

Solution: Dynamic Programming

1. Solve Dynamic Programming Problem and Compare with Naïve approach

You are playing a puzzle. A random number N will be given, you have blocks of length 1 unit and 2 units. You need to arrange the blocks back to back such that you get a total length of N units.

In how many distinct ways can you arrange the blocks for given N.

- Write a description of approach to solve it using Dynamic Programming paradigm
- Implement a function **blockpuzzle_dp(N)** that solves this problem using Dynamic Programming using either top-down or bottom-up approach
- Write pseudocode for the brute force approach
- Compare the time complexity of both the approaches
- Write the recurrence formula for the problem

Name your file **BlockPuzzle.py**

Example 1:

Input: N=2, Result: 2

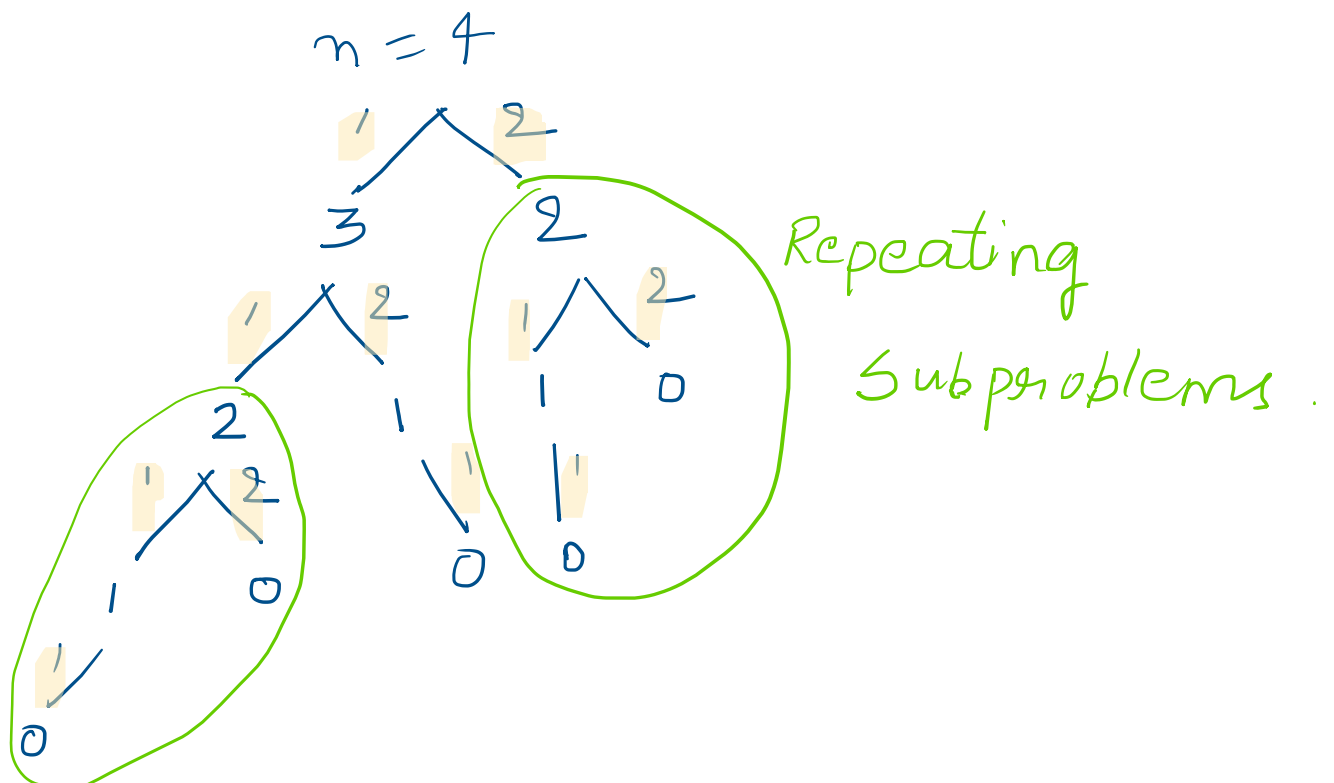
Explanation: There are two ways. $(1+1, 2)$

Example 2:

Input: N=3, Result: 3

Explanation: There are three ways (1+1+1, 1+2, 2+1)

a. We have repeating subproblems as shown in the figure below. We can cache these use the dynamic programming technique to find the result.



b. Following is code for bottom-up approach

```
def blockpuzzle_dp(n):  
    if n==1: return 1  
  
    memo = [0]*(n+1)  
    memo[1] = 1  
    memo[2] = 2  
    for i in range(3, n+1):  
        memo[i] = memo[i-1] + memo[i-2]  
  
    return memo[n]
```

c. brute force approach

```
def blockpuzzle_bruteforce(n):  
    if (n <= 0):  
        return 0  
    if (n == 1):  
        return 1  
    if (n == 2):  
        return 2  
    return blockpuzzle_bruteforce(n-1) + blockpuzzle_bruteforce(n-2)
```

d. Brute force approach has time complexity of $\theta(2^n)$ (If we draw recursion tree: it will have n levels and at level k 2^k amount of work will be done)

The dynamic programming approach has time complexity of $\theta(n)$.

The dynamic programming approach solves this problem extremely faster compared to the brute force approach.

e. Recurrence Formula:

$F(n) = 0$ for $n = 0$

$F(n) = 1$ for $n = 1$

$F(n) = 2$ for $n = 2$

$F(n) = F(n-1) + F(n-2)$ for all values of n

2. Solve a problem using top-down and bottom-up approaches of Dynamic Programming technique

Two players A & B are playing a game. The rules of the game are:

At the start one number N will be given. The player who starts would have to pick a number i such that $0 < i < N$, the condition is that $N \% i == 0$. The second player would pick a number j from $N-i$, satisfying the condition $0 < j < (N-i)$. And the game goes on until there is no

more possibility of making any selection. Each player would play in turns, and A always starts the game. Assume both players play optimally.

Given a number N return if A would win the game or not.

Example 1:

Input: N=2, Result: True

Explanation: A choses 1, and B has no more numbers to chose

Example 2:

Input: N=3, Result: False

Explanation: A choses 1, B choses 1, and A has no more numbers to chose

- a. Implement a solution to this problem using Top-down Approach of Dynamic Programming, name your function **game_topdown(N)**

```
def game_topdown(n):
    memo = [None]*(n+1)
    return game_topdown_helper(n, memo)

def game_topdown_helper(n, memo):
    print('solving for: ', n)
    if n<=1:
        return False

    if memo[n] is not None:
        return memo[n]

    for x in range(1, n):
        if (n%x==0):
            if not game_topdown_helper(n-x, memo):
                memo[n] = True
                return True
    memo[n] = False

    return False
```

- b. Implement a solution to this problem using Bottom-up Approach of Dynamic Programming, name your function **game_bottomup(N)**

Memo[i] holds the value of the game result for A for n=i. 'A' will check for all factors and see if there is any factor that will make 'B' lose.

```
def game_bottomup(n):
    memo = [False]*(n+1)

    memo[0] = False
    memo[1] = False

    for i in range(2, n+1):
        for j in range(1,i):
```

```

        if(i%j == 0 and memo[i-j] == False):
            memo[i] = True
            break

print(memo)
return memo[n]

```

- c. Explain your approach to solve this problem. How is your top-down approach different from the bottom-up approach?

The bottom-up approach solves the problem from $i=2$ onwards until the value of n .
`game_topdown_helper()` recursively solves the problem for n until base case is reached following top down approach.

- d. What is the time complexity and Space complexity using Top-down Approach

Time: $O(n^2)$ [x goes from 1 to n and for each x helper function is called at most n times]

Space: $\theta(n)$

- e. What is the time complexity and Space complexity using Bottom-up Approach

Time: $O(n^2)$

Space: $\theta(n)$

- f. Write the subproblem and recurrence formula for your approach

Subproblem: $f(i)$ can be given as: will A win the game for a value of $N=i$.

Recurrence Formula:

$F(n) = \text{False}$ for $n = 0$ or 1

$F(n) = !F(n-x)$ for, $0 < x < N$ such that $n \% x = 0$ for $n > 1$