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1.
A - Recurrance Formula:
                                     for 0 <= n <= 2
        { c
t(n) = {n + max(t(n-2), t(n-3)}
                                     for value n < 3
        \{\max(t(n-2), t(n-3)\}
                                     for value n < 0}
B - Pseudocode:
def max_independent_set(array):
    memo = []
    three_back = []
    two_back = []
    one_back = []
                                #set variables to hold returns and tracking
    if len(array) < 3
        if length = 0
            return empty memo
        if length = 1
            append array[0] to memo and return memo
            or if array[0] < 0 just return empty memo
        if length = 2
            append max (array[0],array[1]) to memo
            or if both negative, return empty memo
                                # all arrays length 3+
    else:
        if array[2] is negative
            append three back to array[0]
            append two back to array[1]
            append one back to array[0]
                                            #since 2 is negative and skipped
            add correct current sums to memo
        if array[0] is negative
            append three_back to array[0]
            append two back to array[1]
            append one back to array[2]
            add correct current sums to memo
        if array[1] is negative, and array[0] and array[2] are not:
            append three_back to array[0]
            append two back to array[1]
            append one back to array[2]
            add correct current sums to memo
    for i in range 3-array length
        if array[i] is not negative
            append memo with max(memo[i-3],memo[i-2])
            if memo[i-3] greater
                store three back in array temp
                move all back variables one step back
                one back becomes temp
            else memo[i-2] greater
                store two_back in array temp
                move all back variables one step back
                one_back becomes temp
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else if array[i] is positive
            append memo with (array[i] + max(memo[i-3],memo[i-2])
            if memo[i-3] greater
                store three back in array temp
                move all back variables one step back
                one_back becomes temp
            else memo[i-2] greater
                store two back in array temp
                move all back variables one step back
                one back becomes temp
    #memo stores the max sum at each location in the array
    #we only need to check the last 2 for what is max, and then return
    #one back or two back
    if memo[len(memo-1)] > memo[len(memo-2)]
        return one back
    else:
        return two back
C - Time Complexity: O(n)
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The code uses one for loop that goes the length of the array passed to it. everything else is if statements running in constant time.

2.

B - Time Complexity O(2^n) Every additional element of the set makes it run expentionally longer.

Debrief:

- 1. About 15 hours this week on the course. The assignment was 8 hours and about 7 on the explorations and third party resources to try and understand the material. I ended up not using the knapsack technique for the first problem because I was not very comfortable with the way it was explained.
- 2. Moderate+
- 3. 50% maybe. I understand this specific instance of it, but I dont feel confident in the subject overall.
- 4. I really with there was some more examples for the problems we cover in the explorations. I feel that one simple look at the problem from a theoretical approach and then one specific example thats more in depth would be helpful.