#### **CSE 102**

# Introduction to Analysis of Algorithms Master Theorem Practice Problems

#### **Master Theorem**

Let  $a \ge 1$ , b > 1, f(n) be asymptotically positive, and let T(n) be defined by T(n) = aT(n/b) + f(n). Then we have three cases:

- 1. If  $f(n) = O(n^{\log_b(a) \varepsilon})$  for some  $\varepsilon > 0$ , then  $T(n) = O(n^{\log_b(a)})$ .
- 2. If  $f(n) = \Theta(n^{\log_b(a)})$ , then  $T(n) = \Theta(n^{\log_b(a)} \cdot \log(n))$ .
- 3. If  $f(n) = \Omega(n^{\log_b(a) + \varepsilon})$  for some  $\varepsilon > 0$ , and if  $af(n/b) \le cf(n)$  for some c in the range 0 < c < 1 and for all sufficiently large n, then  $T(n) = \Theta(f(n))$ .

## **Practice Problems**

For each of the following recurrences, if the Master Theorem can be applied, give a tight asymptotic bound on the solution T(n). Otherwise, indicate that (and explain why) the Master Theorem does not apply.

- 1.  $T(n) = 3T(n/2) + n^2$
- 2.  $T(n) = 4T(n/2) + n^2$
- 3.  $T(n) = T(n/2) + 2^n$
- 4.  $T(n) = 2^n T(n/2) + n^n$
- 5. T(n) = 16T(n/4) + n
- 6.  $T(n) = 2T(n/2) + n \log(n)$
- 7.  $T(n) = 2T(n/2) + n/\log(n)$
- 8.  $T(n) = 2T(n/4) + n^{0.51}$
- 9. T(n) = (0.5)T(n/2) + 1/n
- 10. T(n) = 16T(n/4) + n!
- 11.  $T(n) = \sqrt{2} T(n/2) + \log(n)$
- 12. T(n) = 3T(n/2) + n
- 13.  $T(n) = 3T(n/3) + \sqrt{n}$
- 14. T(n) = 4T(n/2) + cn
- 15.  $T(n) = 3T(n/4) + n \log(n)$
- 16. T(n) = 3T(n/3) + n/2
- 17.  $T(n) = 6T(n/3) + n^2 \log(n)$
- 18.  $T(n) = 4T(n/2) + n/\log(n)$
- 19.  $T(n) = 64T(n/8) n^2 \log(n)$
- $20. T(n) = 7T(n/3) + n^2$
- 21.  $T(n) = 4T(n/2) + \log(n)$
- 22.  $T(n) = T(n/2) + n(2 \cos(n))$

### **Answers:**