Winograd Transform in Error Correcting Code theory

Translation and re-adaptation of the Master Thesis
La Trasformata di Winograd nella Teoria dei Codici Correttori
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Introduction

"Exact computations to start knowing all the existing things and all the obscure and mysterious secrets."

- Ahmes, 1600 a.C.

The starting point of this thesis are some *Laboratorio di applicazioni dell'algebra* laboratory held at the University of Turin in 2011 and a pre-print by prof. Umberto Cerruti [8] that I had used to learn several topics regarding Error Correcting Code theory that arose my interest and curiosity.
The main ones reported in this thesis are:

Chapter 1 Decomposition of the algebra $\mathcal{R}_{r,\mathbb{F}} = {\mathbb{F}[x]}\Big/{(x^r-1)}$ in a product of fields. Through the study of the factorisation of x^r-1 through a group isomorphic to the Galois group $Gal(\mathbb{F}(\xi),\mathbb{F}))$ acting on the group generated by x over $\mathcal{R}_{r,\mathbb{F}}$ we build a new algebra over the irriducible factors of x^r-1. For $M^{(v)}(x)$ irriducible factor of x^r-1, then every quotient

$$\mathbb{F}[x]\Big/M^{(v)}(x)$$

is a field, and we can build a new algebra as the product of these fields. This is still isomorphic to $\mathcal{R}_{r,\mathbb{F}}$ and the isomorphism between them is called **Winograd transform**.
Still in Chapter 1 we prove a formula based on the Burnside theorem to determine the cardiality of the set of irriducible factors $M^{(v)}(x)$.

Chapter 2 Study of $\mathcal{R}_{r,q}$ ideals and idempotents: from this chapter onwards we account only for the finite fields, as this is where the applications considered in the subsequent chapters are. After defining some of the operators over $\mathcal{R}_{r,q}$, we present a study over the ideals and idempotents, playing a fundamental role in the Error Correcting Code theory. Given the field's ideals and idempotents simplicity when represented in the new algebra defined in the previous chapter, the analysis is not carried forward over $\mathcal{R}_{r,q}$.

Chpater 3 Winograd transform as linear transform between $\mathcal{R}_{r,q}$ and its splitting product: we dive into the definition of Winograd transform, obtaining its transformation matrix and its inverse. We also present some of its more relevant properties. To define the transformation matrix in the most direct way, we start from the definition of **interlinked circulant vectors**. This

is still isomorphic to the ones already provided, and it keeps the structure of product of fields, though with a simple structure to allow to talk about transformation matrix. In this chapter too there can be found several numerical examples.

Chapter 4 and 5 Error correcting code theory introduction: in this interlude we present the error correcting code theory from its basics, to define linear codes, cyclic codes and BCH codes. In this chapter we will be using most of the results provided in chapters 1 and 2 and we prepare the terrain for presenting the applications of Winograd transform in error correcting code theory, aim of the thesis.

capitolo 6 Applications of the study of $\mathcal{R}_{r,q}$ and the Winograd transform in the error correcting code theory: the first applicatoin we see that a choice of Winograd transform blocks defines a matrix, that can work as control matrix as well as generating matrix. The second application is a system to encode a message diminishing the quantity of information, detecting in each word subvectors whose information contained can be neglected.

Originally, Winograd transform was discovered as a tool to decrease the computational complexity of convolution product, and as alternative to the Discrete Fourier Transform (DFT).
It had been introduced in the Shmuel Winograd's paper *On Computing the Discrete Fourier Transform* [31].
In this research I omitted the connection between the Winograd transform and DFT, we omitted any reference to the spectral code theory, and the application of the Winograd transofrm to error correcting codes discovered by Miller Truong and reed is nowhere to be found (*Efficient Program for decoding the* $(255, 223)$ *Reed-Solomon Code over* $GF(2^8)$ [22]).
I had instead considered the Winograd transform as a linear transform between two spaces, and the consequent direct applications. Main sources for this research, other than the already cited Cerruti's pre-print are *Theory and Practice of Error Control Codes*, di Richard E. Blahut [4] and *Algebra e teoria dei codici correttori* di Luigia Berardi [2].

Contents

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Chapter 0

Appendix

0.1 Recipe: transform a sentence into a list of binary numbers and backwards

Given an alphabet $\mathcal{A}$, that is an array of symbols (typically letters, though any kind of symbols would work), we want to transform any sequence of its elements (or a sentence) into an array of binary numbers.

Consider

$$\mathcal{A} = [\text{“ ”}, \text{“}n\text{”}, \text{“}a\text{”}, \text{“}u\text{”}, \text{“}y\text{”}, \text{“}l\text{”}, \text{“}a\text{”}, \text{“}f\text{”}, \text{“}m\text{”}] \tag{1}$$

where the first element is the empty space, and consider the sentence

$$\text{“}my\ funny\ val\text{”}.$$

We first replace the sentence with the sequence of the numbers corresponding to the positions of the letters in the list $\mathcal{A}$:

$$\text{“}my\ funny\ val\text{”} \longrightarrow [8, 4, 0, 7, 3, 1, 1, 4, 0, 6, 2, 5].$$

Then we chose a length to group the numbers in the array together. Here we chose 5, and we pad the last two numbers with zeros to have all the sub-lists of the same length:

$$[8, 4, 0, 7, 3, 1, 1, 4, 0, 6, 2, 5] \longrightarrow [[8, 4, 0, 7, 3], [1, 1, 4, 0, 6], [2, 5, 0, 0, 0]].$$

Each one of the sub-list is called here a *worm*, and its length *worm length*. Now we must convert a worm into an integer, that will then be represented into its binary form. To this end we use the length of the alphabet as the base for the number:

$$[8, 4, 0, 7, 3] \longrightarrow 8 + 4|\mathcal{A}| + 0|\mathcal{A}|^2 + 7|\mathcal{A}|^3 + 3|\mathcal{A}|^4 = 24830$$

And in general:

$$[w_1, w_2, \ldots, w_{\mathrm{w}}] \longrightarrow \sum_{j=0}^{\mathrm{w}-1} w_j |\mathcal{A}|^j \tag{2}$$

We can then convert all the worms into:

$$[[8,4,0,7,3],[1,1,4,0,6],[2,5,0,0,0]] \longrightarrow [24830,39700,47].$$

And then all the integers into their binary representation:

$$[24830,39700,47] \longrightarrow [0110000011111110,1001101100010100,0000000000101111]$$

Where the binaries numbers have been padded with zeros so to have all the same length (the sought length is the length of the binary representation of the largest integer $|\mathcal{A}|^{\mathrm{w}} - 1$).

To go backward, the only critical passage is the inverse of 2, transforming an integer n into the corresponding worm w. This can be done recursively with:

$$\begin{aligned} &w[0] := n \bmod |\mathcal{A}| \\ &\text{for } i = 1 \ldots \mathrm{w} - 1 \\ &\quad n := \frac{n - w[i-1]}{|\mathcal{A}|} \\ &\quad w[i] := n \bmod |\mathcal{A}| \end{aligned}$$

where $n \bmod |\mathcal{A}|$ is the reminder of n modulo n $|\mathcal{A}|$.

0.2 Recipe: find the minimal polynomials over a finite field

To find the minimal polynomial over a finite field there is a procedure that does not involve the research for the group H isomorphic to the Galois group $Gal(\mathbb{F}(\xi),\mathbb{F})$ that we show in one of the chapters of the Italian version. This method is a better candidate to be coded[1]. It is underpinned by the following definition, reformulating the meaning of conjugate elements, cyclotomic classes and minimal polynomial.

Definition 0.2.1. *For an element ξ of the fields $\mathbb{F}_{q^m}$ there is an integer t, the smallest one, so that $\xi^{p^t} = \xi$. We call the set*

$$C(\xi) = \{\xi, \xi^p, \xi^{p^2}, \ldots, \xi^{p^{t-1}}\}$$

cyclotomic class *and two elements of this set are called* **conjugates**. *The* **minimal polynomial of** ξ *is the smallest polynomial in $\mathbb{F}_q[x]$ having ξ as its root.*

With this definition at hand, it is then possible to prove[2] that the minimal polynomial of a non zero element of $\mathbb{F}_q$ is the smallest polynomial having as its roots all the element of the same cyclotomic class. Let's prove it formally:

Theorem 0.2.1. *Let ξ be a non zero element of $\mathbb{F}_q$, then its minimal polynomial, $M_\xi(x)$ is irruducible in $\mathbb{F}_q[x]$ and it is defined as*

$$M_\xi(x) = \prod_{\beta \in C(\xi)} (x - \beta)$$

[1] As also underlined in [6] pag.83.

[2] From [6], theorems 4.36 e 4.38.

Proof. Let ad absurdum $M_\xi(x)$ be not irriducible. Then it can be written as the product of two polynomials in $\mathbb{F}_q[x]$: $M_\xi(x) = f_1(x)f_2(x)$ for $0 \leq deg(f_j(x)) < deg(M_\xi(x))$. Since $M_\xi(\xi) = f_1(\xi)f_2(\xi) = 0$ and since $\mathbb{F}_q$ is a field, then one of the two polynomials in the facorisation of $M_\xi(x)$ admits ξ as its root, in contradiction with the minimality of $M_\xi(x)$.
To complete our knowledge upon the minimal polynomial (as well as to complete the proof) we have to show that $M_\xi(x) = \prod_{\beta \in C(\xi)}(x - \beta)$.
Let $M_\xi(x) = \sum_{j=0}^{t} m_j x^j$ be defined as the minimal polynomial of ξ. Since ξ is one of its roots then also ξ^p must be one of its roots:

$$M_\xi(\xi^p) = \sum_{j=0}^{t} m_j(\xi^p)^j = \sum_{j=0}^{t} m_j^p(\xi^p)^j \tag{3}$$

$$= \sum_{j=0}^{t}(m_j\xi^j)^p = (\sum_{j=0}^{t} m_j\xi^j)^p \tag{4}$$

$$= (M_\xi(\xi))^p = 0 \tag{5}$$

Moreover, thanks to the minimality $M_\xi(x)$ does not have other roots. Therefore:

$$M_\xi(x) = \prod_{\beta \in C(\xi)} (x - \beta)$$

Equation 3 holds only if we can say that $m_j^p = m_j$ which happens only if $m_j \in \mathbb{F}_q$.
On one hand:

$$(M_\xi(x))^p = \prod_{\beta \in C(\xi)} (x - \beta)^p = \prod_{\beta \in C(\xi)} (x^p - \beta^p)$$

$$= \prod_{\beta \in C(\xi)} (x^p - \beta) = (M_\xi(x^p)) = \sum_{j=0}^{t} m_j x^{jp}$$

And on the other hand:

$$(M_\xi(x))^p = \sum_{j=0}^{t}(m_j x^j)^p = \sum_{j=0}^{t} m_j^p x^{jp}$$

Therefore $m_j^p = m_j$ and $m_j \in \mathbb{F}_q$. □

Therefore obtaining the minimal polynomials is a matter of generating the cyclotomic classes starting from the elements of the field.

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