

Lax Pairs and Double Bracket Flows

1 Definition of a Lax Pair

Two operator A and B are said to form a lax pair if they satisfy the equation

$$\frac{dA(t)}{dt} = [B(t), A(t)] \quad (1)$$

2 Unitary Nature of the Flow

It can be shown that this defines a unitary time evolution on $A(t)$, in the following manner. Let $U(t, t_0)$ be the unitary operator that carries this evolution through. We then need to construct a $U(t, t_0)$.

$$A(t) = U(t, t_0)A(t_0)U^\dagger(t, t_0) \quad (2)$$

where $A(t_0)$ is the operator A at a particular time t_0 . The time change of A can then be written as

$$\begin{aligned} \frac{dA(t)}{dt} &= \frac{dU(t, t_0)}{dt}A(t_0)U^\dagger(t, t_0) + U(t, t_0)A(t_0)\frac{dU^\dagger(t, t_0)}{dt} \\ &= \frac{dU(t, t_0)}{dt}U^\dagger(t, t_0)A(t) + A(t)U(t, t_0)\frac{dU^\dagger(t, t_0)}{dt} \quad [A(t) = UAU^\dagger] \\ &= \frac{dU(t, t_0)}{dt}U^\dagger(t, t_0)A(t) - A(t)\frac{dU(t, t_0)}{dt}U^\dagger(t, t_0) \quad [UU^\dagger = 1] \\ &= \left[\frac{dU(t, t_0)}{dt}U^\dagger(t, t_0), A(t) \right] \end{aligned} \quad (3)$$

Looking at the definition of a lax pair, we can now make the connection

$$B(t) = \frac{dU(t, t_0)}{dt}U^\dagger(t, t_0) \quad (4)$$

The equation of motion characterised by the lax pair eq. 1 can thus be said to generate a family of unitarily connected operators $A(t)$, related by the unitaries defined by eq. 4. A direct corollary is that the spectrum of $A(t)$ is preserved during this evolution.

3 Double Bracket Flow

The double bracket flows correspond to a special choice of the operator $B(t)$: $B(t) \equiv [A(t), C]$. A consequence of this choice is that the lax pair evolution then serves to minimize the commutator $[A(t), C]$. To see how, we first write down a function

$$\chi \equiv \text{Tr} \left([A(t) - C]^2 \right) = \text{Tr} [A(t)^2 + C^2 - A(t)C - CA(t)] \quad (5)$$

Since $A^2(t) = U A^2 U^\dagger$, we get $\text{Tr}(A^2(t)) = \text{Tr}(A^2)$. Also, from the cyclic nature of trace, we can write $\text{Tr}(A(t)C) = \text{Tr}(CA(t))$. These considerations (and the fact that C does not depend on t) allows us to write

$$\frac{d\chi}{dt} = -2\text{Tr} \left(\frac{dA(t)}{dt} C \right) = -2\text{Tr} ([B(t), A(t)] C) \quad (6)$$

Using the cyclic property of trace, this becomes

$$\begin{aligned} \text{Tr} ([B(t), A(t)] C) &= \text{Tr} (B(t)A(t)C - A(t)B(t)C) \\ &= \text{Tr} (B(t)A(t)C - B(t)A(t)C) \\ &= \text{Tr} (B(t) [A(t), C]) \end{aligned} \quad (7)$$

If we now substitute the choice of $B(t)$ we made above, we get

$$\frac{d\chi}{dt} = -2\text{Tr} \left([A(t), C]^2 \right) \leq 0 \quad (8)$$

Since χ , the way it is defined, must necessarily be positive semi-definite for all t , the derivative $\frac{d\chi}{dt}$ must vanish in the limit $t \rightarrow \infty$, otherwise $\chi(t)$ will become negative. This gives the result

$$\lim_{t \rightarrow \infty} \frac{d\chi}{dt} = -2 \lim_{t \rightarrow \infty} \text{Tr} \left([A(t), C]^2 \right) = 0 \implies \lim_{t \rightarrow \infty} [A(t), C] = 0 \quad (9)$$

In other words, **the lax pair evolution of $A(t)$ against $[A(t), C]$ leads to the diagonalization of $A(t)$ with respect to C .** This can be used as an iterative algorithm to diagonalize a general matrix with respect to another matrix:

- Define matrices A and B, A being the one we want to diagonalize w.r.t B
- Iteratively run the next two steps until a desired accuracy is reached
- Compute a new matrix $C = A*B - B*A$
- Change A as follows: $A = A + C*A - A*c$