CSQ 102 6-4-24

- Final Exam! Th 6/13 9:30 - 11:00 am

· SETS: Sun. 6/9 11:59 PM

I will add 0.5% to overall scares.

Theorem

at least (2) edge Prober most be Pento-med (in Worst case) to determine connectedness of a g-aph on n Vertices.

Proof

Consider an algorithm for this Problem and run it against the tollowing adversary.

Strategy: answer no to all edse Prober, unless that answer would Prove the g-aph is disconnected. more Pracisely:

maintain Two edge sok A, B SE(Kn) and B= ϕ .

Where instially A = E(Kn) and $B = \phi$.

Adversary Performs following algorithm

each time an edge is Probed

P-obe(e)

1. it A-e is connected

2. A=A-e

3. answer No

4. else

5. B=B+e

6. answer Yes

note: always

· BSA

· A-B = frot yet Probad edger?

all the adversary's answers.

The following invariants are maintained over any Seq. of edge Probes

1a) A in always connected

(b) It A containe a cycle, then none of ite edges belong to B.

P-00+

deleting an edge from a cycle in A would leave A connected. Such an edge would never be added to (e) Bis acyclic. This follows from (b)

(d) It A + B, then Bin disconnected Pool.

Assume that Bix connected.

Then, being acyclic, Bratree.

Since A # B and B EA there exists e EA-R. It e were added to B it would form a cycle

with some other edger of B.

(see treeness thun.) Since BEA

that cycle would also be in A.

: A contains a cycle consisting of e together with other edges in B,

contradicting (6). .. B is discourt.

Suppos the algorithm halts after les than (2) Probes. Then at least one egge of Kn war not Probed, hence A-B + 0 ... A + B. (d) sare Bir disconnected, while (a) says A'n connected. Both graphs are consistent with the seq. of answers given by adversary.

$$= \times$$
. $N = 4$. $\binom{n}{2} = \frac{4.3}{2} = 6$

Probe edges

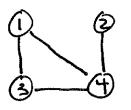
3 9

<u>(1)</u>

11,29 NO

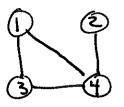
0 0

12,34 no



(1)

h2,4} yes



①

(3)

2

L1,43 no

(

(2)

 Θ

{3,4} yes