

Propositional Logic: An Introduction

**Type Theory and Mechanized Reasoning
Lecture 5**

Introduction

Administrivia

- Assignment 1 is out. Another assignment will be out on Thursday.
- The course has been approved for the semester.
- Details about the final project will be available next week.

Objectives

1. Introduce `propositional logic`, a simple form of logic for reasoning about Boolean connectives.
2. Use Agda to `implement` propositional logic.
3. Use propositional logic as a setting for learning `important concepts and terms` in logic.

Unicode cheatsheet

\rightarrow is `\rightarrow`

\vee is `\or`

\mathbb{N} is `\bN`

\neg is `\neg`

\times is `\times`

$_p$ is `\^p`

\wedge is `\and`

Practice Problem

*Write a function **down** which, given a natural number **n**, returns a vector with the values **n** through **1**.*

*Write a function **up**, which return a vector with the number from **1** to **n**.*

Agda Tutorial: Interaction

Interaction: *At a High Level*

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Agda is a functional language, which means a lot of **pattern matching**.

Agda is pure, there are no print statements, so we need to know how to **compute values**.

The Primary Tool: Holes

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Within a hole, we can determine the `types of everything` in the environment.

We can even `try to fill` in a value (which can have it's own holes).

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5. Rinse and repeat.

Interaction Cheat Sheet

C-c C-l	load file
C-c C-,	check type in hole
C-c C-c	pattern match within hole
C-c C-SPACE	try to fill in hole
C-c C-n	compute value of term

demo

Propositional Logic: Motivation

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This means we need to specify two things:

syntax	what things can I write down ?
semantics	what do those things mean ?

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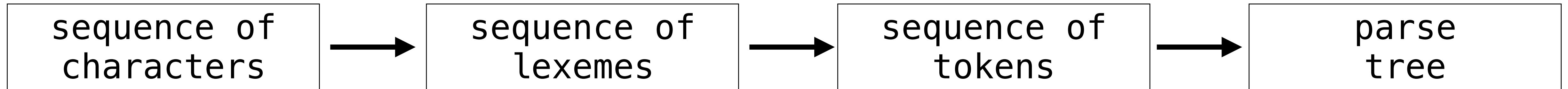
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This means we need to specify two things:

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(this is also what you need for a programming language)

What does syntax mean?



We're not going to focus on *syntax or parsing* in this course.

We will almost always presume we already have a parse tree.

What does syntax mean?

sequence of
characters



sequence of
lexemes



sequence of
tokens



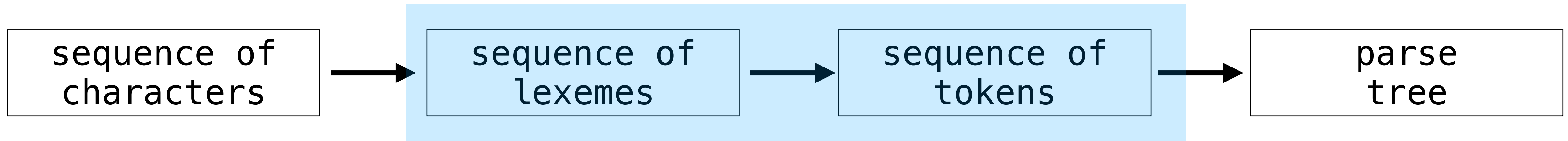
parse
tree

[(, X , □ , a , n , d , □ , Y ,)]

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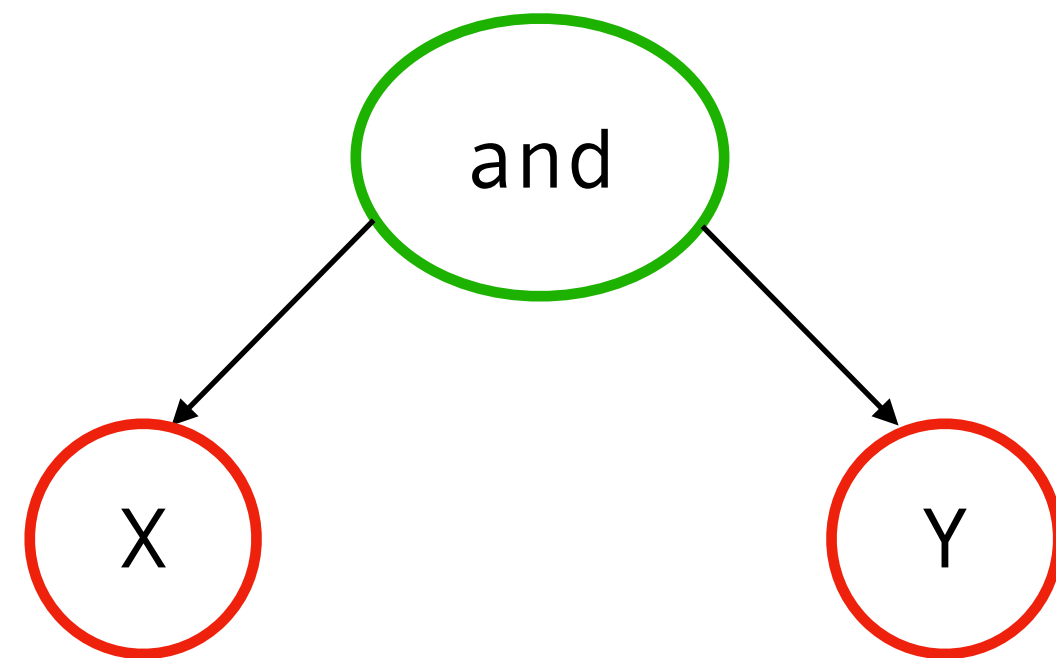
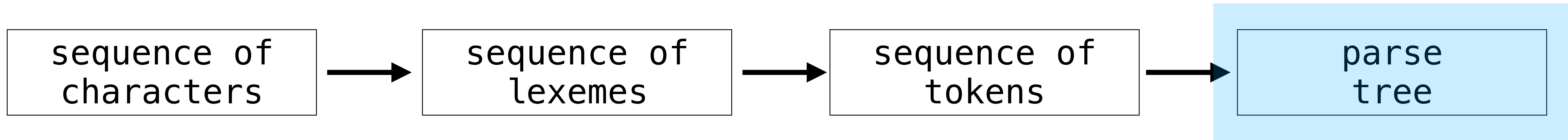


[*sym* '(', *var* 'X', *con* 'and', *var* 'Y', *sym* ')']

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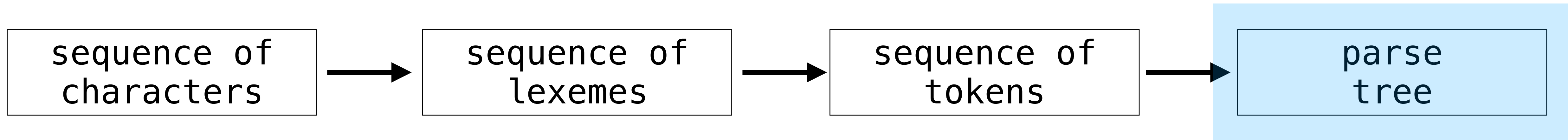
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Very interesting question from a philosophical perspective.

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Very interesting question from a philosophical perspective.

In classical logic (what we will consider today) meaning will mean ***truth***.

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Very interesting question from a philosophical perspective.

In classical logic (what we will consider today) meaning will mean ***truth***.

In intuitionistic logic, meaning will be something different.

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connective proposition connective proposition connective proposition

Why is this?

Boolean Connectives

$$x \mathbf{\wedge} y$$

conjunction (and)

Propositional logic is the study of logical (Boolean) **connectives**.

- » *What do connectives mean? How do they affect truth?*
- » *What connectives exist? How many do we need?*
- » *How do connectives interact?*

Propositional Logic and Programming Conditionals

```
if is_raining and not is_warm:  
    # some code
```

Propositional logic is the study of **Bools**.

- » *When can I replace one conditional with another?*
- » *Why is there only **and**, **or**, and **not**?*
- » *Why can conditionals short-circuit?*

Propositional Logic: Syntax

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We think of these variables in the same way as we think of **variables in algebra**.*

*We will assume a countable number of variable symbols like in algebra

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Definition. A formula is defined inductively as follows:

- A propositional variable x is a formula.
- If P is a formula then so is $\neg P$.*
- If P and Q are formulas then so are $P \wedge Q$ and $P \vee Q$ and $P \rightarrow Q$.

* P is a *meta-variable*. It refers to an arbitrary formula. It is not the same as a propositional variable.

Examples

$$P \wedge (Q \wedge \neg Z)$$

$$A \rightarrow (B \rightarrow (\neg C \vee B))$$

$$A \wedge (A \wedge (A \wedge A))$$

Note. Parentheses are *meta-syntactical*. Remember that these formulas are shorthand for **trees**.

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P is a variable so P is a formula.

P is a formula and $Q \wedge \neg Z$ is a formula so $P \wedge (Q \wedge \neg Z)$ is a formula.

Informal Meaning

$\neg \equiv$ "not"

$\wedge \equiv$ "and"

$\vee \equiv$ "or"

$\rightarrow \equiv$ "implies"

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R It's raining

C It's cold

$(R \wedge C) \rightarrow R$ If it's raining and cold then it's
raining

Understanding Check:

English to Formula

Syntax: In Agda

```
data Formula : Set where
  _p      : String → Formula
  ¬p _    : Formula → Formula
  _p∧_    : Formula → Formula → Formula
  _p∨_    : Formula → Formula → Formula
  _p→_    : Formula → Formula → Formula
```

The **tree structure** of formulas is implicit in it being an ADT.

Syntax: What's Next?

Remember. We haven't actually given formulas *meaning*. That is, we haven't given a **semantics**.

But we can ask about:

- » Function on formulas
- » Transformations of formulas

Example: Depth (In Mathematical English)

Definition. The **depth** of a formula is defined as the depth of its corresponding tree:

$$d(x) = 0$$

$$d(\neg P) = 1 + d(P)$$

$$d(P \square Q) = \max(d(P), d(Q)) + 1$$

where \square is any one of ' \wedge ' or ' \vee ' or ' \rightarrow '.

Example: Depth (In Agda)

Let's do a demo.

Propositional Logic: Semantics

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Given a *formula* $\neg(x \wedge y) \vee (x \rightarrow \neg z)$ if we know the values of x , y , and z , we can compute the value of the the formula.

We think of a propositional variable as having the values of either **true** or **false**.

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A valuation v is like a **state of affairs**.

***The idea.** If we know the state of affairs, we can determine the truth or falsity of any statement.*

Partial Valuations

$$v(x) = \text{true}$$

$$v(y) = \text{false}$$

$$v(z) = \text{false}$$

$$v(_) = \text{false}$$

$$(x \wedge y) \vee z$$

Note. We will typically only care about a **small collection** of variables.

We can assume all **unspecified variables** are assigned to be false.

Valuation (In Agda)

Let's do a demo.

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We have to say what we want the truth of a statement to be based on its constituent parts.

Evaluation: In Agda

Let's do a demo.