Proving the Regular Value Theorem in Lean

Nicolò Cavalleri, Michela Barbieri, Sebastian Monnet

LSGNT

2021

• Computer software that lets us state and prove theorems rigorously.

- Computer software that lets us state and prove theorems rigorously.
- Verifies that each step in a proof is fully justified.

- Computer software that lets us state and prove theorems rigorously.
- Verifies that each step in a proof is fully justified.
- Recently got a lot of attention for helping Peter Scholze check his work.

- Computer software that lets us state and prove theorems rigorously.
- Verifies that each step in a proof is fully justified.
- Recently got a lot of attention for helping Peter Scholze check his work.
- Based on a foundational system called type theory.

• A type is a collection of objects, called terms.

• A **type** is a collection of objects, called **terms**.

Example

The type nat has terms $0, 1, 2, \ldots$

• A type is a collection of objects, called terms.

Example

The type nat has terms $0, 1, 2, \ldots$

• For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".

• A **type** is a collection of objects, called **terms**.

Example

The type nat has terms $0, 1, 2, \ldots$

- For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".
- There is a type called Prop.

• A **type** is a collection of objects, called **terms**.

Example

The type nat has terms $0, 1, 2, \ldots$

- For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".
- There is a type called Prop.
- For each p : Prop, there is a type Proof(p) of "proofs of p", which has at most one term.

• A type is a collection of objects, called terms.

Example

The type nat has terms $0, 1, 2, \ldots$

- For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".
- There is a type called Prop.
- For each p : Prop, there is a type Proof(p) of "proofs of p", which has at most one term.
- We say that p is true if and only if the type Proof(p) is inhabited.

• A **type** is a collection of objects, called **terms**.

Example

The type nat has terms $0, 1, 2, \ldots$

- For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".
- There is a type called Prop.
- For each p : Prop, there is a type Proof(p) of "proofs of p", which has at most one term.
- We say that p is true if and only if the type Proof(p) is inhabited.
- For any propositions p, q, the type $Proof(p) \to Proof(q)$ is called "p implies q".

• A **type** is a collection of objects, called **terms**.

Example

The type nat has terms $0, 1, 2, \ldots$

- For any pair of types α, β , there is another type $\alpha \to \beta$ of "functions from α to β ".
- There is a type called Prop.
- For each p : Prop, there is a type Proof(p) of "proofs of p", which has at most one term.
- We say that p is true if and only if the type Proof(p) is inhabited.
- For any propositions p, q, the type Proof(p) → Proof(q) is called "p implies q".

Example

The implication $Proof(p) \to Proof(p)$ is true for any proposition p, since we can take the identity function as our term.



 Our original goal was to prove that a smooth variety (in the sense of scheme theory) over a normed field K is always a smooth manifold over K.

- Our original goal was to prove that a smooth variety (in the sense of scheme theory) over a normed field K is always a smooth manifold over K.
- This proved too hard for the time constraints, so instead we decided to prove a key step in the theorem, the Regular Value Theorem.

- Our original goal was to prove that a smooth variety (in the sense of scheme theory) over a normed field K is always a smooth manifold over K.
- This proved too hard for the time constraints, so instead we decided to prove a key step in the theorem, the Regular Value Theorem.

Theorem (Regular Value Theorem)

Let $f: M \to N$ be a smooth map from a smooth m-manifold to a smooth n-manifold, where m > n. Let $c \in N$ be a regular value of f with nonempty preimage.

Then $f^{-1}(c) \subseteq M$ is a smooth submanifold of M of dimension m-n.

- Our original goal was to prove that a smooth variety (in the sense of scheme theory) over a normed field K is always a smooth manifold over K.
- This proved too hard for the time constraints, so instead we decided to prove a key step in the theorem, the Regular Value Theorem.

Theorem (Regular Value Theorem)

Let $f: M \to N$ be a smooth map from a smooth m-manifold to a smooth n-manifold, where m > n. Let $c \in N$ be a regular value of f with nonempty preimage.

Then $f^{-1}(c) \subseteq M$ is a smooth submanifold of M of dimension m-n.

This also proved too hard, so we decided to prove...



Our real Objective

• Geometry uses intuition.

- Geometry uses intuition.
- It often leaves technical details up to the imagination.

- Geometry uses intuition.
- It often leaves technical details up to the imagination.
- Since we have such strong pictures in our minds, there is usually a disconnect between our ideas and our axioms.

- Geometry uses intuition.
- It often leaves technical details up to the imagination.
- Since we have such strong pictures in our minds, there is usually a disconnect between our ideas and our axioms.
- When working with Lean, we need to be totally explicit about everything.

• Suppose that we have a map $f: M \to N$ from an m-manifold to an n-manifold, and let $p \in M$ be a point. Consider the statement

"Let $p \in U$ and $f(p) \in V$ be charts, and without loss of generality assume that p and f(p) lie at the origin."

• Suppose that we have a map $f: M \to N$ from an m-manifold to an n-manifold, and let $p \in M$ be a point. Consider the statement

"Let $p \in U$ and $f(p) \in V$ be charts, and without loss of generality assume that p and f(p) lie at the origin."

- Formally, this means
 - "There exist open subsets $U \subseteq M$ and $V \subseteq N$ such that $p \in U$ and $f(p) \in V$, along with diffeomorphisms $\varphi : U \to \mathbb{R}^m$ and $\psi : V \to \mathbb{R}^n$ such that $\varphi(p) = 0$ and $\psi(f(p)) = 0$."

- Suppose that we have a map $f: M \to N$ from an m-manifold to an n-manifold, and let $p \in M$ be a point. Consider the statement
 - "Let $p \in U$ and $f(p) \in V$ be charts, and without loss of generality assume that p and f(p) lie at the origin."
- Formally, this means
 - "There exist open subsets $U \subseteq M$ and $V \subseteq N$ such that $p \in U$ and $f(p) \in V$, along with diffeomorphisms $\varphi : U \to \mathbb{R}^m$ and $\psi : V \to \mathbb{R}^n$ such that $\varphi(p) = 0$ and $\psi(f(p)) = 0$."
- The existence of such charts is not guaranteed directly by the definition of a manifold. To prove it, you have to take charts around the points, and then explicitly define automorphisms of \mathbb{R}^m (respectively \mathbb{R}^n) that translate the relevant points to 0.

• Suppose that we have a map $f: M \to N$ from an m-manifold to an n-manifold, and let $p \in M$ be a point. Consider the statement

"Let $p \in U$ and $f(p) \in V$ be charts, and without loss of generality assume that p and f(p) lie at the origin."

- Formally, this means
 - "There exist open subsets $U \subseteq M$ and $V \subseteq N$ such that $p \in U$ and $f(p) \in V$, along with diffeomorphisms $\varphi : U \to \mathbb{R}^m$ and $\psi : V \to \mathbb{R}^n$ such that $\varphi(p) = 0$ and $\psi(f(p)) = 0$."
- The existence of such charts is not guaranteed directly by the definition of a manifold. To prove it, you have to take charts around the points, and then explicitly define automorphisms of \mathbb{R}^m (respectively \mathbb{R}^n) that translate the relevant points to 0.
- This is not difficult in itself, but it gives some idea of the considerations involved in formalising geometry.