

Team Name: Team 4

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CS3202 Software Engineering Project

Final Report

# Project Story

CS3202 has been an eventful project based module. In the continuation from CS3201, we have implemented the full scope of the Static Program Analyzer (SPA), alongside some extensions. We ensured that we improved on the weaknesses from our previous project and strengthened the base to build our SPA upon. We also had the addition of a new member with, alongside one member of our group opting to change groups. This switch of members actually improved the dynamics of our group, allowing us to work better with one another. The remainder of this report shall discuss in detail how we went about implementing the various components of the SPA. This discussion shall include a summary of our main achievements (Section 1), project plans (Section 2), UML diagrams (Section 3), design decisions (section 4), coding standards and experiences (Section 5), query processing (Section 6), testing (Section 7) and finally end off with a concluding discussion (Section 8).

# Summary of Main Achievements

## Basic SPA Functionality

For the purposes of the CS3202 development of the SPA, we have implemented the full SPA as described in the Project Handbook. This includes the implementation of the components:

* Parser
* Design Extractor
* Program Knowledge Base (PKB)
* Query Processor (QP).

The PKB stores the design abstractions implemented:

* Abstract Syntax Tree (AST)
* Follows/Follows\*
* Parent/Parent\*
* Modifies
* Uses
* Calls/Calls\*
* Next/Next\*
* Affects/Affects\*

The QP handles the processing of queries involving the aforementioned design abstractions alongside a combination of “with”, “such that” and “pattern” clauses. It has been implemented to also return tuple results. The QP also includes components that handle the optimization of query evaluation.

The description above just highlights the main functionality implemented. Overall, all of the required functions from iteration 1-3 have been implemented. The details of their implementation shall be discussed in the later sections of this report.

## Highlights of System

In the implementation of the functionality, defined by the handbook, we have ensured that aspects of our software standout from the norm. This is in the way that we have implemented some functions and also the addition of certain components. The main highlights of our project includes:

* Polymorphism of data structure for PKB (MapTable & ListTable)
* Next\* Implementation
* Addition of Query Representator (QR) and Query Optimizer (QO) in the QP
* Query Evaluator (QE) in the QP [really?]

These highlights will be further elaborated on in Section 4 and Section 6.

## Extension for Bonus Points

We have implemented the first extension for the extended code pattern. This includes the relationships Contains, Contains\* and Siblings. Details entailing the design decisions for this extension shall be discussed in Section 4. The test cases for this extension has been included alongside the other test cases for the SPA.

# Project Plans

## Project Schedule

The tables in this section show how we distributed the work into various tasks throughout the 4 main iterations.

### Iteration 1

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| **Team Member** | **Testing** | **Writing Test Cases** | **Revamp of PKB Tables** | **Refractoring QP** | **Working on QP** | **Extending Parser Functionality** |
| Azima | \* | \* | \* |  |  |  |
| Saima | \* |  |  |  | \* |  |
| Saloni | \* |  | \* |  |  |  |
| Sean | \* |  |  |  |  | \* |
| Tho | \* |  |  | \* | \* |  |

Table 2.1: Iteration 1 Work Distribution

Since this was the starting iteration, we had mainly focused on revamping the design of the whole system, from the Revision of the Prototype iteration. As a result, we were unable to finish implementing all of the required functionality defined in iteration 1.

### Iteration 2

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| **Team Member** | **Testing** | **Writing Test Cases** | **Revamp of TNODE** | **Next for PKB** | **Working on QP** | **Fixing issues with Parser** | **Report** |
| Azima | \* | \* |  |  |  |  |  |
| Saima | \* | \* |  |  | \* |  | \* |
| Saloni | \* | \* |  | \* |  |  |  |
| Sean | \* |  |  |  |  | \* |  |
| Tho | \* |  | \* |  | \* |  | \* |

Table 2.2: Iteration 2 Work Distribution

This iteration saw us complete all of the requirements from iteration 1 and iteration 2. The main setback in this iteration was in the design of the TNode data structure. This class is used to build various trees used for data storage. Previously we had implemented the TNode without the use of pointers. This means that every time we would pass data, the TNode would create a new copy instead of passing the original data from one function to another. This would slow down the SPA process, since each time we would add a new node, the existing tree would be copied another time to attach to the new node. This was seen to be a waste of processing time and storage. It slowed our system down considerably.

We overcame the aforementioned problem by creating a new TNode and assigning a pointer to its address. Basically, creating a pointer of type TNode.

Old way:

TNode node();

TNode \* pointer = &node;

New way:

TNode \* pointer = new TNode();

### Iteration 3

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| **Team Member** | **Testing** | **Writing Test Cases** | **Revamp of Relationships Data Structure** | **Affects** | **Working on QP** | **Fixing Issues with QE** | **Report** |
| Azima | \* | \* |  |  |  |  | \* |
| Saima | \* |  |  |  | \* |  | \* |
| Saloni | \* | \* |  |  |  |  | \* |
| Sean | \* |  | \* | \* |  |  |  |
| Tho | \* |  |  |  | \* | \* |  |

Table 2.3: Iteration 3 Work Distribution

Iteration 3 saw us having a number of issues with the implementation of the Next \* and Affects relationships. Also we had to improve on the managing of temporary results in the Query Evaluator, for multiple clauses. What these issues were and how we fixed them would be discussed in Section 4 and 6 respectively. This slowed us down considerably in this iteration and we could carry out aggressive testing, as we had initially planned to.

### Iteration 4

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| **Team Member** | **Testing** | **Writing Test Cases** | **Extension Relationships** | **Fixing Issues with System** | **Report** |
| Azima | \* | \* |  | \* | \* |
| Saima | \* | \* | \* | \* | \* |
| Saloni | \* | \* | \* | \* | \* |
| Sean | \* |  |  | \* | \* |
| Tho | \* |  |  | \* | \* |

Table 2.4: Iteration 4 Work Distribution

The emphasis placed on this iteration was to implement the extension relations, Siblings and Contains, and also on vigorous testing on the system.

## Organization of Meetings

As this project was implemented in a total of 4 iterations, we measured our progress based on the state of our SPA in the previous iteration and the requirements for the current iteration. To make sure that we stayed consistent, we met, on average, at least 3 times a week, including every Tuesday which was our consultation time slot. In the first meeting of the week we would set an agenda for the things to be achieved in that specific meeting, and for that week. The subsequent meetings would be organized around our tutor’s feedback and on improving our system based on the critique we would have received. Fig 2.1 below shows an excerpt from one of our meeting’s agenda.

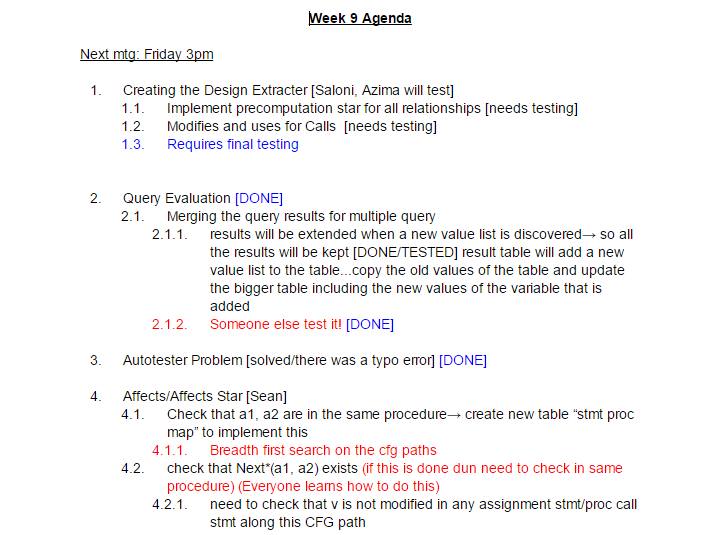


Figure 2.1: Agenda Excerpt

# UML Diagrams

The UML sequence diagrams presented in this section display how the SPA program flow works between the Parser, PKB and QP. These diagrams allowed us to visualize the various component interaction of the SPA and thus aided in the project planning.

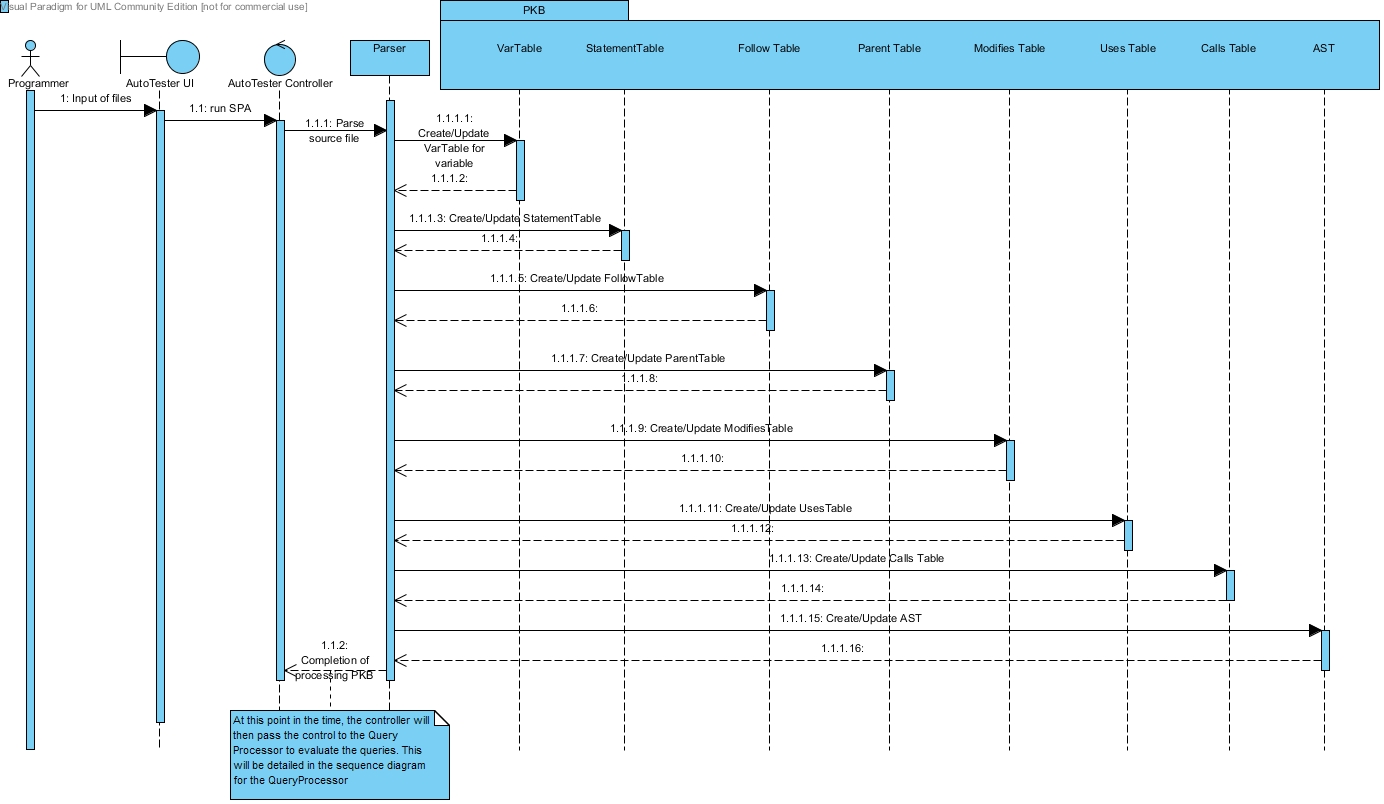


Figure 3.1: Sequence Diagram for Processing PKB

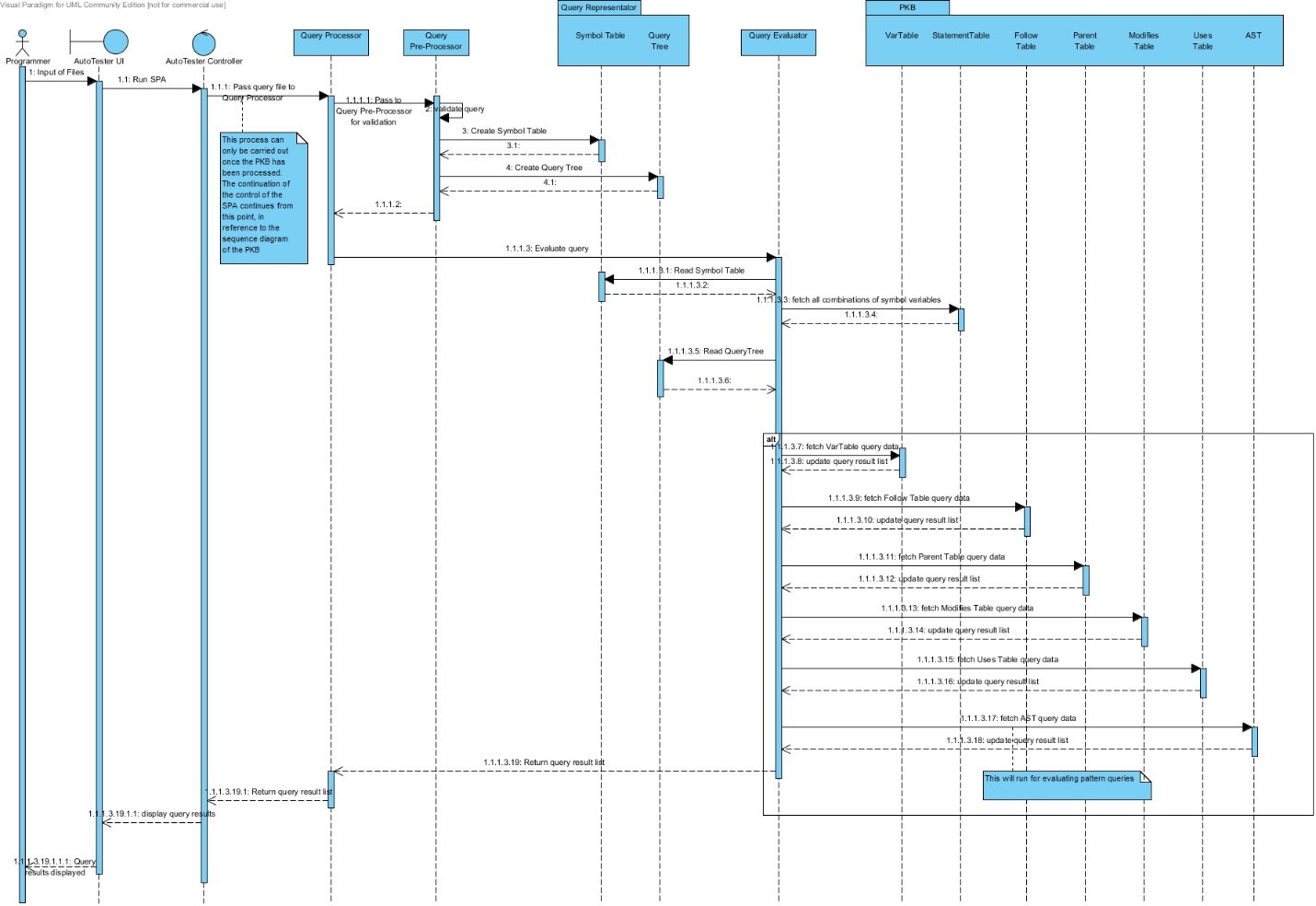


Figure 3.2: Sequence Diagram for Query Processor Flow

These 2 diagrams shown in Fig 3.1 and Fig 3.2, allowed us to understand the system properly before we went to code it out. In making these UML diagrams, we are able to separate the various components and abstract out the most important things. This enabled us to look at each decomposed aspect of the system and come up with an in-depth system architecture. This in turn allowed us to see what the basic requirements for each component were.

# Design Decisions

////talk abt separation of concerns and virtues of simplicity and standardization

//system decomposition with information hiding {section 10.1}

//show software architecture (emphasis on the differences in the QP)

## PKB Data Structure Representation

- discuss the design problem

-identify goals to be met by relevant design solutions

-discuss alternatives, previous implementations

-analyze strengths and weaknesses [include big o notation analysis]

-discuss implemented design based on the weakness of the previous design decisions

-analyze using big o notation

-give example for insertion, searching and getting for one relationship

## PKB API

## SPA Relationships

### AST

### Modifies/Uses [for procedure calls]

### Parent/Follows

### Calls

### Next/Next Star

### Affects/Affects Star

### Contains/Contains Star

### Siblings

## Design Patterns

# Coding Standards & Experiences

In terms of the coding standards, our group has decided to adopt the following naming conventions described in this section.

## Naming Conventions

### General Rules

* Do not use underscore, hyphen or any other non-alphabet characters.
* Any name should has all the first letters of internal words capitalized, e.g. getProcName()
* Avoid using abbreviations. Some words are acceptable in short forms, including: *Var*, *Proc*, *Stmt*, *AST*. Other words such as *Children*, *Number* should be fully spelled out.

### Specific Rules

* API Name:
  + API names should be nouns, in mixed case with the first letter of each internal word capitalized.
* Method:
  + Method names should be in the form of a verb. With method names containing more than one word, use mixed case with the first letter of each internal word capitalized.
  + Name of some specific methods:

1. Methods to insert new records to the database should have the form insertXXX().
2. Methods with return value type BOOLEAN should have the form: isXXX() e.g. isExist(), isMatchVar().
3. Methods with return types of other values should have the form getXXX() e.g. getVarName()
4. Methods that return the number of records inside a table/ list should have the form getSize().
5. Methods that change the values or status of an object should have the form: setXXX()
6. Methods that return values from star queries, such as Calls\* and Next\*, should have the form getXXXStar().

In all of the above examples, the “XXX” is used in place of the specific name which the method will adopt.

To keep the abstract and concrete PKB API in sync, we created a variable table, statement table and procedure tables. These tables are vectors mapping variable names to indexes of the vectors. So that a API method like BOOLEAN isModifies(STMT\_NUM s1, INDEX varIndex)understands that INDEX is the mapped value of a certain variable name, where INDEX is just an integer in C++ type.

# Query Processing

The query processing, the second major component of the SPA, handles the management and evaluation of various queries based on the Simple source parsed by the Parser. As is seen in the system architecture diagram in Fig [num], the QP contains components that we have added on top of the suggested components from the handbook. This would be the Query Representator (QR) and the Query Optimizer (QO). The reasoning for their addition and their functionality shal be discussed under the section for Query Evaluation (Section 6.2).

In this section we will go through how queries are first of all validated and then evaluated. For the discussion we will be referring to the sample query, shown below, to illustrate the process.

assign a; call c; while w; variable v;

Select <c.procName, a> such that Parent\*(w, a) and Uses(a, v) pattern a(\_, \_”2\*y + 3”\_) with c.procName = v.varName

## Query Validation

The query validation process is handled by the QueryPreprocessor (QPP). The QPP receives each query in a string form, and tries to read it from left to right, until the end of the string. During reading, if it meets a block of symbols which fits a PQL grammar rule, the QPP breaks the block apart from the string and validates that block using the corresponding grammar rule. This reading and breaking validating process continues until no more new block of symbols can be found. This is an example of how we have used a top-down approach to tackle this component.

For each undivided block of symbols, QPP will validate it using the grammar rules shown below. If any rule is violated, the QPP will send back information to the QR to handle this error. Otherwise, the information of the block will be saved into the respective components of the QR.

Grammar rules to be checked, line by line:

- On finding a line containing a declaration (e.g. stmt s;), QPP will send this part to the preprocessDeclaration() method.

- On finding a line containing a query part (e.g. Select s such that Follows(s, 1)), QPP will send this part to preprocessQueryPart() method.

- While the preprocessQueryPart() method runs, if the QPP finds a new clause of query, it will call the corresponding method preprocessClause() (e.g. preprocessSuchThatCondtion() ) for this clause. Currently, we have provided the methods for “*such that*” clause and “*pattern*” clause.

Based on the sample query shown above, the validation process is shown below:

In each iteration as a new block of symbols are found, they are enclosed by a square bracket. From iteration 2 onwards, the specific blocks have been colored black and red alternatively to show the distinction of the discovered blocks. This process continues on until everything has been validated. For each block of symbols, QPP will call the function corresponding to the grammar rule for this block, such as preprocessDeclaration() for [assign a;] or preprocessAttrRef() for [c.procName].

## Query Evaluation

### Data Representation (QR)

As was mentioned above, we have implemented a new component called the QR. This component basically stores all of the necessary information of the query, from the validation process, for the evaluation process. We decided to include a QR to deal with the practice of information hiding. With the abstraction of a new component, handling all of this data storage, it would be easier to separate the various concerns at each part of the query processing. Since a query contains two parts: a list of symbol declarations and the main query itself, the QR saves that information into 2 components of the QR:

1. SymbolTable: Used for the storing of all the declarations in the query. A declaration is separated and saved into 2 parts: the declared entity and symbol name.
2. QueryTree: Used for storing of the query itself. Starting from the keyword: “Select” until the end of the query. Each symbols of the query part is stored into a node, TNode and linked with others to form a query tree.
3. In addition to these 2 data structures, the QR also saves a BOOLEAN value to indicate whether the query is free of grammar errors or not. Later the QE will check this value first to decide if it should ignore the query evaluation in case the query is grammatically incorrect.

Shown below is how the sample query would look like in the QR.

SymbolTable:

|  |  |  |
| --- | --- | --- |
| Id | Type | Name |
| 0 | assign | a |
| 1 | call | c |
| 2 | while | w |
| 3 | variable | v |

Figure 6.1: SymbolTable for Sample Query

QueryTree:

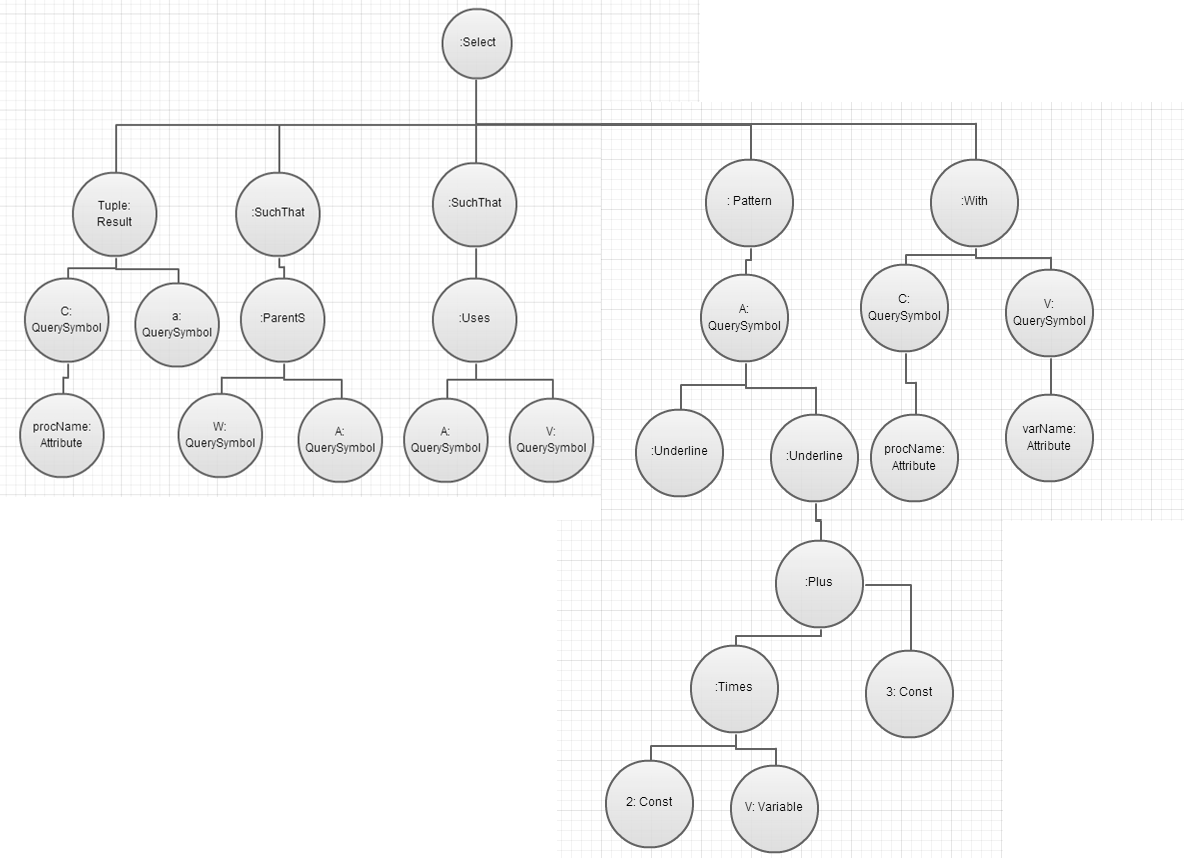


Figure 6.2: Query Tree of sample query

Once all of this information has been stored inside the QR, we can move on with the actual query evaluation, which is handled by the QE.

### Basic Query Evaluation

#### Manage the temporary results

It was noted in the old version of the SPA, in CS3201, that we did not have to worry about dealing with managing temporary results, since the queries did not contain multiple clauses. As a result, the main concern that we had was on how to extend the QE to deal with the management of temporary results.

###### Initial Implementation

The first method that we had implemented to manage temporary results was through a temporary vector inside of the QE. During the evaluation process, this vector would continuously be updated, for each clause until there is no more unsolved clause or an unsatisfiable clause is met. The QE would then use this vector to update the final result, and destroy this temporary result vector after that.

This naïve solution was sufficient to get us started on evaluating simple multiple clauses. However as more clauses were considered in a query, and as each query would get more complex, it was obvious that this solution was clearly insufficient. The main limitation of this solution was that it does not allow for a review of the temporary results, as they are discarded. This is difficult for programmers during debugging and updating the program. Moreover, as the QE only uses one vector to store temporary data and solving the various all clauses, it raises the complexity of the evaluation process. Also not every data from each temporary result is needed in solving a certain clause.

Using the sample query above, we figured out that the time complexity, using this implementation, would be O(N4 P) where N is the number of values for a query symbol and P is the time for a PKB’s operation. It is obvious that this complexity is highly costly and not efficient.

###### Actual Implementation

Given the aforementioned issues, we came up with a new solution that included the use of using tables to store and manage temporary results: a dynamic ResultTable and ResultManager. The ResultTable keeps the list of values for symbols used in the query and the ResultManager maintains the list of ResultTable. The ResultManager interacts with the QE during evaluation to record and return data when the QE demands. The main features of the ResultManager are:

* Returning of concise and non-duplicated data to the QE. Before evaluating a clause, the QE will ask the ResultManager for data relating to that clause. The ResultManager, after receiving a list of symbols used in the clause, will extract the data required from all ResultTable in the list. This extracted data is of the given symbols only, for a particular clause, and no data is duplicated.
* Minimize the space complexity as much as possible: After evaluation, a new ResultTable will be sent from QE to ResultManager. Before saving it to the ResultTable list, ResultManager will try to merge the existing tables with the new one by checking for shared symbols. If there are symbols shared between an existed table and the new table, ResultManager will merge those two tables together. Otherwise, ResultManager will do nothing. In this way, each symbol will only be saved once in the table, and non-related symbols (symbols which are not used in the same clause) will be kept in different tables. This approach cuts down the time and space needed to maintain and extract data from ResultManager.

With this new approach in mind, the symbols’ data are extracted from the ResultManager before evaluation. This means that the cost of solving each clause is O(N2P) at most. After that, when we insert the new ResultTable back to the ResultManager, the time complexity is O(N4M) for merging, where M is the merging time for 2 rows. However, since the number of values in a row is usually small, we can safely assume that M<<P. Thus, in comparison with the old approach, this new approach is definitely faster and more efficient.

Based on the sample query, the temporary result management is illustrated below. The illustration starts from evaluating the pattern clause in the query, assuming that the previous clauses have been evaluated, with some temporary results stored in the ResultManager

2. To solve the pattern clause, QE asks ResultManager to extract the values of symbol “a”. ResultManager will return the following table:

1. Evaluation of query is now at “pattern a(\_, \_”2\*y + 3”\_)”. At that time, ResultManager contains 1 ResultTable.

|  |  |  |  |
| --- | --- | --- | --- |
| **Id** | **w** | **a** | **v** |
| 0 | 1 | 2 | x |
| 1 | 1 | 4 | t |
| 2 | 1 | 8 | y |
| 3 | 3 | 4 | t |
| 4 | 3 | 8 | y |
| 5 | 5 | 8 | y |

|  |  |
| --- | --- |
| **Id** | **a** |
| 0 | 2 |
| 1 | 4 |
| 2 | 8 |

Table 6.1: Table returned to QE

3. After evaluation, the table now contains only data satisfying this pattern clause.

Table 6.2: ResutTable in ResultManager

|  |  |
| --- | --- |
| **Id** | **a** |
| 0 | 8 |

*(Continued on the next page)*

Table 6.3: Updated Table for values of "a"

4. The QE then asks the ResultManager to insert this new table into its table list. ResultManager must merge the existing table, in 1, with this new one, in 3. The table below is the final merged table.

|  |  |  |  |
| --- | --- | --- | --- |
| **Id** | **w** | **a** | **v** |
| 0 | 1 | 8 | y |
| 1 | 3 | 8 | y |
| 2 | 5 | 8 | y |

Table 6.4: Final Result after evaluating Pattern

### Optimization

The optimization process for the evaluation takes place after the validation in the QPP and before the evaluation in the QE. Since this part of the optimization is separated from the other functionalities, we decided to implement a Query Optimizer (QO). After all of the information has been stored inside of the QR, the QO is called to manipulate the results in the SymbolTable and QueryTree for optimization. The QO is in charge of the following 2 duties:

1. Rank all clauses of the given query based on 2 features:
   1. The clause’s type
   2. The number of query symbols used in each clause.

In our program, we give a higher rank to “with” clause, followed by the “such that” and “pattern” clauses. For the second feature, the higher the number of query symbols, the lower a rank the clause will have. For instance, a clause using 2 query symbols will have a lower ranking than another with 2 underlines and 1 query symbol.

1. Sort the query tree again using the ranking of the clauses. We assume that the clauses with a higher ranking will have a better time cost during evaluation, and thus we try to solve those clauses first. After the sorting, the query tree will have its clause nodes re-arranged, with the higher-ranked clause node in the front, following by the lower ones.

Referring to Fig 6.2 of the query tree, the order of the nodes, after optimization would be as follows:

Such-that - > such-that -> pattern-> with [Before Optimization]

Pattern -> with -> such-that -> such-that [After Optimization]

The pattern clause, while having a low ranking for its type, is placed first since it only contains 1 query symbol. For the rest of the clauses, each of them uses 2 query symbols, thus the order is mainly based on their types.

### Design Decisions

This section outlines how we have implemented each of the components of the QP. It mainly shows what data structure we have used.

###### Query Representator

The SymbolTree and QueryTree have been implemented by inheriting from the parent class Tree, which is also used for building the AST. By using the same data structure, we can compare 2 trees easily. This is useful for when we have to solve the pattern clause.

###### Query Optimizer

No specific data structure has been used. The data required is taken from the QPP after validation. There are 2 main functions in QO:

* RankTree() : Gives a ranking score to any clause node met during travelling the tree. The method applies a depth first search.
* SortTree() : Based on the rankings, QO will sort the clauses node from a high ranking to lower ones. This method uses merge sort for sorting.

###### Temporary Result Storage

* ResultTable

Saves a list of temporary results in a table form. The ResultTable provides basic methods to insert/ get/ delete data of the table, and one more method call extractTable(vector<string> symbols) to extract data of symbols in the input.

* ResultManager

Saves a list of pointers to different ResultTable. It provides 2 methods to insert and extract data from its tables.

# Testing

## Testing Experience

The testing of the SPA was carried out in three stages. Firstly, for each component of the PKB and PQL, we carried out unit testing. Unit testing allowed us to test the internal functions of each component. After unit testing was completed, we carried out integration testing of the different components. This ensured that the different components work properly together, for instance, the Parser and the Query Evaluator. Once we tested specific components, we then tested the system as a whole in validation testing. Keeping in mind that his project was implemented in iterations, we carried out this 3 stage testing procedure each time a new functionality was implemented. Even if it was for a new function in a preexisting component.

The most vigorous type of testing that was carried out was validation testing. We tested the system with hundreds of test cases and many complex source codes. The queries that we used to test covered a range of possibilities. From the most basic, to boundary cases, to where the thing to be returned could not be found in any of the tables.

## Examples of Test Cases

### Unit Testing

The aim of unit testing, would be to discover any logical errors present in the code. This allowed us to pinpoint specific errors, saving us the hassle of running into such errors during further stages of testing. Unit testing was done by manually inserting values and asserting that the function outputs were as expected.

The following example illustrates the unit testing of the ListTable component.

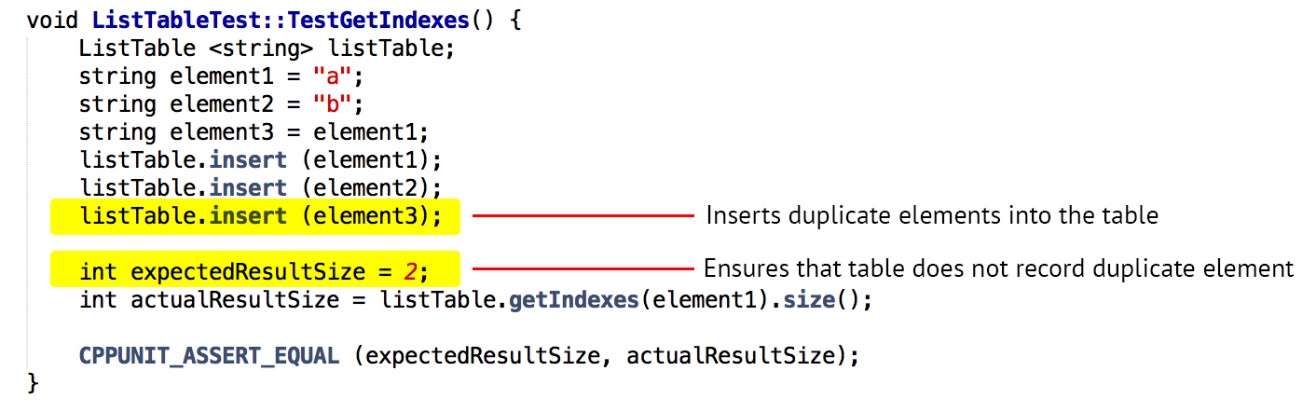


Figure 7.1: Unit Test Example

### Integration Testing

Many of the components work well when tested individually. However when they are integrated with other components, unforeseen errors may appear. This is the aim of integration testing, to identify such errors between the interactions of components. Such errors, in our case, were attributed to the fact thtt the PKB, Parser and PQL were written by different members. Hence the components had slightly different methods of implementation and different expected inputs and outputs. During integration testing, these flaws between the components became apparent to us. The Parser parses the simple source code, which is provided as the input file, and the PKB is built from it. In the following example, the integration between these components, including the Design Extractor, is illustrated when testing the Calls relationship.

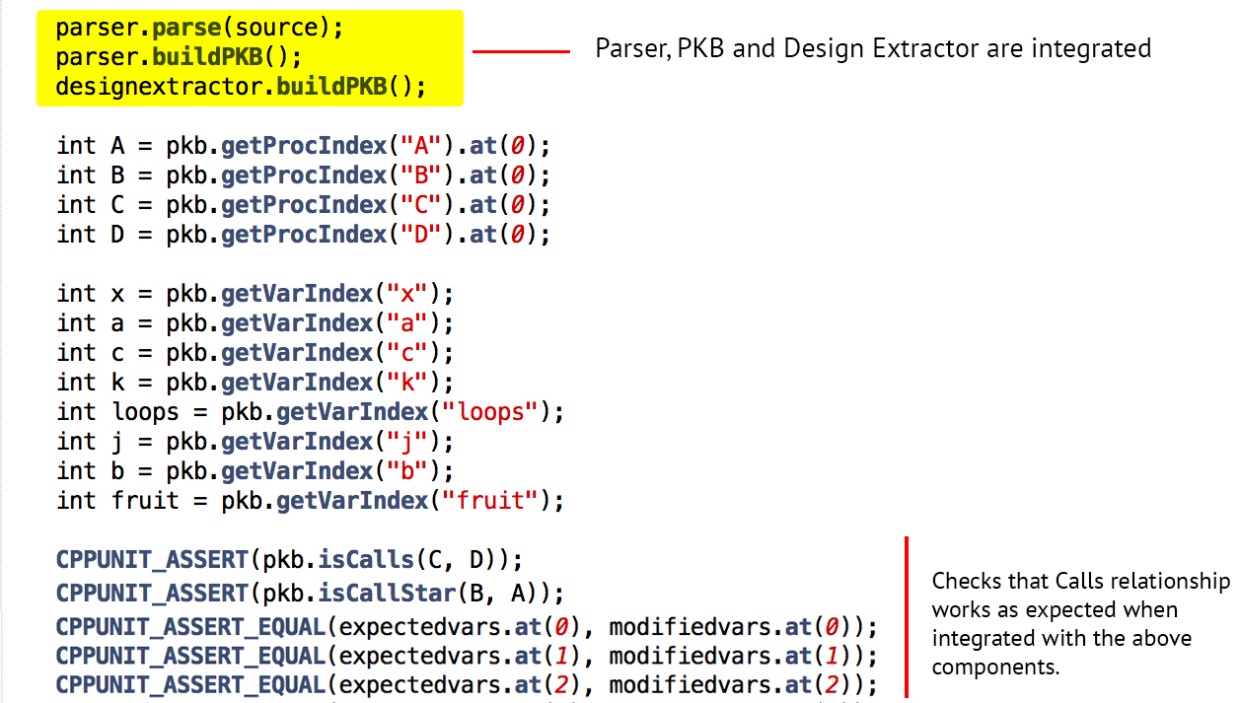


Figure 7.2: Integration Test Example

### Validation Testing

Validation testing, also known as System testing, is crucial to test the SPA system as a whole. There are many logical loopholes that might have been missed during unit and system testing, which can be spotted during extensive system testing. Thus, it is important to have a variety of test cases and source codes. The aim, in our case, was to ensure that the system can handle queries of multiple complexities and parse hundreds of lines of source code, including those haphazardly formatted.

The following Simple source code aims to test complex pattern queries by using multiple expressions and variables, the heavy use of brackets is also to be noted.

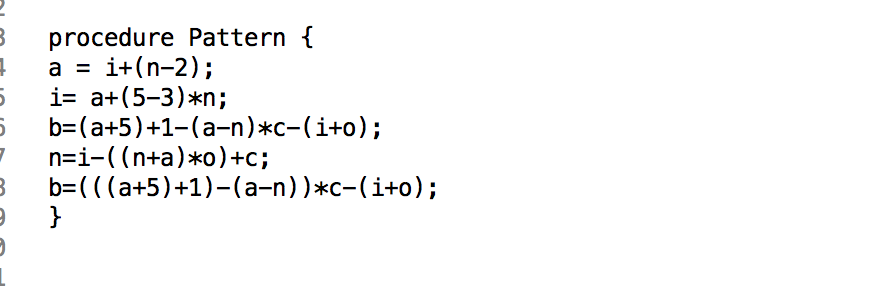


Figure 7.3: Simple Source for Validation Testing

The figure below shows two query examples. The first query tests the combination of multiple relationship clauses. This tests that the evaluated results for the different clauses are merged correctly to output only the correct answers that satisfy the whole query. This is essential for evaluation of any query that has more than one clause. The second query focuses on selecting variables by their attributes, such as ‘stmt#’ and ‘varName’. This is a commonly used selection method in queries, hence it’s essential to ensure that all such possible attribute selections have been addressed by the system.

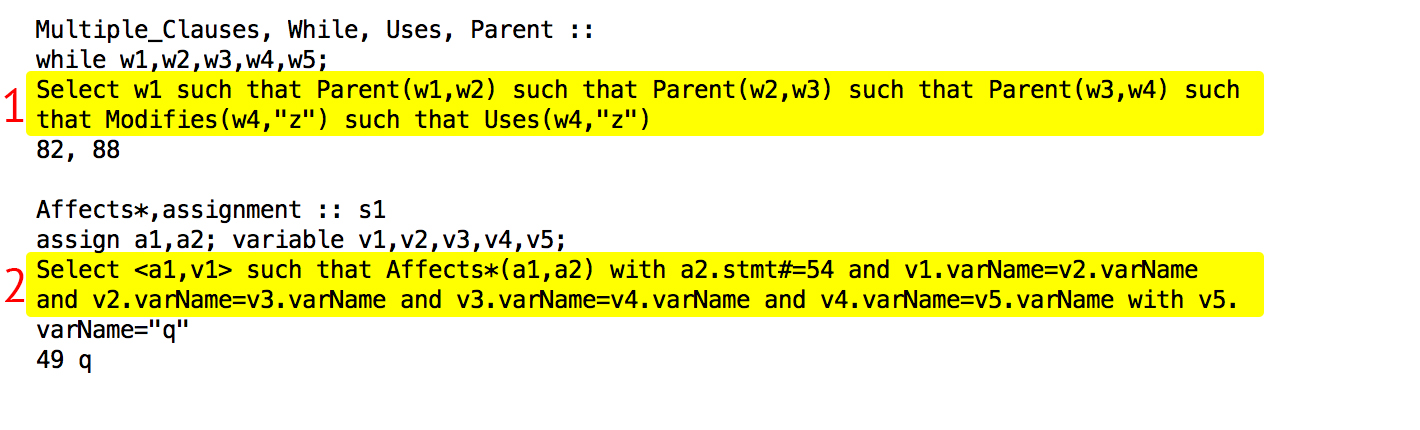


Figure 7.4: Queries for Validation testing

# Discussion

## Possible Improvements to Project

## Project Experience

### Team Work

### Incremental Development

### SPA Complexity & Proj solution

### Things learnt

## Tools used

-usefulness of recommended tools

-other tools used

…..

# Appendix A: Abstract PKB API

# Appendix B: Comments on Handbook

Overall, we found the handbook very helpful in giving us ideas for the implementation of the SPA. One thing that could have been better was if there had been more examples and perhaps some sample exercises for us to practice our concepts on. The extensions such as Contains and Siblings could also be explained briefly in the handbook instead of introducing them later on. This means we would have rough idea of how these relationships work beforehand and when it is time to consider whether or not to implement them, we can focus on finding the best way to integrate them rather than spending time understanding and grasping the new concepts of these relationships, which reduces crucial time which could be spent on the actual implementation.